DNA Sequencing and Data Analysis

Prof Noam Shomron Hadas Volkov

Lecture 7, December 16, 2022

DNA Sequencing and Data Analysis

Friday 8:45 AM to 11:15 AM Arazi-Ofer Building, C.LO3

<u>nshomron@gmail.com</u> <u>hadas.volkov@post.runi.ac.il</u>

DNA Sequencing and Data Analysis

Sequence Mapping and Alignment

Class	Title	Content/assignments	Activity, location
1, 4.11	Introduction to Cells and DNA	Basic knowledge of biology	In the lecture hall, Noam
2, 11.11	DNA Sequencing past and present	Basic knowledge of molecular DNA	In the lecture hall, Noam
3, 18.11	Genomics technologies	DNA, RNA, technologies	In the lecture hall, Noam
4, 25.11	Introduction to Bioinformatics challenges in reading DNA	Focus on three methods: WES/WGS, RNA-seq, cell-free DNA	In the lecture hall, Noam
5, 2.12	Modern DNA Sequencing, 2nd wave File Formats, tools.	Analysis approaches for WES/WGS, RNA-seq, cell-free DNA	In the lecture hall, Hadas and Noam
6, 9.12	De novo Shotgun Assembly	The algorithms and methods behind the assembly problem	In computer class, Hadas and Noam
7, 16.12	Sequence Mapping and Alignment	The algorithms behind mapping and alignment, fast and heuristics	In computer class, Hadas and Noam
8, 23.12	Variant Calling and Somatic Variant Analysis	The bioinformatics behind discovery of novel mutations in cancer	In computer class, Hadas and Noam
9, 30.12	Nanopore data analysis introduction and class activity	The bioinformatics behind Nanopore analysis, activity on computers	In computer class, Hadas and Noam
10, 6.1	Practice molecular biology techniques	Pipetting, transferring small amounts of fluids, running a dry Nanopore experiment	In biology class, Meitar and Noam
11, 13.1	Nanopore DNA sequencing	Nanopore DNA sequencing, experimental run	In biology class, Meitar, Hadas, Assaf
12, 20.1	Nanopore data analysis	Nanopore DNA analysis, experimental run	In computer class, Hadas and Noam
13, 27.1	Nanopore data analysis and presentations	Groups present their results	In the lecture hall, Hadas and Noam

Why Do We Need Sequence Mapping?

Determine the origin of an unknown sequence

Find homologous sequences

Determine genomic position of a sequence

Identify genomic variants between samples (variant calling)

Determine the function of a sequence (annotation)

