

Recursion - 1

Content

- Recursion ?
- How to write recursive code / tracing
- TC/SC of recursive codes → next class

Why recursion ?

- Merge Sort / Quick Sort
- Binary tree / BST / segment trees / Tries / Balanced BST
- Dynamic Programming (DP)
- Backtracking
- Graphs

Recursion: function calling itself

↳ solving a problem, using smaller instance of the same problem.

⇓
sub-problem

eg $Sum(N) = 1 + 2 + 3 + \dots + N-1 + N$

↓
 $Sum(N-1)$

$$Sum(N) = \underbrace{Sum(N-1)}_{\text{sub-problem}} + N$$

$$\begin{aligned}
 Sum(4) &= Sum(3) + 4 &= 10 \\
 &\downarrow \\
 &Sum(2) + 3 \\
 &\downarrow \\
 &Sum(1) + 2 \\
 &\downarrow \\
 &1
 \end{aligned}$$

How to write recursive codes?

Assumption: fix what your function should do

Main logic: Solving assumption using sub-problem

Base condition: Inputs for which you want to stop recursion.

Question 1 : sum of N natural no.s) constraint: $N \geq 1$

```
int sum(N) { Ass.: Given N, calculate & return  
              sum of first N natural no.  
    if (N==1)  
        return 1  
    return sum(N-1) + N  
}
```

Question 2 :

$$\text{fact}(3) = 3 \times 2 \times 1 = 6$$

$$\text{fact}(5) = 24 \times 5 = 120$$

$$\text{fact}(4) = 4 \times 3 \times 2 \times 1 = 24$$

constraint: $N \geq 1$

```
int fact(N) { Ass.: Given N,  
              calculate & return  
              N!  
    if (N==1)  
        return 1  
    return fact(N-1) * N  
}
```

$$\text{fact}(N) = 1 \times 2 \times 3 \dots \times N-1 \times N$$

↓

$$\text{fact}(N) = \text{fact}(N-1) \times N$$

Base condition is must, otherwise

↓
MLE : Stack overflow

Stack Tracing

```
int add(N, m) {  
    return N+m  
}
```

```
int mul(N, m) {  
    return N*m  
}
```

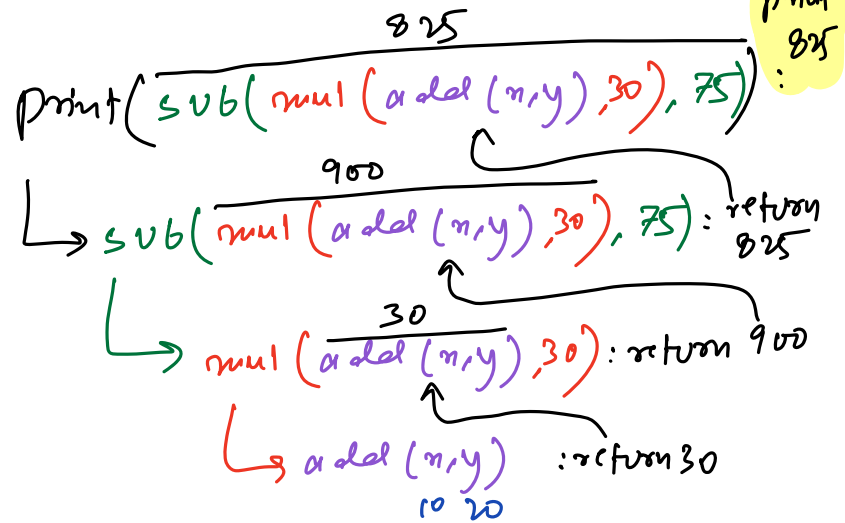
```
int sub(N, m) {  
    return N-m  
}
```

```
int main() {
```

```
    x=10, y=20
```

```
    print(sub(mul(x, y), 30), 75)
```

```
}
```

