

JAVA & DSA Series

# Stacks

Part – 2

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# Today's checklist

- Interview questions related to stacks

# Q. Balanced brackets

Check whether a given bracket sequence is balanced or not

String str = "()"(( )) True

str = " (( )()(" False

str = " ((() ))" True

str = " )(( ))" False

str = " (( ))( )" False

str = "(()())" True

st  
char



### Rules

- 1) Opening  $\rightarrow$  push
- 2) Closing
  - a) st top  $\rightarrow$  C  $\rightarrow$  pop

End case:

If the st is empty,  
they are balanced  $\rightarrow$  true

str = " ( ) () ( )"

i

st



### Rules

1) Opening  $\rightarrow$  push

2) Closing

a) st top  $\rightarrow$  C  $\rightarrow$  pop

if ( st.size() != 0 )  
return false

str = " ((()())"

st



### Rules

- 1) Opening  $\rightarrow$  push
- 2) Closing
  - a) st top  $\rightarrow$  C  $\rightarrow$  pop
  - b) if st is empty -  
return false

Practice → Follow

Homework

Find the minimum number of brackets that we need to remove to make the given bracket sequence balanced.

str = " (( ) ) (" False

str = " > ( ) ( ) " False

str = " ( ( ) ( ) ) " False

# Practice – LeetCode (LeetCode - 20)

Check whether a given bracket sequence is valid or not

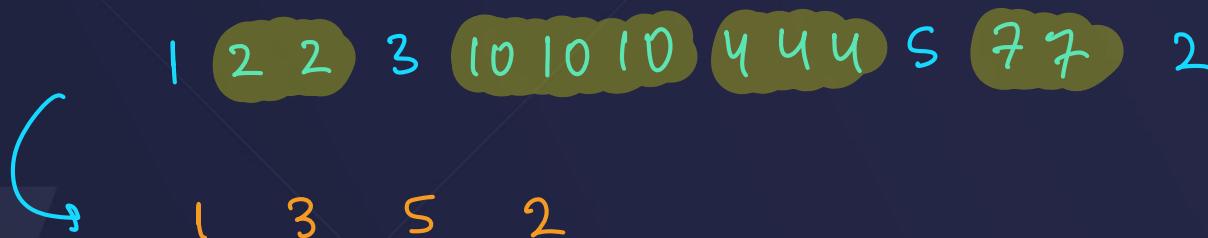
( [ ] ) { }      Valid

( [ ) ]      Invalid

push → ' ( ' || ' { ' || ' [ '

# Q. Remove consecutive subsequences

Given a sequence of numbers. Remove all the consecutive subsequences of length greater than or equal to 2 that contains the same element.



Arr | 1 2 2 3 10 10 10 4 4 4 5 7 7 2



0	1	2	3
1	3	5	2

```

if(st.size() == i) push
if(st.peek() != arr[i]) push
if(st.peek() == arr[i]){
    if(arr[i] != arr[i+1])
        st.pop();
    if(arr[i] == arr[i+1])
        do nothing
}
  
```

int[] res = new int[st.size()];

```

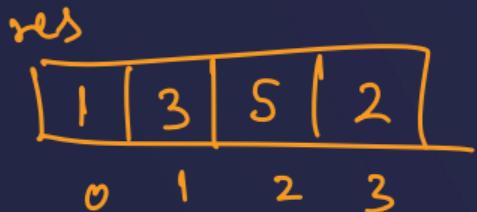
public static void main(String[] args) {
    int[] arr = {1,2,2,3,10,10,10,4,4,4,5,7,7,2};
    int[] res = remove(arr);
    for(int i=0;i<res.length;i++){
        System.out.print(res[i]+" ");
    }
}

public static int[] remove(int[] arr){
    int n = arr.length;
    Stack<Integer> st = new Stack<>();
    for(int i=0;i<n;i++){
        if(st.size()==0 || st.peek()!=arr[i])
            st.push(arr[i]);
        else if(arr[i]==st.peek()){
            if(i==n-1 || arr[i]!=arr[i+1]) st.pop();
        }
    }
    int[] res = new int[st.size()];
    int m = res.length;
    for(int i=m-1;i>=0;i--){
        arr[i] = st.pop();
    }
    return res;
}

