

* Implementation of Undirected Connectivity

- The approach for this problem is also the same as BFS, the only change whenever we reach the end of one part we will select a different leader node and iterate through that node.
- General Algorithm

BFS(G, S)

Mark S as 'explored'

$Q \leftarrow$ queue initialised by S .

while $Q \neq \emptyset$

 remove from front of Q the vertex ' v '

 for each edge (v, w)

 if w is unexplored

 mark w as explored

 Add w to Q

for i to n

 if i is unexplored

 BFS(G, i) source

- Complexity Analysis

We are using two arrays which are both traversed only once. Hence the program will run in linear time.

$$O(m+n)$$