

Agenda :

1. Properties of a good solution to data sync. problem.
2. Solution \rightarrow Mutex.
3. Producer / consumer Problem.
4. Semaphores.

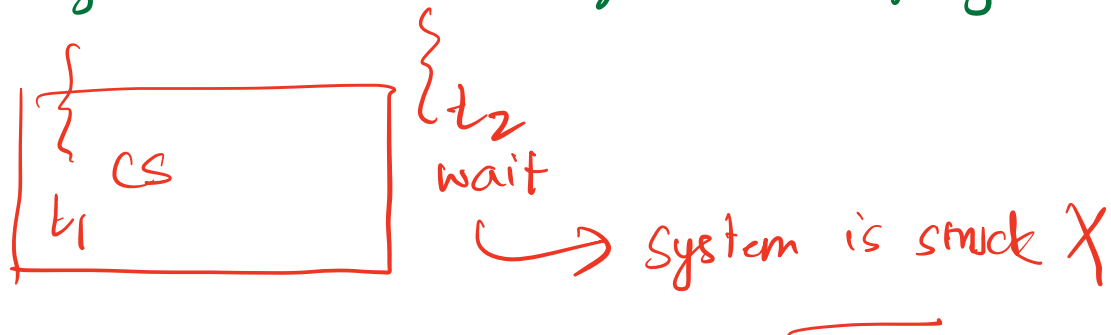
Properties of Good Solution to Data Sync Problem

1. Mutual Exclusion :

- One thread to work on the critical section of a variable at a time.
- Restrict the number of thread to 1 for a critical section of a variable.

2. Progress :

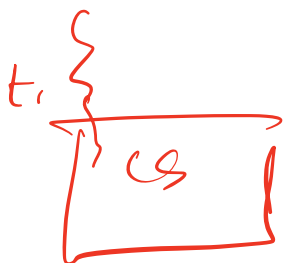
- Overall system must always make progress.



addQuestion () {

deleteQuestion () }

CS { \uparrow List (Question) q ; \downarrow } CS { \uparrow List (Question) q ; \downarrow }



t_2 has to wait

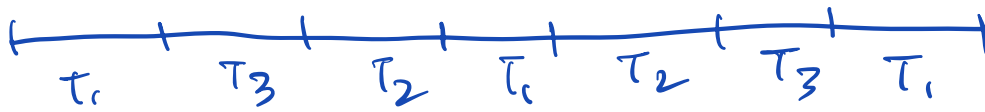
But overall system shouldn't

wait or feel like stuck.

3. Bounded Waiting -

- No thread should wait indefinitely.

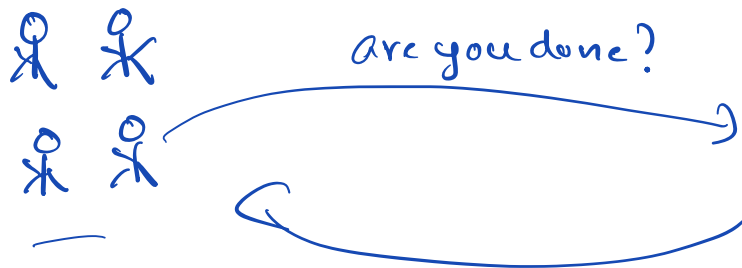
$T_1, T_2, T_3, (T_4)$



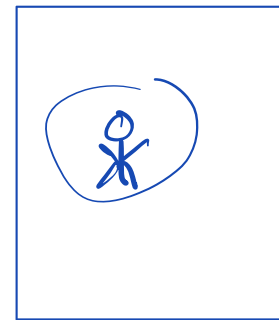
$\{ T_4 \}$
waiting
indefinitely.

- Every thread should get equal amount of CPU time

4. No Busy waiting:



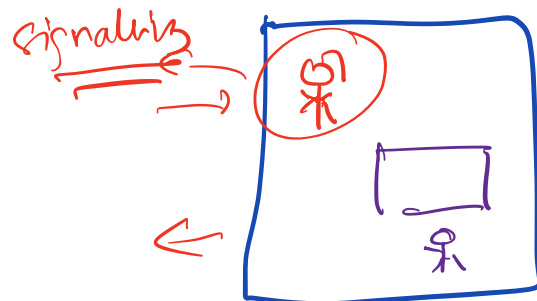
Bathroom.



Sitting room



Doctor's chamber.



