

CS39002 Operating Systems Laboratory Spring 2025

Lab Assignment: 2 Date of submission: 15-Jan -2025

Inter-process communication using signals

Think of n children $C_1, C_2, C_3, \dots, C_n$ standing in a circle, and playing a game with their parent P standing at the center of the circle. P throws a ball to the children in the circular order. If the child (say, C_i) to which the ball is thrown can catch the ball, then C_i continues to play. If C_i misses the ball, then C_i goes out of the game. After each throw, the ball comes back to P who then throws the ball to the next (in the circular order) child who is not yet out of the game. Eventually, $n - 1$ children miss and go out of the game. The remaining child wins the game.

You need to implement this game as a multi-process application where the processes communicate with one another by signals. Write two programs *parent.c* and *child.c* to simulate the working of the parent process and of each child process. Suppose that these two programs are compiled to the executable files *parent* and *child*. The program *parent* (which simulates P) is run with one command-line argument n (the number of child processes). P creates n child processes C_i which exec *child* i for $i = 1, 2, 3, \dots, n$. P also writes, in a text file *childpid.txt*, the child count n and the PIDs of the n child processes created by P . Each child first waits for some time (like one second) for P to finish writing to *childpid.txt*. After this wait, each child reads n and the n PIDs from the text file *childpid.txt*. After the child creation, P waits for some time (like two seconds) so that each child process gets time to read *childpid.txt*. These waits may be implemented by *sleep()* or *usleep()*, but after this, no waits based on these functions will be allowed.

The parent P then enters a loop which continues until only one child is left as the player. Each child C_i , on the other hand, enters an infinite loop. The body of each loop should contain the single system call *pause()* which lets the calling process wait until it receives a signal (this way, a busy wait is avoided). The game of throwing balls and catching/missing throws will be implemented by sending signals. In this assignment, we use the three signals SIGUSR1, SIGUSR2, and SIGINT only. P starts the game by sending SIGUSR2 to C_1 .

For a child process C_i , receiving SIGUSR2 implies that a throw is made to it. It then randomly decides whether it catches the ball (with probability 0.8) or misses the ball (with probability 0.2). If C_i catches the ball, it sends SIGUSR1 to P . If C_i fails to catch the ball, it sends SIGUSR2 to P . Depending on the type of the signal received from the child C_i (to which the throw is made), the parent knows whether that child continues to play the game or is out of the game. P records this information. The child too records its own status.

After the outcome of a throw is recorded as explained above, P initiates a printing of the current status of all the n players. Since P has all the necessary information, it can do that printing itself. However, as a mandatory part of this assignment, this printing should be done collaboratively by the child processes. This is achieved by sending SIGUSR1 to the child processes in turn. Recall that to a child process, receiving SIGUSR2 means that a throw is made to it. On the other hand, the reception of SIGUSR1 instructs that child to print its current status. The possible status of a child are PLAYING (written as \dots), CATCHMADE (written as CATCH), CATCHMISSED (written as MISS), and OUTOFGAME (written as blank). See the sample at the end to know the format of printing. Strictly follow this format.

P initiates the printing by sending SIGUSR1 to C_1 . For each $i < n$, C_i prints its status, and then sends SIGUSR1 to the next child C_{i+1} . The last child C_n prints its status, but does not send SIGUSR1 to any other process. However, C_n takes part in the synchronization activity in a different manner. Until all the child processes finish writing their status, the parent P must wait before it can make the throw to the next playing child. Currently, you know only a few synchronization primitives, so let this wait be accomplished by *waitpid()*. Before sending SIGUSR1 to C_1 , P forks a dummy child process D , and writes the PID of D in a text file *dummycpid.txt*. After sending SIGUSR1 to C_1 , P waits until D exits. When C_n is done printing its status, it reads the PID of D from the file *dummycpid.txt*, and sends SIGINT to D . Write *dummy.c* (the code

for D) that enters an infinite loop of `pause()` at the beginning of its *main()*. It is not meant to do any useful work except putting an end to the wait of the parent P .

When D exits, P wakes up, and works out the next playing child process C_{next} to which the throw is to be made. Recall that P maintains the information of the status of all child processes. P should also keep track of the child process to which the last throw is made. So P can determine C_{next} easily. P then sends `SIGUSR2` to C_{next} , and the game continues as explained above.

After $n - 1$ child processes miss throws, the parent sends `SIGINT` one by one to all of the n child processes. Only the last playing child process prints a happy message, and exits (see the format in the sample). The other processes exit without printing anything.

The sequencing of the throw-and-print cycle must be implemented only by signals (and by `waitpid()` in one situation). No other synchronization mechanism is allowed. Use `fflush(stdout)`; to avoid garbled output. But do not use any `sleep` or `usleep` calls (except only at the very beginning, that is, before the game starts).

You may use the following *makefile*.

```
all:
    gcc -Wall -o parent parent.c
    gcc -Wall -o child child.c
    gcc -Wall -o dummy dummy.c

run: all
    ./parent 10

clean:
    -rm -f parent child dummy childpid.txt dummycpid.txt
```

Submit a zip/tar/tgz archive containing the files *parent.c*, *child.c*, *dummy.c*, and *makefile*.

Sample Output

```
$ make run
gcc -Wall -o parent parent.c
gcc -Wall -o child child.c
gcc -Wall -o dummy dummy.c
./parent 10
Parent: 10 child processes created
Parent: Waiting for child processes to read child database
```

[illegible]