```

arr

1 2 2 3 10 10 10 4 4 4 5 7 7 2  
i h-1



st

m=4

# \* Q. Next greater element → (+ve)

input	1	3	2	1	8	6	3	4
res	3	8	8	8	-1	-1	4	-1

Without Extra Space → MI → T.C. →  $O(n^2)$     S.C. →  $O(1)$

$O(n) \rightarrow$  Solve → using stacks

# Q. Next greater element

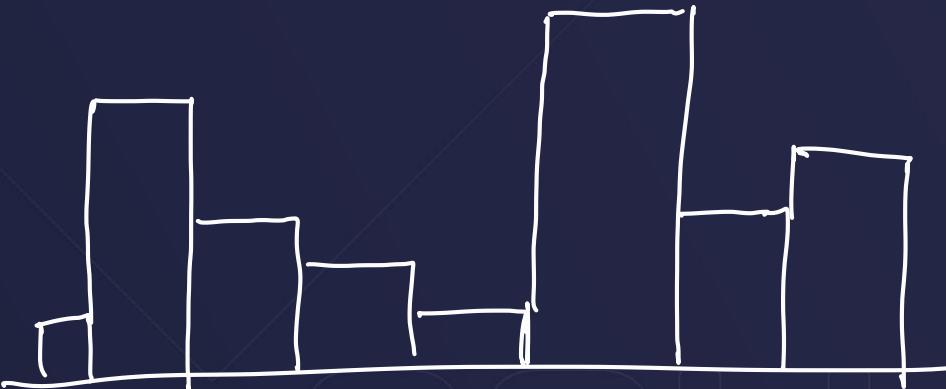
i	
input	1 5 3 2 1 6 3 4
res	5 6 6 6 6 -1 4 -1

```
for(int i=n-2; i>=0; i--){  
    while ( st.peek() < arr[i] && st.size()>0)  
        st.pop();  
    if (st.size() == 0) res[i] = -1;  
    else res[i] = st.peek();  
    st.push(arr[i]);  
}
```



