# Surface Simplification using Quadric Error Metrics

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## **Surface Simplification**

- Triangulated surface model
- Several types of simplification techniques
  - Vertex Decimation
  - Vertex Clustering
  - Iterative Edge Contraction
  - Pair Contraction (QEM)

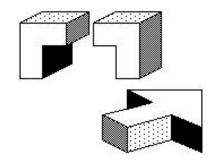
## **Surface Simplification**

- Limitation of other methods:
  - Vertex Decimation: maintain mesh topology, assume manifold model geometry
  - Vertex Clustering: poor control on simplification process, low quality
  - Edge Contraction: no aggregation, assume manifold model geometry.

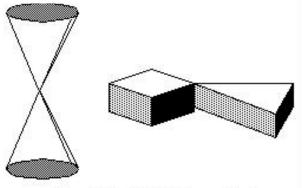
## **Surface Simplification**

#### Manifold Surface

- Points on the manifold surface is isomorphic to a disk
- Non-Manifold Surface
  - vertices with less than 3 adjoining faces
  - edges with more or less than two adjoining faces



Manifold Parts



Non-Manifold / Open Parts

Image source: http://claymore.engineer.gvsu.edu/~jackh/eod/design/design-160.html

## **Model Boundary Representation**

#### Manifold

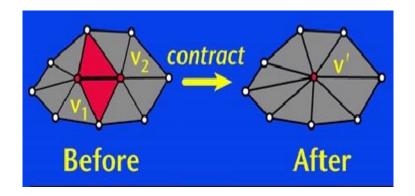
- Half-Edge data structure (directional edge)
- Winged-Edge data structure (every edge connects to exactly two faces, one on each side.)

#### Non-Manifold

 Radial-Edge data structure (edges can connect to arbitrarily many faces)

## **Edge Contraction vs Pair Contraction**

- Edge Contraction
  - Contract v1, v2 to v' when there is an edge connecting v1 and v2.
- Pair Contraction



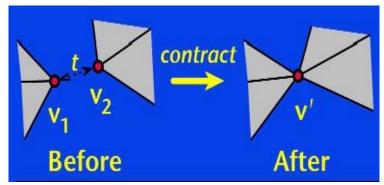


Image source: Garland's presentation in SIGGRAPH97

## **Aggregation**

#### Example

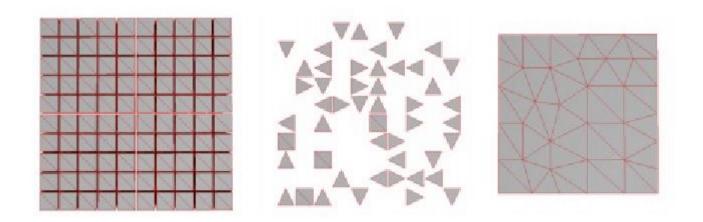


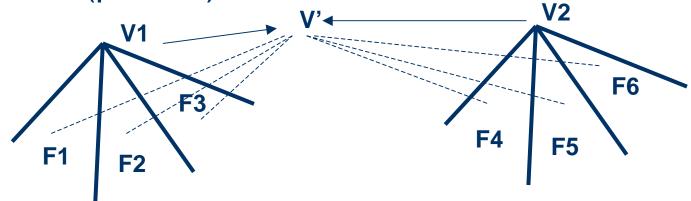
Image source: original QEM paper

#### **Pair Selection**

- Considering contracting every possible pair is not applicable.
- Select a pair v1,v2
  - (v1, v2) is an edge
  - ||v1 v2|| < t, t is a threshold

## Vertex as intersection of planes

 Vertex can be defined by the set of adjacent faces (planes)



#### **Cost of Pair Contraction**

- The cost of contracting a pair (v1,v2) to v' is the error of replacing v1 and v2 by the new created vertex v'.
- For a given point, the error is the sum of square distance to the associated set of planes.
- The cost of the contraction is the sum of square distance to the union of two sets of faces adjacent to v1 and v2. (F1, F2, F3, F4, F5, F6)

- In order to calculate the cost contraction, each vertex must keep track of a set of planes.
- After the contraction, the newly created vertex should keep track of the union of two sets of planes associated with the previous two vertices.
- Explicitly storing these planes takes more storage.