Two Stages of Sequence Mapping

1. SEARCH -

Roughly find the position of the query in the DB

ACCTGAGGATCGTATACAAGTTA

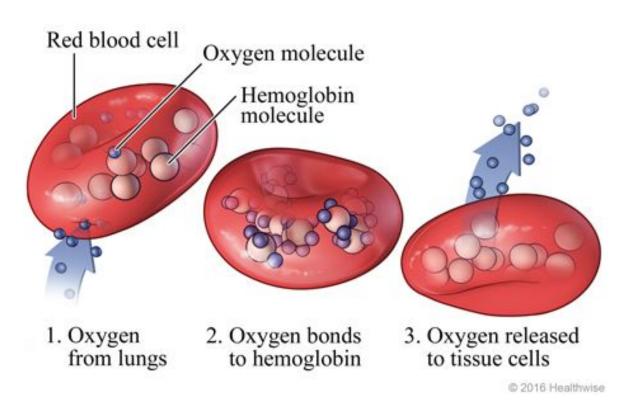
GTGTACAG

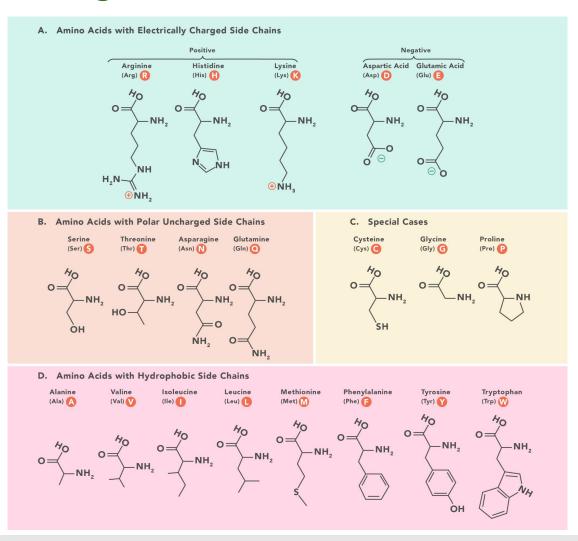
2. ALIGN -

Find the exact pairwise alignment of the query and the DB sequences

G	Т	G	Т	Α	С	Α	_	G
G	Т	Α	Т	Α	С	Α	Α	G

Hemoglobin Homologous





Hemoglobin Homologous

The Hamming Distance

Hamming distance is the number of symbols or positions of two strings at which their corresponding characters are different

```
def hamming_distance(string1, string2):
    if (len(string1) != len(string2)):
        raise Exception('Strings must be of equal length.')
    dist_counter = 0
    for n in range(len(string1)):
        if string1[n] != string2[n]:
            dist_counter += 1
    return dist_counter / len(string1)
```

The Hamming distance between r1 and q1 is: 0.1690 The Hamming distance between r2 and q1 is: 0.3169

The Hamming Distance

```
# NCBI Reference Sequence: XP 028905054.1 (platypus hemoglobin subunit A);
q2 = skbio.Protein("MLTDAEKKEVTALWGKAAGHGEEYGAEALERLFQAFPTTKTYFSHFDLSHGSAQIKAHGKKVADA\
LSTAAGHFDDMDSALSALSDLHAHKLRVDPVNFKLLAHCILVVLARHCPGEFTPSAHAAMDKFLSKVATVLTSKYR")
q2
Protein
Stats:
    length: (141
    has gaps: False
    has degenerates: False
    has definites: True
    has stops: False
   MLTDAEKKEV TALWGKAAGH GEEYGAEALE RLFQAFPTTK TYFSHFDLSH GSAQIKAHGK
   KVADALSTAA GHFDDMDSAL SALSDLHAHK LRVDPVNFKL LAHCILVVLA RHCPGEFTPS
120 AHAAMDKELS KVATVLTSKY R
```

q2 = skbio.Protein("MLTDAEKKEVTALWGKAAGHGEEYGAEALERLFQAFPTTKTYFSHFDLSHGSAQIKAHGKKVADA\LSTAAGHFDDMDSALSALSDLHAHKLRVDPVNFKLLAHCILVVLARHCPGEFTPSAHAAMDKFLSKVATVLTSKYK-")

The Hamming Distance

The Hamming distance between r1 and q2 is: 0.90845

The Hamming distance between r2 and q2 is: 0.92254

MVLSPADKTNVKAAWGKVGAHAGEYGAEALERMFLSFPTTKTYFPHFDLSHGSAQVKGHGKKVADALTNAVAHVDDMPNALSALSDLHAHKLRVDP\
M-LTDAEKKEVTALWGKAAGHGEEYGAEALERLFQAFPTTKTYFSHFDLSHGSAQIKAHGKKVADALSTAAGHFDDMDSALSALSDLHAHKLRVDP\

The Hamming distance between r1 and q2_aligned is: 0.27465

The Hamming distance between r2 and q2_aligned is: 0.34507

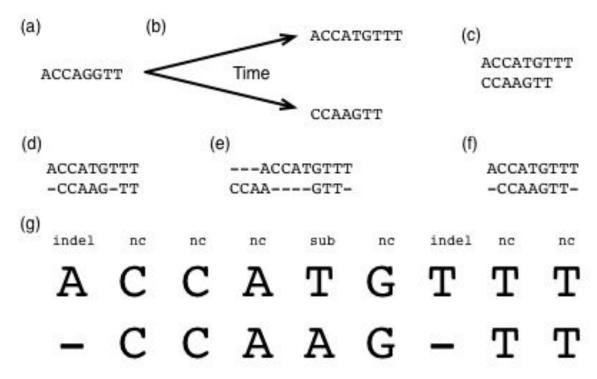
What Is Sequence Alignment?

Mutations:

Substitutions, where one DNA base is replaced with another

Insertions, where one or more contiguous DNA bases are inserted into a sequence

Deletions, where one or more contiguous DNA bases are deleted from a sequence.



