Subarrays

Subarray → continuous part of an array
→ single elements / complete array
→ empty {} is not subarray
→ [i, j] : length = j - i + 1
↳ {i, i+1, i+2, ... j}

What are total possible subarrays ?

a[4] = 2 6 3 9
0 1 2 3

[0,0] [1,1] [2,2] [3,3]
[0,1] [1,2] [2,3]
[0,2] [1,3]
[0,3]

4 3 2 1

= 4+3+2+1 = 10

What are total subarrays for array of size N?

a[n] : a_0 a_1 a_2 . . — . . a_{n-2} a_{n-1}

0 1 2 n-2 n-1

[0,0]
[0,1]
[0,2]
.
.
.
[0, n-1]

count = n-1-0+1 = n

[1,1]
[1,2]
.
.
.
[1, n-1]

count = n-1-1+1 = n-1

[n-2, n-2]
[n-2, n-1]

2

[n-1, n-1]

1

[a,b] = b-a+1

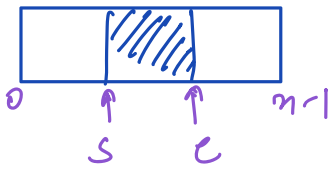
total subarrays = $n + n-1 + \dots + 2 + 1$

$= \frac{n(n+1)}{2}$

Question 1

Given a[N], s & e integers. $0 <= s <= e < n$

Print subarray from [s, e]



eg a[5] = 2 4 5 9 8 (indices 0 1 2 3 4)

s=1 e=3

ans = 4 5 9

Code

```
for(i=s; i<=e; ++i)
    print(a[i])
```

Question 2

Given N array elements, print each & every subarray.

Note: Do it without extra space.

eg a[4] = 6 8 -1 7 (indices 0 1 2 3)

[0,0] → {6}

[0,1] → {6,8}

[0,2] → {6,8,-1}

[0,3] → {6,8,-1,7}

[1,1] → {8}

[1,2] → {8,-1}

[1,3] → {8,-1,7}

[2,2] → {-1}

[2,3] → {1,7}

[3,3] → {7}

total subarrays = $\frac{N(N+1)}{2}$

$\approx O(N^2)$

To print all subarrays, I have to print $O(N^2)$ lines

worst case to print subarray = $O(N)$

Time taken to print all subarrays

$\Rightarrow O(N^2) * O(N)$

$= O(N^3)$

Code

```
def printAll(a[]) {
    n = a.length
    for(i=0; i<n; ++i) {
        for(j=i; j<n; ++j) {
            // [i,j] subarray
            print(a[i:j])
```

→

```
for(k=i; k<=j; ++k)
    print(a[k])
```

```
    }
}
```

3 3

Question 3

Given N array elements, print each subarray sum.

eg a[4] = 6 8 -1 7
0 1 2 3

[0,0] → 6
[0,1] → 6+8 = 14
[0,2] → 6+8+-1 = 13
[0,3] → 6+8+-1+7 = 20
[1,1] → 8
[1,2] → 7
[1,3] → 14
[2,2] → -1
[2,3] → 6
[3,3] → 7

print $O(N^2)$ values

```
def printSum(a[]) {
    n = a.length
    for(i=0; i<n; ++i) {
        for(j=i; j<n; ++j) {
            // [i,j] subarray
            sum=0
            for(K=i; K<=j; ++K)
                sum += a[K]
            print(sum)
    }
}
```

TC: O(N^3)

SC: O(1)

⇓ OPTIMISE using Prefix Sum

Code

```
void printSum(a[]) {
    n = a.length
    pf[n] → TODO → TC O(N), SC O(N)
    for(i=0; i<n; ++i) {
        for(j=i; j<n; ++j) {
            if(i==0) print(pf[j])
            else     print(pf[j] - pf[i-1])
    }
}
```

→ TC: O(N^2)
SC: O(1)

TC: O(N^2)
SC: O(N)

Question 4

Given a[N], print all subarray sums starting at index 3.

eg a[8] = 3 8 4 7 9 4 3 2

(indices 0 1 2 3 4 5 6 7)

sum

7

16

20

23

25

Carry forward

sum = 0 + 7

sum = 7 + 9

sum = 16 + 4

sum = 20 + 3

sum = 23 + 2

25

```
sum = 0
for (j = 3; j < n; ++j) {
    sum += a[j]
    print(sum)
}
```

Question 4 code

```
for (i = 0; i < n; ++i) {
    sum = 0
    for (j = i; j < n; ++j) {
        sum += a[j]
        print(sum)
    }
}
```

TC : O(N²)

SC : O(1)

Optimized code of Question 3

Dry run

a[4] = 6 8 -1 7

0 1 2 3

i=0

sum=0 j : [0,3]

sum = 0 + a[0] = 6

sum = 6 + a[1] = 14

sum = 14 + a[2] = 13

sum = 13 + a[3] = 20

print(6)

14

13

20

i=1

sum=0 j : [1,3]

sum = 0 + a[1] = 8

sum = 8 + a[2] = 7

8

7

⋮

i=2

i=3

BREAK : 8:05 - 8:15

Question 5

Given N array elements, return sum of all subarray sums.

eg a[4] = 6 8 -1 7

(0) (1) (2) (3)

[0,0] → 6

[0,1] → 6+8 = 14

[0,2] → 6+8+-1 = 13

[0,3] → 6+8+-1+7 = 20

[1,1] → 8

[1,2] → 7

[1,3] → 14

[2,2] → -1

[2,3] → 6

[3,3] → 7

ans = 6+14+13+20
+ 8+7+14
+ -1 + 6
+ 7

⇒ 94

print O(1) values.