# Q. Next greater element

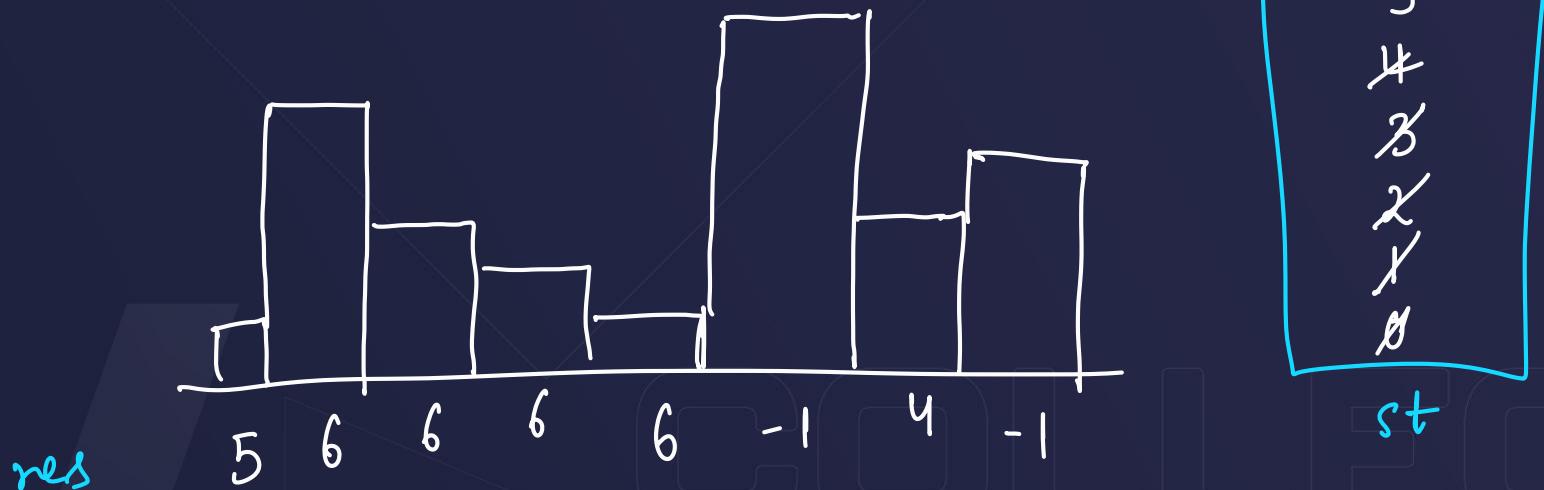
input 1 5 3 2 1 6 3 4



pop , and mark , push

# Q. Next greater element

0 1 2 3 4 5 6 7  
input 1 5 3 2 1 6 3 4

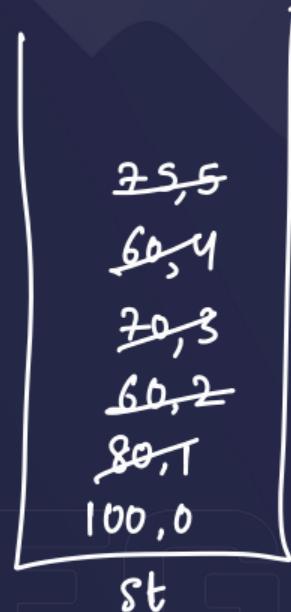


2<sup>nd</sup> approach

# Practice

GFG → Stock span problem

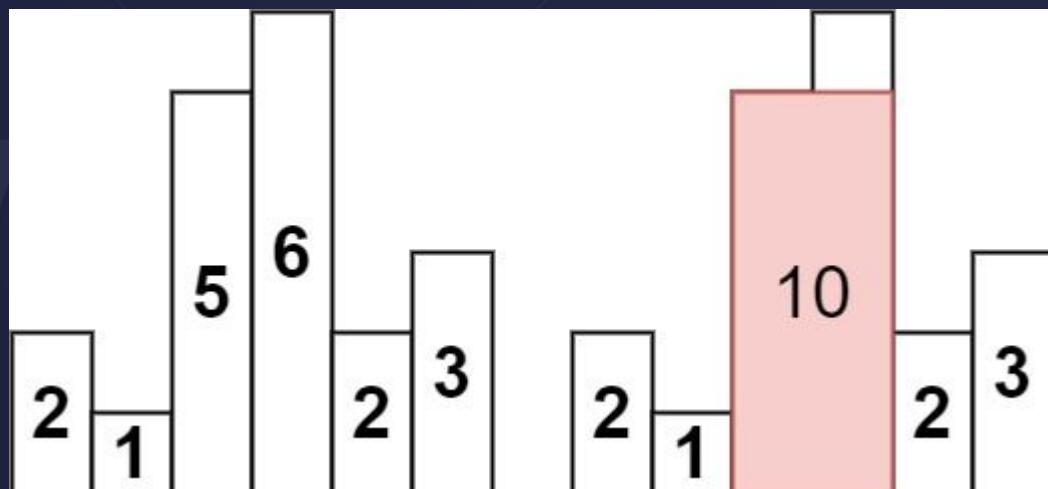
Given a series of N daily price quotes for a stock, we need to calculate the span of the stock's price for all N days. The span of the stock's price in one day is the maximum number of consecutive days (starting from that day and going backward) for which the stock price was less than or equal to the price of that day.