ACCATGTTT CCAAGTT

ACCATGTTT

C

(

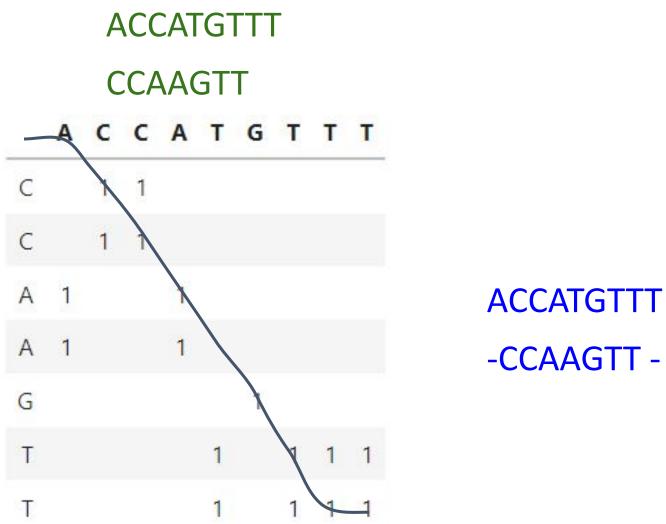
A

A

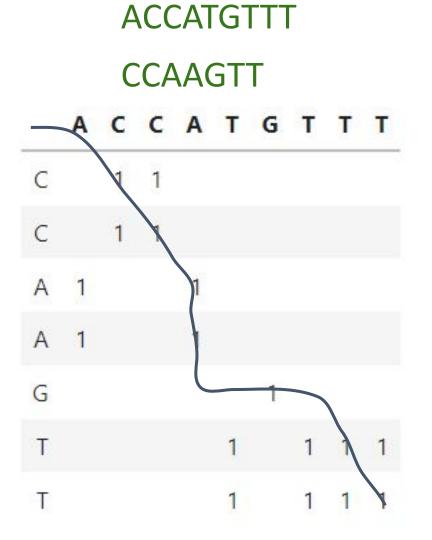
0

1

1



ACCATGTTT



ACCA--TGTTT
-CCAAG---TT

ACCATGTTT

CCAAGTT

ACCATGTTT

-CCAAGTT -

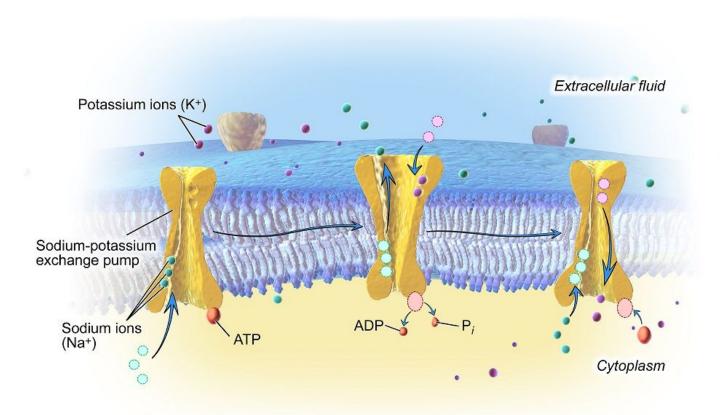
ACCA--TGTTT

-CCAAG- - -TT

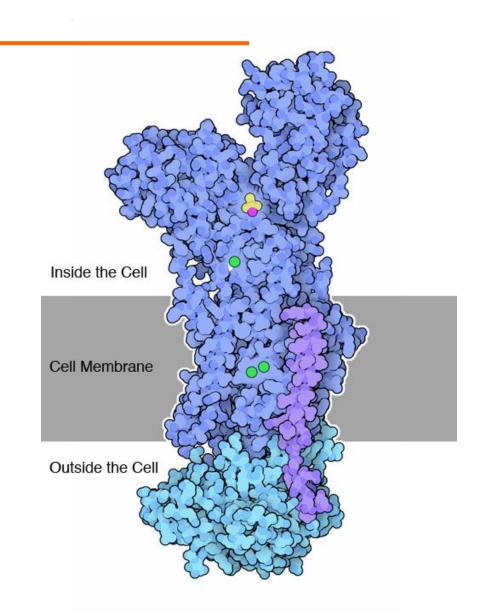
$$S = -1+1+1-1+1+1-1=4$$

$$S = -1+1+1+1-1-1-1-1-1+1+1=-1$$

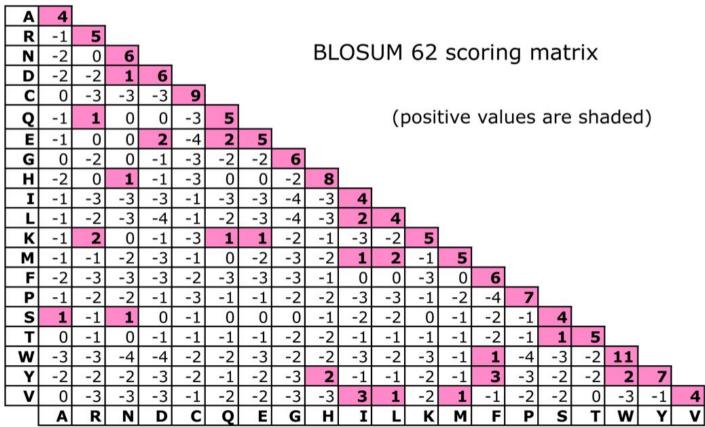
