

# IEOR 4500

## Robust Arbitrage Model

Notation:

- $K$  = number of scenarios
- $0, 1, 2, \dots, n$ : indices for assets (cash is 0)
- $r$  = risk-free interest rate
- $\bar{\pi}_{0j}$  = today's price for asset  $j$ ,
- $\bar{\pi}_{kj}$  = expected value of  $\pi_{kj}$ , the price for asset  $j$  in scenario  $k \geq 1$  ( $0 \leq j \leq K$ )
- $\sigma_{kj}$  = maximum deviation of price  $\pi_{kj}$  away from its expectation.

The robust problem we want to solve is:

$$W^* \doteq \min \sum_{j=0}^n \pi_{0j} x_j, \quad (1)$$

*s.t.*

$$\sum_{j=0}^n \pi_{kj} x_j \geq 0, \quad \text{for all } k, \text{ and} \quad (2)$$

for all  $\pi_{kj}$  such that  $\bar{\pi}_{kj} - \sigma_{kj} \leq \pi_{kj} \leq \bar{\pi}_{kj} + \sigma_{kj}$

$$-1 \leq x_j \leq 1, \quad 0 \leq j \leq n. \quad (3)$$

**Lemma 1.**  $W^* < 0$  if and only if a type A arbitrage exists.

Lemma 1 casts the problem of testing for a type A arbitrage into a Linear Program with an infinite number of constraints. Next we will render this into a useful formulation. Given a vector  $x$ , we say that  $x$  is *feasible for scenario*  $k$  ( $1 \leq k \leq K$ ) if

$$\sum_{j=0}^n \pi_{kj} x_j \geq 0,$$

for every choice of prices  $\pi_{kj}$  for scenario  $k$  such that  $\bar{\pi}_{kj} - \sigma_{kj} \leq \pi_{kj} \leq \bar{\pi}_{kj} + \sigma_{kj}$  for every asset  $j$ . So the robust problem can be rewritten as

$$W^* \doteq \min \sum_{j=0}^n \pi_{kj} x_j, \quad (4)$$

*s.t.*

$$x \text{ is feasible for scenario } k, \text{ for every } 1 \leq k \leq K \quad (5)$$

$$-1 \leq x_j \leq 1, \quad 0 \leq j \leq n. \quad (6)$$

Constraint (5) describes an infinite number of conditions. We can make this into a compact statement as follows. Consider a given, fixed asset vector  $\hat{x}$ . Note that  **$\hat{x}$  is feasible for scenario  $k$  if and only if  $V_k^*(\hat{x}) \geq 0$** , where

$$V_k^*(\hat{x}) \doteq \min \sum_{j=0}^n \hat{x}_j \pi_{kj}, \quad (7)$$

$$s.t. \quad \bar{\pi}_{kj} - \sigma_{kj} \leq \pi_{kj} \leq \bar{\pi}_{kj} + \sigma_{kj} \quad 0 \leq j \leq n. \quad (8)$$

In this LP, the  $\hat{x}$  vector is given data, and the  $\pi_{kj}$  are the variables. We can rewrite this LP as follows:

$$V_k^*(\hat{x}) \doteq \min \sum_{j=0}^n \hat{x}_j \pi_{kj}, \quad (9)$$

$$s.t. \quad (10)$$

$$\pi_{kj} \geq \bar{\pi}_{kj} - \sigma_{kj}, \quad 0 \leq j \leq n, \quad (11)$$

$$-\pi_{kj} \geq -\bar{\pi}_{kj} - \sigma_{kj}, \quad 0 \leq j \leq n. \quad (12)$$

The value of an LP is equal to that of its dual. The dual of (9)-(12) is a *maximization* problem. Thus, it follows that  $V_k^*(\hat{x}) \geq 0$  if and only if there exists a vector of duals feasible for the dual of (9)-(12), of nonnegative value. Denoting the dual of constraint (11) by  $u_{kj}$ , and the dual of constraint (12) by  $v_{kj}$ , we have that  $V_k^*(\hat{x}) \geq 0$  (i.e., vector  $\hat{x}$  is feasible for scenario  $k$ ) if there exist values  $u_{kj}, v_{kj}$  ( $1 \leq j \leq n$ ) such that

$$\sum_{j=0}^n (\bar{\pi}_{kj} - \sigma_{kj}) u_{kj} + (-\bar{\pi}_{kj} - \sigma_{kj}) v_{kj} \geq 0, \quad (13)$$

$$u_{kj} - v_{kj} = \hat{x}_j, \quad 0 \leq j \leq n, \quad (14)$$

$$u_{kj} \geq 0, v_{kj} \geq 0, \quad 0 \leq j \leq n, \quad (15)$$

In summary, our robust arbitrage testing LP is

$$W^* \doteq \min \sum_{j=0}^n \pi_{0j} x_j, \quad (16)$$

*s.t.*

$$\sum_{j=0}^n (\bar{\pi}_{kj} - \sigma_{kj}) u_{kj} + (-\bar{\pi}_{kj} - \sigma_{kj}) v_{kj} \geq 0, \quad 1 \leq k \leq K, \quad (17)$$

$$u_{kj} - v_{kj} - x_j = 0, \quad 0 \leq j \leq n, \quad 1 \leq k \leq K, \quad (18)$$

$$u_{kj} \geq 0, v_{kj} \geq 0, \quad 0 \leq j \leq n, \quad 1 \leq k \leq K, \quad (19)$$

$$-1 \leq x_j \leq 1, \quad 0 \leq j \leq n. \quad (20)$$

In this linear program, the variables are the  $x$ , the  $u$  and the  $v$ . We can think of this process as a 'game': we pick the vector  $x$  and there is an adversary who in each scenario will try to manipulate prices so that the value of the portfolio in that scenario becomes negative. By picking  $x$  so that constraints (18) and (19) hold, we make the adversary's task **impossible**: no (legal) choice of the prices will make the portfolio value negative in any scenario.