- We need the set of planes to calculate the cost of contraction, which is the sum of square distances to the set of planes.
- As long as the sum of square distance to the set of planes can be calculated, we do not need to store a set of plane equations.
- The set of plane equations associated with each vertex can be replaced by a 4x4 symmetric matrix Q.

 The distance between a point (x0,y0,z0,1) and a plane ax+by+cz+d=0:

- Let p = [a b c d], 
$$v = [x0, y0, z0, 1]$$
  $(p^Tv)$ 

- Distance formula 
$$D = \frac{ax_0 + by_0 + cz_0 + d}{\sqrt{a^2 + b^2 + c^2}}$$

The sum of square distances:

$$\Delta(\mathbf{v}) = \Delta([v_x \ v_y \ v_z \ 1]^\mathsf{T}) = \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} (\mathbf{p}^\mathsf{T} \mathbf{v})^2$$

#### Sum of Square Distances

$$\begin{split} \Delta(\mathbf{v}) &= \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} (\mathbf{v}^\mathsf{T} \mathbf{p}) (\mathbf{p}^\mathsf{T} \mathbf{v}) \\ &= \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} \mathbf{v}^\mathsf{T} (\mathbf{p} \mathbf{p}^\mathsf{T}) \mathbf{v} \\ &= \mathbf{v}^\mathsf{T} \left( \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} \mathbf{K}_{\mathbf{p}} \right) \mathbf{v} \qquad \mathbf{K}_{\mathbf{p}} = \mathbf{p} \mathbf{p}^\mathsf{T} = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix} \end{split}$$

Sum of all quadric matrices

Quadric Surface

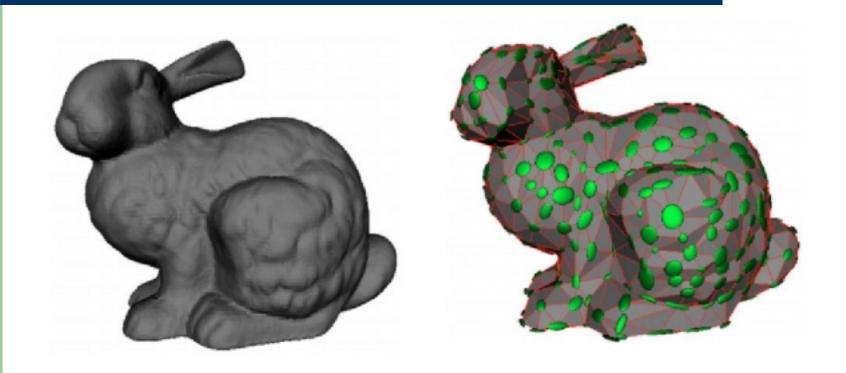
$$\mathbf{v}^{\mathsf{T}}\mathbf{Q}\mathbf{v} = \begin{bmatrix} x & y & z & 1 \end{bmatrix} \begin{bmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{22} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ q_{14} & q_{24} & q_{34} & q_{44} \end{bmatrix} \begin{bmatrix} x \\ y \\ z \\ 1 \end{bmatrix}$$

$$v^{T}Qv = q_{11}x^{2} + 2q_{12}xy + 2q_{13}xz + 2q_{14}x$$
  
  $+ q_{22}y^{2} + 2q_{23}yz + 2q_{24}$   
  $+ q_{33}z^{2} + 2q_{34}z + q_{44}$ 

#### **Quadric Surface**

- When the sum of square distances is fixed, the equation is an equation of quadric surface: (ellipsoid, hyperboloid, etc)
- The ellipsoid is centered at the vertex.
- Every point on the ellipsoid surface has the same error.
- The ellipsoid captures the local shape of the surface.

# **Example of quadric surface**



## Optimal position for the new vertex

- Now we know how to evaluate the cost of contraction, so to determine the cost of contraction, we need to figure out the position of the new vertex after contraction.
- Solution: find the vertex location that minimize the cost of the contraction.

