

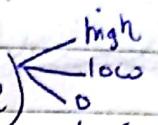
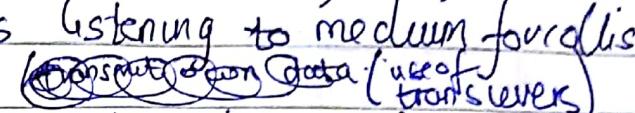
18/9/17

→ observed power greater than transmitted power of own signal, means collision occurs.

## Multiple Access (continued ...)

### • CSM A / CD:

#### - Carrier Sense Multiple Access with Collision Detection

- Reviewing signal strength (RSS - value) 
- while transmitting, sender is listening to medium for collision.
- Transmit + Receiving 
- If collision occurs, defer transmission (ie stop transmitting, wait for a random time)
  - (a) Abort transmission
  - (b) Transmit jam signal to inform others to stop as collision occurring
  - (c) back OFF waiting time

when transmitting without error/collision

transmit own data and receive own data

when collision occurs

RSS value rises

8 - transmitting own data

but receiving extra data.

### Actual Procedure & (same as flow diagram in ppt slide 20)

1. Use one of the persistence process (non-persistent, 1-persistent or p-persistent). Basically, sense channel before transmission by using one of the persistence processes.
2. Transmission & collision detection are simultaneous processes if collision detected then
  - (a) Abort transmission
  - (b) transmit a jam signal (18bit) to notify other stations and make them aware of collision. also notifies to discard the transmitted frame.
  - (c) After sending jam signal back off (wait) for a random time
  - (d) transmit again

~~CSMA/CD~~

- Through put of CSMA/CD is better than Aloha & Slotted Aloha.

1. Persistence

2. Transmission & Reception

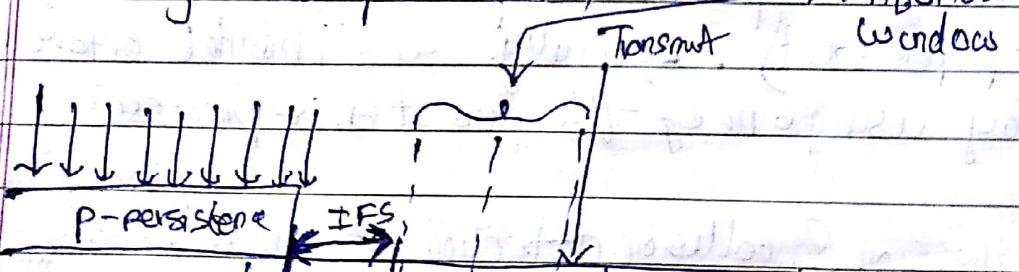
3. If collision detected. - CSMA/CD is used for bus topology LAN's (Ethernet)

~~Transmission & Reception~~ ~~Collision detection~~ ~~ACK~~ ~~now retransmit~~

- CSMA/CA : (Carrier Sense Multiple Access With Collision Avoidance)
  - specifically applicable for wireless networks
  - avoid collision before it happens

1. Apply any one of the three persistence method  
(sensing)

2. Interframe space (IFS)



If S is a time interval used to avoid collision

we wait for IFS time even if the channel is idle we don't transmit

3. ~~Contention~~ contention window:

$$K=2$$

{0, 1, 2, 3}

(Backoff method)

↑  
Random selection

now wait for 2 time slots (contention window)  
and then transmit blindly

Contention window is an amount of time divided into slots. A station that is ready to send chooses a random number of slots as its waiting time.

Date	/ /
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- Transmit & receive & wait for acknowledgement.

~~Collision~~ - positive ack and timeout timer help guarantee that the receiver has received the frame.

CSMA/CD is implemented by IEEE 802.3 ethernet LAN  
CSMA/CA is ~~impl~~ applicable on WIRELESS LAN  
IEEE 802.11

## Controlled Access Protocol

→ 3 methods

1. Reservation

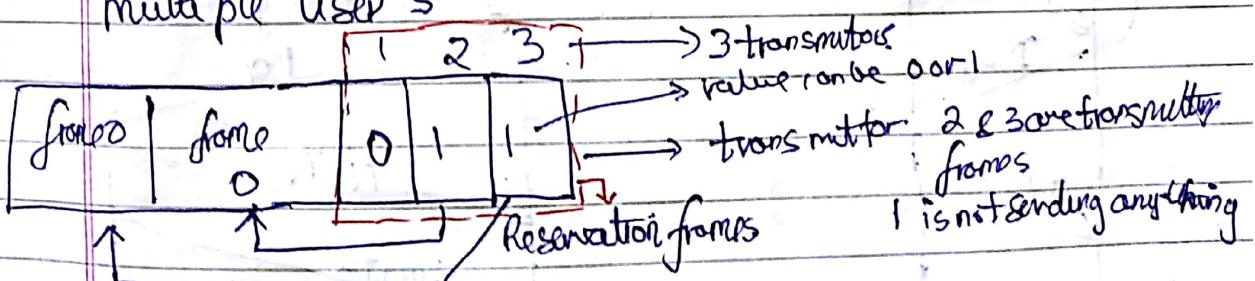
2. Polling

3. Token passing

→ In controlled Access, stations consult one another to find which station has the right to send. A station cannot send until it is authorized by others.