Too Simplistic



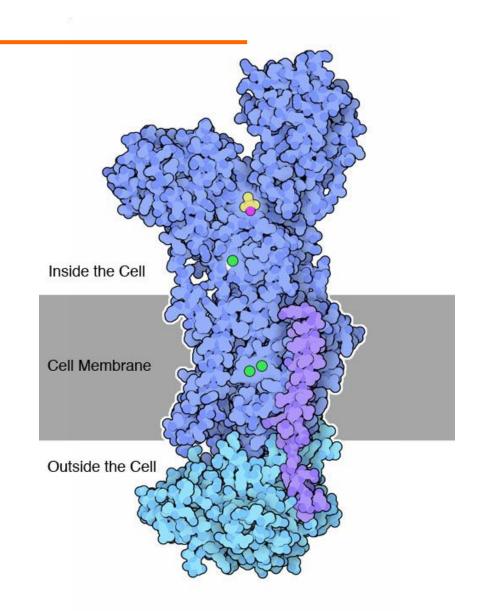
The Sodium-Potassium Exchange Pump

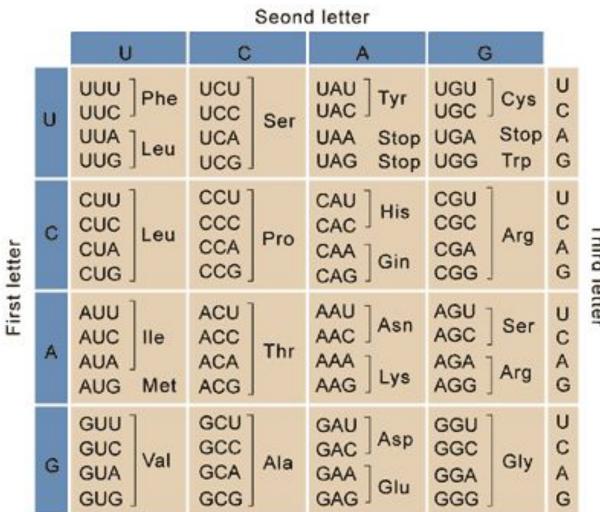


Too Simplistic

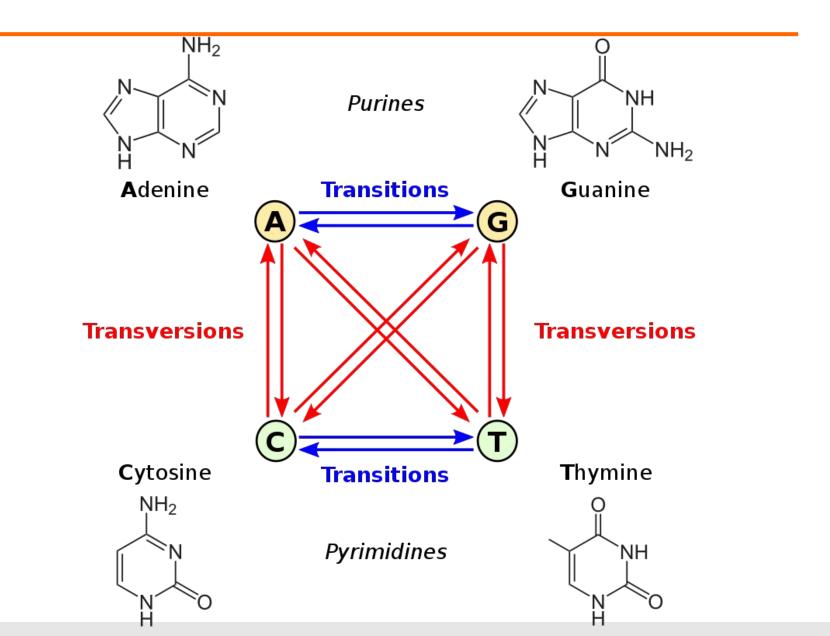


The values for amino acid substitutions were obtained from Henikoff S & Henikoff JG (1992) Amino acid substitutions matrices from protein blocks. *Proc. Natl. Acad. Sci.* **89**: 10915-10919.





