

# Operating Systems



# Module Introduction

- CCAT Exam
  - Operating Systems: 5 Questions
  - Computer Fundamentals + & Concepts of Programming: 10 Questions
- Quiz
  - Operating Systems: 4 Quizzes (10 marks)
  - Computer Fundamentals: 1 Quiz (10 marks)
  - Operating Systems: Module end quiz (20 marks)
- GitLab Repository
  - OS Galvin Slides
  - Practice MCQ
  - Notes/Diagrams (Daily)
  - Reference Book:  
Operating System Concepts - Galvin



# Operating System Concepts

- **Introduction**

- Introduction to Operating System, What is OS, Booting the System

- **System Architecture Design of OS**

- System Calls, Dual Mode Operation: System mode and Kernel mode

- **Process Management**

- What is Process & PCB?
- States of the process
- CPU scheduling & CPU scheduling algorithms
- Inter Process Communication: Shared Memory Model & Message Passing Model
- Process Synchronization/Co-ordination
- Deadlocks & deadlock handling methods



# Operating System Concepts

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## \* **Memory Management**

- Swapping
- Memory Allocation Methods
- Segmentation
- Paging
- Virtual Memory Management



# Operating System Concepts

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## **File Management**

- What is file?
- What is filesystem & filesystem structure?
- Disk space allocation methods
- Disk scheduling algorithms





→ End User .

Editor

IDE

Web Browser

Media Player

→ appln software .

① Interface .

② Resource allocate/  
manage .

Operating System (kernel) .

③ Control Prog .

④ CD/DVD → core OS . + appln software + utilities .

⑤ Kernel → Core OS

CPU

RAM

HDD

KB D

Monitor

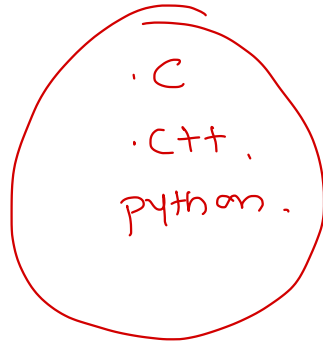
→ Computer Hardware

Program → Set of Instr<sup>n</sup>. given to computer system  
(executable file).

Software → Set of Program.

Utility → Prog with OS.

IDE → Integrated Development Environment.



# Operating System Concepts

## Introduction:

- Why there is need of an OS?
- What is an OS?
- History of OS- Multi-threading, multiprocessing etc
- Functions of an OS
- What is Process?
- States of Process





# Operating System Concepts

## Q. Why there is a need of an OS?

- Computer is a machine/hardware does different tasks efficiently & accurately.
- Basic functions of computer:
  1. Data Storage
  2. Data Processing
  3. Data Movement
  4. Control
- As any user cannot communicates/interacts directly with computer hardware to do different tasks, and hence there is need of some interface between user and hardware.



# Operating System Concepts

## Q. What is an Operating System?

- An OS is a **system software** (i.e. collection of system programs) which acts as an interface between the user and hardware.
- An OS also acts as an interface between programs and hardware.
- An OS allocates resources like main memory, CPU time, i/o devices access etc... to all running programs, hence it is also called a **resource allocator**.
- An OS controls the execution of all programs and it also controls hardware devices which are connected to the computer system hence it is also called a **control program**.
- An OS manages limited available resources among all running programs, hence it is also called a **resource manager**.



# Operating System Concepts

- From End User: An OS is software (i.e. collection of programs) that comes either in CD/DVD and has the following main components:

**1. Kernel:** It is a core program/part of an OS which runs continuously into the main memory and does basic minimal functionalities of it.

e.g. Linux: vmlinuz, Windows: ntoskrnl.exe

**2. Utility Software:** e.g. disk manager, Control panel, windows firewall, anti-virus software etc...

**3. Application Software:** e.g. google chrome, Shell, notepad, MS Office etc.



# Operating System Concepts

## Q. What is a Software?

-Software is a collection of programs.

## Q. What is a Program?

- Program is a set of instructions written in any programming language (either low-level or high-level programming language) given to the machine to do a specific task.

- Three types of programs are there:

1. “user programs”: programs defined by the programmer user/developers

e.g. main.c, hello.java, addition.cpp etc....

2. “application programs”: programs which come with an OS/can be installed later

e.g. MS Office, Notepad, Compiler, IDE's, Google Chrome, Mozilla Firefox, Calculator, Games etc....

3. “System Programs”: programs which are inbuilt into an OS/part of an OS.

e.g. Kernel, Loader, Scheduler, Memory Manager etc...



