

1)

a) $F(n) = F_{n+1} - 1.$

$$T_F(0) = 0$$

$$T_F(1) = 0$$

Assume the below is true.

$$T_F(k) = F(k+1) - 1.$$

Substitute $n = k+1$,

$$T_F(k+1) = 1 + T_F(k) + T_F(k-1)$$

$$= 1 + F(k+1) - 1 + F(k) - 1$$

$$= F(n+2) - 1$$

b)

2) $f(0; a, b) = a$

$$f(1; a, b) = b$$

$$f(n; a, b) = f(n-1; b, a+b)$$

$$T_F(0) = 0$$

$$T_F(1) = 0$$

$$T_F(n) = 1 + T_F(n-1)$$

$$= 1 + 1 + T_F(n-2)$$

$$= 1 + 1 + 1 + T_F(n-3)$$

Put $k = n$,

$$T_F(n) = k + T_F(n-k)$$

$$T_F(n) = k + T_F(0)$$

$$T_F(n) = n$$

The time complexity of fibItHelper function is $O(n)$.

3) For $n=1$,

$$L(a, b) = (b, a+b)$$

$$\text{i.e., } (f(1; a, b), f(2; a, b)) = (b, f(1; b, a+b)) \\ = (b, a+b)$$

Assume true for $n=k$,

$$L^k(a, b) = (f(k; a, b), f(k+1; a, b))$$

So for $n = k+1$,

Substituting $k+1$ we get,

$$L^{(k+1)} = L(L^k(a, b))$$

$$= L(f(k; a, b), f(k+1; a, b))$$

$$= (f(k+1; a, b), f(k; a, b) + f(k+1; a, b))$$

$$= (f(k+1; a, b), f(k+2; a, b))$$

For any n we can prove $\in \mathbb{N}$, $L^n(a, b) = (f(n; a, b), f(n+1; a, b))$.

5)

- a. False, Fib is not a pseudo-polynomial. This is because time complexity grows exponentially with increasing input value. For instance, the amount of time needed to calculate the recursive functions $\text{fib}(n-1)$ and $\text{fib}(n-2)$ will be same.
- b. True, Yes, fibIt is pseudo-polynomial. This is due to the fact that, in the worst situation, computing the n th number in the Fibonacci series will require $O(n)$.
- c. True, FibPow is indeed a pseudo-polynomial. In the worst case, it will take $O(n \log(n))$ because each time fibPow calls itself, the value of n is half.

6)

- a)

$$\begin{aligned}
 T(n+1) &= T(n) + 5 \\
 T(0) &= 5 \\
 T(1) &= T(0) + 5 \\
 &= 5 \\
 T(n) &= T(n-1) + 5 \\
 &= T(n-2) + 5 + 5 \\
 &= T(n-3) + 5 + 5 + 5 \\
 &= T(n-4) + (5 \cdot 4) \\
 \text{Using } k \\
 &= T(n-k) + (5 \cdot k) \\
 \text{Putting } k \text{ in } n \\
 n &= k \\
 &= T(k-k) + (5 \cdot k) \\
 &= T(0) + 5 \cdot k \\
 &= 5k \\
 T(n) &= 5n
 \end{aligned}$$
- b)

$$\begin{aligned}
 T(n+1) &= n + T(n) \\
 T(0) &= 0 \\
 T(n) &= (n-1) + T(n-1) + (n-1) \\
 T(n) &= (n-2) + T(n-2) + (n-1) \\
 T(n) &= (n-3) + T(n-3) + (n-1) \\
 T(n) &= (n-k) + T(n-k) + (n-1) \\
 n &= k \\
 T(k) &= 0 + T(0) + (k-1) \\
 T(n) &= T(0) + 0 + 0 + 1 + 2 + 3 \dots \\
 i = 0 \quad \sum n &= i - n \\
 &= n(n+1)/2 - n \\
 &= n^2 - n/2
 \end{aligned}$$

7)

- a) $T(n+1) = 2T(n)$

$$\begin{aligned}
T(0) &= 1 \\
T(1) &= 2T(n-1) \\
T(1) &= 2T(0) \\
T(1) &= 0 \\
T(n) &= 2T(n-1) \\
T(n) &= 2(2T(n-2)) \\
T(n) &= 2(2(2T(n-3))) \\
&\text{Put } k \\
T(n) &= 2^k T(n-k) \\
&n = k \\
T(n) &= 2^n T(0) \\
T(n) &= 2^n \cdot 1 \\
T(n) &= 2^n \\
&\text{Hence Proved}
\end{aligned}$$

b)

$$\begin{aligned}
T(n+1) &= 2n+1 + T(n) \\
T(n) &= 2n + T(n-1) \\
T(n) &= 2n + 2n-1 + T(n-2) \\
T(n) &= 2n + 2n-1 + 2n-2 + T(n-3) \\
T(n) &= 2n + 2n-1 + 2n-2 + \dots + 2n-(k-1) + T(n-k) \\
&\text{Let } k = n \\
T(n) &= 2n + 2n-1 + 2n-2 + \dots + 2n-1 + 2n \\
&= (2n+1-1)/(2-1) \\
&= 2n+1 - 1
\end{aligned}$$