

Assignment - 11

- 1) Big Omega Notation Prove that $g(n) = n^3 + 2n^2 + 4n$ is $\Omega(n^3)$

Solution:

To prove that exist Positive constant c and n_0 such that for all $n > n_0$ to show $g(n) \geq c \cdot n^3$

$g(n) = n^3 + 2n^2 + 4n$ is $\Omega(n^3)$ we can choose $c=1$ and $n_0 = 1$ then for all $n \geq 1$

$$n^3 + 2n^2 + 4n \geq 1 \cdot n^3$$

$$g(n) = n^3 + 2n^2 + 4n \text{ is } \Omega(n^3)$$

- 2) Big theta Notation. Determine whether $g(n) = 4n^3 + 3n$ is $\Theta(n^3)$ or not.

Solution:

To determine whether $f(n) = 4n^3 + 3n$ is $\Theta(n^3)$ we need to check if there exist positive constant c_1 , c_2 and n_0 such that for all $n \geq n_0$

$$c_1 \cdot n^3 \leq h(n) \leq c_2 \cdot n^3$$

$$h(n) = 4n^3 + 3n \leq c_2 \cdot n^3$$

$$h(n) = 4n^3 + 3n \leq 4n^3 + 3n^2 \leq 7n^3$$

$$c_2 = 7$$

$$h(n) = 4n^2 + 3n \leq 7n^2$$

$$c_1 = 4$$

$$4n^2 \leq nn^2 + 3n \leq 7n^2$$

$$h(n) = 4n^2 + 3n \text{ is } O(n^2)$$

3. Let $f(n) = n^3 - 2n^2 + n$ and $g(n) = n^2$ show that $f(n) = \Omega(g(n))$ is true or false justify your answer?

$$f(n) \geq c \cdot g(n)$$

Substituting

$f(n)$ and $g(n)$ into this inequality

we get

$$n^3 - 2n^2 + n \geq (c - n^2)$$

$$n^3 - 2n^2 + n \geq -cn^2$$

$$n^3 - 2n^2 + n + cn^2 \geq 0$$

$$n^3 + (c-2)n^2 + n \geq 0$$

$$n^3 + (c-2)n^2 + n \geq 0 \quad (n^3 \geq 0)$$

$$n^3 + (1-2)n^2 + n = n^3 - n^2 + n \geq 0$$

$$f(n) = n^3 - 2n^2 + n \text{ is } \Omega(g(n) = \sqrt{2}(n^2))$$

$$f(n) = \Omega(g(n)) \text{ is True}$$

Determine whether $h(n) = n \log n + n$ is $O(n \log n)$. Prove a rigorous proof for your conclusion

$$c_1 n \log n < h(n) \leq c_2 n \log n$$

$$h(n) \leq c_2 n \log n$$

$$h(n) \leq n \log n + n$$

$$n \log n + n \leq c_2 n \log n$$

$$1 + \frac{n}{n \log n} \leq 2$$

$$1 + \frac{1}{\log n} \leq 2$$

$$h(n) \text{ is } O(n \log n)$$

$$h(n) \geq c_1 A_1$$

$$h(n) \geq n \log n + n$$

$$n \log n + n \geq c_1 n \log n$$

$$1 + \frac{n}{n \log n} \geq c_1$$

$$\frac{1}{\log n} \geq 0$$

$$1 + \frac{1}{\log n} \geq c_1$$

$$h(n) \text{ is } \Omega(n \log n)$$

$$h(n) \leq n \log n + n \text{ is } O(n \log n)$$

$$1 + \frac{1}{\log n} \geq 1$$

5.

Solve the following recurrence relations and find the order of growth of solutions

$$T(n) = 2T\left(\frac{n}{2}\right) + n^2, \quad T(1) = 1$$

$$T(n) = 4T\left(\frac{n}{2}\right) + n^2, \quad T(1) = 1$$

$$T(n) = 4T\left(\frac{n}{2}\right) + f(n)$$

$$a = 4, \quad b = 2, \quad f(n) = n^2$$

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

$$f(n) = O(n^{\log_b a - c})$$

$$f(n) = O(n^{\log_b a}) \quad T(n) = O(n^{\log_b a} \log n)$$

$$f(n) = O(n^{\log_b a + \epsilon}) \quad \text{then } T(n) = f(n)$$

$\log_b a$:

$$\log_b a = \log_2 n = 2$$

$$f(n) = n^2 = O(n^2)$$

$$f(n) = O(n^2) = O(n^{\log_b a})$$

$$T(n) = n + (n/2) + n^2$$

$$T(n) = O(n^{\log_b a} \log n) = O(n^2 \log n)$$

$$T(n) = nT(n/2) + n^2 \quad \text{with } T(1) = 1 \quad \text{is } O(n^2 \log n)$$