# Operating System Concepts

## # Functions of an OS:

### Basic minimal functionalities/Kernel functionalities:

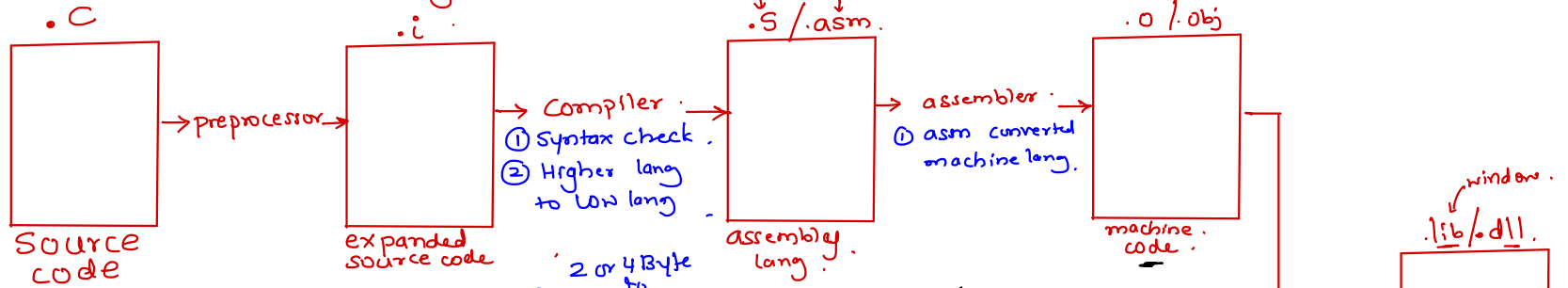
1. Process Management
  2. Memory Management
  3. Hardware Abstraction
  4. CPU Scheduling
  5. File & IO Management
- } kernel,

### Extra utility functionalities/optional:

6. Protection & Security
  7. User Interfacing
  8. Networking
- }



# Process is Prog under Execution



identify the file format

eg :-

- .bmp → BM
- .exe → MZ
- .out → ? ELF

magic no

addr of entry point function

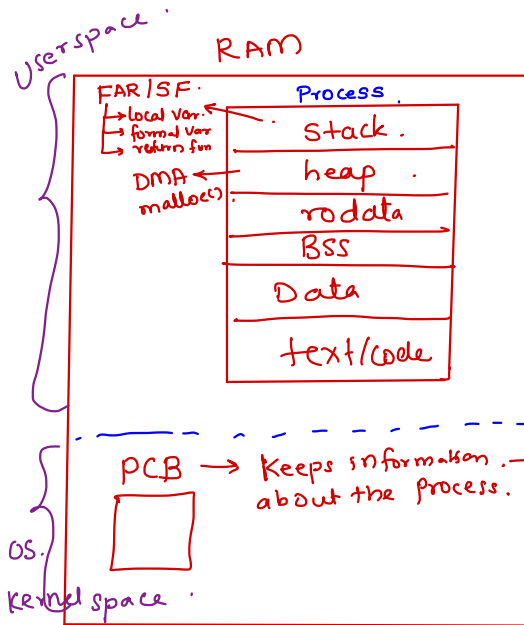
adds of all the section.

ELF: Executable & Linking Format

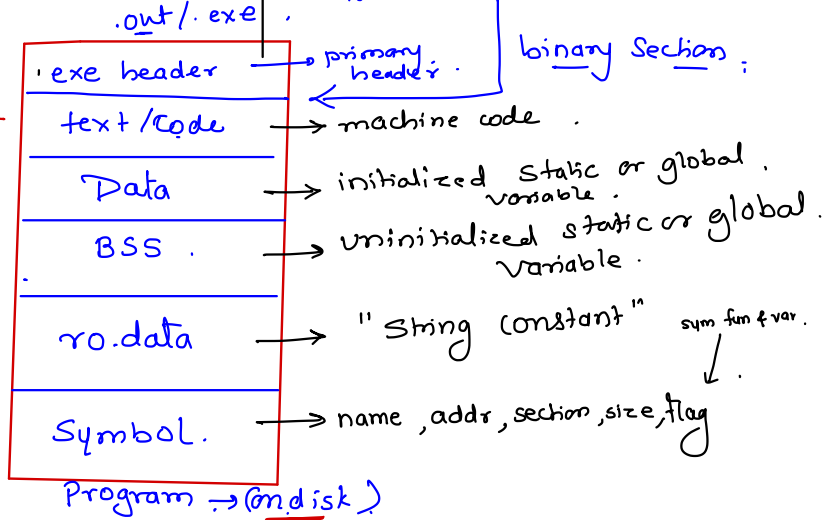
PE: Portable Executable

.out / .exe

.a / .so → Linux



- ① Pid.
- ② Sch info (state, Prior, alog)
- ③ Meminfo.
- ④ Fileinfo.
- ⑤ IPC info
- ⑥ Execut context
- ⑦ Kernel stack
- ⑧ Exit



```
int main()  
{  
    _____  
    _____  
    _____  
    _____  
    return 0;  
}
```

05 → success.

-1 → unsuccessful.

# Operating System Concepts

## Q. What is an IDE (Integrated Development Environment) ?

- It is an application software i.e. collection of tools/programs like **source code editor, preprocessor, compiler, linker, debugger** etc... required for **faster software development**.  
e.g. VS code editor, MS Visual Studio, NetBeans, Android Studio, Turbo C etc....

1. **"Editor"**: it is an application program used for to write a source code.  
e.g. notepad, vi editor, gedit etc...