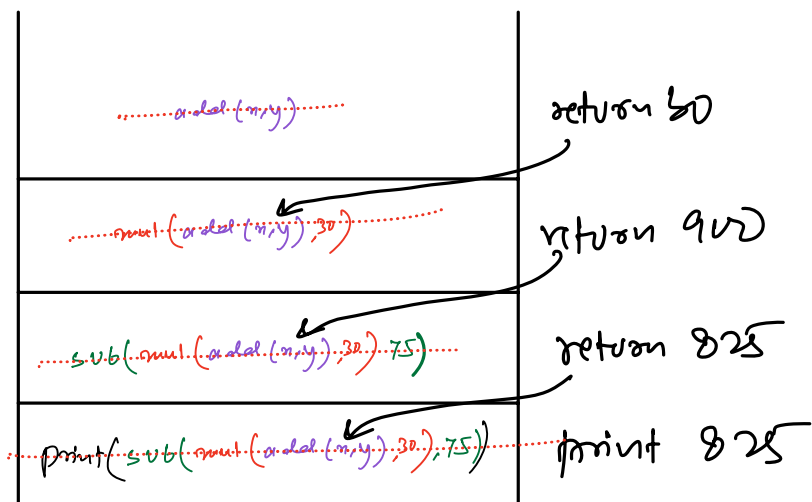


Data Structure for Stack Tracing?

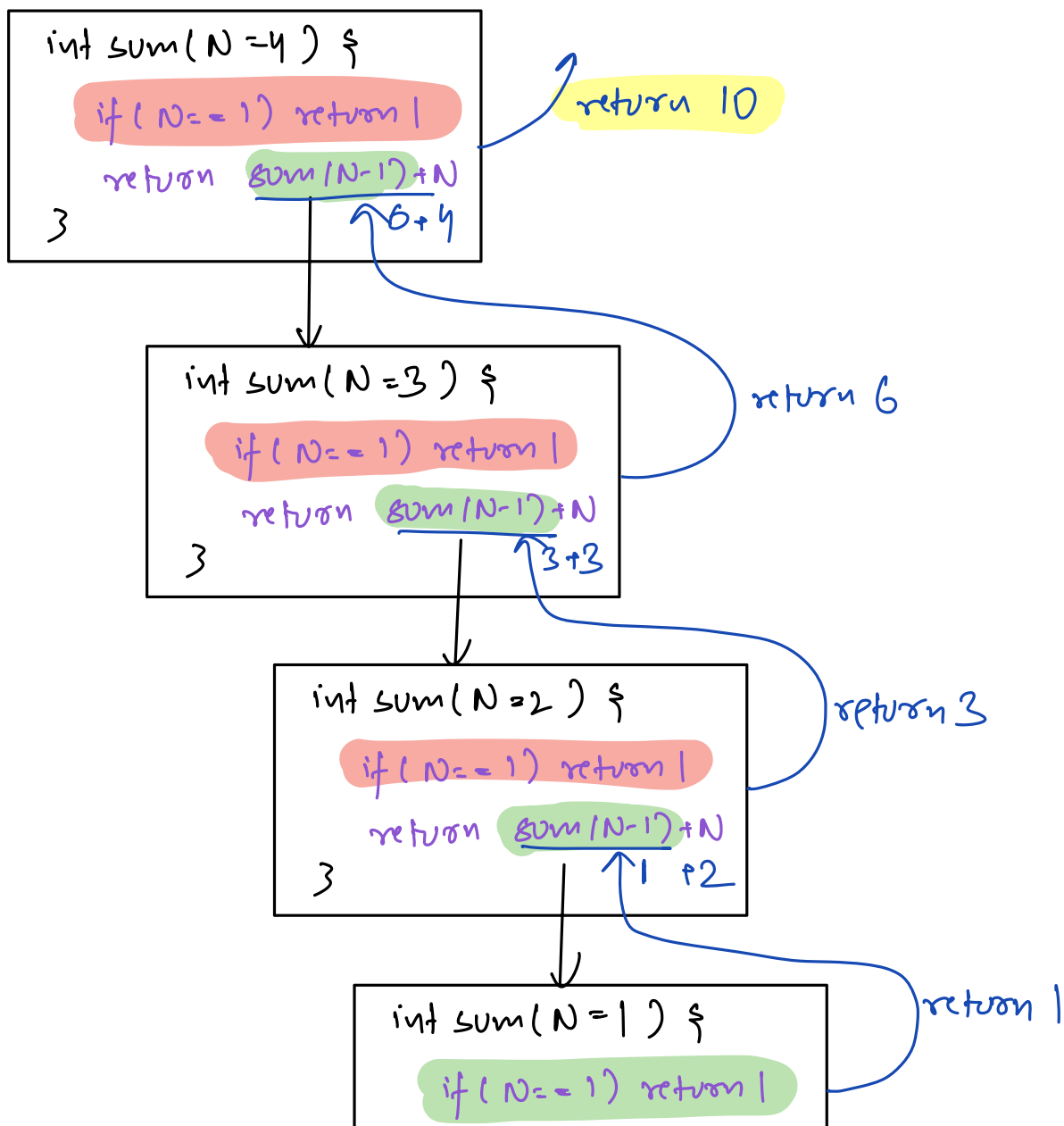
Observation 1: Whenever a function call happens, we add the function call at top.