# Q. Largest rectangle in Histogram

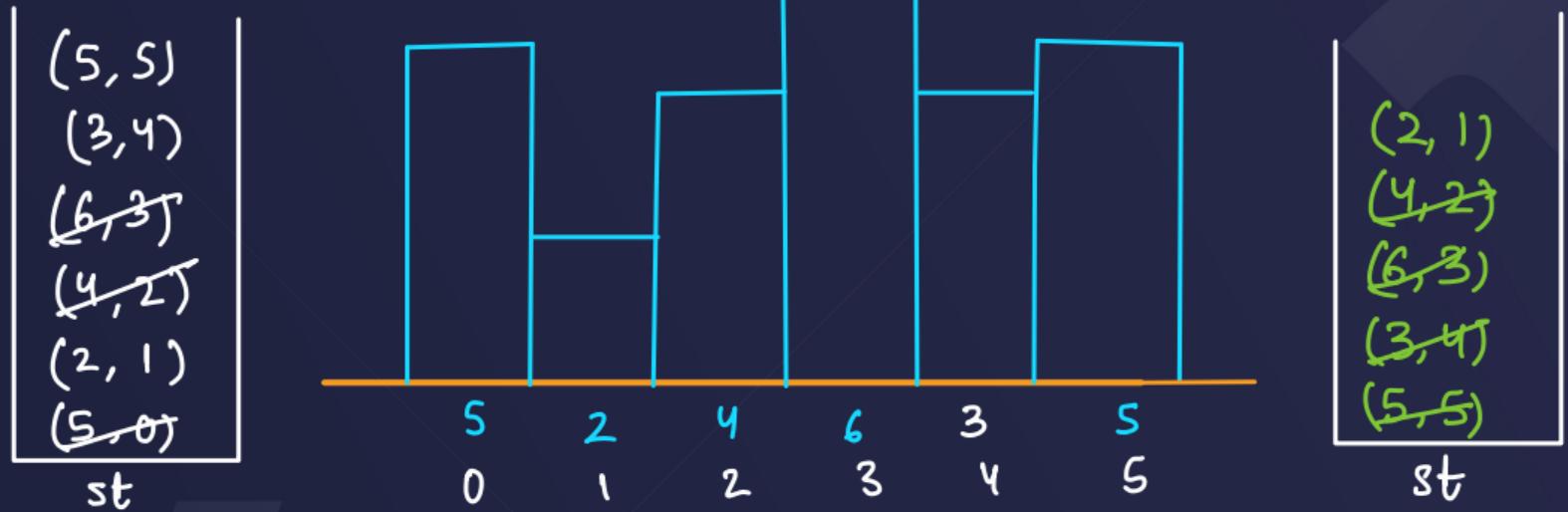
 $\Theta(n)$ 

Given an array of integer heights representing the histogram's bar height where the width of each bar is 1, return the area of the largest rectangle in the histogram.





PSE & NSE



nse[1]

1 6 4 4 6 6

pse[1]

-1 -1 1 2 1 4

$$(heights[4])^*(nse[4] - pse[4] - 1)$$

# Next Smaller Element

h 2, 1, 5, 6, 2, 3  
nse 6 6

```
int n = heights.length;
Stack<Integer> st = new Stack<>();
int[] nse = new int[n];
int[] pse = new int[n];
// calculate nse[]
st.push(n-1); // index
nse[n-1] = n;
for(int i=n-2;i>=0;i--){
    while(st.size()>0 && heights[st.peek()]>heights[i]){
        st.pop();
    }
    if(st.size()==0) nse[i] = n;
    else nse[i] = st.peek();
}
```

n-1

```
int n = heights.length;
Stack<Integer> st = new Stack<>();
int[] nse = new int[n];
int[] pse = new int[n];
// calculate nse[]
st.push(n-1); // index
nse[n-1] = n;
for(int i=n-2;i>=0;i--){
    while(st.size()>0 && heights[st.peek()]>heights[i]){
        st.pop();
    }
    if(st.size()==0) nse[i] = n;
    else nse[i] = st.peek();
}
```

heights      1, 1  
              0 1  
  
nse          2



# Q. Min Stack [Special stack implementation]

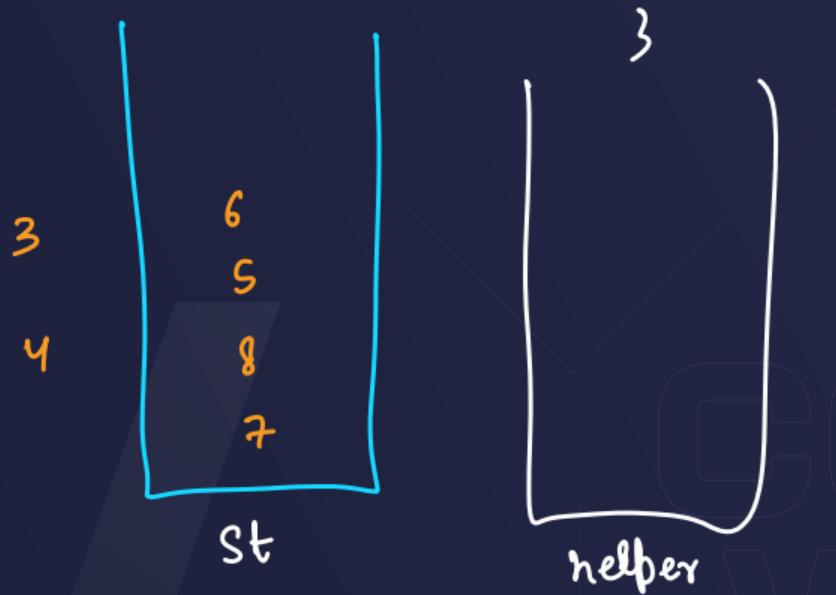
st. peek() → O(1)

st. push() → O(1)

st. pop() → O(1)

st. getMin() → O(1)

```
int getmin() {  
    stack helper;  
    int min = max;
```