### 1. Reservation Access Method

- identifying the sequence of frames transmitted by multiple user's



↑  
Data Word

→ In reservation method, a station needs to make a reservation before sending data. Reservation frames precede data frames in each interval slot. → 5 stations, ∴ 5 mini slots inside the reservation frame example → in first interval station 1, 3, 4 have reservation. In second interval only station 1 has reservation.

→ It is network specific  
can't be used everywhere.

Q. Polling Protocol (has 2 function's)

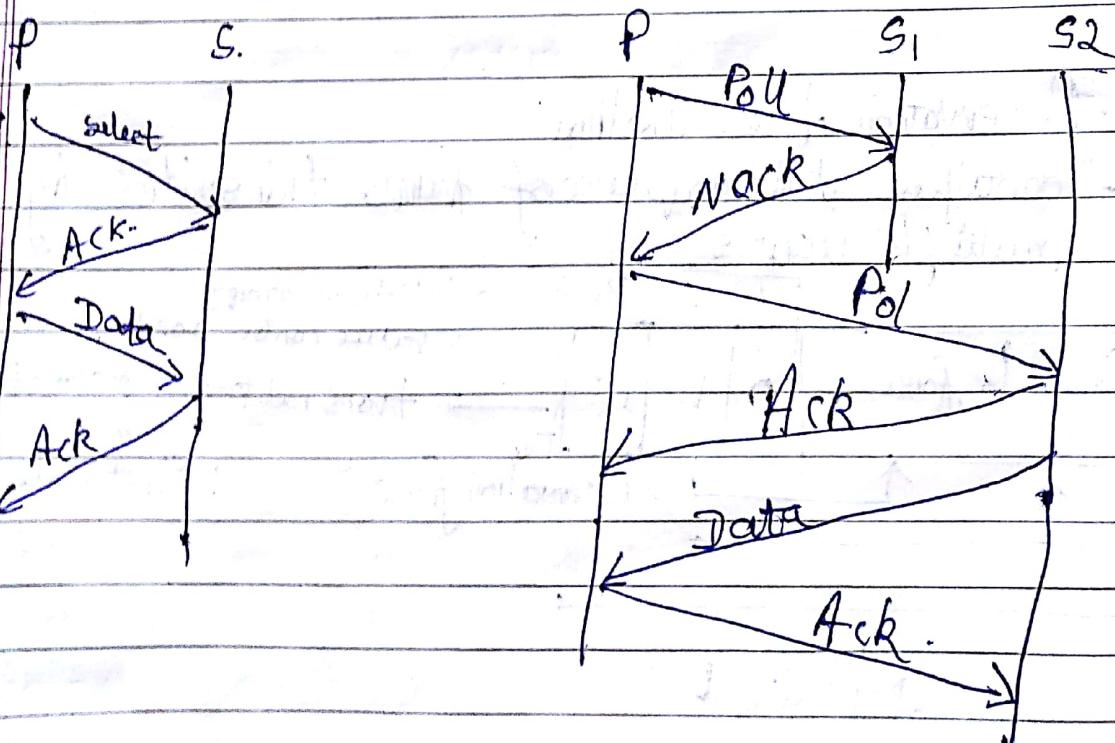
- for topologies where one device as primary device and others as secondary device
- All data exchanges must be made through primary device.

### Select

- function is used when primary device need to send the data.
- Select & Ack signal is used to reserve channel

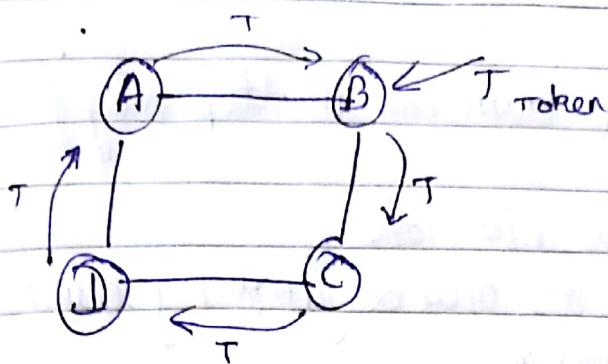
### Poll

- function is used when primary device wants to receive data.
- Ask each & every secondary device if they want to send any data
- Nack means no.
- +ve ack means yes and then we receive data.



→ Draw back  
If Primary device fails then system fails.

### 3. Token Passing Protocol:

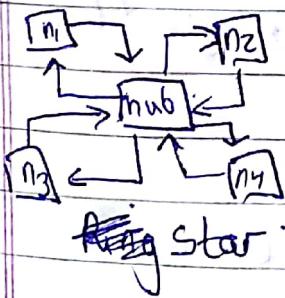


token is allocated to B.  
it sends some data  
for a very short  
while and  
then B  
passes the  
token to C.

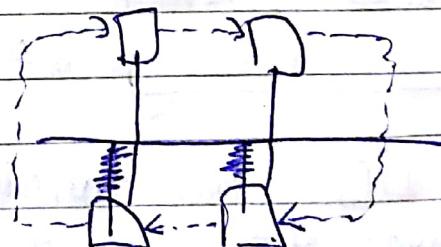
C transmits for  
a short time  
and then token is  
passed to D.

