

CSC110 Lecture 9: Programming and Proofs

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Announcements and Today's Plan

Assignment 1 due tomorrow (noon Eastern Time)

- Assignment 1 has been posted!
 - Check the [FAQ \(+ corrections\)](#) page
 - Additional TA office hours ([schedule on Quercus](#))
 - **(NEW)** [Academic Integrity in CSC110 advice page](#)
 - **(NEW)** Review the [Assignment Policies](#) page, including [grace credit policy](#)
- For your submission:
 - Submit [all](#) required files (including the `a1.tex` file!)
 - Make sure your Python files all run (right-click -> "Run File in Python Console")
 - Run `doctest` and `python_ta` on your Python files (using provided code)
 - You can submit on MarkUs multiple times!

Other announcements

- Join a [Recognized Study Group](#)
- Term Test 1 info has been posted!
 - Info on [Quercus](#):
 - Test [time](#) and [location](#) (not MY 150!)
 - Test [coverage](#)
 - Advice for advice
 - Review provided [reference sheet](#)

Story so far

Data: data types, literals, basic operators, comprehensions

Functions: using built-in functions, methods; defining our own (top-level) functions

Logic: translating boolean expressions, filtering comprehensions, if statements

Our code is getting more complex, and will only continue to do so in the coming weeks.

How can we be confident that our code is correct?

How can we be confident that our code is correct?

1. **Testing:** in the Python console; `doctest`; `pytest`; property-based tests (`hypothesis`)
2. **Code reuse:** reuse functions that we've already defined (and are confident is correct), rather than copying the same code in multiple places.
3. **Simplification:** translate more complex code structures into simpler ones (e.g., simplifying if statements)
4. **(TODAY) Proof:** Write a formal proof that a function's implementation is correct for all valid inputs.

Today you'll learn to...

1. Understand unfamiliar **mathematical definitions**, and create new predicates based on these definitions.
 - Problem domain: **basic number theory** (divisibility and prime numbers)
2. Use definitions to simplify and expand statements in predicate logic.
3. Write **proofs** of simple statements using these definitions.
4. Connect mathematical statements to improve/develop **algorithms**.

Working with Definitions

Definitions are our way of taking a problem domain and breaking it down precisely into understandable parts.

A definition

Definition. Let $n, d \in \mathbb{Z}$. We say that d **divides** n when there exists a $k \in \mathbb{Z}$ such that $n = dk$.

Note: using this definition, 0 divides 0.

(In fact, every number divides 0!)

A definition (broken down)

Let $n, d \in \mathbb{Z}$.

The domain of the definition

We say that d **divides** n

The term being defined

when

Connection word

there exists a $k \in \mathbb{Z}$ such that
 $n = dk$

The body of the definition

Translating a definition into a predicate

Every definition can be turned into a predicate that asks, “Does this definition apply to these values?”

Let $n, d \in \mathbb{Z}$. We say that d **divides** n when there exists a $k \in \mathbb{Z}$ such that $n = dk$.

$Divides(d, n) : \exists k \in \mathbb{Z}, n = dk$ where $d, n \in \mathbb{Z}$

$Divides(d, n)$ is often written as $d \mid n$.

Note: $d \mid n$ vs. $n \div d$

\mid is a **predicate**: $d \mid n$ evaluates to a **boolean**.

\div is an **arithmetic operator**: $n \div d$ evaluates to a **number**.

Expanding definitions

“For every integer x , if x divides 10, then it also divides 100.”

$$\forall x \in \mathbb{Z}, x \mid 10 \Rightarrow x \mid 100$$

$$\forall x \in \mathbb{Z}, \underline{(\exists k_1 \in \mathbb{Z}, 10 = k_1 x)} \Rightarrow \underline{(\exists k_2 \in \mathbb{Z}, 100 = k_2 x)}$$

Proofs

A **proof** is a structured argument for convincing someone else why a statement is True.

A **program** is a structured text for telling a computer what to compute.

Just as we learn a **programming language** to write programs, we learn a **proof language** to write proofs.

Two things can make proofs challenging:

1. Proofs use both English and mathematical notation, and the language is less rigid than a programming language.
 - It's up to us to write our proofs in a structured way.
2. There's no "proof interpreter" we can run to check our proofs.
 - It's up to us to check our proofs carefully.

Proof example 1

Prove the following statement: every integer is divisible by 1.