Observation 2: When function returns, we remove it from top.

Last In First Out (LIFO)
↓
Stack



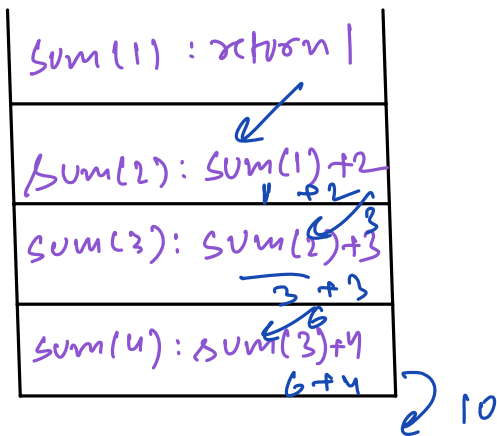
Sum(N) forcing



```

    return sum(N-1)+N
}

```



Question 3

$N \geq 0$

Input(N) :	0	1	2	3	4	5	6	7	8
fib()	0	1	1	2	3	5	8	13	21

N^{th} fibonacci no. = $(N-1)^{\text{th}}$ fibo no. + $(N-2)^{\text{th}}$ fibo no.

```

int fib(N) {
    // Acc.: calculate & return Nth fibo no.
    if(N==0) return 0;
    if(N==1) return 1;
    return fib(N-1)+fib(N-2);
}