2. **"Preprocessor"**: it is an application program gets executes before compilation and does two jobs - it executes all preprocessor directives and removes all comments from the source code.  
e.g. cpp

3. **"Compiler"**: it is an application program which convert high level programming language code into low level programming language code i.e. human understandable language code into the machine understandable language code.  
e.g. gcc, tc, visual c etc...





# Operating System Concepts

4. "**Assembler**": it is an application program which converts assembly language code into machine language code/object code.

e.g. masm, tasm etc...

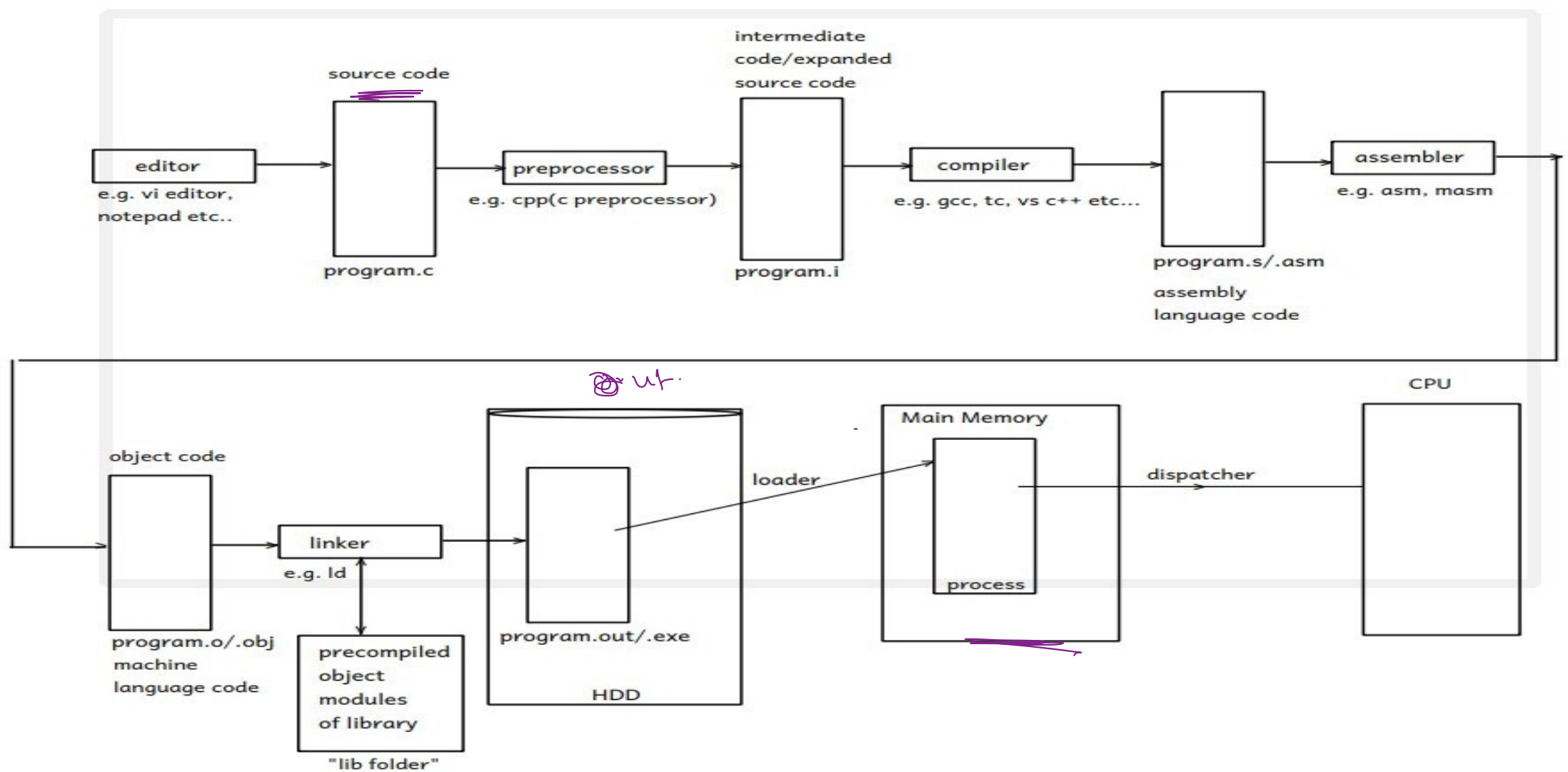
- Program written in any programming language is called as a "**source code**".

5. "**Linker**": it is an application program which links object file/s in a program with precompiled object modules of library functions exists in a lib folder and creates final single executable file.

e.g. ld. link editor in Linux.

