

ASSIGNMENT - 2

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DESIGN AND ANALYSIS
OF ALGORITHM

⑫ Demonstrate Binary Search method to search key=23 from the array $a[] = \{2, 5, 8, 12, 16, 23, 38, 56, 72, 91\}$

$$mid = \frac{l+h}{2}$$

$$a[mid] > key$$

$$h = mid - 1$$

$$a[mid] < key$$

$$l = mid + 1$$

$$l = 0$$

$$a[] = \{2, 5, 8, 12, 16, 23, 38, 56, 72, 91\}$$

$$h = 9$$

$$mid = \frac{0+9}{2} = \frac{9}{2} = 4.5$$

$$a[4] < key$$

~~mid~~ = l = 4 + 1

$$l = 5$$

$$mid = \frac{5+9}{2} = \frac{14}{2} = 7$$

$$a[7] = 56$$

$$56 > key$$

$$l = mid - 1$$

$$= 7 - 1$$

$$h = 6$$

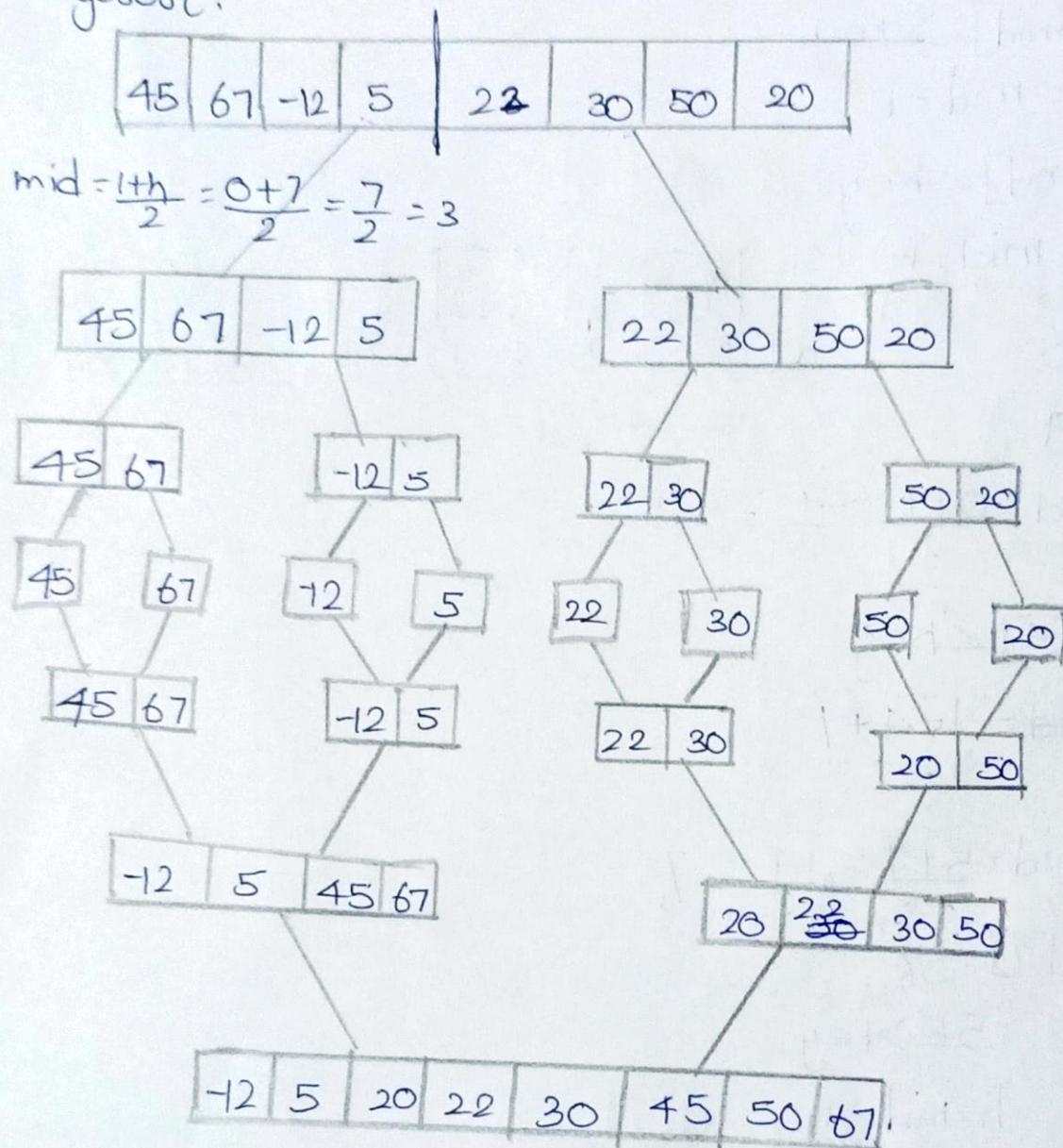
$$mid = \frac{5+6}{2} = \frac{11}{2} = 5.5 \text{ or } 5$$

