Distance Vector Algorithm – Problems, Solutions and a Standard

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Recap

- Nodes exchange with their neighbors their current routing table information (destination, estimated cost)

- On receipt of a message, nodes update cost to destination based on Bellman-ford equation

- Messages sent periodically as well as when table changes

Reference Node C

Example

D	C	H
A	5	A
B	3	B
D	4	D

Routing Table of C

(1)

To	A
A	0
B	1
C	5

Message from A
C to A: C = 5

D	C	H
A	5	A
B	3	B
D	4	D

Routing Table of C

D	C	H
A	5	A
B	3	B
D	4	D

Routing Table of C

(2)

To	B
A	1
B	0
C	3
E	9

Message from B
C to B: C = 3

D	C	H
A	4	B
B	3	B
D	4	D
E	12	B

Routing Table of C

D	C	H
A	4	B
B	3	B
D	4	D
E	12	B
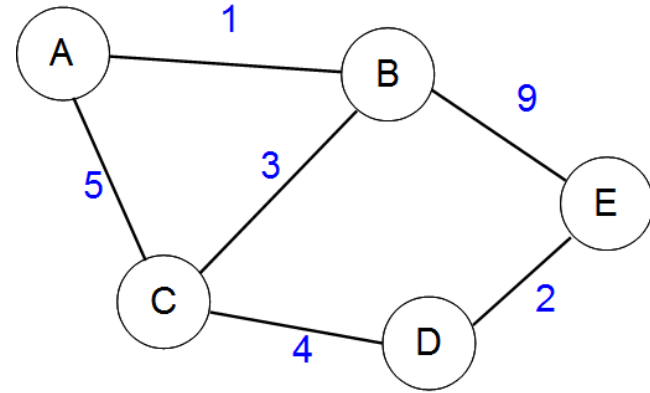
Routing Table of C

(3)

To	D
C	4
D	0
E	2

Message from D
C to D: C = 4

D	C	H
A	4	B
B	3	B
D	4	D
E	6	D

Routing Table of C

Every path has its puddle!

Counting to Infinity

Distance to Node D

Mesg.	A	B	C
	2,B	∞,-	2,B
B→A	∞,-	∞,-	2,B
C→A	3,C	∞,-	2,B
B→C	3,C	∞,-	∞,-
A→B	3,C	4,A	∞,-
C→A	∞,-	4,A	∞,-
B→C	∞,-	4,A	5,B
A→B	∞,-	∞,-	5,B
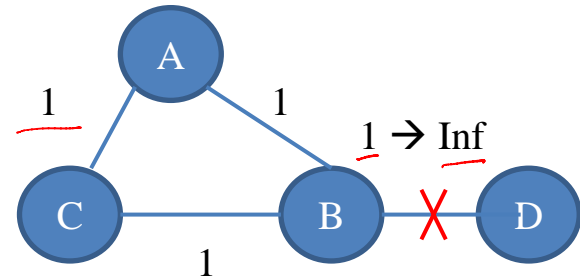


State maintained by nodes A,B and C

Partial Solutions

- ## Make infinity small

  → hop count

  - Use for example 16 to represent infinity (assumes max no of hops under 16)