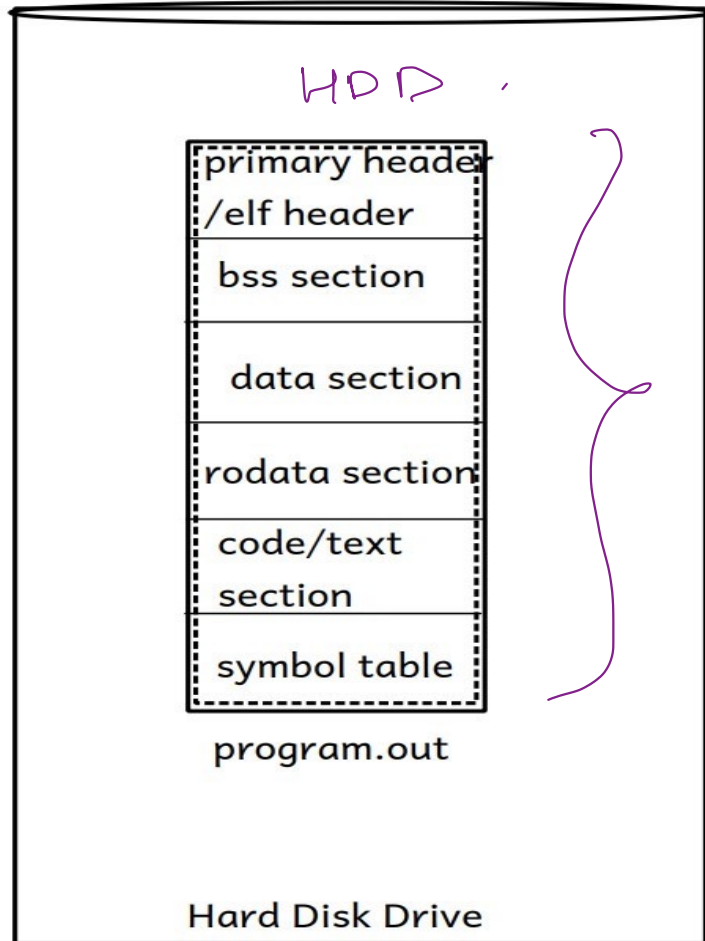


# Operating System Concepts



# Operating System Concepts

Structure of an executable file  
ELF file format in Linux

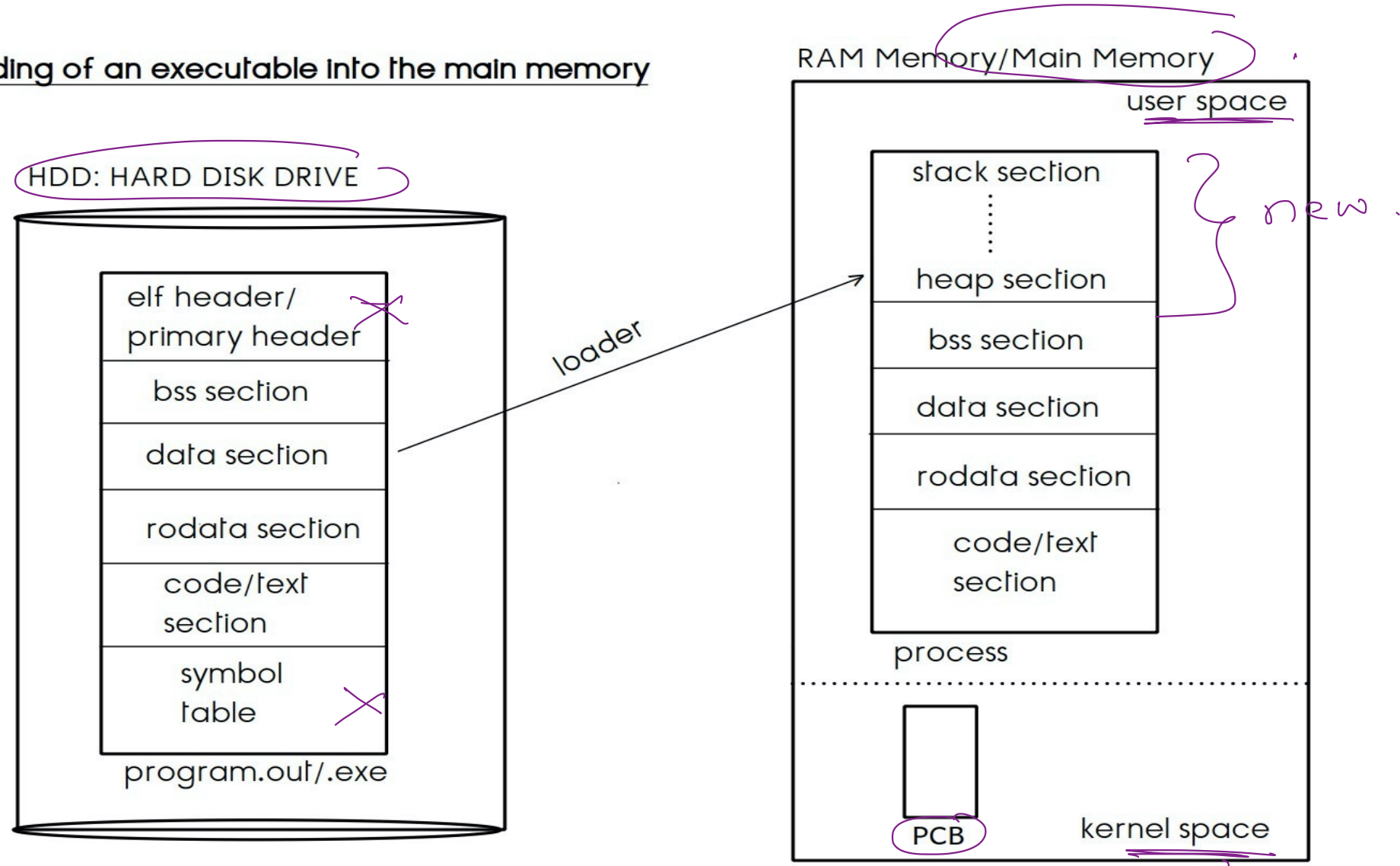


- 1. primary header/exe header:** it contains information which is required to start an execution of the program.  
e.g. - addr of an entry point function --> main() function  
- **magic number**: it is a constant number generated by the compiler which is file format specific.
  - magic number in Linux starts with ELF in its eq **hexadecimal format**.
  - info about remaining sections.
- 2. bss(block started by symbol) section:** it contains uninitialized global & static vars
- 3. data section:** it contains initialized global & static vars
- 4. rodata (readonly data) section:** it contains string literals and constants.
- 5. code/text section:** it contains executable instructions
- 6. symbol table:** it contains info about functions and its vars in a tabular format.



# Operating System Concepts

## Loading of an executable into the main memory



# Operating System Concepts

## Process

- Program under execution.
- Process execute in RAM. (PCB)
- The process control block contains information about the process (required for the execution of the process).
  - Process id
  - Exit status
  - Scheduling information (State, Priority, Sched algorithm, Time, ...)
  - Memory information (Base & Limit, Segment table, or Page table)
  - File information (Open files, Current directory, ...)
  - IPC information (Signals, ...)
  - Execution context
  - Kernel stack
- PCB is also called process descriptor (PD), uarea (UNIX), or task\_struct (Linux).



## Q. What is a Process? User

### view:

- A program in execution is called a process.
- Running a program is called a process.
- When a program gets loaded into the main memory it is referred to as a process.
- Running an instance of a program is referred to as a process.

### System view:

- The process is a file loaded into the main memory which has got bss section, rodata section, code section, and two new sections gets added for the process:
- **stack section**: contains function activation records of called functions.
- **heap section**: dynamically allocated memory