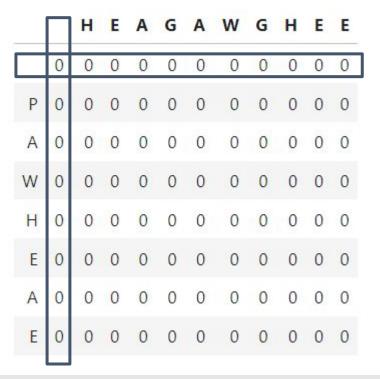
Third letter



HEAGAWGHEE PAWHEAE

F - The Dynamic Programming Matrix

T - The Traceback Matrix



$$F(0,0)=0$$

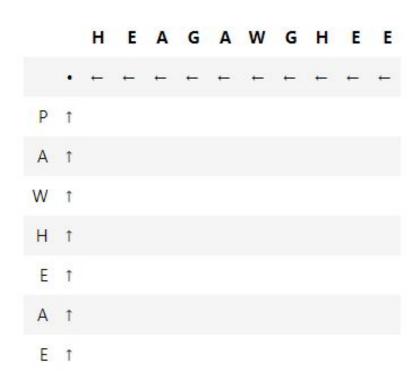
$$F(i,0)=F(i-1,0)-d$$

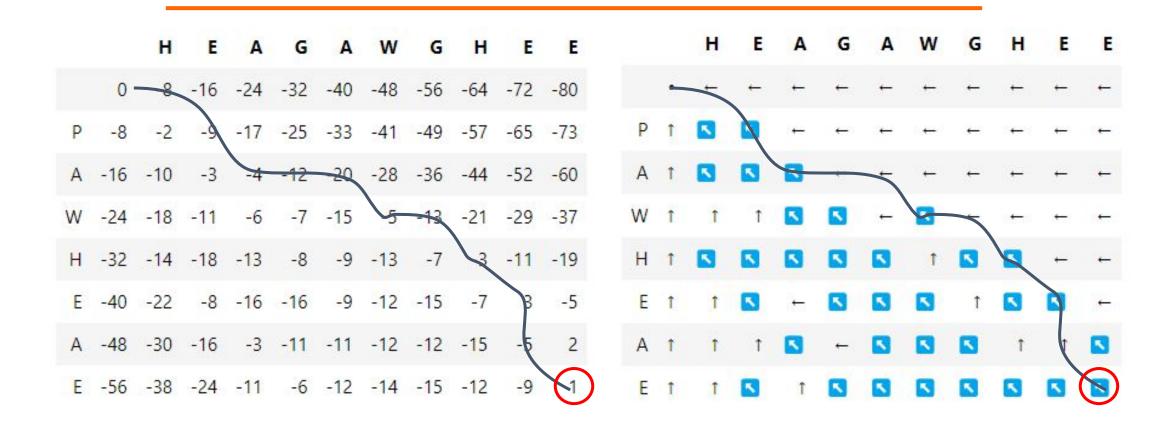
$$F(0,j)=F(0,j-1)-d$$

$$T(0,0)=.$$

$$T(i,0)=\leftarrow$$

$$T(0,j)=\uparrow$$





HEAGAWGHE - E - PA - -W - HEAE Score = 1

Local vs. Global Alignment

Global alignment - try to match entire sequences

Useful for closely-related sequences of similar size

Local alignment - allow partial matching

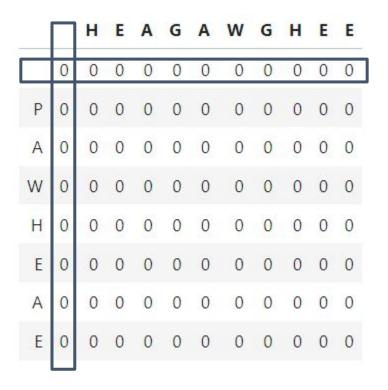
Useful for sequences expected to contain some similarity regions