$$a[5] = 23$$

$$a[mid] = key$$

$$\text{Index} = 5$$

(13) Apply mergesort and order ^{the} list of 8 elements. Data $d = (45, 67, -12, 5, 22, 30, 50, 20)$. Setup a recurrence relation for the number of key comparison made by mergesort.



$T(n)$ is the number of comparisons

Recurrence Relation

$$T(n) = 2T\left(\frac{n}{2}\right) + n - 1$$

When $n=1$ no comparison needed to sort the array.

The list is divided into two halves so, $2T\left(\frac{n}{2}\right)$ and for merging $(n-1)$

⑭ Find the no. of times to perform swapping for selection sort. Also estimate the time complexity for the order of notation Set S(12, 7, 5, -2, 18, 6, 13, 4)

12	7	5	-2	18	6	13	4
----	---	---	----	----	---	----	---

min = -2

Swap 12 and -2

-2	7	5	12	18	6	13	4
----	---	---	----	----	---	----	---

min = 4

Swap 7 and 4

-2	4	5	12	18	6	13	7
----	---	---	----	----	---	----	---

min = 5

No swap

-2	4	5	12	18	6	13	7
----	---	---	----	----	---	----	---

min = 6

Swap 12 and 6

-2	4	5	6	18	12	13	7
----	---	---	---	----	----	----	---

min = 7

Swap 18 and 7

-2	4	5	6	7	12	13	18
----	---	---	---	---	----	----	----

Sorted Array

Total Comparisons

$$(n-1) + (n-2) + \dots + 1 = \frac{n(n-1)}{2} = O(n^2)$$

n=8 Swaps=4 times

Time complexity $O(n^2)$

- (15) Find the index of the target value 10 using binary search from the following list of elements
[2, 4, 6, 8, 10, 12, 14, 16, 18, 20]