```

→ if $(N \leq 1)$ return N

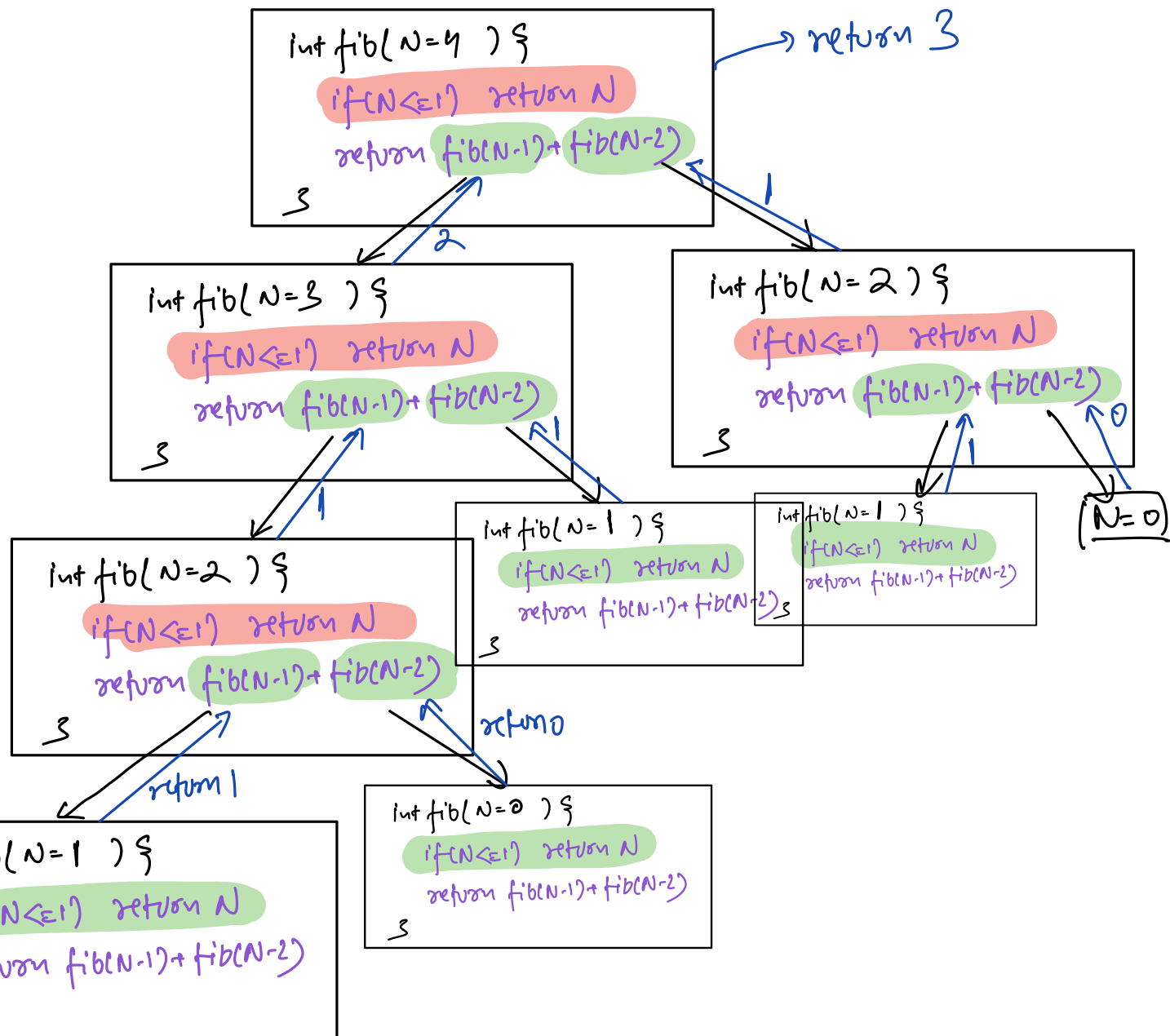
3

Note: How to properly figure out base conditions
 → for which valid input, expression is invalid
 ↳ main logic

$\text{fib}(0) = \text{fib}(-1) + \text{fib}(-2)$ ✗

$\text{fib}(1) = \text{fib}(0) + \text{fib}(-1)$ ✗

$\text{fib}(2) = \text{fib}(1) + \text{fib}(0)$ ✓✓



TODD: Try Stack Tracing with Stack

Question 4

Given N , print all no. from $1-N$ in increasing order.

