CS372 Assignment 4

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March 13, 2012

P9 Show that the maximum efficiency of pure ALOHA is 1/(2e). *Note*: This problem is easy if you have completed the problem above!

The probability that a node in pure ALOHA has a successful transmission is

$$\lim_{N \to \infty} p(1-p)^{2(N-1)}$$

To find the maximum efficiency of pure ALOHA and show that is 1/(2e) we take the dirivitive

$$f(n) = Np(1-p)^{2(N-1)}$$

$$f'(n) = N(1-p)^{2(N-1)} - Np2(N-1)(1-p)^{2(N-3)}$$

= $N(1-p)^{2(N-3)}((1-p) - p2(N-1))$

- **P10** Consider two nodes, A and B, that use the slotted ALOHA protocol to contend for a channel. Suppose node A has more data to transmit than node B, and node A's retransmission probability p_A is greater than node B's retransmission probability, p_B .
 - a. Provide a formula for node A's average throughput. What is the total efficiency of the protocol with these two nodes?
 - b. If $p_A = 2p_B$, is node A's average throughput twice as large as that of node B? Why or why not? If not, how can you choose p_A and p_B to make that happen?
 - c. In general, suppose there are N nodes, among which node A has retransmission probability 2p and all other nodes have retransmission probability p. Provide expressions to compute the average throughputs of node A and of any other node.
- P13 Consider a broadcast channel with N nodes and a transmission rate of R bps. Supose the braodcast channel uses polling (with an additional polling node) for multiple access. Suppose the amount of time from when a node completes transmission until the subsequent node is permitted to transmit (that is, the polling delay) is d_p . Suppose that within a polling round, a given node is allowed to transmit at most Q bits. What is the maximum throughput of the broadcast channel?

- P14 Consider three LANs interconnected by two routers, as shown in Figure 5.38.
 - a. Assign IP addresses to all the interfaces. For Subnet 1 use addresses of the form 192.168.1.xxx; for Subnet 2 use address of the form 192.168.2.xxx; and for Subnet 3 use addresses of the form 192.168.3.xxx.
 - b. Assign MAC addresses to all the adapters
 - c. Consider sending an IP datagram from Host E to Host b. Suppose all of the ARP tables are up to date. Enumerate all the steps, as done for the single-router example in Section 5.4.2.
 - d. Repeat (c), now assuming that the ARP table in the sending host is empty (and the other tables are up to date).

P19 Explain why a minimum frame size is required for Ethernet. For example, 10Base Ethernet imposes a minimum frame size constraint of 64 bytes. (If you have done the previous problem, you might have relized the reason). Now suppose that the distance between two ends of an Ethernet LAN is d. Can you derive a formula to find the minimum frame size needed for an Ethernet packet? Based on your reasoning, what is the minimum required packet size for an Ethernet that spans 2 kilometers?

P24 In this problem you will derive the efficiency of a CSMA/CD-like multiple access protocol. In this protocol, time is slotted adn all adapters are synchronized to the slots. Unlick slotted ALOHA, however, the length of a slot (in seconds) is much less than a frame time (the time to transmit a frame). Let S be the length of a slot. Suppose all frames are of constant length L = kRS, where R is the transmission rate of the channel and k is a large integer. Suppose there are N nodes, each with an infinite number of frames to send. We also assume that $d_{prop} < S$, so that all nodes can detect a collision before the end of a slot time. The protocol is as follows:

- If, for a given slot, no nod has possesion of the channel, all nodes contend for the channel; in particular, each node transmits in the slot with probability p. If exactly one node transmits in the slot, that node takes possession of the channel for the subsequent k-1 slots and transmits its entire frame.
- If some node has possession of the channel, all other nodes refrain from transmitting until the node that possesses the channel has finished transmitting its frame. Once this node has transmitted its frame, all nodes contend for the channel.

Note that the channel alternates between two states: the productive state, which lasts exactly k slots, and the nonproductive state, which lasts for a random number of slots. Clearly, the channel efficiency is the ratio of k/(k+x), where x is the expected number of consecutive unproductive slots.

- a. For fixed N and p, determine the efficiency of this protocol.
- b. For fixed N, determine the p that maximizes the efficiency.

c. Using the p (which is a function of N) found in (b), determine the efficiency as N approaches infinity.

P25 Suppose two nodes, A and B, are attached to opposite ends of an 800 m cable, and that they each have one frame of 1,500 bits (including all headers and preambles) to send to each other. Both nodes attempt to transmit at time t=0. Suppose there are four repeaters between A and B, each inserting at 20-bit delay. Assume the transmission rate is 100 Mbps, and CSMA/CD with backoff intervals of multiples of 512 bits is used. After the first collision, A draws K=0 and B draws K=1 in the exponential backoff protocol. Ignore the jam signal and the 96-bit time delay.

- a. What is the one-way propagation delay (including repeater delays) between A and B in seconds? Assume that the signal propagation speed is $2 * 10^8$ m/sec.
- b. At what time (in seconds) is A's packet completely delivered at B?
- c. Now suppose that only A has a packet to send and that the repeaters are replaced with switches. Suppose that each switch has a 20-bit processing delay in addition to a store-and-forward delay. At what time, in seconds, is A's packet delivered at B?

P27 Consider Figure 5.38 in problem P14. Provide MAC addresses and IP addresses for the interfaces at Host A, both routers, and Host F. Suppose Host A sends a datagram to Host F. Give the source and destination MAC addresses in the fram encapsulating this IP datagram as the frame is transmitted (i) from A to the left router, (ii) from the left router to the right router, (iii) from the right router to F. Also give the source and destination IP addresses in the IP datagram encapsulated within the frame at each of these points in time.