  - Bounds time it takes to count to infinity

- ## Split horizon

  - Don't send routes learnt from a neighbor back to it

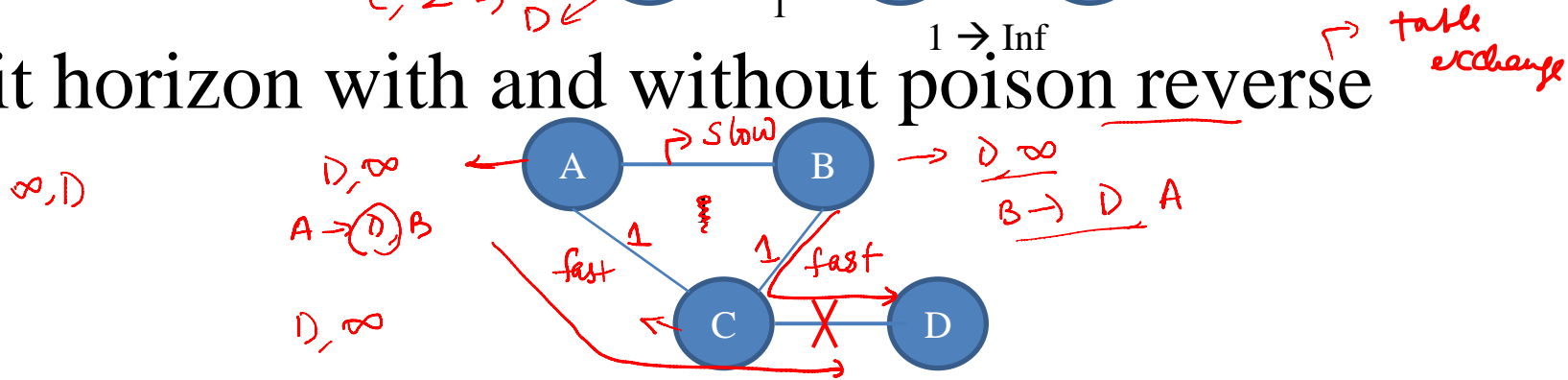A,

C, B, 2

2

A —— B ✕ C

1

1 → Inf
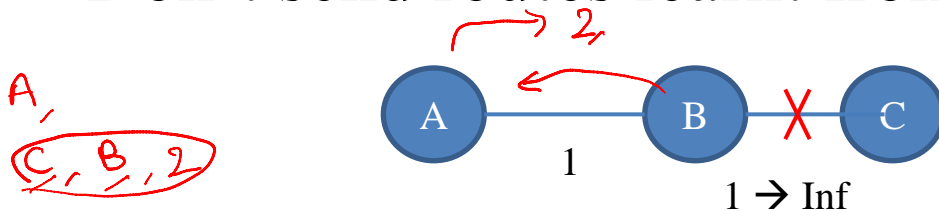
Partial Solutions

- Split horizon with poison reverse

  - Send routes learnt from a neighbor back to it but with infinite cost
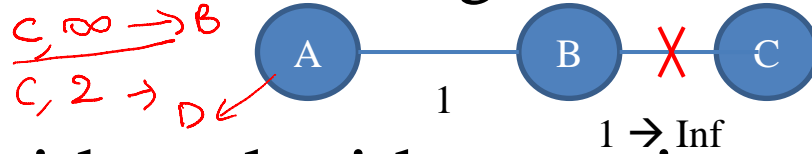

- Split horizon with and without poison reverse

- Both don't work for loops with more than 2 nodes

Partial Solutions

- Hold-Down Timer: Wait some time before propagating link failure

    - Slows down convergence

- Path-vector routing is a variation of distance-vector

    - Each node sends to its neighbors not just the cost, but the entire path to the destination

    - Avoids the looping problem of DV but more overhead
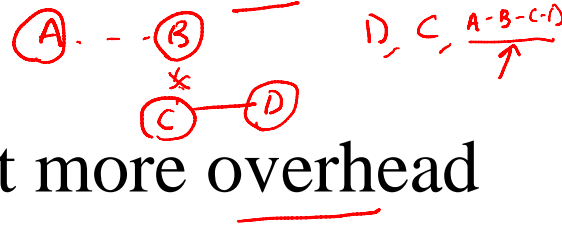
RIP

- Routing Information Protocol (RIP) is a standard that implements DV routing

- One of the oldest DV based protocol

  - Popular once, not used much due to convergence problems

    ↳ size

RIP Features

- Uses UDP and work over reserved port 520

- Period updates sent every 30 sec

- Supports multiple address families

  IP

- Cost of a link is 1 (finds minimum hop route)

- 16 represents infinity

  Split horizon, hold-down timer

- RIP can run only on very small networks

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RIPv1 Packet Format

$\rightarrow R_1 \quad C_1 \quad$ (NH1)

0   31

request/reply

Command	Version	Reserved
Family of Net 1 ⟷ IP		Reserved
Address of Net 1 (IP Address)		
Reserved		
Reserved		
Distance to Net 1		
Family of Net 2		Reserved
Address of Net 2 (IP Address)		
Reserved		
Reserved		
Distance to Net 1		

$\rightarrow$ IP Prefix (NH) Route

5×4   classful

1 to 25 sets of entries, each entry is 20 bytes

classless

Reserved fields are set to all zero

IPF_1) $\{ IPF_2$

$R \quad \underset{I}{} \quad R_1$

$R \quad 0 \quad IPF_3$

$IPF_1 \quad 0$

$IPF_2 \quad 0$

$IPF_3 \quad 1$

Summary

- Distance vector is a distributed, dynamic algorithm

- Exchanges information locally to determine routes

- Suffers from poor convergence, routing loops

- RIP is a standard that implements the DV protocol
  - Handles above problems via split horizon, hold-down timer and using a value of 16 to represent infinity

- Better approach: Link-state routing