- Token is the controller to govern communication.
- The network is organised in a logical ring. For each station there is a predecessor and a successor in the ring. The current station has the right to access the channel as it holds a special packet called Token which circulates through the ring.
- When station has data to send, it waits for the token to come to it

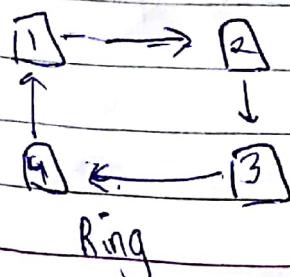
Q. Is star topology can implement token passing protocol?  
yes



~~Ring~~ Star



Bus



Ring

is a multiple access method in which available bandwidth of a link is shared in time, frequency or code among different stations

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## • Channelization

FDMA

TDMA

CMA

### 1. FDM: (frequency division Multiplexing)

- frequency is divided into slots
- we use guard bands b/w 2 frequency bands.  
to avoid interference
- A single subband can be used TDD or FDD for communication

### 2. TDM (Time division multiplexing)

- is a digital process that allows several connection's to share high bandwidth of a link. They share time. There is no division of BW
- each connection occupies a portion of time

### 1. FDD (frequency division duplexing)

- A single frequency band is divided into a forward channel and reverse channel

10 kHz

transmission      reception

4 kHz      4 kHz.

forward      reverse

2 kHz

some frequency band, say band one.

### 2. TDD (time division duplexing)

- divide the time slot into forward time slot & reverse time slot for transmission & reception respectively.

10 kHz

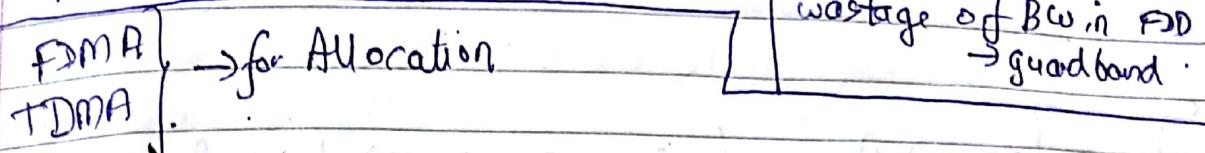
Transmission      Reception

$\uparrow T/2$        $\downarrow T/2$

Reception      Transmission

$\downarrow T/2$        $\uparrow T/2$

- \* - Bandwidth of TDD is more as full channel bandwidth is used. only time division is there.
- \* - Synchronous transmission is given by FDD as no waiting. In TDD you have to wait as alternate transmit & receive takes place.



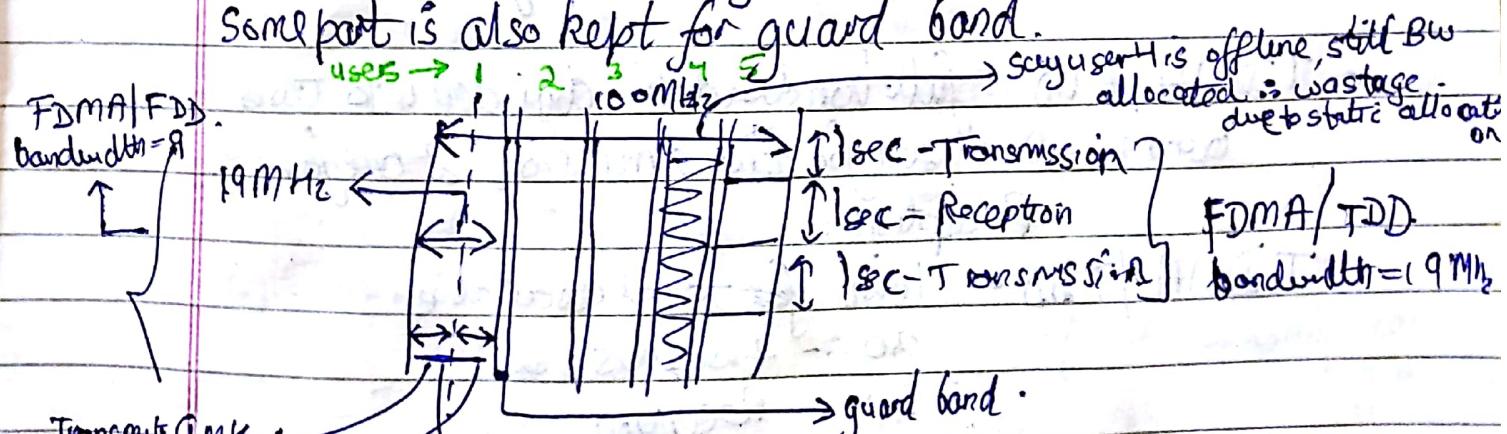
{ FDM - physical layer concept where actual signal is generated.  
 FDM is ~~one to one~~ ~~one to many~~ is access method in Data Link Layer.

