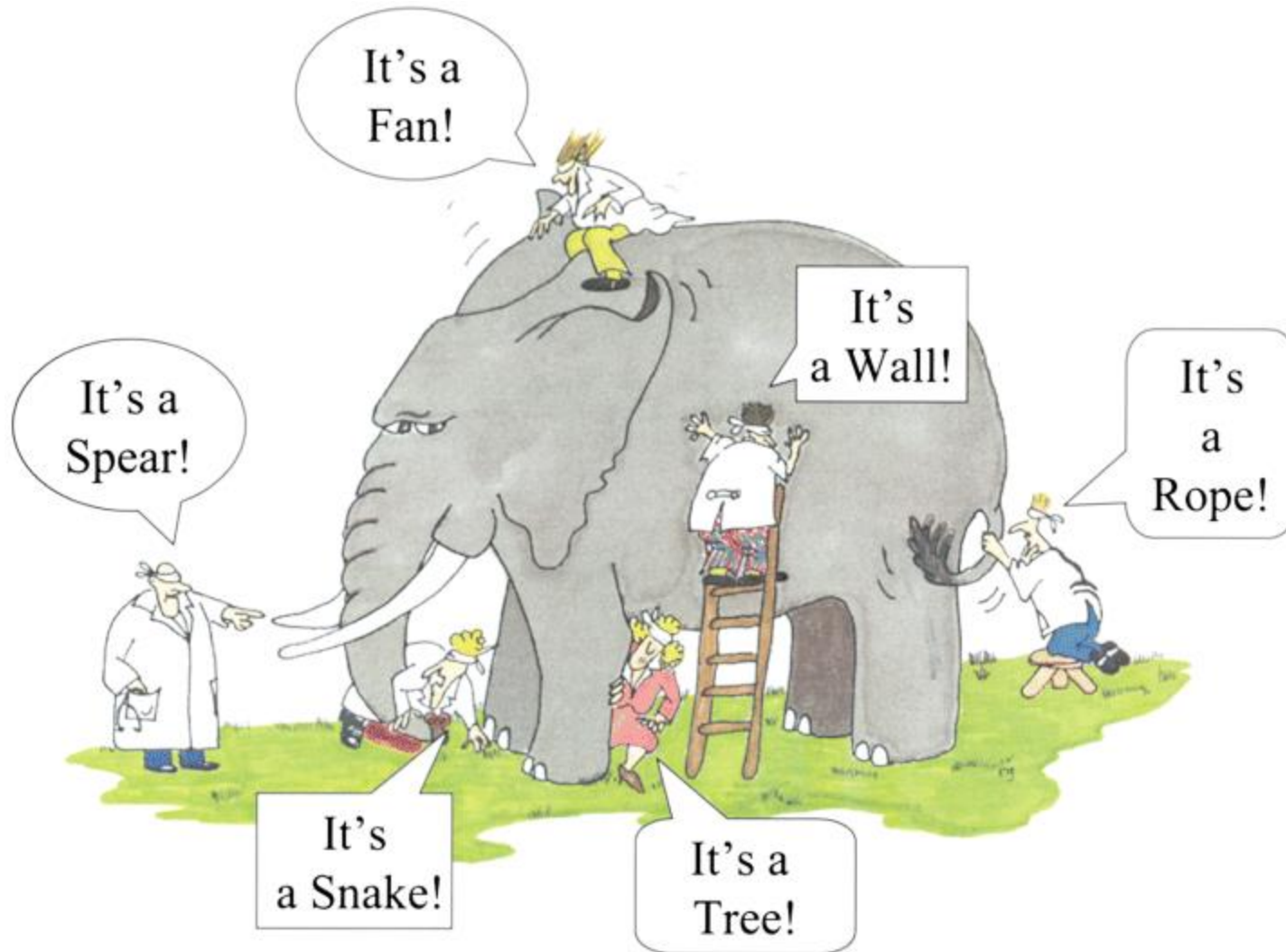




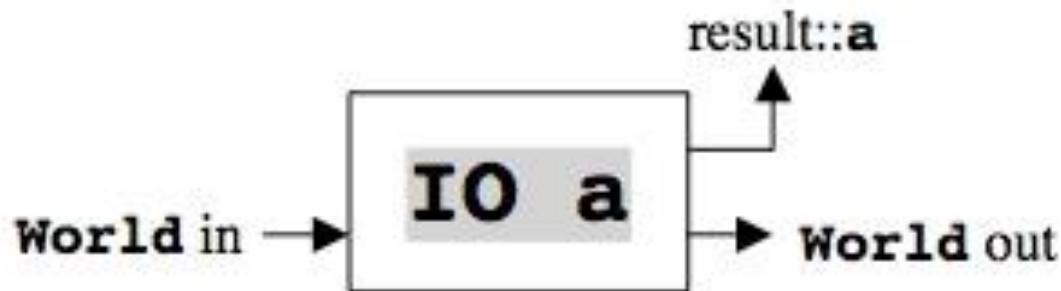
So, what's a monad?





