### **Algorithm Analysis and Design**

# Linear Time Sorting (Radix Sort and Bucket Sort)



**Lecture -22** 

# **Linear Time Sorting** (Radix Sort)

#### Overview

- Running time of counting sort is  $\Theta(d(n+k))$
- Required extra space for sorting.
- Is a stable sorting.



- Radix sort is non comparative sorting method
- Two classifications of radix sorts are least significant digit (LSD) radix sorts and most significant digit (MSD) radix sorts.
- LSD radix sorts process the integer representations starting from the least digit and move towards the most significant digit. MSD radix sorts work the other way around.

# Radix Sort (Algorithm)

```
Radix_Sort(A,d)

for i \leftarrow d down to 1

Use a stable sort to sort the array A on digit i

(i.e. Counting Sort)
```

• In input array A, each element is a number of d digit.  $Radix\_Sort(A, d)$ 

```
for i ← 1 tod
  do "use a stable sort to sort array A on digit i;
329
457
657
839
436
720
```

355

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457	355	329
657	436	436
839	457	839
436	657	355
720	329	457
355	839	657
	<b>^</b>	<b>^</b>

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457	355	329	355
657	436	436	436
839	457	839	457
436	657	355	657
720	329	457	720
355	839	657	839
	<b>^</b>	<u> </u>	<b>^</b>

# Radix Sort (Analysis)

```
Radix_Sort(A,d)

for i \leftarrow d down to 1

Use a stable sort to sort the array A on digit i

(i.e. Counting Sort)
```

- Here Counting Sort execute for d times.
- The running time of Counting Sort is  $\Theta(n+k)$
- Hence the running time complexity of Radix Sort is O(d(n+k))

# **Linear Time Sorting** (Bucket Sort)

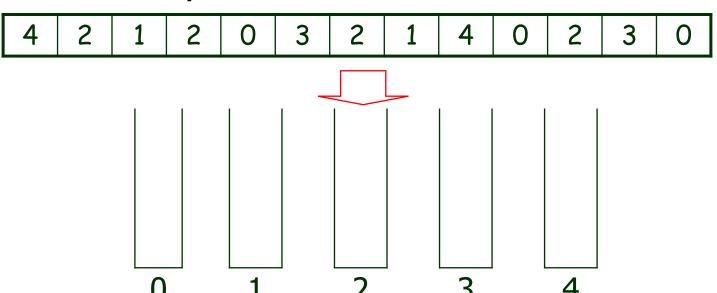
#### Overview

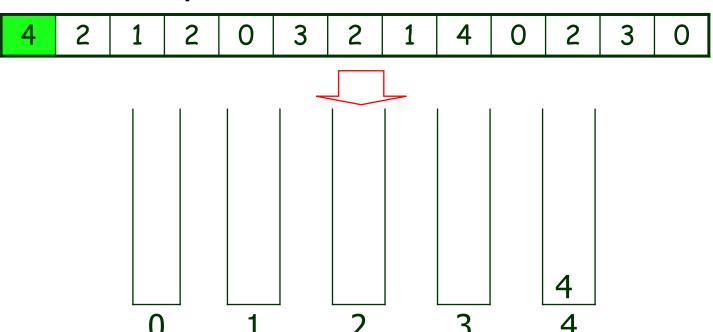
- The average time complexity is O(n + k).
- The worst time complexity is  $O(n^2)$ .
- Required extra space for sorting.
- Is a stable sorting.