$$a = [2, 4, 6, 8, 10, 12, 14, 16, 18, 20]$$

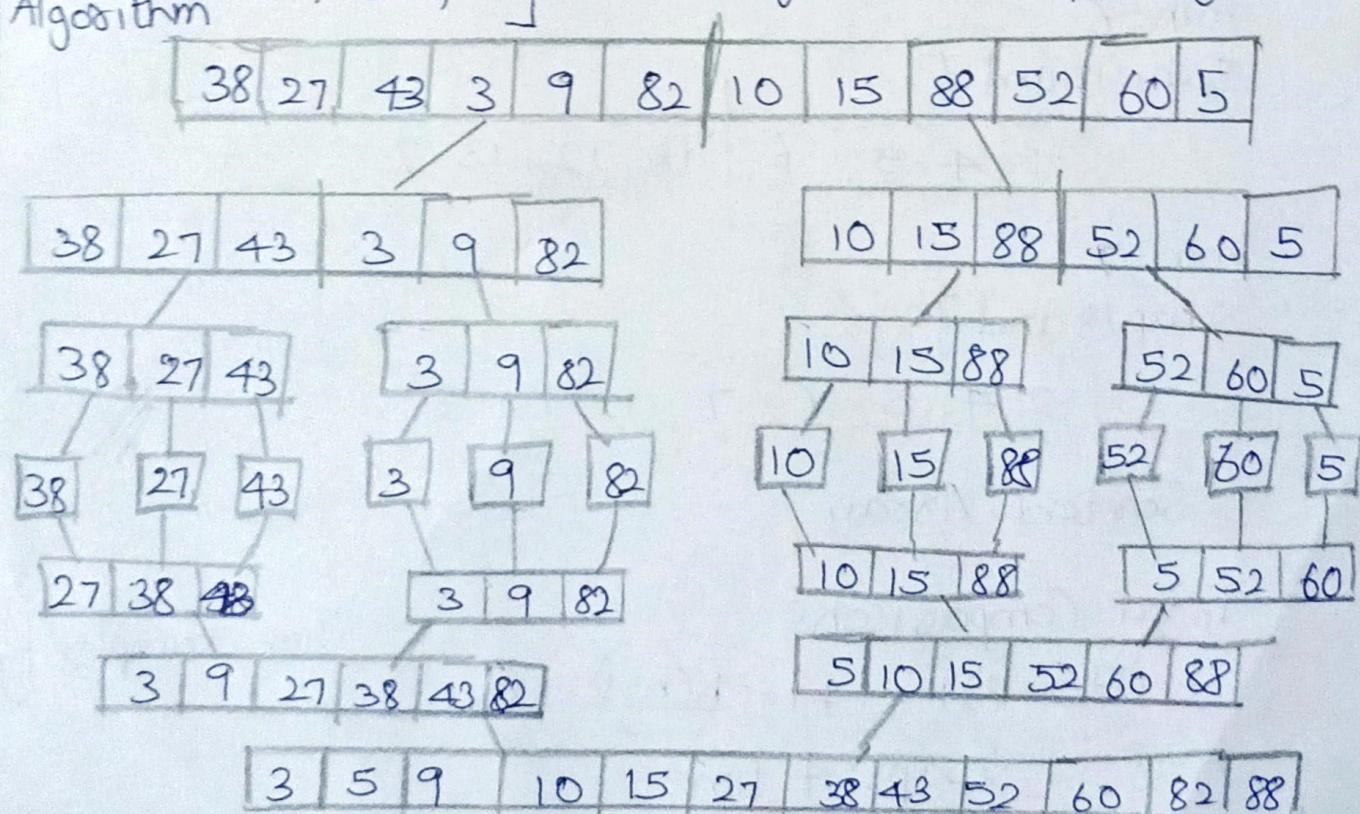
$$\text{mid} = \frac{0+9}{2} = \frac{9}{2} = 4.5 \text{ or } 4$$

$$a[\text{mid}] = 10$$

$$a[\text{mid}] = \text{Target}$$

$$\text{Index} = 4$$

- (16) Sort the following elements using merge sort divide and conquer strategy [38, 27, 43, 3, 9, 82, 10, 15, 88, 52, 60, 5] and analyze the complexity of Algorithm



⑯ Sort the array 64, 34, 25, 12, 22, 11, 90 using bubble sort

64	34	25	12	22	11	90
----	----	----	----	----	----	----

Best Case: $O(n^2)$

34	64	25	12	22	11	90
----	----	----	----	----	----	----

Average Case: $O(n^2)$

34	25	64	12	22	11	90
----	----	----	----	----	----	----

Worst Case: $O(n^2)$

Pass I

34	25	12	64	22	11	90
----	----	----	----	----	----	----

34	25	12	22	64	11	90
----	----	----	----	----	----	----

34	25	12	22	11	64	90
----	----	----	----	----	----	----

25	34	12	22	11	64	90
----	----	----	----	----	----	----

Pass II

25	12	34	22	11	64	90
----	----	----	----	----	----	----

25	12	22	34	11	64	90
----	----	----	----	----	----	----

25	12	22	11	34	64	90
----	----	----	----	----	----	----

12	25	22	11	34	64	90
----	----	----	----	----	----	----

Pass III

12	22	25	11	34	64	90
----	----	----	----	----	----	----

12	22	11	25	34	64	90
----	----	----	----	----	----	----

Pass IV

12	22	11	25	34	64	90
----	----	----	----	----	----	----

12	11	22	25	34	64	90
----	----	----	----	----	----	----

Pass V

11	12	22	25	34	64	90
----	----	----	----	----	----	----

⑯ Sort the array 64, 34, 25, 12, 22, 11 using Selection Sort

64	34	25	12	22	11
----	----	----	----	----	----

min = 11

Swap 64 and 11

11	34	25	12	22	64
----	----	----	----	----	----

min = 12

Swap 34 and 12

11	12	25	34	22	64
----	----	----	----	----	----

min = 22

Swap 25 and 22

11	12	22	34	25	64
----	----	----	----	----	----

min = 25

Swap 34 and 25

11	12	22	25	34	64
----	----	----	----	----	----

Sorted Array

Time Complexity

Best Case: $O(n^2)$

Average Case: $O(n^2)$

Worst Case: $O(n^2)$

⑯ Given an array of [4, -2, 5, 3, 10, -5, 2, 8, -3, 6, 7, -4, 1, 9, -1, 0, -6, -8, 11, -9] integers sort the following elements using insertion sort using brute force approach.

