	SHRADDHA BAHUGUNA	
	CST-SPL1 30	Date.———
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	DESIGN AND ANALYSIS OF ALGORI	THMS
	DESIGNATION NAMED IN THE PROPERTY OF THE PROPE	.,,,,,
	TUTORIAL-2	
0:	1 What is the Time Complexity of below a	da 7
7	so the time complexity of Bellion a	<u> </u>
	void fun (int n)	
	int j=1, i=0; while (i <n)< td=""><td></td></n)<>	
	while (i< n)	
	î=i+j*,	
	j++,,	
	7	
	D 0	
and the second s	J D	
	2 3	
	3 6	
	4 10	
	5 15	
	T. T. T.	<u> </u>
	6 = D+1+3+6+10+15+Tk	(2)
	also S = 0+1+3+6+Tk-1+Tk	
	0=1+2+3+4+k-Tk	
	T. = 1+2+3+4 + k	
	$T_{k} = J k (k+1)$	
	2	Moraddhe
		Moradon

for k iteration

1+2+3+ ... k <n

 $\frac{2}{\sqrt{k^2 + k}} < \sqrt{n}$ 

k = 0(m)

$$T(n) = O(\sqrt{n})$$

Q2 Wink recurence relation for recursive function that prints Libonacci series. Solve the recurerence relation to get time complexity of the program what will be the space complexity of this program and why?

0 1 1 2 3 5 ... n

int fib(int n)

if (n <=1)

return n, 1/0(1)return fib(n-1) + fib(n-2) - T(n-1) + T(n-2)

T(n) = T(n-1) + T(n-2) + 1

Sheraddha

$$fib(4)$$
 $fib(3)$ 
 $fib(2)$ 
 $fib(2)$ 
 $fib(1)$ 
 $fib(0)$ 
 $fib(0)$ 
 $fib(0)$ 
 $fib(0)$ 
 $fib(0)$ 
 $fib(0)$ 
 $fib(0)$ 

$$T(n) = 1 + 2 + 4 + 8 + \dots + n$$

$$= \underbrace{1(2^{n+1}-1)}_{1}$$

$$T(n) = 2^n \cdot 2^{-1}$$

$$T(n) = O(2^m)$$

Space complexity -> 0(1)

As recursive implementation doesn't store any values and calculates every value from scratch so, as complexity of I call is o(1)

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Q3 Program which have complexity.

1) (nlog n)

for (i=1; i<=n; i=i\*2) // log n
for (==1; j<=n; j++) // n

unt s=1;

O(nlog n)

3) log(log n)

for (int 1=2; 1<n; ≈i=paw(1,2)

cont < "Hi";

Dwaddho-

$$\frac{4}{3} T(n) = T(n/4) + T(n/2) + cn^2$$

by neglecting lower order term 
$$T(n|4)$$

$$T(n) = T(n|2) + cn^2$$

$$a = 1$$
,  $b = 2$ 

$$c = \log 1$$

$$= 0^{2}$$

$$n^{c} = n^{o} = 1 < cn^{2}$$

$$T(n) = \Theta(n^2)$$

for 
$$i=1$$

$$j=1+2+3+4+5+6+...$$
for  $i=2$ 

$$j=1+3+3+4+5+6+...$$
for  $i=3$ 

$$j=1+3+3+4+5+6+...$$

$$T(n) = n + n + n + n + \dots 1$$

$$2 \quad 3 \quad 4$$

$$n \left(1 + 1 + 1 + 1 + \dots 1\right)$$

$$= n \int \frac{1}{x}$$

Shuddhe

6 Time Complexity

where k is constant

first iteration 
$$i=2$$
second  $i=2^k$ 
third  $i=(2^k)^k=2^{k^2}$ 

with  $i=(2^k)^{\ell}=n$ 

$$n = 2^{k^2} = n$$

applying log -> log n = log ki = ki

applying log again

$$i = \log \log (n)$$

 $T(n) = \log \log n$ 

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7 Given algo dévides averay in 99% and 1% paint

$$T(n) = t(n-1) + o(1)$$

$$\begin{array}{c|c}
 & n \\
 & 1 \\
 & 1
\end{array}$$

$$n = \left[T(n-1) + T(n-2) + \dots T(1) + O(1)\right] \times n$$

$$T(n) = O(n^2)$$

howest height = 2 hoighest height = n

$$\frac{1}{n} \left[ \frac{diff}{diff} = n-2 \right] \quad (n > 1)$$

The given also provides lineau result

Smaddha

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8	Avrange the following in increasing order of growth rate.
a)	Averange the following in increasing order of growth rate. $n, n!$ , log $n$ , log log $m$ , root $(n)$ , log $(n!)$ , $nlog n$ , $log^2 n$ , $2^n$ , $2^n$ , $4^n$ , $100$
	$\frac{100 < \log\log n < \log n < (\log n)^2 < \ln < n < n\log n < \log(n)}{< n^2 < 2^n < 2^n < 4^n < 2^{2^n}}$
b)	$2(2^n)$ , $4n$ , $2n$ , $1$ , $\log n$ , $\log \log n$ ), $\log n$ , $\log 2n$ , $2\log n$ , $n$ , $\log (n!)$ , $n!$ , $n^2$ , $n\log(n)$
	$1 < \log \log n < (\log n) < (\log 2n) < 2\log n < n < n \log n < 2n < 2\log n < n < n < n < n < n < n < n < n < n $
د)	8 <sup>2</sup> n, log n, nlog n, nlog n, log n, nl, log n, 1=96,8 n <sup>2</sup> , 7n <sup>3</sup> ,5
	$96 < \log n < \log 2n < 5n < n \log n < n \log n < \log(n) < 8n^{2} < 7n^{3} < n   < 8^{2}   < 7n^{3} < n   < 8^{2}  $
	Munaddha

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