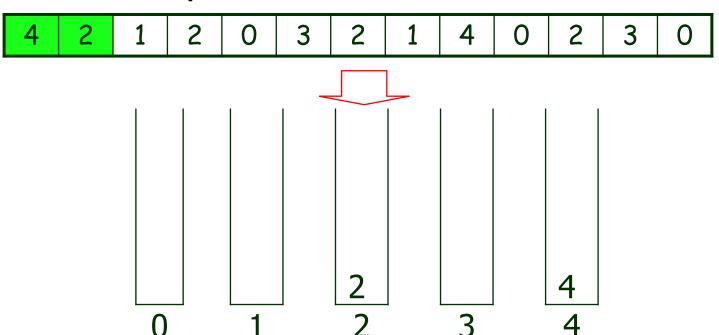
- Bucket sort is a comparison sort algorithm that operate on elements by dividing them into different bucket and return the result.
- Buckets are assigned based on each element's search key.
- A the time of returning the result, First concatenate each bucket one by one and then return the result in a single array.

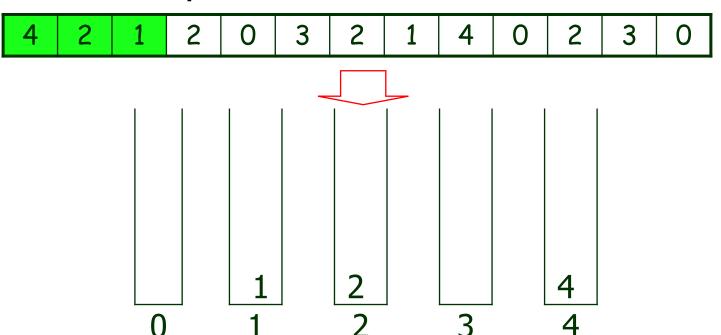
- Some variations
  - Make enough buckets so that each will only hold one element, use a count for duplicates.
  - Use fewer buckets and then sort the contents of each bucket.
- The more buckets you use, the faster the algorithm will run but it uses more memory.

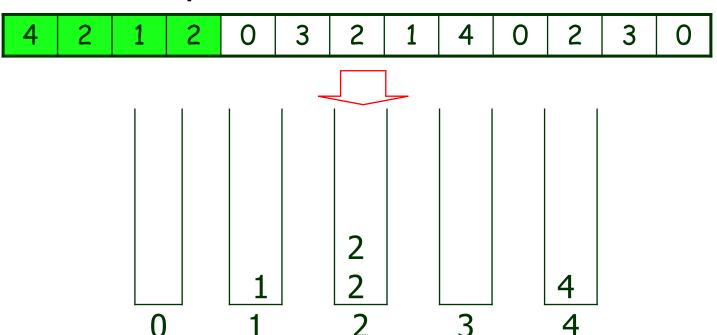
- Time complexity is reduced when the number of items per bucket is evenly distributed and it is closed to one item per bucket.
- As buckets require extra space, This algorithm trading increased space consumption for a lower time complexity.
- In general, Bucket Sort beats all other sorting techniques in time complexity but can require a huge of space.