#### Optimal position for the new vertex

Taking partial derivatives of the cost function:

$$\partial \Delta / \partial x = \partial \Delta / \partial y = \partial \Delta / \partial z = 0.$$

$$\begin{bmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{22} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ 0 & 0 & 0 & 1 \end{bmatrix} \bar{\mathbf{v}} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

$$\bar{\mathbf{v}} = \begin{bmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{22} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ 0 & 0 & 0 & 1 \end{bmatrix}^{-1} \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

Position for the new Vertex

## Optimal position for the new vertex

 If the matrix is not invertible, find the optimal vertex on the segment joining the vertex pair, or choosing among the two vertices and the midpoint of the two vertices.

## **Simplification Algorithm**

- 1. Compute the **Q** matrices for all the initial vertices.
- 2. Select all valid pairs.
- 3. Compute the optimal contraction target  $\bar{\mathbf{v}}$  for each valid pair  $(\mathbf{v}_1, \mathbf{v}_2)$ . The error  $\bar{\mathbf{v}}^{\mathsf{T}}(\mathbf{Q}_1 + \mathbf{Q}_2)\bar{\mathbf{v}}$  of this target vertex becomes the *cost* of contracting that pair. Union of two sets of planes
- Place all the pairs in a heap keyed on cost with the minimum cost pair at the top.
- 5. Iteratively remove the pair  $(\mathbf{v}_1, \mathbf{v}_2)$  of least cost from the heap, contract this pair, and update the costs of all valid pairs involving  $\mathbf{v}_1$ .

# **Simplification Result**

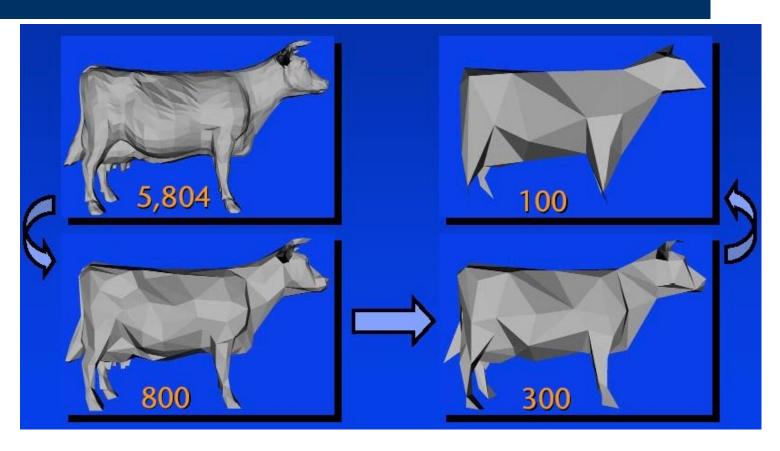


Image source: Garland's presentation in SIGGRAPH97

## **Simplification Result**

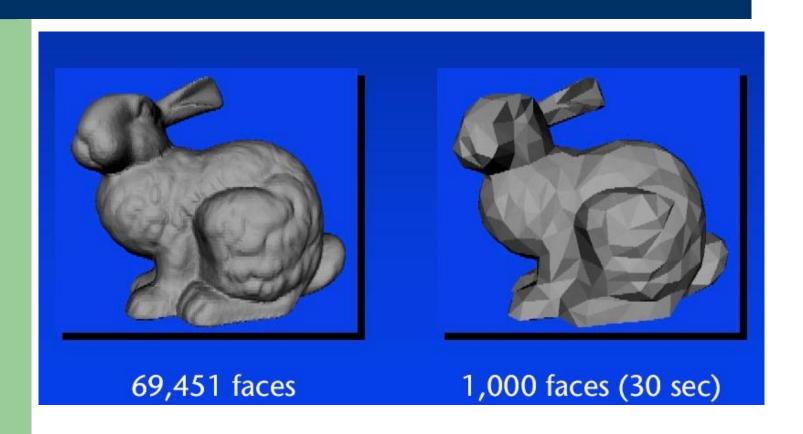


Image source: Garland's presentation in SIGGRAPH97

#### Conclusion

- Fast Algorithm and good approximation
- Compact storage
- General surface handling

#### Reference

- K. Weiler. The Radial Edge data structure: A topological representation for non-manifold geometric boundary modeling. In J. L. Encarnacao, M. J. Wozny, H. W. McLaughlin, editors, Geometric Modeling for CAD Applications, pages 3--36. Elsevier Science Publishers B. V. (North--Holland), Amsterdam, 1988.
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