# Operating System Concepts

- file format of an executable file in Windows is **PE** (Portable Executable), whereas the file format of an executable file in Linux is **ELF (Executable & Linkable Format)**.
- The file format is a specific way to store data & instructions of a program inside an executable file, and it is different in diff OS.
- ELF file format divides an executable file logically into sections and inside each section specific contents can be kept in an organized manner:
  1. elf header
  2. bss section (block started by symbol)
  3. data section
  4. rodata (read-only data )section
  5. code/text section
  6. symbol table



# Memory Layout of Program and Process

## Program Consist of

exe header/primary header

Block started by symbol (bss) section (un initialized static / global variables)

Data section (initialized static / global variables)

Rodata Section(Constant/literals)

code/text section (contains executable instructions)

Symbol Table

## Process Consist of

Skipped

Block started by symbol (bss) section (un initialized static / global variables)

Data section (initialized static / global variables)

Rodata Section(Constant/literals)

code/text section (contains executable instructions)

Skipped

Stack Section

Heap Section





# Process and Program

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- A **process** is an instance of a program in execution.
- Running a program is also known as **Process**.
- When a program gets loaded into memory is also known as **Process**.
- A **Program** is a set of instructions given to the machine to do a specific task.



# Interaction with an OS : Two Types of Interface (CUI and GUI)

## 1. CUI/CLI : Command User Interface/Command Line Interface

- By using this kind of interface user can interact with an OS by means entering commands onto the terminal/command line in a text format.

e.g. In Windows name of the program which provide CUI => cmd.exe command prompt

In Linux name of an application program which provides CUI => shell/terminal

In MSDOS name of the program which provides CUI => command.com (MicroSoft Disk Operating System).

## 2. GUI : Graphical User Interface

- by using this kind of interface user can interact with an OS by means making an events like click on buttons, left click/right click/double click, menu bar, menu list etc.....