### • 3 Types of channelization

#### 1. FDMA (Frequency Division Multiple Access)

- Identify number of users available. // say 5
- spectrum / Bandwidth is also identified. // say 100 MHz
- divide given frequency band into the number of users and each user is assigned a band.

Some part is also kept for guard band.



Transmitting = 9 MHz  
 Reception = 9 MHz  
 guard band = 1 MHz.

Now this is full bandwidth for user One  
 (All other users shared it).

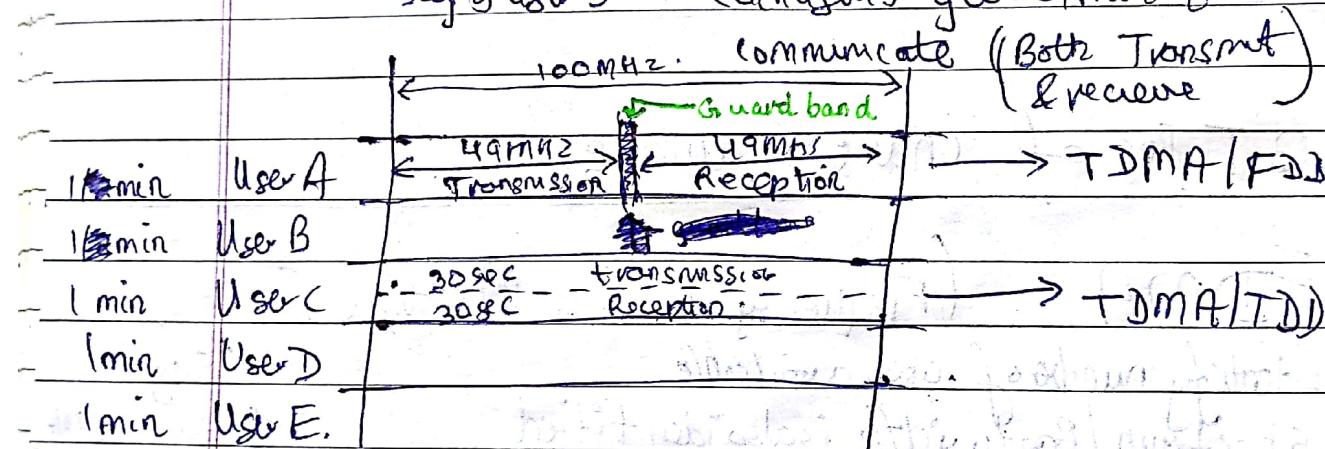
transmission as well as Reception. Inside this frequency band now we divide it into half using FDD or TDD.

there is

- static division of frequency, based on number of users
- wastage of Bandwidth.
- All users will get their dedicated frequency band.

## 2. TDMA (Time Division Multiple Access)

- Superframe of limited time say 1 hr
- say 5 user's - each user is given 1 min to communicate (Both Transmit & receive)



- each user is given full bandwidth for 1 min

- TDMA/FDD - full bandwidth is divided into two bands - One used for transmission & one for Reception.

- TDMA/TDD - 1 min for total available 30s - transmission  
30s - Reception

Channell or occupies entire BW

CDMA differs from TDMA as all stations can send data simultaneously; there is no time slotting.

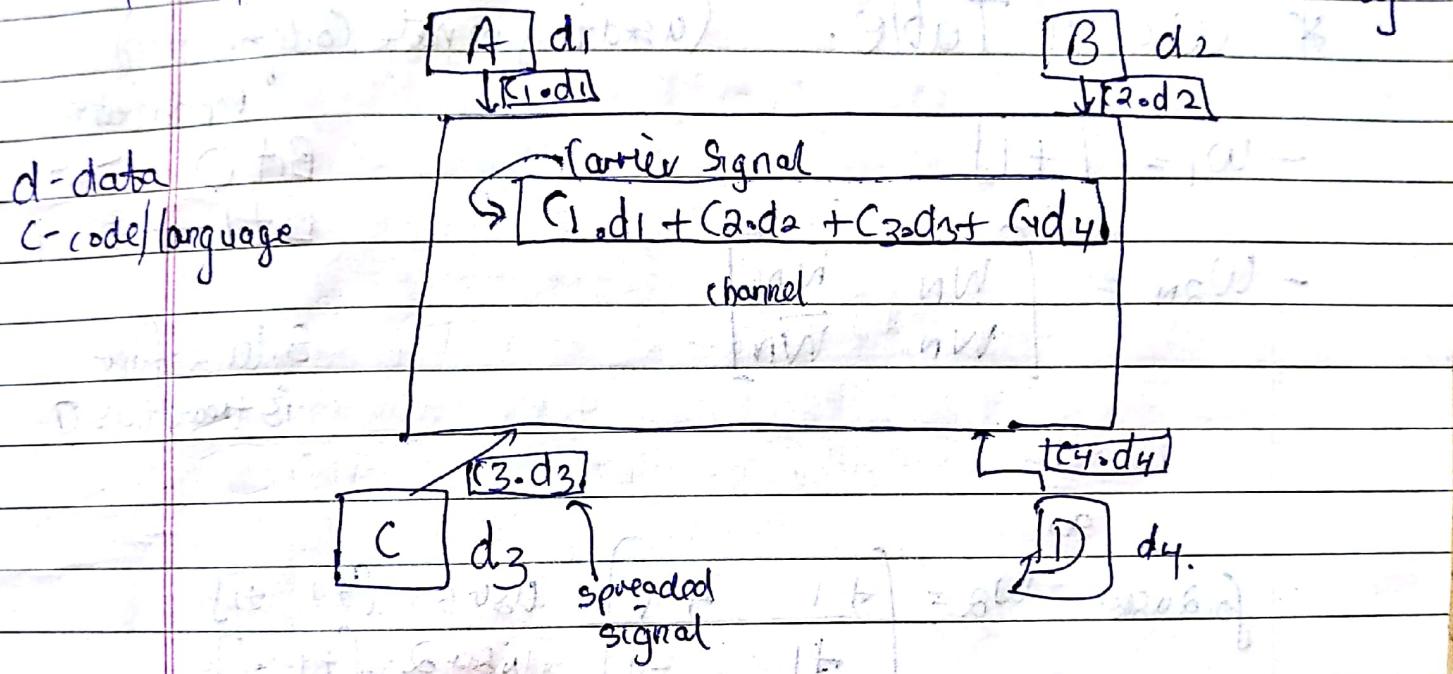
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### 3. CDMA (Code Division Multiple Access)