4	-2	5	3	10	-5	2	8	3	6	7	-4	1	9	-1	0	-6	-8	11	-9
---	----	---	---	----	----	---	---	---	---	---	----	---	---	----	---	----	----	----	----

key = 4

-2	4	5	3	10	-5	2	8	-3	6	7	-4	1	9	-1	0	-6	-8	11	-9
----	---	---	---	----	----	---	---	----	---	---	----	---	---	----	---	----	----	----	----

key = 5

-2	4	5	3	10	-5	2	8	-3	6	7	-4	1	9	-1	0	-6	-8	11	-9
----	---	---	---	----	----	---	---	----	---	---	----	---	---	----	---	----	----	----	----

key = 3

-2	3	4	5	10	-5	2	8	-3	6	7	-4	1	9	-1	0	-6	-8	11	-9
----	---	---	---	----	----	---	---	----	---	---	----	---	---	----	---	----	----	----	----

key = -5

1	5	-2	3	4	5	10	2	8	-3	6	7	-4	1	9	-1	0	-6	-8	11	-9
---	---	----	---	---	---	----	---	---	----	---	---	----	---	---	----	---	----	----	----	----

key = 2

-5	-2	2	3	4	5	10	8	-3	6	7	-4	1	9	-1	0	-6	-8	11	-9
----	----	---	---	---	---	----	---	----	---	---	----	---	---	----	---	----	----	----	----

key = 8

-5	-2	2	3	4	5	8	10	-3	6	7	-4	1	9	-1	0	-6	-8	11	-9
----	----	---	---	---	---	---	----	----	---	---	----	---	---	----	---	----	----	----	----

key = -3

-5	-3	-2	2	3	4	5	8	10	6	7	-4	1	9	-1	0	-6	-8	11	-9
----	----	----	---	---	---	---	---	----	---	---	----	---	---	----	---	----	----	----	----

key = 6

-5	-3	-2	2	3	4	5	6	8	10	7	-4	1	9	-1	0	-6	-8	11	-9
----	----	----	---	---	---	---	---	---	----	---	----	---	---	----	---	----	----	----	----

key = 7

-5	-3	-2	2	3	4	5	6	7	8	10	-4	1	9	-1	0	-6	-8	11	-9
----	----	----	---	---	---	---	---	---	---	----	----	---	---	----	---	----	----	----	----

key = -4

3	-5	-4	-3	-2	2	4	5	6	7	8	10	1	9	-1	0	-6	-8	11	-9
---	----	----	----	----	---	---	---	---	---	---	----	---	---	----	---	----	----	----	----

key = 1

- ② Sort the following elements using insertion sort using brute force approach strategy [38, 27, 43, 3, 9, 82, 10, 15, 88, 52, 60, 5]

38	27	43	3	9	82	10	15	88	52	60	5
----	----	----	---	---	----	----	----	----	----	----	---

Key = 38

27	38	43	3	9	82	10	15	88	52	60	5
----	----	----	---	---	----	----	----	----	----	----	---

Key = 43

3	27	38	43	9	82	10	15	88	52	60	5
---	----	----	----	---	----	----	----	----	----	----	---

Key = 9

3	9	27	38	43	82	10	15	88	52	60	5
---	---	----	----	----	----	----	----	----	----	----	---

Key = 10

3	9	10	27	38	43	82	15	88	52	60	5
---	---	----	----	----	----	----	----	----	----	----	---

key = 15

3	9	10	15	27	38	43	82	88	52	60	5
---	---	----	----	----	----	----	----	----	----	----	---

key = 52

3	9	10	15	27	38	43	52	82	88	60	5
---	---	----	----	----	----	----	----	----	----	----	---

key = 60

3	9	10	15	27	38	43	52	60	82	88	5
---	---	----	----	----	----	----	----	----	----	----	---

key = 5

3	5	9	10	15	27	38	43	52	60	82	8
---	---	---	----	----	----	----	----	----	----	----	---