Translation:

$$\forall n \in \mathbb{Z}, 1 \mid n$$

$$\forall n \in \mathbb{Z}, \underline{\exists k \in \mathbb{Z}, n = 1 \cdot k}$$

Proof example 1, continued

$$\forall n \in \mathbb{Z}, \exists k \in \mathbb{Z}, n = 1 \cdot k$$

Proof (structure only).

Let $n \in \mathbb{Z}$.

Let $k = \underline{\hspace{2cm}}$.

We want to show that $n = 1 \cdot k$

.

...

Proof.

Let $n \in \mathbb{Z}$.

Let $k = n$.

We want to show that $n = 1 \cdot k$

.

That's just a calculation:

$1 \cdot k = 1 \cdot n = n$, since $k = n$.

Proof example 2

Prove the following statement: for all **positive** integers n and d , if $d \mid n$ then $d \leq n$.

Translation:

$$\forall n, d \in \mathbb{Z}^+, d \mid n \Rightarrow d \leq n$$

$$\forall n, d \in \mathbb{Z}^+, \underline{(\exists k \in \mathbb{Z}, n = d \cdot k)} \Rightarrow d \leq n$$

Proof example 2, continued

$$\forall n, d \in \mathbb{Z}^+, \quad \underline{(\exists k \in \mathbb{Z}, n = d \cdot k)} \Rightarrow d \leq n$$

Proof.

Let $n, d \in \mathbb{Z}^+$.

Assume there exists $k \in \mathbb{Z}$ such that $n = d \cdot k$.

We want to show that $d \leq n$.

...

Proof example 2, continued

$$\forall n, d \in \mathbb{Z}^+, \quad \underline{(\exists k \in \mathbb{Z}, n = d \cdot k) \Rightarrow d \leq n}$$

Proof. Let $n, d \in \mathbb{Z}^+$. Assume there exists $k \in \mathbb{Z}$ such that $n = d \cdot k$. We want to show that $d \leq n$.

Since $n > 0$ and $d > 0$ and $k = \frac{n}{d}$, we know $k > 0$.

Since $k \in \mathbb{Z}$ and $k > 0$, we know $k \geq 1$.

Then $d \leq d \cdot k$, since $k \geq 1$.

By our assumption, $n = d \cdot k$, and so $d \leq n$.

Proof example 2 (alternate visualization 1)

$$\forall n, d \in \mathbb{Z}^+, \quad \underline{(\exists k \in \mathbb{Z}, n = d \cdot k)} \Rightarrow d \leq n$$

Proof. Let $n, d \in \mathbb{Z}^+$. Assume there exists $k \in \mathbb{Z}$ such that $n = d \cdot k$. We want to show that $d \leq n$. Since $n > 0$ and $d > 0$ and $k = \frac{n}{d}$, we know $k > 0$. Since $k \in \mathbb{Z}$ and $k > 0$, we know $k \geq 1$. Then $d \leq d \cdot k$, since $k \geq 1$. By our assumption, $n = d \cdot k$, and so $d \leq n$.

Proof example 2 (alternate visualization 2)

$$\forall n, d \in \mathbb{Z}^+, \quad \underline{(\exists k \in \mathbb{Z}, n = d \cdot k)} \Rightarrow d \leq n$$

Proof. Let $n, d \in \mathbb{Z}^+$. Assume there exists $k \in \mathbb{Z}$ such that $n = d \cdot k$.
We want to show that $d \leq n$.

Then:

Conclusion	Justification
$k > 0$	$n > 0$ and $d > 0$ and $k = \frac{n}{d} >$
$k \geq 1$	$k \in \mathbb{Z}$ and $k > 0$
$d \leq d \cdot k$	$k \geq 1$
$d \leq n$	$d \leq d \cdot k$ and $n = d \cdot k$ (assumption)

Exercise 1: Practice with proofs

Divisibility and the Quotient- Remainder Theorem

```
def divides(d: int, n: int) -> bool:
    """Return whether d divides n.
    """
    if d == 0:
        return n == 0
    else:
        return n % d == 0
```

Why does this work?

- If branch ($d == 0$):
 - in this case, $0 \mid n \Leftrightarrow n = 0$.
- Else branch ($d \neq 0$):
 - in this case, $d \mid n \Leftrightarrow n \% d = 0$

Quotient-Remainder Theorem

Theorem (Quotient-Remainder Theorem). For all $n, d \in \mathbb{Z}$, if $d \neq 0$ there exist unique integers $q \in \mathbb{Z}$ and $r \in \mathbb{N}$ such that $n = qd + r$ and $0 \leq r < |d|$.

