* i = Arr[i] problem

- Here we will ignore -ve numbers and zero as they are out of the equation. Rest approach will be same, eliminate one of the halves to save on time complexity.

Recurrence Relation:

$$T(n) = T(n/2) + O(1)$$

$$a = 1,\ b = 2,\ d = 0$$

$$a = b^d \Rightarrow O(\log n)$$

Algorithm:

$$mid = \left\lfloor \frac{p+q}{2} \right\rfloor$$

A[mid] = mid
A[mid] < mid ⇒ L.dis
A[mid] > mid ⇒ R.dis

```
if (arr[mid] == mid)
  return mid
if (arr(mid < mid)
  return indequl (arr, mid+1, right)
else
  return indequl (arr, left, mid-1)
```