$N \geq 1$

$\text{Inc}(N) : \underbrace{1, 2, 3, \dots, N-1, N}$
↓

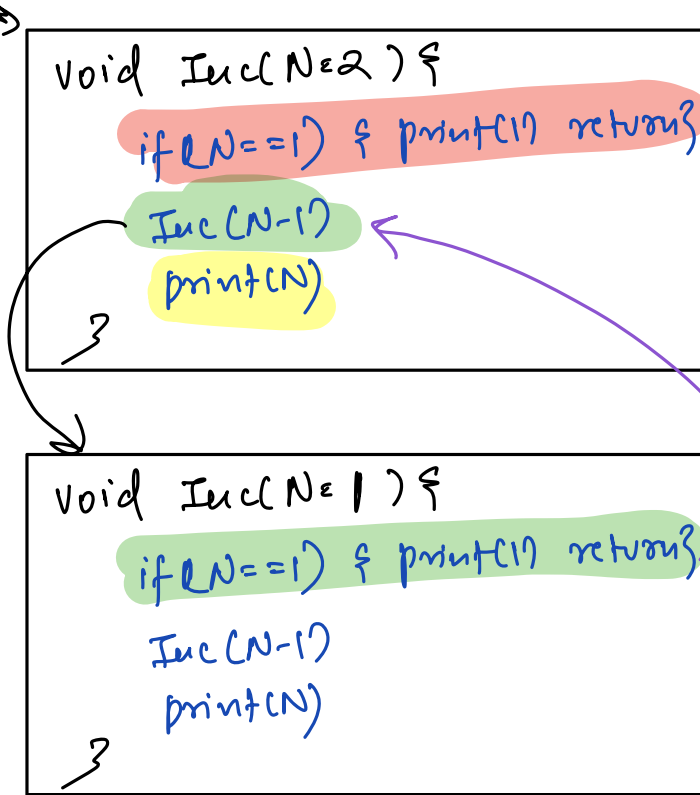
$\text{Inc}(N) : \text{Inc}(N-1) ; \text{print}(N)$

```
void Inc(N) {  
    if(N==1) {  
        print(1) return  
    }  
    Inc(N-1) → 1 2 3 ... N-1 N  
    print(N)
```

}

```
void Inc(N=4) {  
    if(N==1) { print(1) return }  
    Inc(N-1)  
    print(N)  
}
```

```
void Inc(N=3) {  
    if(N==1) { print(1) return }  
    Inc(N-1)  
    print(N)  
}
```

Output :

<code>print(1)</code>	1
<code>print(2)</code>	2
<code>print(3)</code>	3
<code>print(4)</code>	4

A downward arrow is positioned between the first and last rows of the output table.

Note :

1. Even for void return type, we can return inside function.
2. Once a function completes it will go back to its parent function.

Homework: Print in decreasing order

`Dec(N) : N N-1 3 2 1`

Question 5:

Given a substring, check if its palindrome or not?