Approach - 2 :     $\text{getmin}() \rightarrow O(1)$

S.C.  $\rightarrow O(n)$

push, pop, peek,  
 $\text{getmin} \rightarrow O(1)$

7 8 5 6 3 4

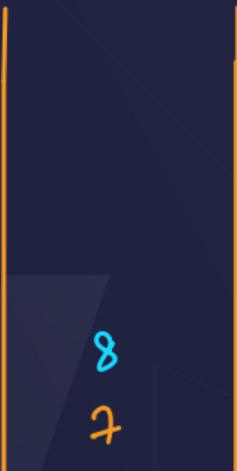


Approach - 3 : without extra stack

variable → one

```
if(val < min)
    st.push(val - min)
```

push 4  
push 3\*  
push 6  
push 5\*  
push 8  
push 7



Val  
min = 7 8 3 9 2  
min - oldmin = st.peek()  
oldmin = min - st.peek()  
min = oldmin

```
getmin()
    return min;
```

```
pop()
    if(st.peek() > min)
        st.pop()
    if(st.peek() < min)
        int old = min - st.peek();
        min = old;
        st.pop()
```

```
void push(int val){  
    if(st.size() == 0){  
        st.push(val);  
        min = val;  
    }  
    else if(val > min)  
        st.push(val);  
    else if(val < min){  
        st.push(val - min);  
        min = val;  
    }  
}
```



$$\min = 7 \neq 3$$

$$2^{\text{nd}} \text{ val} - \min$$

void pop() {

if (st.size == 0) return;

if (st.peek() > min)  
st.pop();

if (st.peek() < min) { // restore

int old = min - st.peek();

min = old;  
st.pop();

min = 7 & 5

6  
-2  
8  
7

st

old = 2 \* min - st.peek();

3

["MinStack","push","push","push","pu  
h","getMin","pop","getMin","pop","get  
Min","pop","getMin"]

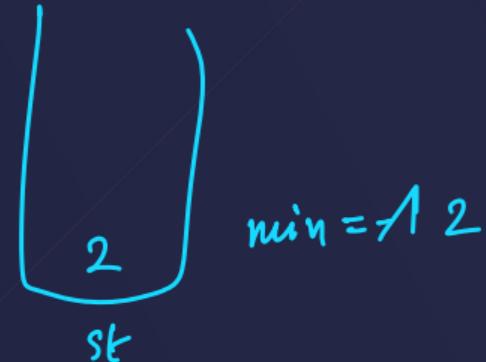
[[],[~~2~~],[0],[3],[0],[],[],[],[],[],  
[],[]]

## Output

[null,null,null,null,null,0,null,0,nu  
ll,2,null,2]

## Expected

[null,null,null,null,null,0,null,0,nu  
ll,0,null,2]



-2

8

7

st

min = 7 & 8 & 5

int getmin(){

| if(st.size() == 0) return -1;

| return min;

3

int top(){  
| if(st.size() == 0) return -1;  
| if(st.peek() > min )  
| | return st.peek();  
| if(st.peek() < min ){ //farzi  
| | return min ;  
| }  
3

COLLEGE  
WALLAH

## Input

["MinStack","push","push","push","getMin","pop","top","getMin"]

[[], [-2], [0], [-3], [], [], [], []]

## Output

[null,null,null,null,-3,null,0,-3]

## Expected

[null,null,null,null,-3,null,0,-2]



val  
-3

$$\min = 1 - 2 - 3$$

$$2^* \text{val} - \min$$

$$2^*(-3) - (-2)$$
$$-6 + 2 = -4$$

-2    0    - 3



$$\min = -2 \quad (-3)$$

$$2^*(-3) - (-2) = -4$$

$$st.\text{peek}() = 2^*\text{newMin} - \text{old min}$$

$$\text{Oldmin} = 2^*\min - st.\text{peek}()$$

$$\text{val} < \min_{-3, -2} \rightarrow \boxed{\text{val} - \min < 0}$$

$$\text{val} - \min \leq \text{val}$$

$$2^* \text{val} - \min < \text{val}$$

↓

$$\text{val} + (\text{val} - \min)$$

$$\text{val} + (-\text{val}) < \text{val}$$