Sorted

Best case: $O(n)$ Worst Case: $O(n^2)$

Average Case: $O(n^2)$

3	-5	4	3	-2	1	2	4	5	6	7	8	10	9	-1	0	-3	11	9	
key=9	-5	4	3	-2	1	2	3	4	5	6	7	8	9	10	-1	0	-6	11	9
key=-1	-5	4	3	-2	-1	1	2	3	4	5	6	7	8	9	10	0	-6	11	9
key=0	-5	4	3	-2	-1	1	2	3	4	5	6	7	8	9	10	0	-6	11	9
key=-6	-5	4	3	-2	-1	0	1	2	3	4	5	6	7	8	9	10	-7	11	9
key=-8	-6	-5	4	-3	-2	-1	0	1	2	3	4	5	6	7	8	9	10	-9	11
key=-9	-8	-6	-5	4	-3	-2	-1	0	1	2	3	4	5	6	7	8	9	10	11

Sorted

Best Case: $O(n)$

Average Case: $O(n^2)$

Worst Case: $O(n^2)$

⑩ Solve the recurrence relation

$$T(n) = 4T\left(\frac{n}{2}\right) + n^2, T(1) = 1 \quad \text{Case 2: } \log_b a = k$$

$$T(n) = aT\left(\frac{n}{b}\right) + f(n) \quad p=1 \quad p > -1 \quad \Theta(n^k \log^{p-1} n)$$

$$a=4 \quad b=2 \quad k=2$$

$$\log_b a = \log_2 4 = 2 \quad \Theta(n^2 \log^{1-1} n)$$

$$\Theta(n^2 \cdot \log n)$$

The order of growth solution to the recurrence relation is $\Theta(n^2 \log n)$

① If $t_1(n) \in O(g_1(n))$ and $t_2(n) \in O(g_2(n))$ prove that $t_1(n)+t_2(n) \in O(\max\{g_1(n), g_2(n)\})$

for any four arbitrary real numbers a, b, a_1, b_1
such that $a_1 \leq b_1$ and $a_2 \leq b_2$

$$\text{we have } a_1 + a_2 \leq 2 \max\{b_1, b_2\}$$

since $t_1(n) \in O(g_1(n))$, then there exists some constant c_1 and non-negative integer n_1 such that

$$t_1(n) \leq c_1 g_1(n) \text{ for all } n \geq n_1$$

since $t_2(n) \in O(g_2(n))$, then there exists integer n_2 such that

$$t_2(n) \leq c_2 g_2(n) \text{ for all } n \geq n_2$$

Let $c_3 = \max\{c_1, c_2\}$ and $n_0 = \max\{n, n_2\}$

$$\begin{aligned} t_1(n) + t_2(n) &\leq c_1 g_1(n) + c_2 g_2(n) \\ &\leq c_3 g_1(n) + c_3 g_2(n) \\ &= c_3 \{g_1(n) + g_2(n)\} \\ &\leq 2c_3 \max\{g_1(n), g_2(n)\} \end{aligned}$$

Hence, $t_1(n) + t_2(n) \in O(\max\{g_1(n), g_2(n)\})$

constants c_3 and n_0

② Find the time complexity of the below recurrence equation

$$③ T(n) = \begin{cases} 2T\left(\frac{n}{2}\right) + 1 & \text{if } n > 1 \\ T(n) = aT\left(\frac{n}{b}\right) + f(n) & \text{otherwise} \end{cases}$$

$$a = 2$$

$$b = 2$$

$$\log_b^a = \log_2^2 = 1$$

$$k = 0$$

$$\log_b^a > k$$

Case: 1

$$\Theta(n \cdot \log_b^a)$$

$$\Theta(n \cdot \log_2^2)$$

$$\Theta(n \cdot 1)$$

$$\Theta(n)$$

$$④ T(n) = \begin{cases} 2T(n-1) & \text{if } n > 0 \\ 1 & \text{otherwise} \end{cases}$$
$$T(n) = 2T(n-1) \rightarrow ① \quad T(0) = 0$$