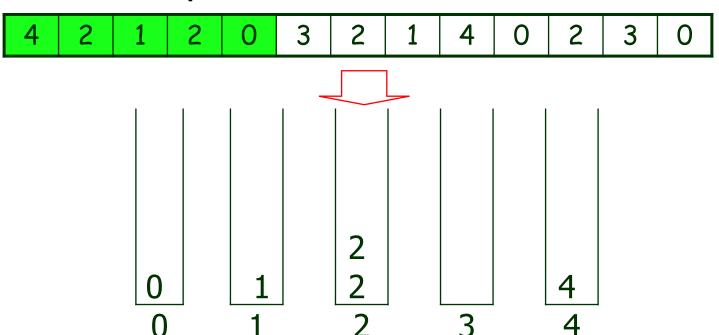


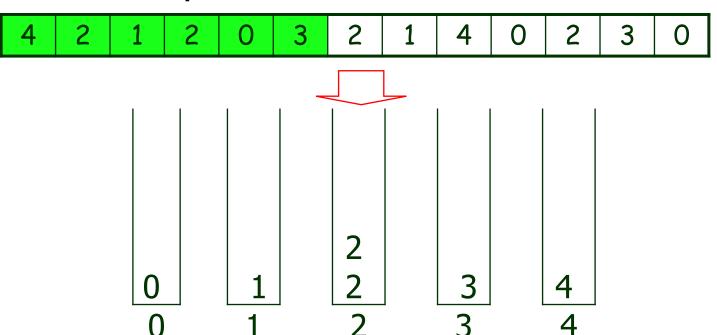


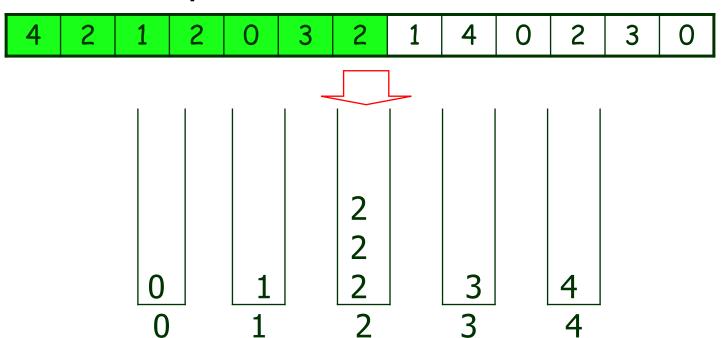


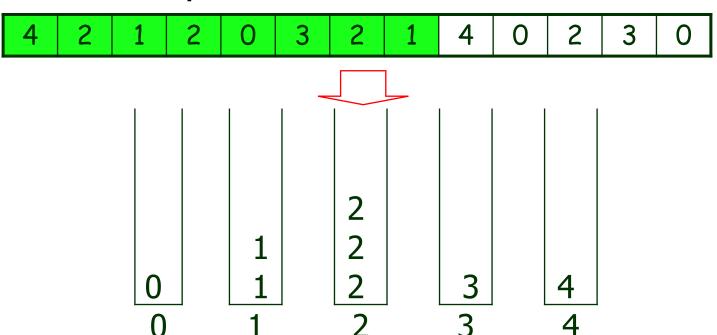


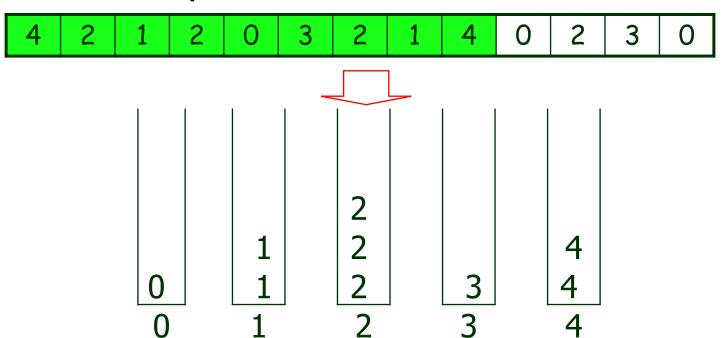


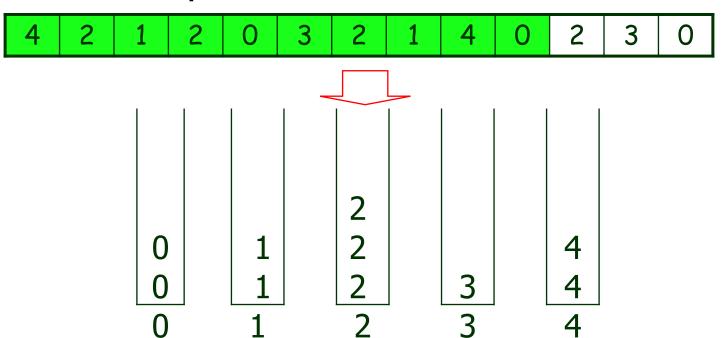


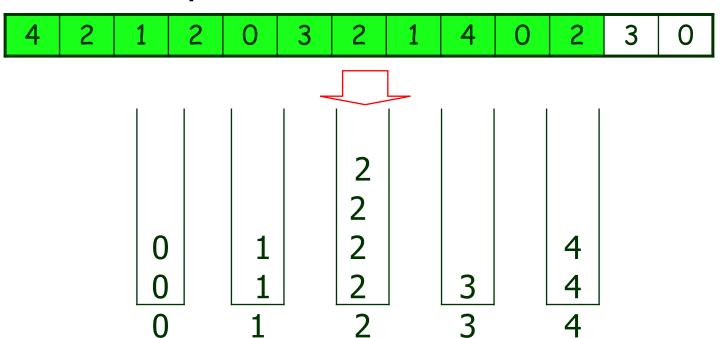


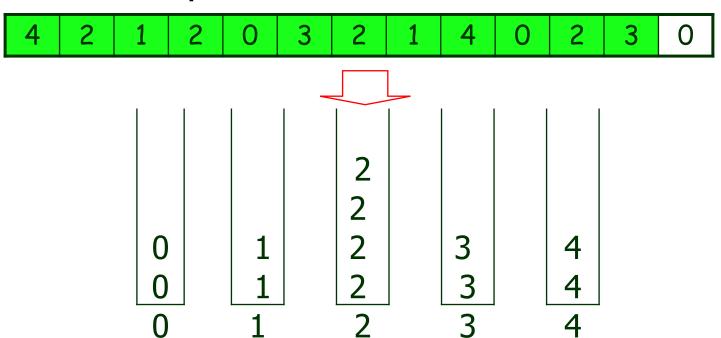


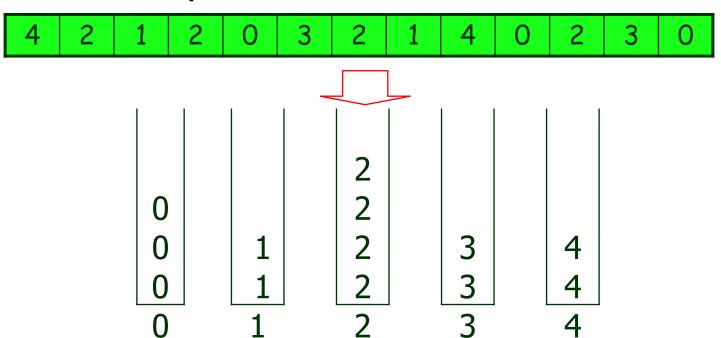


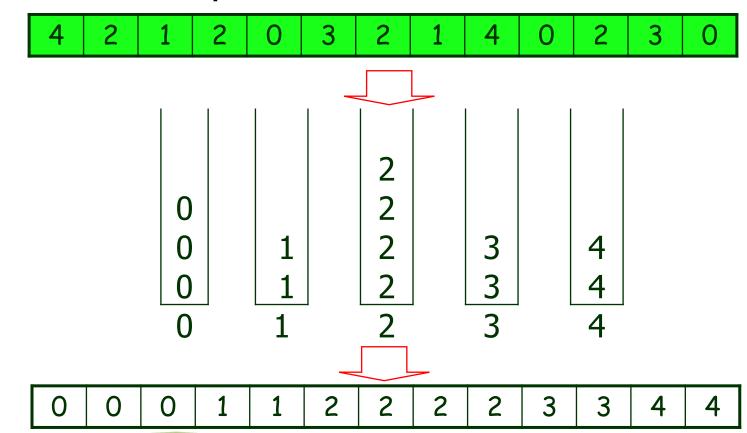








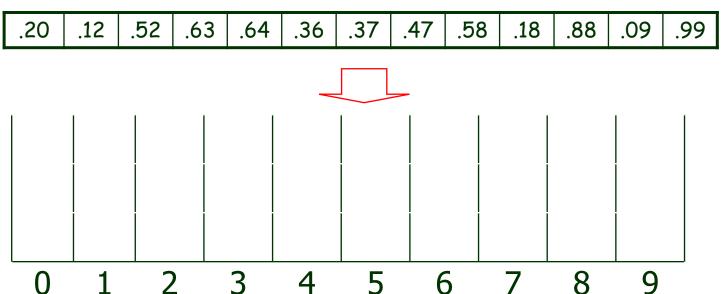