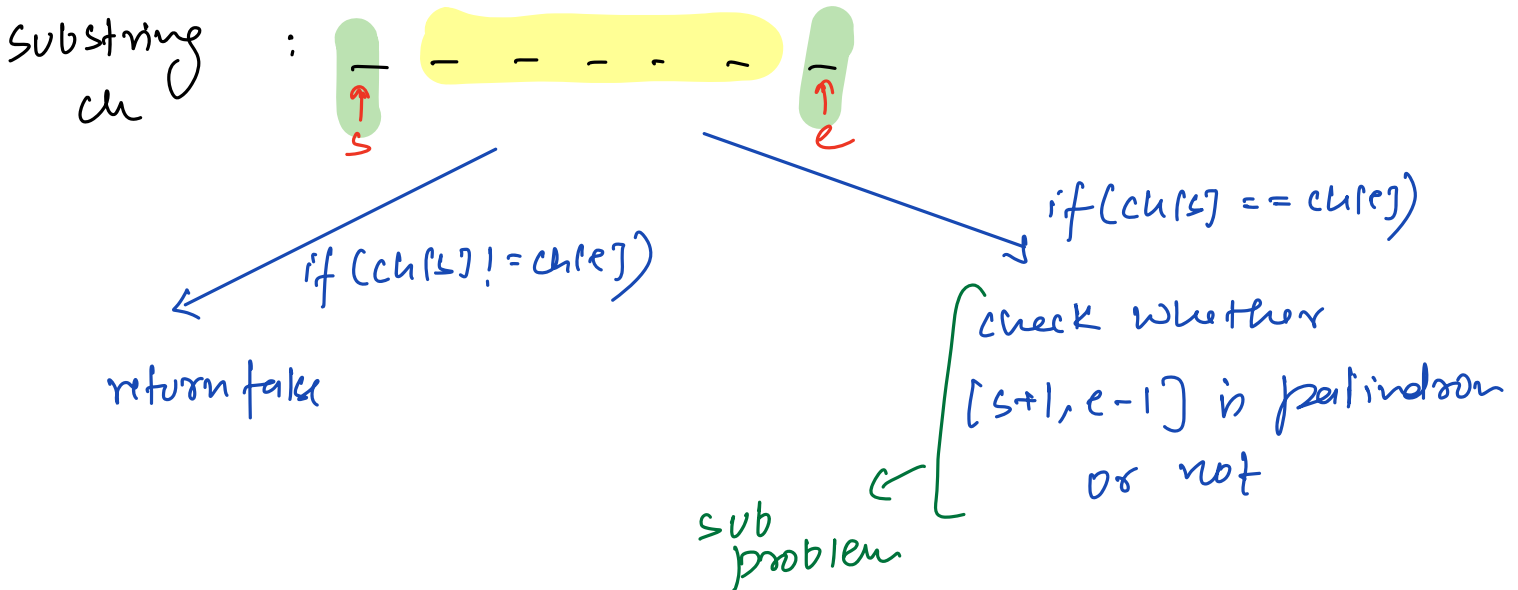
eg g o o d d a d
0 1 2 3 4 5 6

$s=4, e=6 \rightarrow \text{return true}$

$s=2, e=5 \rightarrow \text{return false}$

main
problem

\rightarrow check whether $[s, e]$ is palindrome or not?



```
bool isPalin(char ch[], s, e) {  
    if ( $s > e$ ) { return true }  
    if ( $ch[s] \neq ch[e]$ )  
        return false  
    return isPalin(ch,  $s+1$ ,  $e-1$ )  
}
```

}

Input : m a d d a m
 0 1 2 3 4 5

s=0, e=5

```
bool isPalin(char ch[], s, e) {  
    if(s > e) { return true }  
    if(ch[s] != ch[e]) m=m  
        return false  
    return isPalin(ch, s+1, e-1)  
}
```

return true

true

```
bool isPalin(char ch[], s, e) {  
    if(s > e) { return true }  
    if(ch[s] != ch[e]) a=a  
        return false  
    return isPalin(ch, s+1, e-1)  
}
```

true

```
bool isPalin(char ch[], s, e) {  
    if(s > e) { return true }  
    if(ch[s] != ch[e]) d=d  
        return false  
    return isPalin(ch, s+1, e-1)  
}
```

true

```
bool isPalin(char ch[], s, e) {  
    if(s > e) { return true }  
    if(ch[s] != ch[e])  
        return false  
    return isPalin(ch, s+1, e-1)  
}
```

Input: a n m e t n a
 0 1 2 3 4 5 6

s=0, e=6

```
bool isPalin(char ch[], s, e) {
  if(s > e) { return true; }
  if(ch[s] != ch[e]) a=a
  return false
  return isPalin(ch, s+1, e-1)
}
```

return false

false

```
bool isPalin(char ch[], s, e) {
  if(s > e) { return true; }
  if(ch[s] != ch[e]) n=n
  return false
  return isPalin(ch, s+1, e-1)
}
```

false

```
bool isPalin(char ch[], s, e) {
  if(s > e) { return true; }
  if(ch[s] != ch[e]) m!=t
  return false
  return isPalin(ch, s+1, e-1)
}
```

Doubt

$$a[i] - a[j] = k$$

$$a[j] = a[i] - k$$