Global Alignment

Local Alignment

Smith-Waterman Local Alignment

HEAGAWGHEE PAWHEAE



Smith-Waterman Local Alignment

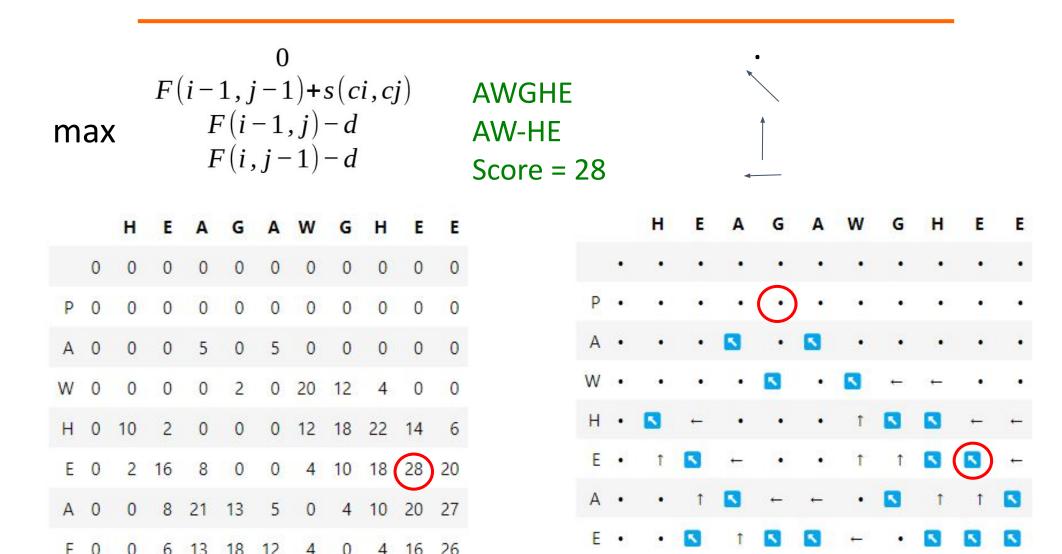
$$F(0,0)=0$$

 $F(i,0)=0$
 $F(0,j)=0$

$$T(0,0)=.$$

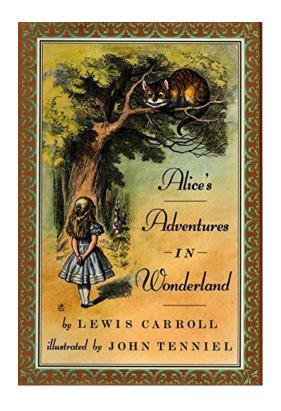
 $T(i,0)=.$
 $T(0,j)=.$

Smith-Waterman Local Alignment



Search

Imagine we have a big book...



... and we want to search it for a specific sentence

It would be 66 so nice if something made sense for a Lewis Carroll change.

Alice in Wonderland

Search

- How can we do it in a timely manner?
 - Brute force
 - Indexing
- Do we allow slight changes?
 - e.g.: "it could be so nice if something made sense"
- Do we allow insertions and deletions?
 - e.g.: "it would be so nice if something made a little sense"
- What if the sentence is repeated in several places in the book?

It would be
66 so nice if
something
made sense
for a

Lewis Carroll
Alice in Wonderland change.

Sequence Mapping Challenges

Large DBs - millions to billions of nucleotides/AAs

Repetition - biological sequences tend to repeat

Noisy - sequencing errors and real biological variants

BLAST - Basic Local Alignment Search Tool

The most popular alignment tool

BLAST finds regions of similarity between biological sequences.

Compares nucleotide or protein sequences to sequence databases

Calculates the statistical significance of DB hits

Allows searching for **imperfect** sequence matches

Uses a **heuristic** algorithm to improve efficiency



BLAST - Algorithm

- 1. Index the DB
- 2. Generate query words
- 3. compute neighbour words
- 4. Search the DB for exact word matches seeds
- 5. Elongate and combine seeds to get final alignment
- 6. Score alignment