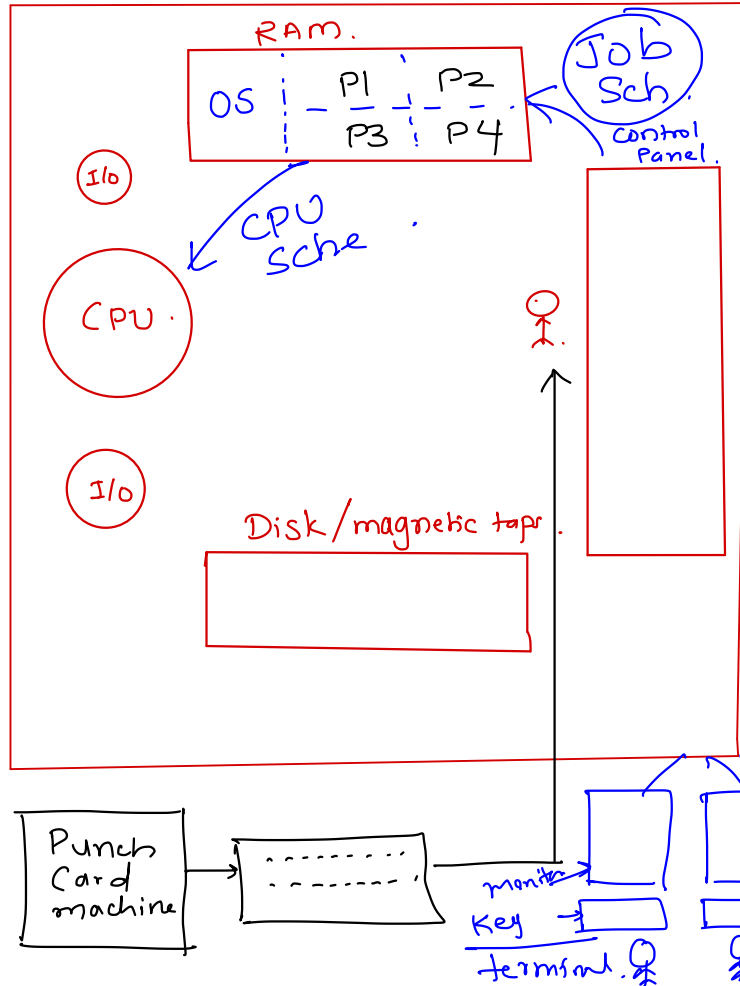
- Windows = User friendly GUI.

e.g. In Windows name of an application program which provides GUI => explorer.exe

In Linux name of an application program which provides GUI => GNOME/KDE (GNU Network Object Model Environment / Common Desktop Environment).



old computer  $\rightarrow$  MainFrame .

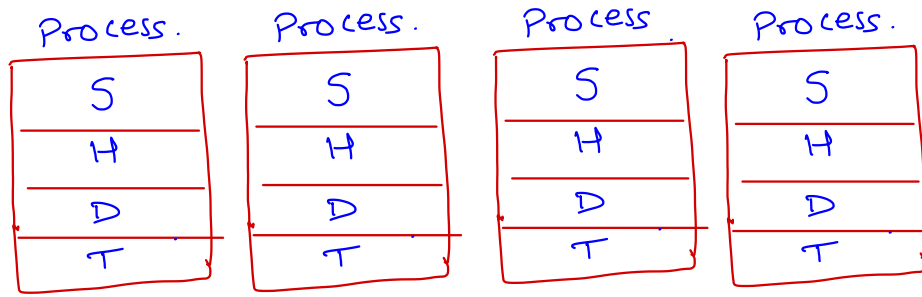


Prog.  $\rightarrow$  CPU Burst time + I/O Burst time .

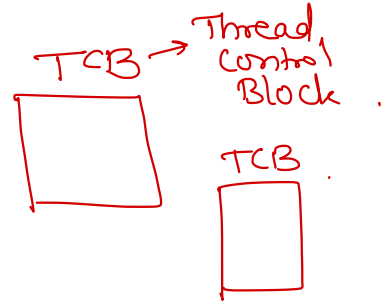
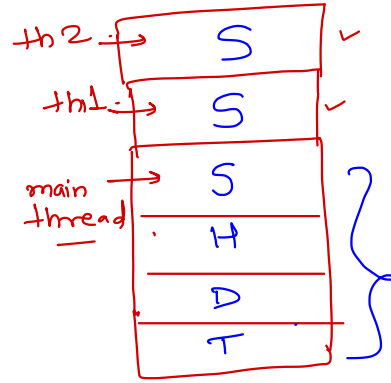
### History of OS:

- ① Resident monitor .
- ② Batch Prog .
- ③ multi-programming
  - $\hookrightarrow$  multiple prog loaded in main mem ,
  - $\hookrightarrow$  Utilization of CPU
- ④ Multi-tasking / time-sharing
  - Process base
  - thread base
  - resp time  $< 1$  sec .
- ⑤ multi-user .
- ⑥ multi-core processing .

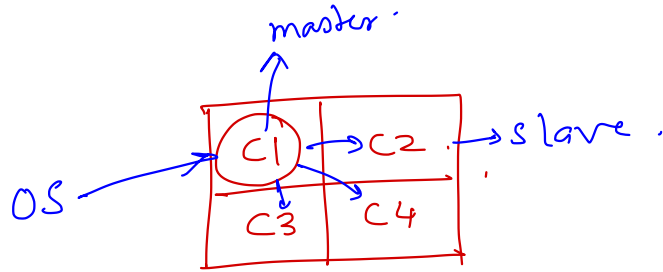
## Process-based



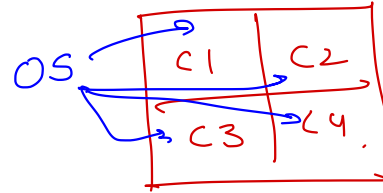
## thread-based → Light weight process.



