DPLL Based SAT-Solver

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1 Key Implementation Ideas

1. **Parser:** By `input().split()`, I could easily split those words. To provide convenient debug way, the *CNF-parser.py* also accepts inputs from terminal. To maintain robustness and alert when input is incorrect, format checkers are also added in it. You can see different error types at the head of *CNFparser.py*.

2. **Data Structure:** The way I chose to express clauses is to have D *a dictionary of variable to the set of equation index it appears*, E *a list of equations which are sets of variables*, and L *a list of current equation index*.

3. **Basic Property I:** Any time in the program running, it is guaranteed that the 3 data structures above only represent current situation, which means eliminated variables and equations will not appear in any of them.

4. **Set Variable:** *Set Variable* is a fundamental function for *BCP*, *Set Pure*, and *Decision*. To keep *Property I*, when using D to go through the equation E_k where variable v appears, I need to check two thing:

   (a) If the equation becomes empty, I need to eliminate it in L.

   (b) If the equation becomes true, I need to eliminate it in L, and remove all the other variables v' in this equation which clears E_k, and get $D[v']$ rid of k.

5. **Basic Property II:** Whenever program goes into *Set Variable*, it is assumed that the assignment will not generate conflicts. Thus *Set Variable* does not check satisfiability. Moves described below will guarantee this.

6. **BCP:** *BCP* needs to go through E to single out those one literal equation. Use a *set* to store them. When one is added, check whether its negation is inside. There is a potential trap in it depending how you implement *BCP*, I will describe it below.

7. **Set Pure:** *Set Pure* needs to go through D to pick those literals the negation of which does not appear in equations. After this, set them to *true*, which will maximumly preserve the satisfiability.

8. **Decision:** *Decision* is what program will do after *BCP* and *Set Pure*. The way to pick what variable to try the assignment is pretty tricky. I tried several strategies, such as *Random Pick, Frequent Pick, Balanced Pick*[1]. It turns out *Pick the most frequently appeared* is about 30% faster than *Random Pick* and about 250% outperforms *Balanced Pick*. *Frequent Pick* and *Balanced Pick* combined together can not provide substantial improvement.

9. **DPLL:** *DPLL* is basically the combination of the algorithms above. To maintain *Property I*, I need to fix change when program recalls. Considering efficiency, I have 2 set of *lists*; one records changes from *BCP* and *Set Pure*, one records changes from *Decision*. When program backtracks from the first assignment trial, undo the *Decision* change; when it returns from the second trial, undo all the change and return to parent function. There is a case which will not be examined during this. I will describe it below.

2 Potential Traps

1. In *DPLL*, my program will not check whether there exists conflict in one clause, which is both X and $\neg X$ in the same clause. So I checked this in my preparation part.

2. In *BCP*, there is a case needs further consider. Assume after you go through all the clauses containing one literal, these literals are set *true* at the same time without further consideration, there will be a potential bug. Considering this case $\{X\}, \{\neg X, \neg Y\}, \{Y\}$. If you set $X = true$ and $Y = true$ without checking the clause $\{\neg X, \neg Y\}$, it will be satisfiable, which is incorrect. So in my version, I set them one by one.

3. Pay attention to Python Import. Name file and function carefully. I was bothered by the strange error log when I have file like *string.py* or *parser.py* with function *parse* in it.

4. I used test data from UBC website[2]. In these data, there is a mysterious % at the end of every file input. So, I had to add a special check for this in my *CNFparser.py*.

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[1] *Balanced Pick* means the difference between it and its negation is as small as possible. It indicates either *true* or *false* assignment is balanced in *DFS*.

[2] `http://www.cs.ubc.ca/~hoos/SATLIB/benchm.html`