We say that q is the **quotient** when n is divided by d , and that r is the **remainder** when n is divided by d .

See Section 4.5 for further details (and more example proofs).

Prime numbers

Definition. Let $p \in \mathbb{Z}$. We say p is **prime** when it is greater than 1 and the only natural numbers that divide it are 1 and itself.

“greater than 1”:

- $p > 1$

“the only natural numbers that divide it are 1 and itself”:

- $\forall d \in \mathbb{N}, d \mid p \Rightarrow d = 1 \vee d = p$

(**Note:** the converse, $d = 1 \vee d = p \Rightarrow d \mid p$, holds for all integers.)

From definition to predicate, prime edition

IsPrime(p) :

$$p > 1 \wedge (\forall d \in \mathbb{N}, d \mid p \Rightarrow d = 1 \vee d = p)$$

where $p \in \mathbb{Z}$

IsPrime(p) :

$$p > 1 \wedge (\forall d \in \mathbb{N}, \underline{(\exists k \in \mathbb{Z}, p = kd)} \Rightarrow d = 1 \vee d = p)$$

where $p \in \mathbb{Z}$

Translating into Python

IsPrime(p) :

$$p > 1 \wedge (\forall d \in \mathbb{N}, d \mid p \Rightarrow d = 1 \vee d = p)$$

where $p \in \mathbb{Z}$

```
def is_prime(p: int) -> bool:
    """Return whether p is prime."""
    possible_divisors = range(1, p + 1) # {1, 2, ..., p}
    return (
        p > 1 and
        all({d == 1 or d == p
             for d in possible_divisors if divides(d, p)})
    )
```

```
def is_prime(p: int) -> bool:
    """Return whether p is prime."""
    possible_divisors = range(1, p + 1) # {1, 2, ..., p}
    return (
        p > 1 and
        all({d == 1 or d == p
             for d in possible_divisors if divides(d, p)})
    )
```

This algorithm checks p possible divisors. Can we do better?


```
from math import floor, sqrt

def is_prime(p: int) -> bool:
    """Return whether p is prime."""
    possible_divisors = range(2, floor(sqrt(p)) + 1)
    return (
        p > 1 and
        all({not divides(d, p) for d in possible_divisors})
    )
```

How do we **know** this version of `is_prime` is still correct?

Prove the following statement:

$$\forall p \in \mathbb{Z}, \text{Prime}(p) \Leftrightarrow (p > 1 \wedge (\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p))$$

This is a bigger statement! Let's break down the structure.

$$\forall p \in \mathbb{Z}, \text{Prime}(p) \Leftrightarrow (p > 1 \wedge (\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p))$$

Proof. Let $p \in \mathbb{Z}$.

Part 1: Prove that IF $\text{Prime}(p)$, THEN $p > 1$ and

$\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p$.

...

Part 2: Prove that IF $p > 1$ and $\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p$, THEN $\text{Prime}(p)$.

...

Proof (Part 1, $\text{Prime}(p) \Rightarrow \dots$)

Part 1: Prove that IF $\text{Prime}(p)$, THEN $p > 1$ and

$\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p$.

Assume p is prime, i.e. (expanding the definition) that

- $p > 1$, and
- $\forall d_1 \in \mathbb{N}, d_1 \mid p \Rightarrow d_1 = 1 \vee d_1 = p$

We want to show that:

- $p > 1$, and
- $\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p$

Proof (Part 1, $\text{Prime}(p) \Rightarrow \dots$)

We want to show that:

- $p > 1$, and
- $\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p$

First, we know $p > 1$, as this is one of our assumptions.

For the second part: let $d \in \mathbb{N}$, and assume $2 \leq d \leq \sqrt{p}$. We want to prove that $d \nmid p$.

Since $d \geq 2$, we know $d \neq 1$.

Since $p > 1$, we know $\sqrt{p} < p$, and so $d \leq \sqrt{p} < p$.

Therefore $d \neq 1$ and $d \neq p$, and so $d \nmid p$ (by the definition of prime).

Proof (Part 2, $\dots \Rightarrow \textit{Prime}(p)$)

Part 2: Prove that IF $p > 1$ and $\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p$, THEN $\textit{Prime}(p)$.