# Multi-core processing ..



Asymmetric multi-core ,



Symmetric multi-core .

# Operating System Concepts

## ■ History of Operating System

### 1. Resident monitor

### 2. Batch System

- The batch/group of similar programs is loaded in the computer, from which OS loads one program in the memory and execute it. The programs are executed one after another.
- In this case, if any process is performing IO, CPU will wait for that process and hence not utilized efficiently.

### 3. Multi-programming

- Better utilization of CPU
- Loading multiple Programs in memory
- Mixed program(CPU bound + IO bound)

### 4. Time-sharing/Multitasking

- Sharing CPU time among multiple process/task present in main memory and ready for execution
- Any process should have response time should be less then 1sec
- Multi-tasking is divided into two types
  - Process based multitasking
  - Thread based multitasking



# Operating System Concepts

- Process based multitasking: Multiple independent processes are executing concurrently. Processes running on multiple processors called as "multi-processing".
- Thread based multi-tasking OR multi-threading: Multiple parts/functions in a process are executing concurrently.

Thread is a light weight process

- When new thread is created a new stack and new TCB is created.
- Thread Share text, data, heap sections with the parent process

## Process vs thread

- In modern OS, process is a container holding resources required for execution, while thread is unit of execution/scheduling.  
Process holds resources like memory, open files, IPC (e.g. signal table, shared memory, pipe, etc.).  
PCB contains resources information like pid, exit status, open files, signals/ipc, memory info, etc.
- CPU time is allocated to the threads. Thread is unit of execution.
- TCB contains execution information like tid, scheduling info (priority, sched algo, time left, ...), Execution context, Kernel stack, etc.
- For each process one thread is created by default it is called as main thread.



## 5. Multi-user system

Multiple users runs multiple programs concurrently

## 6. Multi-processor/Multi-core system

System can run on a machine in which more than one CPU's are connected in a closed circuit.

Multiprocessing Advantage is it increased throughput (amount of work done in unit time)

- There are two types of multiprocessor systems:
- Asymmetric Multi-processing Symmetric Multi-processing

### ▪ Asymmetric Multi-processing

OS treats one of the processor as master processor and schedule task for it. The task is in turn divided into smaller tasks and get them done from other processors.

### ▪ Symmetric Multi-processing

OS considers all processors at same level and schedule tasks on each processor individually. All modern desktop systems are SMP.





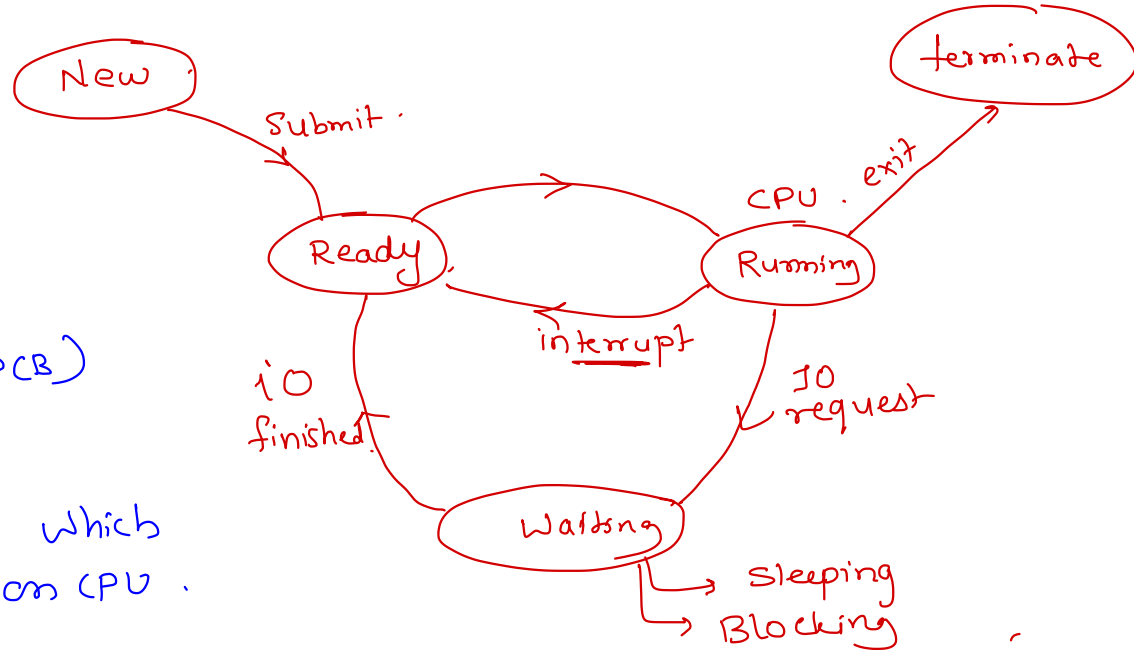
Process - PCB

- ↓
- ① Pid.
- ② Sch inf
- ③ State
- ④ Priority
- ⑤ Alg

OS data structure

- ① Job queue  
list of all process (PCB)
- ② Ready queue  
list of process (PCB) which are ready to run on CPU.
- ③ Waiting queue  
list of PCB which are ready of IO device.