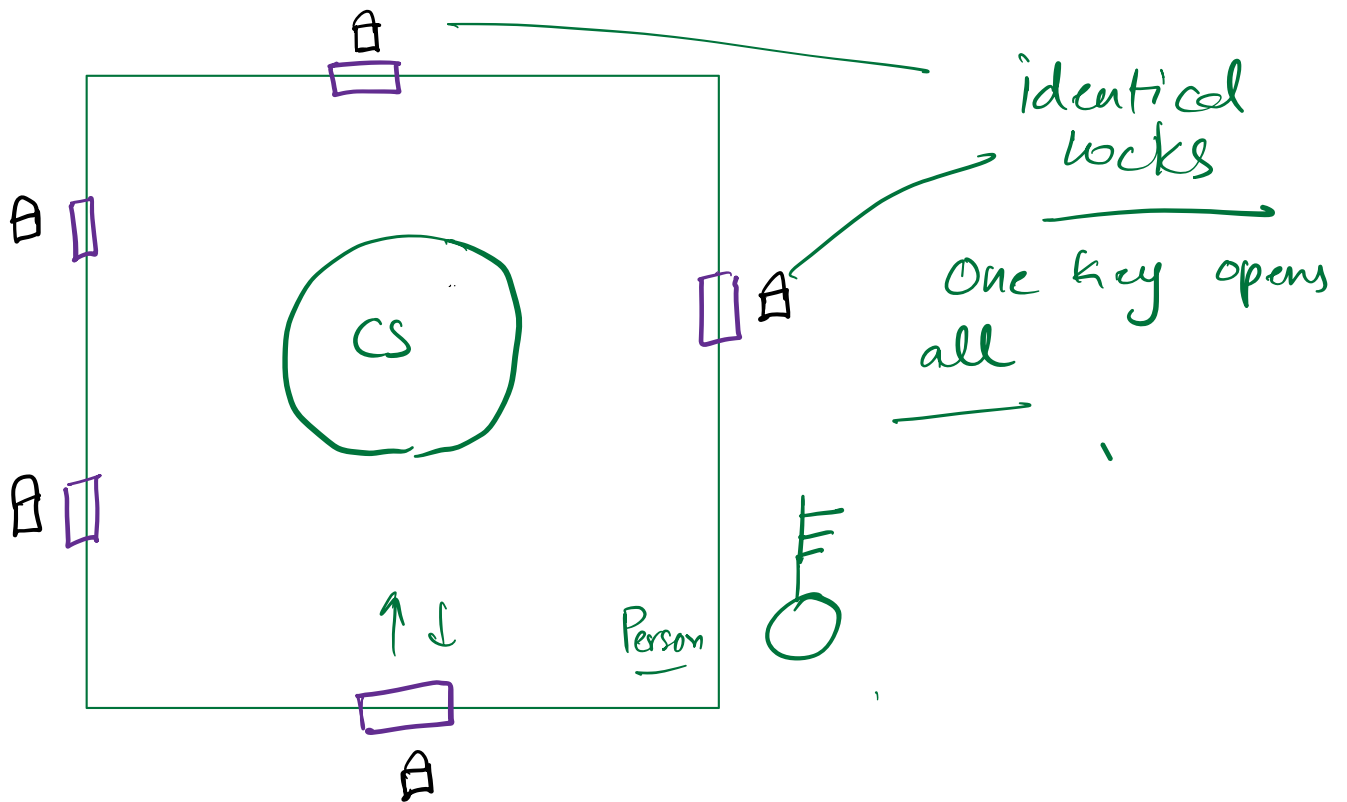
- Threads should not continuously check if CS is free. It is waste of resource.

```
while (true) {  
    if (available) {  
        =  
    }  
}
```

→ Resource
waisting

Mutex → Mutual Exclusion.

↳ Let's us restrict one thread to enter a critical section.



Adder.

$\{ t_1$

lock()

$v.val += 1$

unlock()

Subtractor

$\{ t_2$ wait

lock()

$v.val -= 1$

unlock()

8:10 → Break

Synchronized keyword

1. synchronized method
 2. Block synchronization.
- } Mutual Exclusion

1. Synchronized method.

- Implicit lock (Every object has an implicit lock)

```
synchronized public void add() {  
      
      
}
```

2. Block synchronization:

- when you don't want to take lock over whole function

```
public void add() {  
      
      
      
}
```

```
    synchronized(this) {
```

```
          
    }
```

```
          
    }
```

→ Thread Safe/
synchronised

Class Something {

synchronized void func1() { - }

void func2() { - }

synchronized void func3() { - }

}

Something - new
Something();
Something B = new
Something();

{ A. func1() , { A. func1() }
t₁ → takes lock t₂ → waits

A. func1() , A. func3() }
t₁ → takes lock t₂ → wait.

A. func1() , A. func2() }
t₁ → takes lock t₂ → start running func2
without any lock.

A. func1() ,
t₁ → takes lock

B. func1() }
t₂ → takes lock.
B. func3() ,

$t_2 \rightarrow \text{wait}$

class Account {

sync. phonePay();

sync. gpay();

}

Account a = new Account();