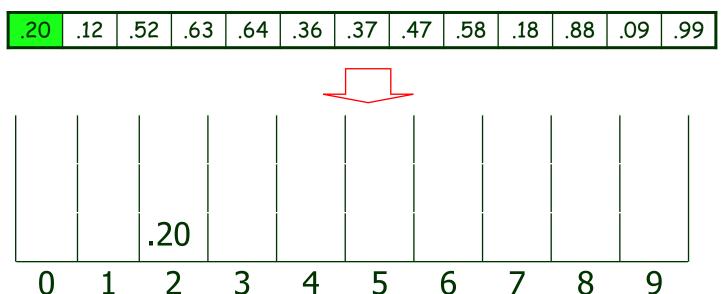


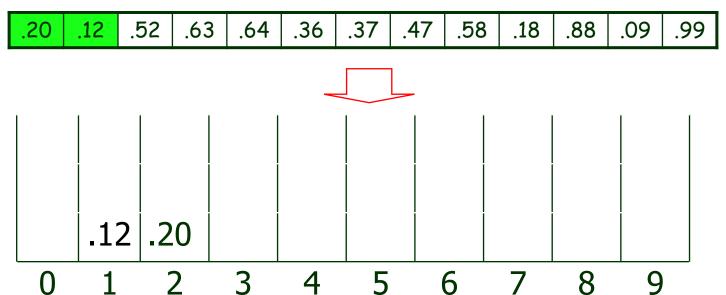
```
• One Value per bucket:
Algorithm BucketSort(S)
(values in S are between 0 and m-1)
for j \leftarrow 0 to m-1 do // initialize m buckets
  b[j] \leftarrow 0
for i \leftarrow 0 to n-1 do // place elements in their
   b[S[i]] \leftarrow b[S[i]] + 1 // appropriate buckets
i \leftarrow 0
for j \leftarrow 0 to m-1 do
                               // place elements in buckets
