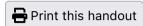
CSC110 Lecture 24: Analyzing Built-In Data Type Operations



You will find the following formula helpful:

def f1(nums: list[int]) -> None:

$$orall n \in \mathbb{N}, \,\, \sum_{i=0}^n i = rac{n(n+1)}{2}$$

Exercise 1: Running time of list operations

Each of the following Python functions takes a list as input. Analyse each one's running time in terms of n the size of its input.

```
Let n be the length of the lost input to f1.

The call to list insert is in senting at the beginning of the list. This cause each of the n elements to shift by one place in memory. ... RT<sub>G1</sub>(n) = n 6-O(n) (or use: list insert(1st, i,...) is O(n-i) and here "i" is 0.)
```

```
def f2(nums: list[int]) -> None:
    for i in range(0, 100):
        list.append(nums, 10000)
```

Let n be the length of the list input to f2.

The for loop iterates 100 times.

Each i teration appends to the list, a $\Theta(i)$ operation that we will count as 1 ste.

3.

Note: the length of nums changes at each iteration, and so the running time of list.insert does a well!

rough analysis:
$$qq$$

RT $(n) = \sum_{i=0}^{n+i} (n+i)$

rept

value install

 $qq = a$ constant

$$= 100n + \frac{\sum_{i=0}^{2} i}{i=0}$$

$$= 100n + \frac{99 \cdot 100}{2}$$

$$\in \Theta(s)$$

rough and lyst?:

RT_{f4}(n) = 1 +
$$\frac{1}{2}$$
 (n+i) it is when it is insultable in a second in a second

Exercise 2: Running-time analysis with multiple parameters

Each of the following functions takes more than one list as input. Analyse their running time in terms of the size of their inputs; do not make any assumptions about the relationships between their sizes.

(Let n_1 be the size of nums1 and n_2 be the size of nums2.)

Rough analysis:
$$RT_{c}(n_{1},n_{3}) = \sum_{i=0}^{n_{2}-1} 1$$

$$= n_{2} \cdot 1$$

$$= n_{2} \cdot 1$$

$$= n_{2} \cdot 1$$

(Let n_1 be the size of nums 1 and n_2 be the size of nums 2.)

Rough analysis:
$$n_{2}$$
-1

RT_f(n_{1},n_{2}) = $\sum_{i=0}^{2} (n_{i} + i)$

$$= n_{1} n_{2} + \sum_{i=0}^{\infty} i$$

$$= n_{1} n_{2} + (n_{2}-1)(n_{2})$$

$$= n_{2} (2n_{1} + n_{2} - 1)$$

$$\in \Theta(n_{1}n_{2} + n_{2}^{2})$$
Can simplify
$$\Theta(n_{2}(n_{1}+n_{2}))$$

$$= \Theta(n_{2} \max(n_{1},n_{2}))$$

Exercise 3: Sets, dictionaries, and data classes

Analyse the running time of each of the following functions.

def f7(nums: set[int]) -> bool:

7.

Let n be the length of the set in put to f7.

Since searching in a set is $\Theta(i)$. $RT_{C7}(n) \in \Theta(i)$.

```
def f8(num_map: dict[int, int]) -> None:
    for num in num_map:
        num_map[num] = num_map[num] + 1
```

Let n be the length of the did input to for.

There are n i teratrons of the forloop body.

Since did indexing and assignment are OCD operations

$$RT_{f8}(u) = n \cdot 1$$

$$= n$$

$$\in \Phi(u)$$

9.

```
def f9(grades: dict[str, list[int]], new_grades: dict[str, int]):
    for course in new_grades:
        if course in grades:
            list.append(grades[course], new_grades[course])
        else:
            grades[course] = [new_grades[course]]
```

Let n be the length of the input new-grades to fq.

The for loop strades in times.

Since key search in a dict, list appelled and assignment to a dict are all $\Theta(i)$ operations,

$$RT_{fq}(n) = n$$

 $e \Leftrightarrow (n)$

10.

```
from dataclasses import dataclass
import math

@dataclass
class Person:
    """Docstring omitted"""
    name: str
    age: int

def f10(people: list[Person]) -> int:
    """Precondition: people != []"""
    max_age_so_far = -math.inf

for person in people:
    if person.age > max_age_so_far:
        max_age_so_far = person.age

return max_age_so_far
```

Let n be the length of the input list to \$10.

Since destaclase attribute lookup and assignment one
$$\Theta(i)$$
 operations,

RT_{fio}(n) = 1 + n·1 + 1

= n+2

 $\in \Theta(n)$.

Additional exercises

Analyse the running time of each of the following functions.

1.

```
def extra1(nums: list[int]) -> None:
    for i in range(0, len(nums)):
        nums[i] = 0
```

2.

```
def extra2(nums: list[int]) -> None:
    for i in range(0, len(nums)):
        list.pop(nums)
```

```
def extra3(nums: list[int]) -> None:
    for i in range(0, len(nums)):
        list.pop(nums, 0)
```

Note: the length of nums changes at each iteration, and so the running time of list.pop does as well!

4.

```
def extra4(nums1: list[int], nums2: list[int]) -> None:
    for i in range(0, len(nums1)):
        for j in range(0, len(nums2)):
            nums1[i] = nums1[i] + nums2[j]
```

(Let n_1 be the size of nums1 and n_2 be the size of nums2.)

5.

(Let n be the length of nested_nums, and assume every inner list has length m.)

6.

```
def extra6(nums: set[int]) -> list[int]:
    new_nums = []

for num in nums:
    list.insert(new_nums, 0, num ** 2)

return new_nums
```



```
def extra7(nums: list[int]) -> dict[int, int]:
    counts_so_far = {}

for num in nums:
    if num in counts_so_far:
        counts_so_far[num] = counts_so_far[num] + 1
    else:
        counts_so_far[num] = 1
return counts_so_far
```