- is a part of DSSS
- code the data in such a way that no one understands then multiplex
- DSSS (Direct spread spectrum scheme) and send



- \* - one channel carries all transmission's simultaneously



- say  $A \xrightarrow{d_1} D$   
→  $D$  receives ~~full~~ carrier signal.  
∴  $D$  must know  $A$ 's code (i.e. transmitter code/language)

\* Properties of Code used in CDMA

- 1. Auto-correlation

$c_i \cdot c_i = \text{High Value}$  (means we will be able to get data)  
 CDMA Rake Receiver does this multiplication as this high value can be differentiated from data.

- 2. Cross Correlation

$c_i \cdot c_j = 0$  (neutralization of others data)

- Receiving at  $\rightarrow (d_1c_1 + d_2c_2 + d_3c_3 + d_4c_4)$

$\rightarrow$  Sender was A  $\therefore$  multiply by C.

$$\rightarrow (d_1c_1 + d_2c_2 + d_3c_3 + d_4c_4) \cdot C$$

$$= \begin{matrix} \checkmark & & & \\ d_1 & + d_2 & C_2 C_1 & + d_3 C_3 C_1 & + d_4 C_4 C_1 \\ \times & & \times & & \times \end{matrix}$$

\* Walsh Table. (used to generate Code for CDMA)  
Representation

$$- W_1 = [+1]$$

Bit 0  $\rightarrow$  -

Bit 1  $\rightarrow$  1

$$- W_{2N} = \begin{bmatrix} W_N & \bar{W}_N \\ \bar{W}_N & \bar{\bar{W}}_N \end{bmatrix}$$

Silence of user  
is treated as 0

for users:  $W_2 = \begin{bmatrix} +1 & +1 \\ +1 & -1 \end{bmatrix}$  User 1 -  $[+1 +1]$   
User 2 -  $[+1 -1]$ .