↓ a.phonePay()
 $t_1 \rightarrow \text{take lock}$

a.gpay(),
 $t_2 \rightarrow \text{wait}$

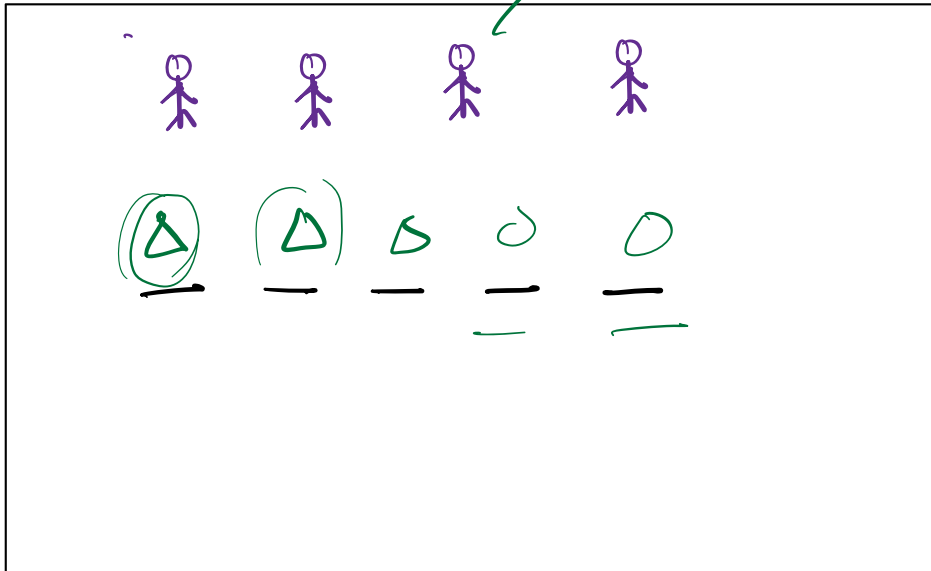
Account b = new Account();

a.phonePay()
 $t_1 \rightarrow \text{take}$

b.phonePay();

Producers / Consumer :

Producers :
they will produce
items when shelf
is empty .



Consumers will
consume item
when it is
available .

Class Store {

int maxSize; → # of shelves.
List<Object> itemList;

getSize() → current size

getMaxSize()
 addItem()
 removeItem()

2

addItem

$\downarrow P_2 \quad \downarrow P_1 = 4$
 $P_1 \rightarrow 4 < 5$
 $P_2 \rightarrow 4 < 5$

if (items.currentSize < maxSize)

{

items.addItem()

}

$\dots 4 \quad 5, 6$

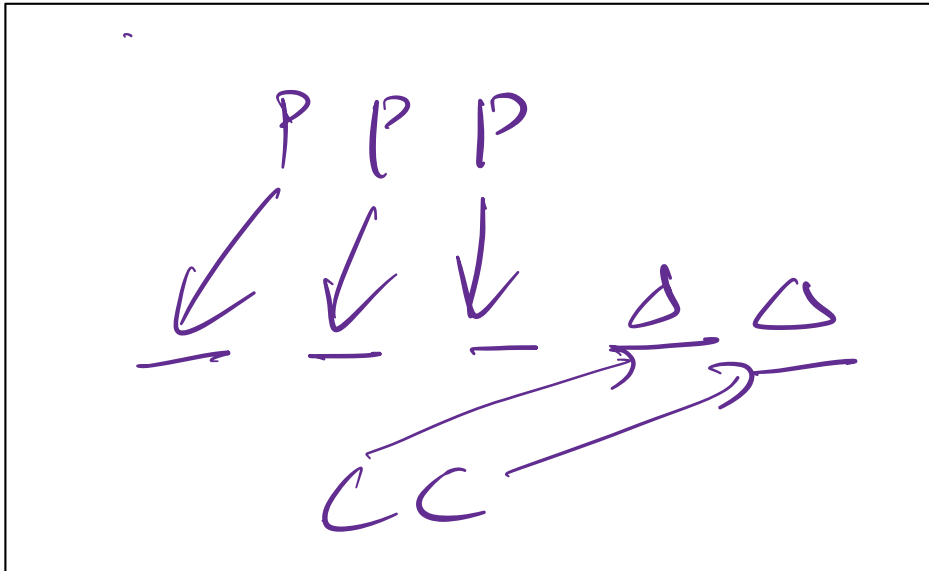
Remove Item

$\downarrow C_1 \quad \downarrow C_2 \rightarrow 1$

if (items.currentSize > 0) {

items.removeItem()

}



no of producer + no. of consumers
 = no. of shelves

- ✓ 2P, 3C ✓
- ✓ 1P, 4C ✓
- ✓ 4P, 1C ✓
- 4P, 3C ✗

Need a capability of having fixed number of threads to work on a critical section.

↳ ? locking + count

Semaphores → It is a
mutex lock with count