Dealing with state

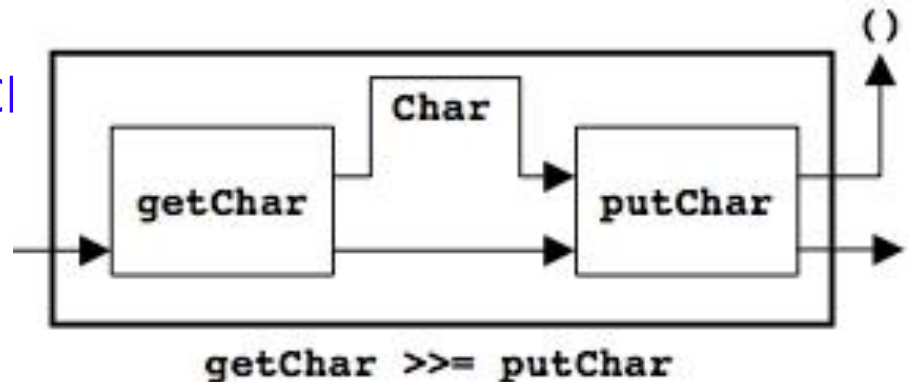
- To have state *and* pure functions, the old state of the world must be passed in as a parameter, and the new state of the world returned as a result
- A **monad** is a way of automatically maintaining state
- **IO a** can be thought of as a function whose type is $\text{World} \rightarrow (a, \text{World})$





The “bind” operator, `>>=`

- We will want to take the “state of the world” resulting from one function, and pass it into the next function
- Suppose we want to read a character and then print it
- Types:
 - `getChar :: IO Char`
 - `putChar :: Char -> IO ()`
- The result of `getChar` isn’t something that can be given to `putChar`
 - The `IO Char` “contains” a `Char` that has to be extracted to be given to `putChar`
 - `(>>=) :: IO a -> (a -> IO b) -> IO b`
- Hence,
- `Prelude> getChar >>= putChar`
`aPrelude>`





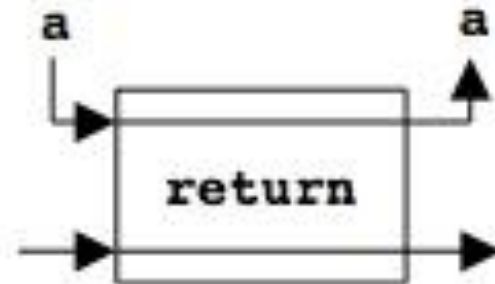
The “then” operator, >>

- The second argument to >>= is a function (such as `putChar`)
 - This is what we need for passing along a result
 - It is convenient to have another function that doesn't demand a function as its second argument
- The “then” operator simply throws away its contents
 - `(>>) :: IO a -> IO b -> IO b`
 - `Prelude> putChar 'a' >> putChar 'b' >>
putChar '\n'
ab
Prelude>`



The return function

- Finally, it is helpful to be able to create a monad container for arbitrary values
- `return :: a -> IO a`
- The action `(return v)` is an action that does no I/O, and immediately returns `v` without having any side effects
- `getTwoChars :: IO (Char,Char)`
`getTwoChars = getChar >>= \ c1 ->`
`getChar >>= \ c2 ->`
`return (c1 c2)`



do notation

- From the last slide:
 - `getTwoChars :: IO (Char,Char)`
`getTwoChars = getChar >>= \ c1 ->`
`getChar >>= \ c2 ->`
`return (c1,c2)`
- That's pretty hard to read
- The `do` provides “syntactic sugar”
 - `get2Chars :: IO (Char,Char)`
`get2Chars = do`
`c1 <- getChar`
`c2 <- getChar`
`return (c1,c2)`
 - The `do` also allows the `let` form (but without `in`)



Formal definition of a monad

- A monad consists of three things:
 - A type constructor M
 - A bind operation, $(\gg=) :: (Monad\ m) \Rightarrow m\ a \rightarrow (a \rightarrow m\ b) \rightarrow m\ b$
 - A return operation, $return :: (Monad\ m) \Rightarrow a \rightarrow m\ a$
- And the operations must obey some simple rules:
 - $return\ x \gg= f = f\ x$
 - $return$ just sends its result to the next function
 - $m \gg= return = m$
 - Returning the result of an action is equivalent to just doing the action
 - $do\ \{x \leftarrow m1;\ y \leftarrow m2;\ m3\} = do\ \{y \leftarrow do\ \{x \leftarrow m1;\ m2\}\ m3\}$
 - $\gg=$ is associative

sequence

- sequence takes a list of I/O actions and produces a list of results
- `sequence :: [IO a] -> IO [a]`
 - `main = do`
 `rs <- sequence [getLine, getLine,`
 `getLine]`
 `print rs`
 - is equivalent to
 - `main = do`
 `a <- getLine`
 `b <- getLine`
 `c <- getLine`
 `print [a,b,c]`