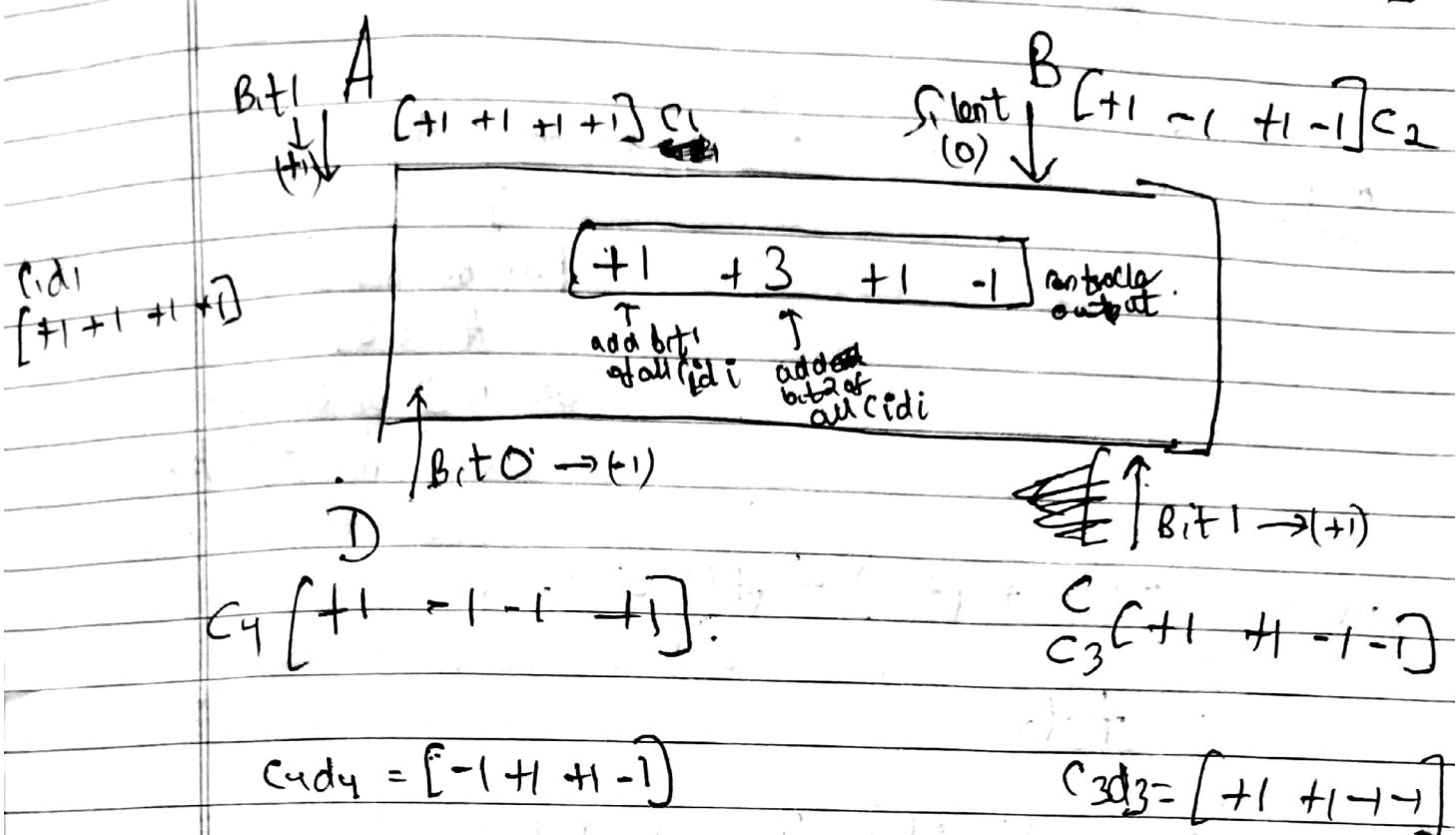
for 4 users	$W_4$ will be	$+1 +1 +1 +1$	User 1 - C <sub>1</sub>
range		$+1 -1 +1 -1$	User 2 - C <sub>2</sub>
calculated		$+1 +1 -1 -1$	User 3 - C <sub>3</sub>
only using		$+1 -1 -1 +1$	User 4 - C <sub>4</sub>

$$\text{User} = [+1 +1 +1 -1 +1]$$

$d_1 = -1$ . ie sending bit 0

$$d_1 \cdot C_1 = (-1 \cdot -1 + 1 \cdot -1) = -1$$

code = [00000]



### \* Components

#### - CDMA transmitter

pb is to multiply data ~~with~~ with code.

#### - CDMA controller

input is all spreaded data from all users.

combine all signals

\* Before transmission a handshake must take b/w sender & receiver to exchange the keys.

#### - CDMA Receiver

A  $\xrightarrow{?} D$

say ~~all nodes~~ Receiver is D

Received signal  $[+1 +3 +1 -1]$   $\rightarrow$  This is received by all nodes

for D calculate sop.

step1  $[+1 \ +3 \ +1 \ -1] \circ [+1 \ +1 \ +1 \ +1]$

key or code of

A pre C,

which D exchanged  
during handshake

step2  $= (+1) + (+3) + (+1) + (-1)$

$$= +4.$$

step3 Now divide the output by number of users; (only for Walsh table Method)

$$\Rightarrow +4/4 = (+1)$$

mean high voltage

$\therefore$  Received bit is 1

example 2

for ~~B~~

$$D \xrightarrow{o} B$$

$$[+1 \quad +3 \quad +1 \quad -1] \cdot [f_1 \quad f_2 \quad f_3 \quad f_4]$$

$$= \underbrace{(+1) + (+3) + (-1) + (-1)}$$

$$= (+1) + (+3) + (-1) + (-1)$$

$$= (-4)$$

$$-4/4 = -1$$

means low voltage

Received bit is 0

book pg 351  
2 numericals

## WIRELESS LAN

- protocols used is CSMA/CA.
- IEEE has defined specifications for a wireless LAN called IEEE 802.11 which covers physical and data link layers.

- There are 2 types of service sets

1. BSS: Basic Service Set
- AP: access point

Category 1 BSS without AP eg Bluetooth.

Category 2 BSS with AP eg Router, cellular network.

- IEEE 802.11 defines BSS as building blocks of a wireless LAN.
- It is made up of a stationary or mobile wireless station and an optional base station called Access Point (AP).
- A BSS without AP is a stand-alone network. It cannot send data to other BSS. It is called ad-hoc architecture.

2. ESS ( Extended service set ).

- many BSS are included.

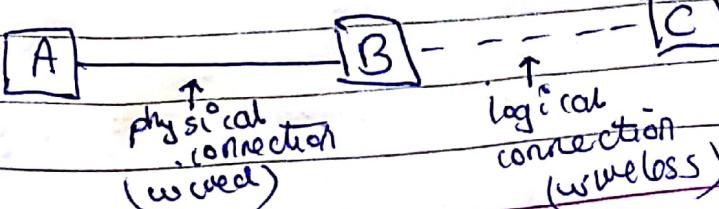
ESS ( 2 or more BSS with AP's )

A BSS with AP is called infrastructure BSS.