BLAST - Indexing

Only needed the first time a DB is used

Mask repetitive and low-complexity regions -

ATATATTTATT → atatatttatt

Break DB sequences into overlapping words of length W

- W=3 for amino acids
- W=11 for nucleotides

Create a lookup table of words with their positions



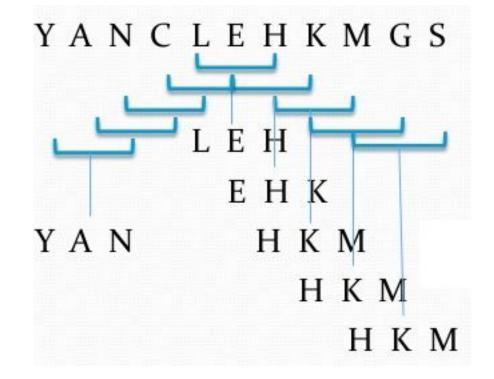


BLAST - Breaking Query to Words

A query of length L produces L-W+1 overlapping words of length W

L = 11

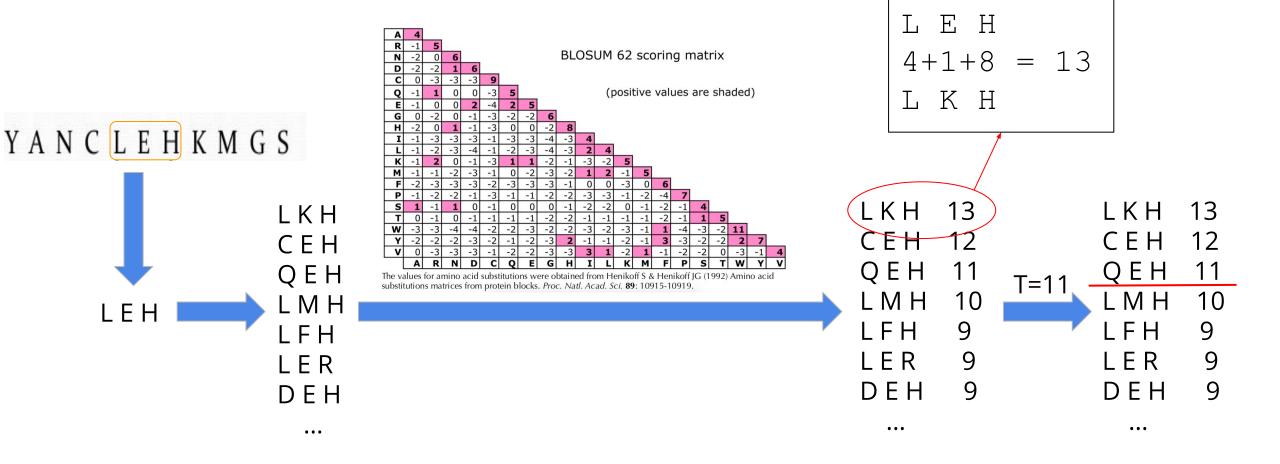
M = 3



BLAST - Finding Neighbour Words

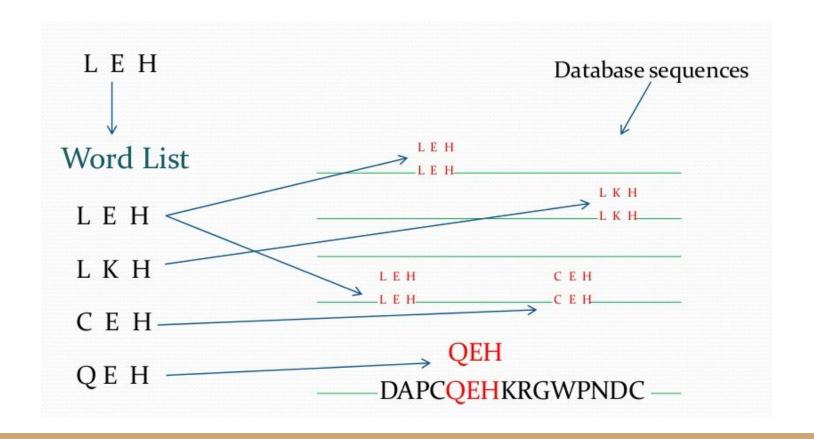
- 1. For each word, find all neighbourhood words
 - = words with one change
- 2. Use a scoring matrix to assign each neighbourhood word a score
- 3. Discard neighbourhood words with score < T

