<u>Distinction in Complexities of Heuristic and Non – Heuristic Approach</u> Using 8-Tile Puzzle Problem

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INDEX

S.No	Topic	Page No.
1	Project Title	3
2	Abstract	3
3	Introduction	3
4	Literature Review	3
5	Objective	4
6	Methodology	4
7	Algorithm	5
8	Result	8
9	Result Analysis	9
10	Conclusion	9
11	References	9
12	Appendix	10

Project Title:

Distinction in complexities of Heuristic and non – Heuristic approach using 8-Tile Puzzle problem

Abstract:

8-tile puzzle problem, normally called 8-puzzle problem is a sliding puzzle problem. It consists of the frame (3x3) in which numbers are placed randomly from (1-8) with a blank tile. The bigger picture of this problem is the N-puzzle problem, which is solved using graph – search technique. 8-puzzle problem is a prominent workbench model for analysing the performance of a heuristic algorithm. To find the optimal solution of an N-puzzle problem is np-Hard. In this investigation, we will be analysing the performance of informed (heuristic) and uninformed (non-heuristic) search algorithm on the 8-puzzle problem to find an optimal solution to this problem. The 8-puzzle problem is converted into a game - search tree and then heuristic/non-heuristic approach is applied to it.

The heuristic search method is an intelligent method, which allows us to take decisions to reach from the current node the goal state in optimal manner. On the other hand, the latter search method looks for all possible solutions without additional information about the search space.

Introduction:

8 puzzle problem uses a 3 by 3 grid board with eight inlaid square sliding pieces marked with values 1 through 8 along with a blank tile. The blank tile can move one step at a time. The aim of this puzzle is to arrange the numbers in the grid according to the desired output with the help of legal moves.

8-puzzle problem is the largest possible N-puzzle problem that could be solved completely. The general extension of this 8-puzzle is the N-puzzle which is an np-hard problem.

☐ Alexander Reinefeld [1] investigated 8-puzzle problem and his results include data on

Literature Review:

the expected solution lengths, the 'easiest' and 'worst' configurations, and the density and distribution of solution nodes in the search tree.
Daniel R. Kunkle [2] applied A* algorithm which guarantees that the best solution (that with the least number of moves) will be found. Heuristics are examined to allow the algorithm to find the optimal solution while examining as few states as possible (maximizing the informedness of the heuristic).
Masuma Sultana et. al. [3] provide a complete solution of 8-puzzle problem using BFS in Compute Unified Device Architecture (CUDA) Environment. Objective of their work is to find the complete solution of 8-puzzle problem i.e. examining all the permutations for solvability. Using Breadth First search (BFS) graph traversal, we can reach the solution for a definite goal. The parallel algorithm is capable of providing us with much faster solution using CUDA. HuaShi [4] conducted a thorough analysis of the search solution for the 8-puzzle problem. The Breadth-first search algorithm, the Depth-first search algorithm and A *algorithm are used to solve the problem.

Objective:

- Comparison between the complexities of heuristic and non-heuristic algorithms such as A*, Branch and Bound (heuristic) and Breadth First Search (non-heuristic).
- Understanding C language and Data-structures

Methodology:

- In general, a typical 8-puzzle problem may take about 20 steps to solve, but it highly depends on the initial state of the problem. The movement of blank tile depends on its position: four possible moves if it is in middle, two possible moves if it is in corner and three along the edge, which makes the branching factor approximately 3.
- The Algorithms which we have used are: BFS (non-heuristic), A*(heuristic) and Branch and Bound (heuristic). It is to show the comparison between heuristic and non-heuristic approaches to solve the problem.
- **BFS** algorithm traverse a tree in a breadth wise motion and uses queue to remember to get the next vertex to start a search. This algorithm selects a single node (initial or source point) and then visits all the nodes adjacent to the selected node. Once the algorithm visits and marks the starting node, then it moves towards the nearest unvisited nodes and analyses them.
- A* Search Algorithm at each step picks the node according to a value **f**, which is a parameter equal to the sum of two other parameters **g** and **h**. At each step, it picks the node having the lowest **f** and process that node. Here **g** is the movement cost to move from the starting point to a given square on the grid, following the path generated to get there and **h** is the estimated movement cost to move from that given square on the grid to the final state.
- Branch and Bound heuristic algorithm depends on upper bound and lower bound regions of search space. There are three types of nodes:
 1. Live node: It is a node that has been generated but whose children have not yet
 - **2. E-node:** It is a live node whose children are currently being explored. In other words, an E-node is a node currently being expanded.
 - **3. Dead node:** It is a generated node, that is not to be expanded or explored any further. All children of a dead node have already been expanded.
- Each node X in the search tree is associated with a cost. The cost function is useful for determining the next E-node. The next E-node is the one with the least cost. The cost function is defined as:

c(x) = f(x) + h(x),

• where f(x) is the length of the path from root to x (the number of moves so far) and h(x) is the number of non-blank tiles not in their goal position (the number of misplaced tiles). There are at least h(x) moves to transform state x to a goal state.

Algorithm:

BFS:

- **STEP 1**: Create a structure NODE. Declare a matrix and variables x, y and level, cost and pointer as next.
- **STEP 2:** Create a structure Queue which has pointers as front and rear.
- **STEP 3:** Create a function Create_Queue which returns q node and initialize the front and rear to NULL.
- STEP 4: Create a function Push to push the node into the Queue.
- **STEP 5:** Create a function Display_Matrix to display the node of the tree.
- **STEP 6:** Create a function Pop to pop the front node from the Queue.
- **STEP 7:** Create a function Swap to swap two nodes.
- **STEP 8:** Create a function NewNode to create any new node.
- **STEP 9**: Create a function Cal_Cost to count no of misplaced tiles from the **goal node.**
- **STEP 10:** Create a function is_safe to check whether the node is safe or not . if the node is safe .
- **STEP 11:** Create a function Display path to print the path from the root node to goal node.
- **STEP 12:** Create a function Solve to solve the puzzle.
- **STEP 13**: In the solve function we would push root node to queue.
- **STEP 14:** pop the node from QUEUE and check whether it is equal to final node. If it is equal to final node then go to step 16 else make it live node to explore its children.
- **STEP 15:** check whether the child node is safe and calculate its cost . If safe , then push the node to queue .
- **STEP 16:** Display the path and matrix from root to goal node.

A-STAR:

STEP 1: Create a structure node_array. Declare a matrix and variables depth, text and cost and pointers as right, left ,top, down, parent and next.

STEP 2:Intialize the front and rear of the Queue to NULL and Front_list and Rear_list to NULL.

STEP3: Input goal and intial matrix.

STEP 4: Calculate heuristic function using Manhattan distance function (mod) and return its value to f variable.

STEP 5: Call the A_star function by passing the root node.

STEP 6: In A_star Function insert the root to list and queue using functions Insert_list and Insert_Queue respectively.

STEP 7: If the current node is equal to the goal node then print the path from root to goal state else go to step 8.

STEP 8: Explore the children of the current node with the help of function next_move.

STEP 9: Push the child nodes into queue and list.

STEP 10: Now prioritize the nodes in the queue according to their minimum costs with the help of the function Arrange.

STEP 11: Go to Step 4

STEP 12: Now print the number of steps taken to reach goal state

BRANCH & BOUND

```
struct list_node
{
    list_node *next;

    // Helps in tracing path when answer is found
    list_node *parent;
    float cost;
}
algorithm LCSearch(list_node *t)
{
    // Search t for an answer node
    // Input: Root node of tree t
    // Output: Path from answer node to root
```

```
if (*t is an answer node)
  {
    print(*t);
    return;
 E = t; // E-node
 Initialize the list of live nodes to be empty;
  while (true)
   for each child x of E
      if x is an answer node
        print the path from x to t;
        return;
      Add (x); // Add x to list of live nodes;
      x->parent = E; // Pointer for path to root
   if there are no more live nodes
     print ("No answer node");
     return;
    }
// Find a live node with least estimated cost
   E = Least();
   // The found node is deleted from the list of
   // live nodes
  }
```

Result:

Final

1	2	3
8	0	4
7	6	5

1	3	4
8	6	2
7	0	5

2	8	1
0	4	3
7	6	5

5	6	7
4	0	8
3	2	1

Easy

Medium

Worst

Algorithm	Difficulty level	Solution Length	Executed or not	Execution time
BFS(non- heuristic)	Easy	5	executed	0
	Medium	12	executed	0.001s
	Worst	-	TLE	-
A- Star(heuristic)	Easy	5	executed	0
	Medium	9	executed	0
	Worst	30	executed	0.88s
Branch&Bound	Easy	5	executed	0
	Medium	14	executed	0.003s
	Worst	-	TLE	-

Result Analysis:

The aim of the project was to show distinction between the complexities for solving the 8 puzzle problem.

Both the heuristic and non heuristic algorithms are implemented to prove the efficiency of the algorithms.

For the verification of the puzzle solvability A*, BFS and Branch and Bound are implemented and verified.

From the above result we can analyze the following points:

- A* algorithm which is heuristic works for all kinds of difficulty level of puzzle.
- For the medium difficulty level three of the algorithms works fine but branch and bound and BFS took more time.
- Also the length of the solution in A* algorithm is much smaller than BFS.

So we can finally assess that A* which is pure heuristic algorithm does a great job in solving this 8 puzzle problem. It provided us optimal solution for every test case with great efficiency.

Conclusion:

As we can see from the results, A* saves more time and space than BFS and Branch and Bound search, which reflects the theoretical comparison between the three algorithms' time complexities. In fact, while A* solved all three test cases, BFS and Branch and Bound didn't due to the memory requirement not being met by our computer. Therefore, we conclude that A* search algorithm is more efficient than BFS and Branch and Bound.

References:

- [1] Alexander Reinfeld, "Complete Solution of the Eight-Puzzle and the Benefit of Node Ordering in IDA*", Paderborn center for parallel computing, wax burger str.100, D-33095 Paderborn, Germany.
- [2] Daniel R. Kunkle, "Solving the 8 Puzzle in a Minimum Number of Moves: An Application of the A* Algorithm", College of Computing and Information Sciences, Rochester Institute of Technology, assessed in October 08, 2001.
- [3] Masuma Sultana, Rathindra Nath Dutta & S. K. Setua "Complete Solution of Eight Puzzle Problem using BFS in CUDA Environment", IEEE International WIE Conference on Electrical and Computer Engineering, 19-20 December, 2015
- [4]Shi & Hua. "Searching Algorithms Implementation and Comparison of Eight-Puzzle Problem", International Conference in Computer Science and Network Technology (ICCSNT), Proceedings, IEEE Xplore, Vol. 2, (pp.1203-1206), 2011.

Appendix:

A*

```
include<stdio.h>
```

```
front_list=nd;
```

```
strcpy(nd->shift_down->text,"move down");
```

For BFS:

```
#include<stdlib.h>
#include<time.h>

//#include "state.h"
//#include "list.h"
//#include "node.h"
//#include "io.h"
//#include "io.h"
//#include blank_CHARACTER '0'
//this enumerates available movements in the game relative to the blank
```

```
else if(move == LEFT && col - 1 >= 0) {
    char temp = newState->board[row][col - 1];
```

```
Move action;
```

```
//if the head node of `toAppend` is to be inserted after
```

```
extern unsigned int nodesGenerated; extern unsigned int solutionLength;
```

```
unsigned int nodesExpanded; //number of expanded nodes
SolutionPath* BFS search(State *, State *);
```

For Branch and Bound

```
Program to print path from root node to destination node for N*N -1 puzzle algorithm using Branch and Bound The solution assumes that instance of puzzle is solvable
Node* newNode(int mat[N][N], int x, int y, int newX,
```

```
// update new blank tile cordinates
node->x = newX;
```

```
printf("\n");
  // Solvable Final configuration
  // Value 0 is used for empty space

int final[N][N];
printf("Enter final matrix");
for(int i=0;i<N;i++)
{
    for(int j=0;j<N;j++)
    {
        cin>>final[i][j];
    }
}
printf("\n");

// Blank tile coordinates in initial
  // configuration
  int x = 1, y = 2;
  solve(initial, x, y, final);
  return 0;
}
```