Idea: for every subarray get sum & add it to total sum.

Approach 1
→ 3 nested loops
$O(N^3)$ $O(1)$

Approach 2
→ prefix sum
$O(N^2)$ $O(N)$

Approach 3
→ carry forward
$O(N^2)$ $O(1)$

Code

```
ans=0
for (i=0; i<n; ++i) {
    sum=0
    for (j=i; j<n; ++j) {
        sum += a[j]
        ans += sum
    }
}
print(ans)
```

TC : $O(N^2)$

SC : $O(1)$

eg A[6] = 3 (0) -2 (1) 4 (2) -1 (3) 2 (4) 6 (5)

In how many subarrays, index 3 is present?

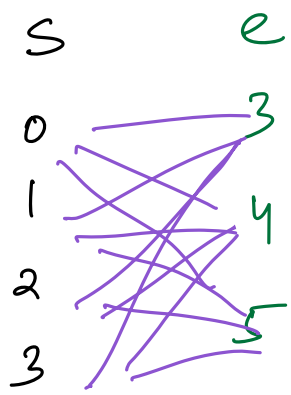
[0,3] [0,4] [0,5]

[1,3] [1,4] [1,5]

[2,3] [2,4] [2,5]

[3,3] [3,4] [3,5]

total = 12

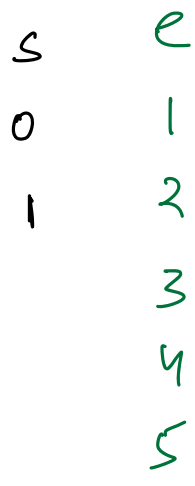


⇒ 4 × 3 = 12

In how many subarrays, index 1 is present?

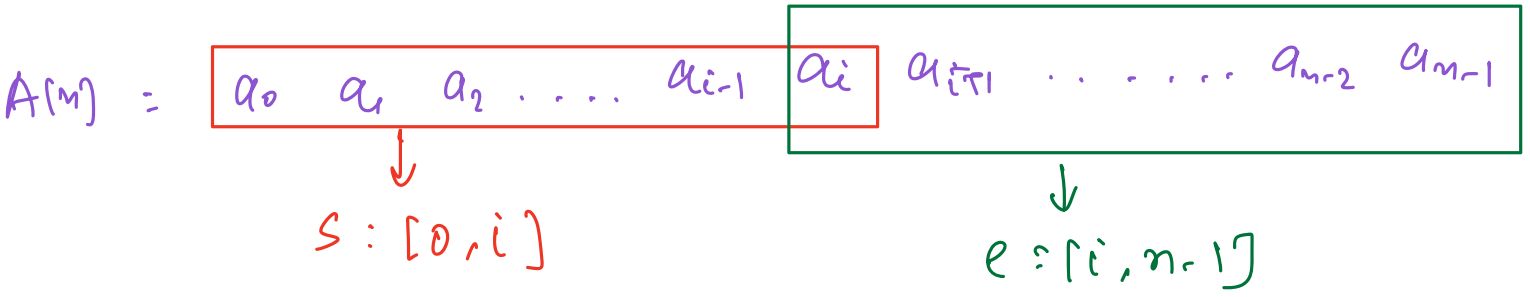
[0,1] [0,2] [0,3] [0,4] [0,5]

[1,1] [1,2] [1,3] [1,4] [1,5]



⇒ 2 × 5 = 10

Generalize: Given N elements, find number of subarrays where ith index is present?



count = i - 0 + 1

= i + 1

count = n - 1 - i + 1

= n - i

total = (i+1) * (n-i)

	0	1	2	3
a[4] =	6	8	-1	7
i+1 :	1	2	3	4
n-i :	4	3	2	1
(i+1)*(n-i) :	4	6	6	4

individual contribution = 24 + 48 + -6 + 28 = 94

Dry run

N=3

	0	1	2
a[3] =	4	3	7
i+1	1	2	3
n-i	3	2	1
(i+1)*(n-i)	3	4	3

sum = 12 + 12 + 21

= 45

subarrays

[0,0] : {4}

[0,1] : {4,3}

[0,2] : {4,3,7}

[1,1] : {3}

[1,2] : {3,7}

[2,2] : {7}

Code

```
ans=0
for(i=0; i<n; ++i) {
    total = (i+1)*(n-i)
    ans += (total * a[i])
}
print(ans)
```

TC: O(N)

SC: O(1)