Process-state



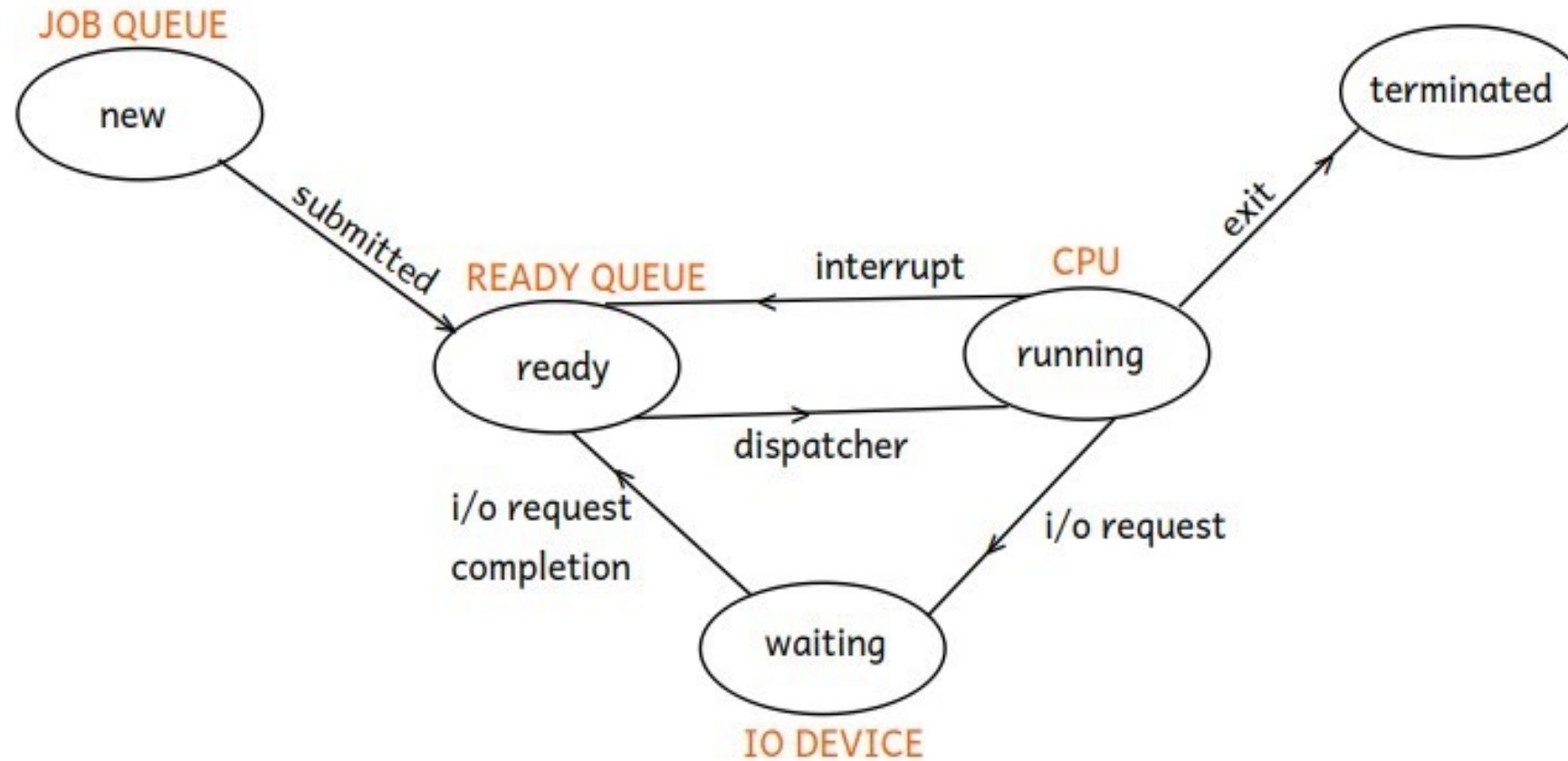
## ▪ **Process States:**

To keep track on all running programs, an OS maintains few **data structures** referred as **OS data Structure**

1. **Job queue:** it contains list of all the processes(PCB).
2. **Ready queue:** it contains list of PCB's of processes which are ready to run on CPU (not for io)
3. **Waiting queue:** it contains list of PCB's of processes waiting for io device or for synchronization



# Five State Process Diagram



PROCESS STATE DIAGRAM

# Five State Process Diagram

## # Process States:

-Throughout execution, process goes through different states out of which at a time it can be only in a one state.

### States of the process:

- 1. New state:** upon submission or when a PCB for a process gets created into the main memory process is in a new state.
- 2. Ready state:** after submission, if process is in the main memory and waiting for the CPU time, it is in a ready state.
- 3. Running state:** if currently the CPU is executing any process then state of that process is considered as a running state.
- 4. Waiting state:** if a process is requesting for any i/o device then state of that process is considered as a waiting state.
- 5. Terminated state:** upon exit, process goes into terminated state and its PCB gets destroyed from the main memory.