Assume that:

- $p > 1$, and
- $\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p$

We want to show that p is prime, i.e., that

- $p > 1$, and
- $\forall d_1 \in \mathbb{N}, d_1 \mid p \Rightarrow d_1 = 1 \vee d_1 = p$

Proof (Part 2, $\dots \Rightarrow \textit{Prime}(p)$)

We want to show that p is prime, i.e., that

- $p > 1$, and
- $\forall d_1 \in \mathbb{N}, d_1 \mid p \Rightarrow d_1 = 1 \vee d_1 = p$

First, we know $p > 1$, as this is one of our assumptions.

For the second part: let $d_1 \in \mathbb{N}$, and **assume** $d_1 \mid p$, i.e., that $\exists k \in \mathbb{Z}$ such that $p = d_1 k$. We want to show that $d_1 = 1 \vee d_1 = p$.

Since $d_1 \mid p$, by our earlier assumption we know $d_1 < 2$ or $d_1 > \sqrt{p}$.

This “or” gives us two **cases** in our proof.

Proof (Part 2, $\dots \Rightarrow \textit{Prime}(p)$)

We want to show that $d_1 = 1 \vee d_1 = p$.

Since $d_1 \mid p$, by our earlier assumption we know $d_1 < 2$ or $d_1 > \sqrt{p}$.

Case 1: assume $d_1 < 2$.

Then since $d_1 \in \mathbb{N}$, we know $d_1 = 1$ or $d_1 = 0$.

Then since $d_1 \mid p$, we know $d_1 \neq 0$, and so $d_1 = 1$.

Proof (Part 2, $\dots \Rightarrow \textit{Prime}(p)$)

We want to show that $d_1 = 1 \vee d_1 = p$.

Since $d_1 \mid p$, by our earlier assumption we know $d_1 < 2$ or $d_1 > \sqrt{p}$.

Case 2: *assume* $d_1 > \sqrt{p}$. Now we calculate:

$$d_1 k = p \quad (\text{by definition of divisibility})$$

$$k = \frac{p}{d_1}$$

$$k < \frac{p}{\sqrt{p}} \quad (\text{since } d_1 > \sqrt{p})$$

$$k < \sqrt{p}$$

We want to show that $d_1 = 1 \vee d_1 = p$.

Case 2: assume $d_1 > \sqrt{p}$.

...

So $k < \sqrt{p}$.

By our earlier assumption again, since $k \mid p$, we know $k < 2$ or $k > \sqrt{p}$.

.

Since $k < \sqrt{p}$, this means $k < 2$.

By the same reasoning as earlier, $k = 1$.

Since $p = d_1 k$, we conclude that $d_1 = p$.

Putting it all together

$$\forall p \in \mathbb{Z}, \text{Prime}(p) \Leftrightarrow (p > 1 \wedge (\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p))$$

$$\forall p \in \mathbb{Z}, \text{Prime}(p) \Leftrightarrow \underline{\text{is_prime}(p) \text{ returns True}}$$

```
def is_prime(p: int) -> bool:
    """Return whether p is prime."""
    possible_divisors = range(2, floor(sqrt(p)) + 1)
    return (
        p > 1 and
        all({not divides(d, p) for d in possible_divisors})
    )
```

```
def is_prime(p: int) -> bool:
    """Return whether p is prime."""
    possible_divisors = range(2, floor(sqrt(p))) # CHANGE
    return (
        p > 1 and
        all({not divides(d, p) for d in possible_divisors})
    )
```

$$\forall p \in \mathbb{Z}, \textit{Prime}(p) \Leftrightarrow (p > 1 \wedge (\forall d \in \mathbb{N}, 2 \leq d \leq \sqrt{p} \Rightarrow d \nmid p))$$

$$\forall p \in \mathbb{Z}, \textit{Prime}(p) \Leftrightarrow (p > 1 \wedge (\forall d \in \mathbb{N}, 2 \leq d < \sqrt{p} \Rightarrow d \nmid p))$$

Summary

Today you learned to...

1. Understand unfamiliar **mathematical definitions**, and create new predicates based on these definitions.
 - Problem domain: **basic number theory** (divisibility and prime numbers)
2. Use definitions to simplify and expand statements in predicate logic.
3. Write **proofs** of simple statements using these definitions.
4. Connect mathematical statements to improve/develop **algorithms**.

Homework

- Readings:
 - Today: 4.6, 4.7
 - Next class: 5.1 (to be posted soon)
- Finish up [Assignment 1](#)
- Review for [Term Test 1](#)