- AP redirect request to one which is connected to server.
- BSS's are connected through distributed system.  
( controller-distributed systems )

so we need protocols

we use  
CSMA/CA



hidden terminal problem → A & B are communicating but C does not know. If no multiple access protocol is used then collision will occur at B.

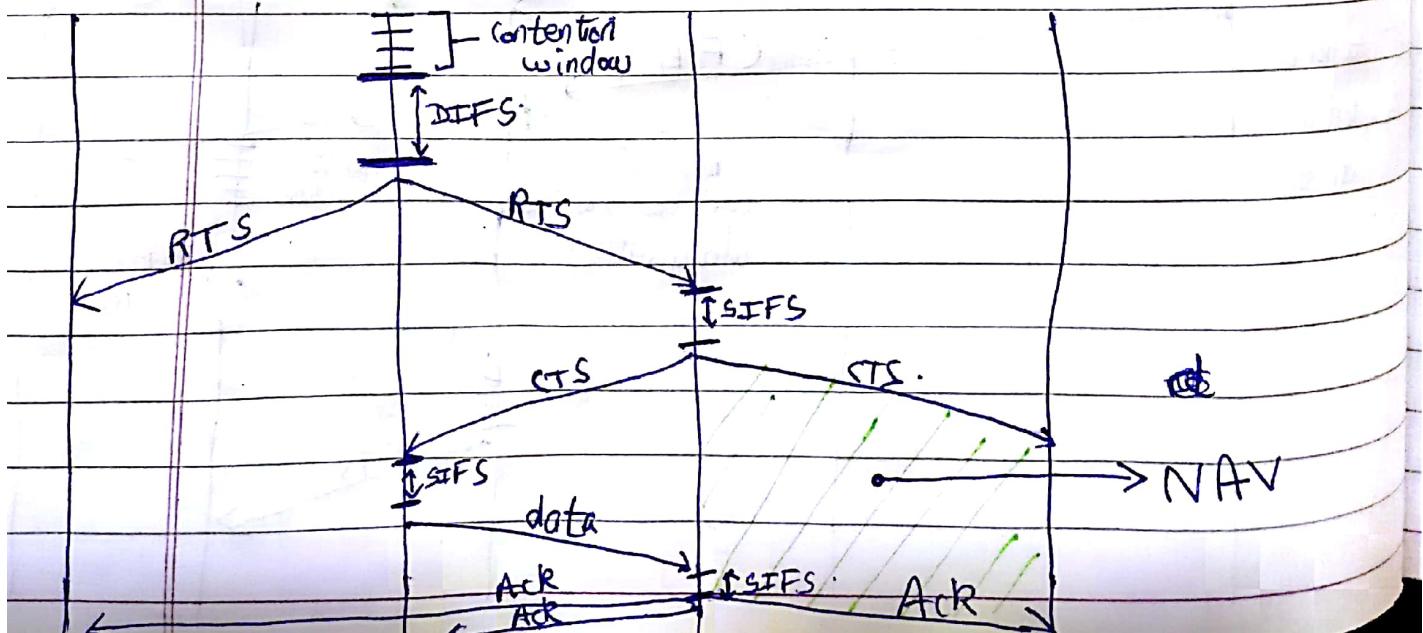
## MAC sublayers:

1. DCF
2. PCF

### 1. DCF (Distributed coordination function)

- uses CSMA/CA as access method
- Before sending frame, source senses the medium
  - (a) channel uses persistence strategy with backoff until channel is idle.
  - (b) After station is found idle, station waits for a period of time called DIFS (distributed interframe space) then it sends a control frame called (Request to send) RTS.
- After receiving RTS and waiting for SIFS (short interframe space) the destination sends a control frame called (Clear to send) CTS to source. It indicates that destination is ready to receive.
- Source sends data after waiting for SIFS time.
- Destination after waiting for SIFS time sends ACK.

other station	Source	Destination	other station
D	A	B	C



SIFS → constant time interval

DIFS → SIFS + 2 \* slot-time

- RTS is sent to all neighbours, informing them that a comm is going to take place. There after all other stations do not transmit or receive.
- CTS is also sent to all neighbours.
- other stations do not perform any communication until sender and they themselves receive acknowledgement.
- Contention Window

Let  $K=3$  Range  $\{0, 1, 2, 3, 4, 5, 6, 7\}$

means transmitter waits for say I choose 4.

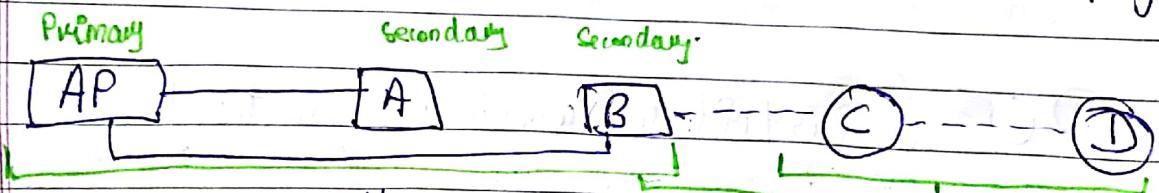
4 time slots.