BLAST - Breaking Query to Words



BLAST - Finding Alignment Seeds in DB

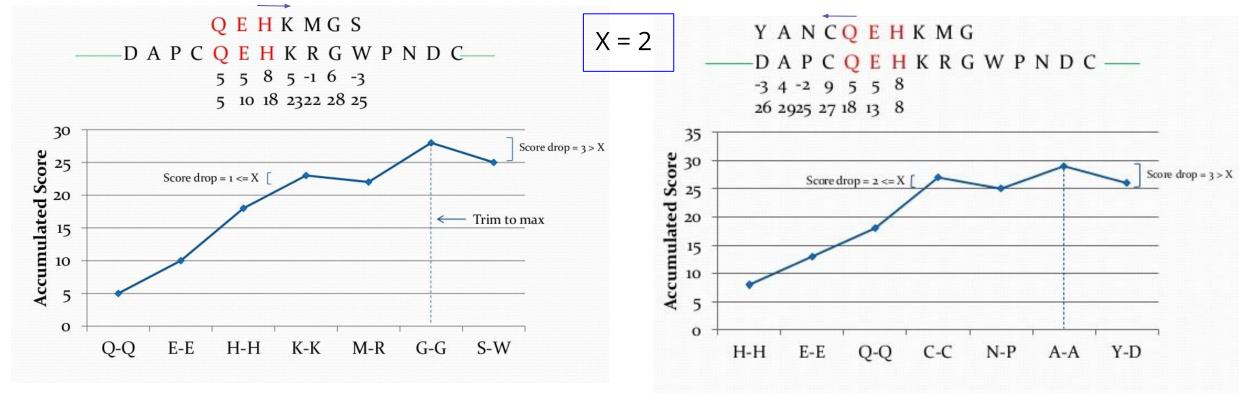
- Look for exact matches of query words with the DB words
- Masked regions are ignored



BLAST - Seed Elongation

Elongate each seed to both directions until a score drop > X is encountered

Query: YANC**LEH**KMGS



BLAST - Scoring the Alignment

Calculate total alignment score



Discard alignments with score < S

Remaining alignments are called High scoring Sequence Pairs - HSPs

BLAST - Scoring the Alignment

- Calculate alignment bit score
 - Independent of query length
 - Independent of DB size

$$S' = \frac{\lambda S - \ln(K)}{\ln(2)}$$

 Calculate E-value - the number of hits with score >= s that one can expect to find in DB by chance

$$E = \frac{L \times N}{2^{S'}}$$

L - query length, N - DB length, S' - bit score

Smaller $E \rightarrow better hit$

BLAST - Scoring the Alignment

W – word size (query and DB)

T – neighborhood words score cutoff

X – allowed score drop during seed elongation

S – HSP score cutoff

What would happen if we **increase** T?

Can We Use Blast in NGS?

Blastn - ~100 reads / sec

Human genome - ~ 3Gb

Assume 100bp reads

How long to map x10 data to the human genome?

Data required:

3 Gb x 10 = 30 Gb

Reads required:

30 Gb / 100 = 300 M reads

Time to map:

300 M reads / (100 reads/sec) = 3M sec = \sim 35 days

Scale and Speed

We need to map millions to hundreds of millions of reads

Can we use Blast?

Blastn - ~100 reads / sec

Human genome - ~ 3Gb

Assume 100bp reads

How long to map x10 data to the human genome?

Hint: how many reads do we need?

BWA - Burrows-Wheeler Aligner

Specifically designed for mapping of short reads

Maps ~2,200 reads / sec (one CPU)

Allows parallel computing

Contains three algorithms - the most useful is BWA-MEM

BWA - Limitations

Only works for nucleotides (usually DNA, not RNA)

Less effective when:

- Queries are very long
- Reads are highly diverged from the reference
- Reads contain lots of sequencing errors

Usually offers a good accuracy-speed balance

BWA - Algorithm Overview

Step 1: Index the reference genome

Step 2: Search for reads

Indexing is based on the Burrows-Wheeler's transformation

Index allows easy searching:

- Quick
- Memory efficient

BWA - The Burrows Wheeler Transform

abracadabra\$

Rotations

abracadabra\$
\$abracadabra
a\$abracadabra\$abracada
abra\$abracada
dabra\$abraca
adabra\$abrac
cadabra\$abra
acadabra\$abr
racadabra\$abr
racadabra\$ab

Sort

\$abracadabra
a\$abracadabra
abra\$abracad
abracadabra\$abra
adabra\$abracada
bracadabra\$a
bracadabra\$a
cadabra\$abra
dabra\$abraca
ra\$abracadab
racadabra\$abraca

BWT(abracadabra\$) = ard\$rcaaaabb

BWA - The Burrows Wheeler Transform

BWT is reversible - we can get back from BWT(G) to G

BWT(G) tends to cluster the same characters together - easy to compress

BWT(abracadabra\$) = ard\$rcaaaabb

Using some additional data structures, BWT(G) can be searched efficiently

BWA - The Burrows Wheeler Transform

1. Create index of reference genome:

Input: reference in fasta format

\$ bwa index genome.fasta

2. Map reads to reference:

Input: reads file or pair (for PE data) in fastq format

\$ bwa mem genome.fasta reads_R1.fq reads_R2.fq -o aln.sam