  for r \leftarrow 1 to b[j] do // back in S (Concatination)
       S[i] \leftarrow j
       i \leftarrow i + 1
```

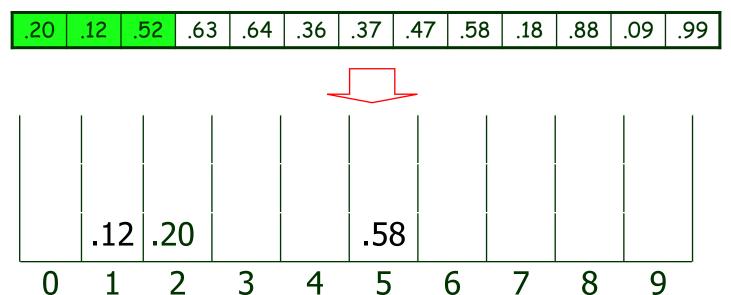
#### One Value per bucket (Analysis)

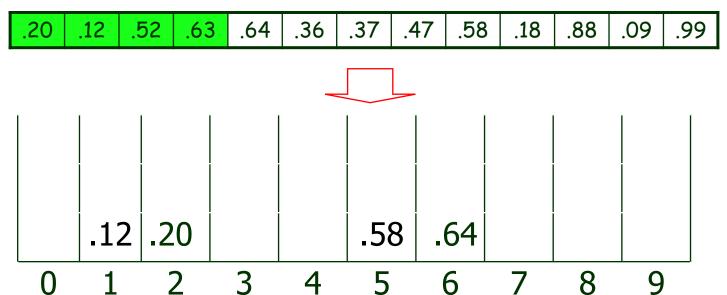
- Bucket initialization: O(m)
- From array to buckets: O(n)
- From buckets to array: O(n)
  - Due to the implementation of dequeue.
- Since m will likely be small compared to n, Bucket sort is O(n)
- Strictly speaking, time complexity is O(n + m)