- NAV (network allocation vector)

when station sends an RTS frame, it includes duration of time that it needs to occupy the channel. The station's that are affected by this transmission create a timer called NAV that shows how much time must pass before those stations are allowed to check channel for idleness.

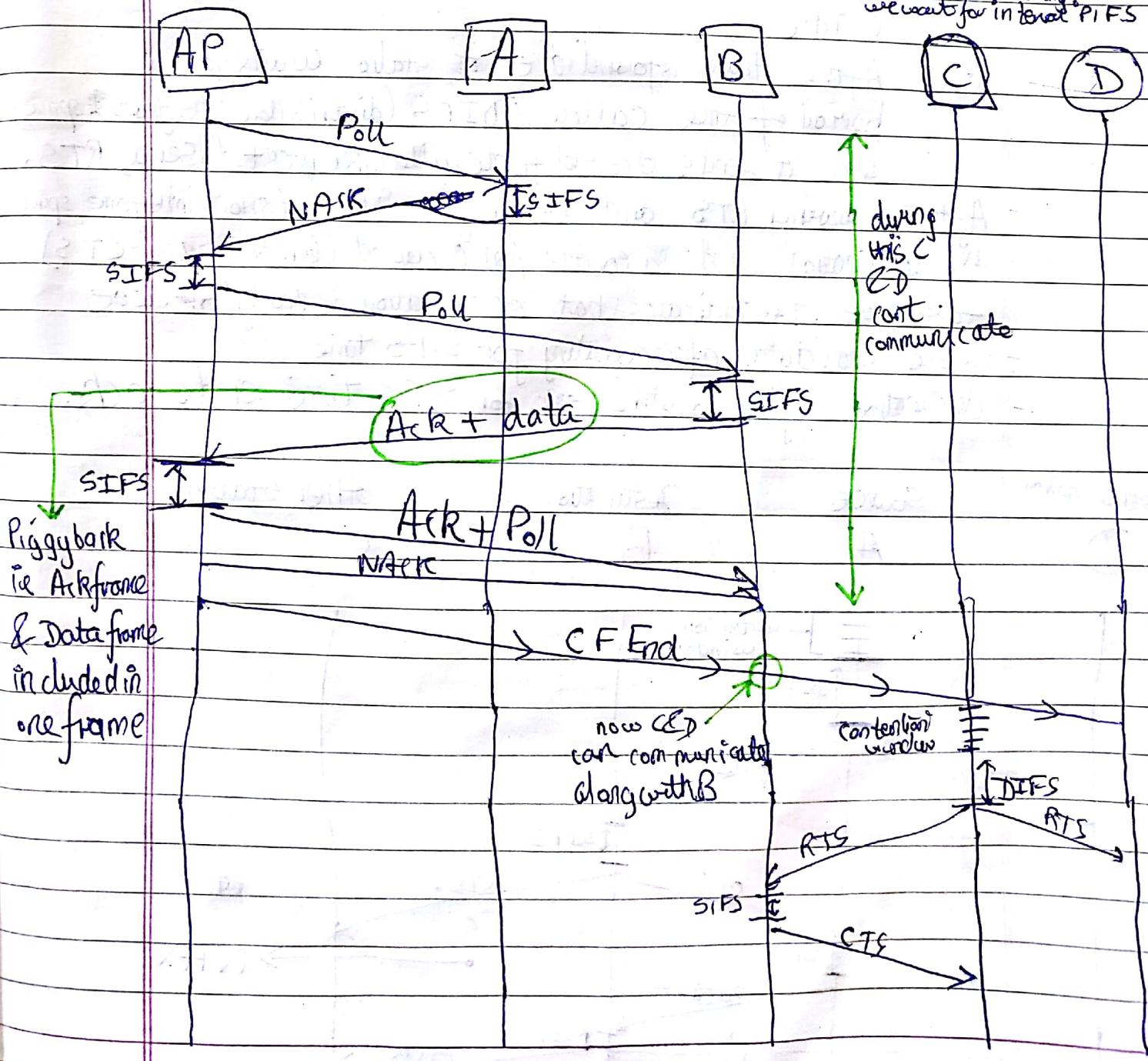
## Q. PCF (Point Coordination Function)

- It is an optional access method that can be implemented only in infrastructure network. It is implemented on top of DCF.



A & B are physically connected to AP  
C & D form an Ad-Hoc network.

a diff types of communication, some need switching.  
we want for instant PIFS



- AP performs polling for stations that are capable of being polled. stations are polled one after another, sending any data they have to AP.
- To give priority to PCF over DCF, interframe space PIFS (PCF IFS) is defined. PIFS is shorter than DIFS.
- If at sometime station wants to use only DCF and an AP wants to use PCF, then the AP has the priority.
- To manage b/w PCF & DCF we use Beacon & CFend signals (frames)
- B: Beacon frame signal
- CF: contention free frame signal.
- Adhoc network is wireless (can do only PCF mode) ie C & D.
- CFend - means contention free slot ended ie closing PCF so now we can switch to DCF mode.