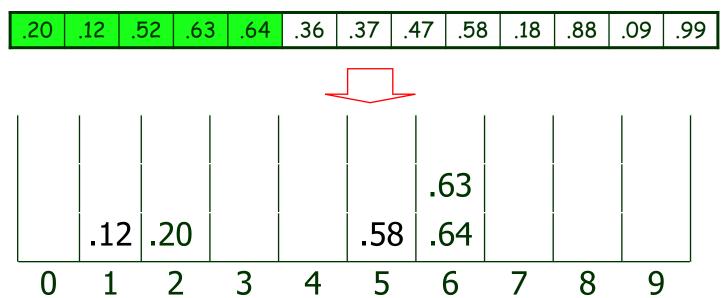


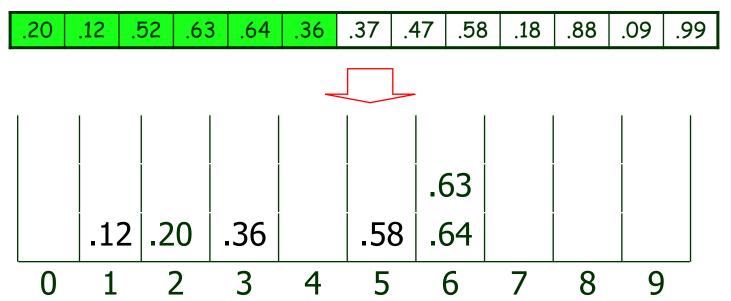


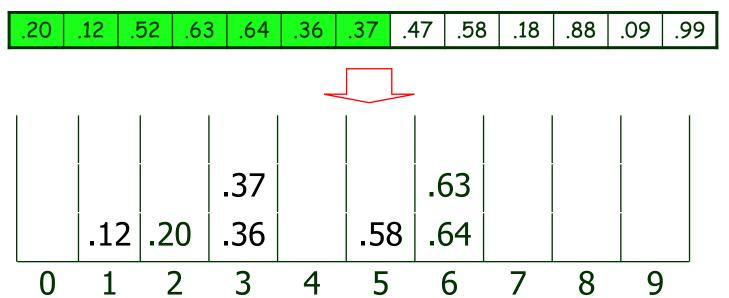


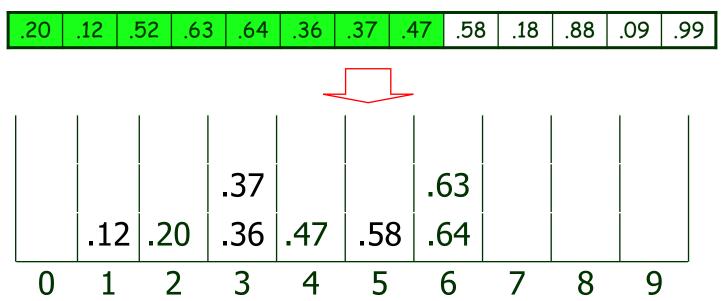


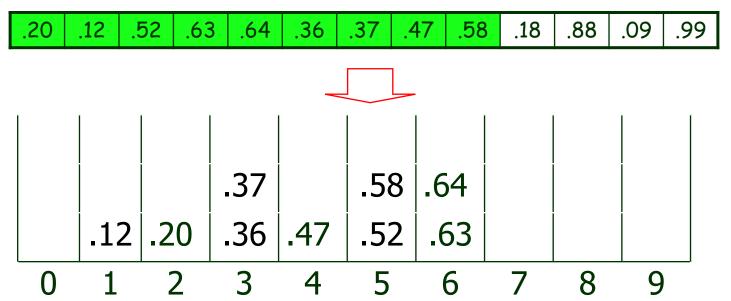


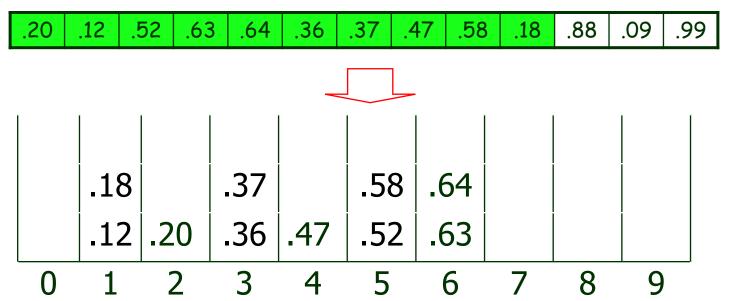


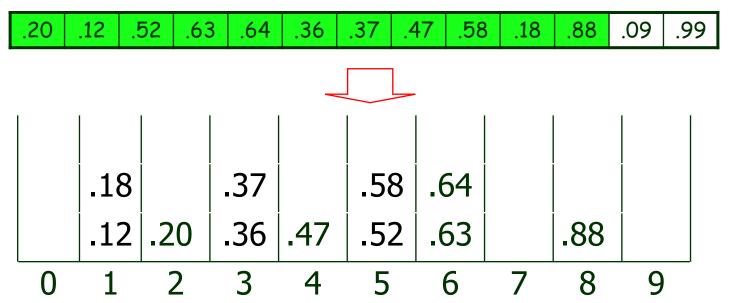


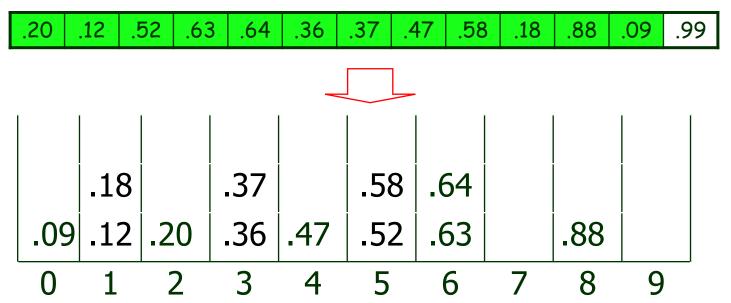


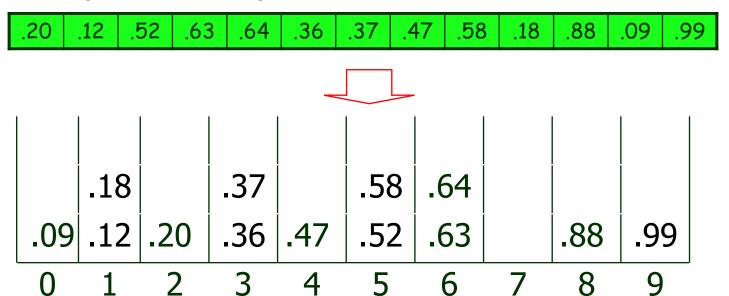


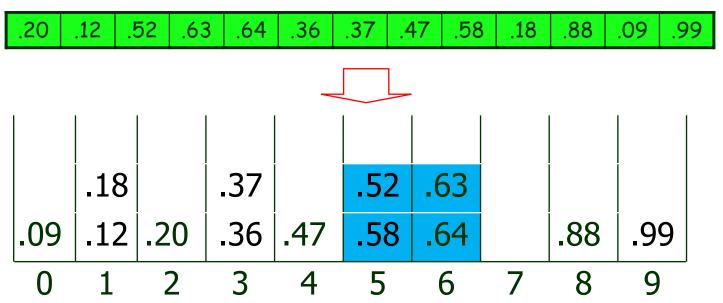




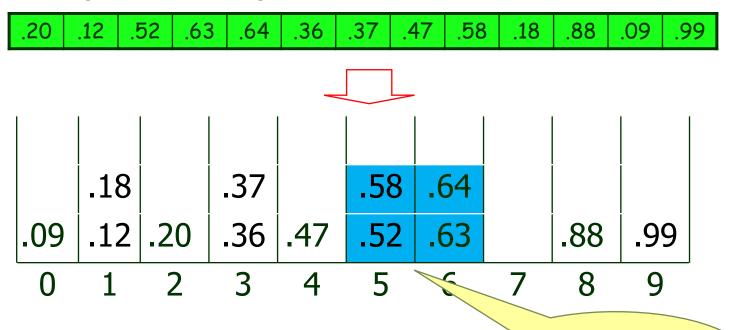




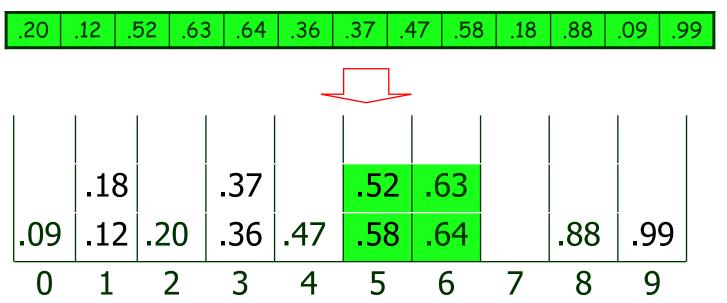


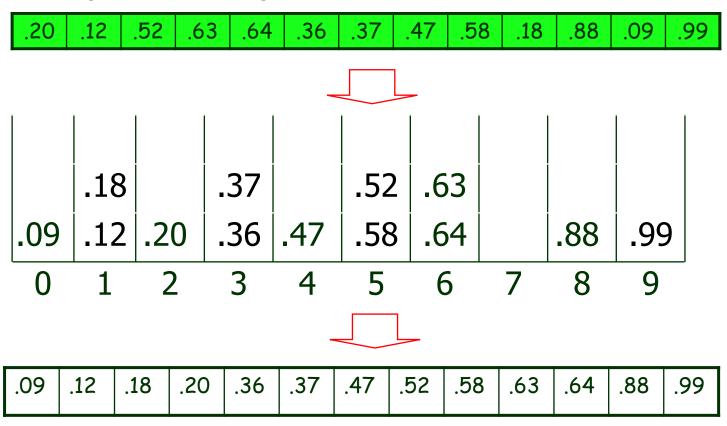


Multiple items per bucket:



Apply Internal sorting(stable)





- Multiple items per bucket:
   Algorithm BucketSort(S)
- 1. Let B[0..(n-1)] be a new array.
- 2.  $n \leftarrow A.length$
- 3. for  $i \leftarrow 0$  to n-1
- 4. make B[i]an empty list
- 5. for  $i \leftarrow \leftarrow 1$  to n
- 6. insert A[i] into list B[n \* A[i]]
- 7. for  $i \leftarrow 0$  to n-1
- 8. sort list B[i] with a stable sorting(insertion sort)
- 9. Concatenate the list  $B[0], B[1], B[2], \dots, B[n-1]$  together in order.

#### Multiple items per bucket (Analysis)

- It was observed that except line no 8 all other lines take O(n) time in worst case.
- Line no. 8 (i.e. insertion sort) takes  $O(n^2)$ .
- The average time complexity for Bucket Sort is O(n + k) in uniform distribution.

#### **Characteristics of Bucket Sort**

- Bucket sort assumes that the input is drawn from a uniform distribution.
- The computational complexity estimates involve the number of buckets.
- Bucket sort can be exceptionally fast because of the way elements are assigned to buckets, typically using an array where the index is the value.

#### **Characteristics of Bucket Sort**

- This means that more auxiliary memory is required for the buckets at the cost of running time than more comparison sorts.
- The average time complexity is O(n + k).
- The worst time complexity is  $O(n^2)$ .
- The space complexity for Bucket Sort is O(n + k).