$$h = n-1$$

$$T(n-1) = 2T(n-1-1) \\ = 2T(n-2) \rightarrow ②$$

Sub ② in ①

$$T(n) = 2[2T(n-2)] \\ = 2^2 T(n-2) \rightarrow ③$$

$$n=n-2$$

$$T(n-2) = 2T(n-2)-1$$

$$= 2T(n-3) \rightarrow ④$$

Sub ④ in ③

$$T(n) = 2^2 [2T(n-3)]$$

$$T(n) = 2^3 T(n-3) \rightarrow ⑤$$

$$n=n-3$$

$$T(n-3) = 2T(n-3)-1$$

$$T(n-3) = 2T(n-4) \rightarrow ⑥$$

Sub ⑥ in ⑤

$$T(n) = 2^3 [2T(n-4)]$$

$$= 2^4 T(n-4) \rightarrow ⑦$$

$$T(n) = 2^k T(n-k)$$

$$n-k=0 \Rightarrow n=k$$

$$\text{if } T(0)=1$$

$$T(n) = 2^k \cdot T(0)$$

$$T(n) = 2^k \cdot 1$$

$$T(n) = 2^k$$

$$\therefore n=k$$

$$T(n) = O(2^n)$$

⑤ Big O Notation: show that $f(n)=n^2+3n+5$ is $O(n^2)$

To prove that $f(n)=n^2+3n+5$ is $O(n^2)$ we need to find constants c and n_0 such that $f(n) \leq c \cdot n^2$ for all $n \geq n_0$

$$f(n) = n^2 + 3n + 5$$

For $n \geq 1, n^2 \geq n \dots$ so on

$$f(n) = n^2 + 3n + 5 \leq n^2 + 3n^2 + 5n^2$$

$$f(n) = n^2 + 3n + 5 \leq 9n^2 \text{ for } n \geq 1$$

so, for $c=9$ and $n_0=1$

$$f(n) \leq c \cdot n^2 \text{ for all } n \geq n_0$$

that proves $f(n)$ is $O(n^2)$

- ⑥ Big omega notation: prove that $g(n) = n^3 + 2n^2 + 4n$ is $\Omega(n^3)$

To prove that $g(n) = n^3 + 2n^2 + 4n$ is $\Omega(n^3)$
we need to find constants c and n_0 such that

$$g(n) \geq c \cdot n^3 \text{ for all } n \geq n_0$$

$$g(n) = n^3 + 2n^2 + 4n$$

For $n \geq 1$,

$$g(n) = n^3 + 2n^2 + 4n \geq n^3$$

since $2n^2$ and $4n$ are both less than n^3 when $n \geq 1$

so, for $c=1$ and $n_0=1$

$$g(n) \geq c \cdot n^3 \text{ for all } n \geq n_0$$

that proves $g(n)$ is $\Omega(n^3)$

⑦ Big Theta notation: Determine whether $h(n) = 4n^2 + 3n$ is $\Theta(n^2)$ or not

1. $h(n) = 4n^2 + 3n$ is $\Theta(n^2)$

For $n \geq 1$, $h(n) \leq 4n^2 + 3n^2$

(Since $3n$ is less than n^2 when $n \geq 1$)

For this simplifies to $h(n) \leq 7n^2$
for $n \geq 1$

Therefore, $h(n)$ is $O(n^2)$

2. $h(n) = 4n^2 + 3n$ is $\Omega(n^2)$

For $n \geq 1$, $h(n) \geq 4n^2$

(Since $3n$ is positive)

Therefore $h(n)$ is $\Omega(n^2)$

Since $h(n)$ is both $O(n^2)$ and $\Omega(n^2)$, it is $\Theta(n^2)$

⑧ Let $f(n) = n^3 - 2n^2 + n$ and $g(n) = -n^2$ show whether $f(n) = -n^2 g(n)$ is true or false and justify your answer

$n=1$

$$\begin{aligned} f(1) &= 1^3 - 2(1)^2 + 1 & g(1) &= (-1)^2 \\ &= 1 - 2 + 1 & &= 1 \\ &= 0 & &= 1 \end{aligned}$$

$n=2$

$$\begin{aligned} f(2) &= 2^3 - 2(2)^2 + 2 & g(2) &= (-2)^2 \\ &= 8 - 8 + 2 & &= 4 \\ &= 2 & & \end{aligned}$$

$n=3$

$$\begin{aligned} f(3) &= 3^3 - 2(3)^2 + 3 & g(3) &= (-3)^2 \\ &= 27 - 18 + 3 & &= 9 \end{aligned}$$

$n=4$

$$\begin{aligned}f(4) &= 4^3 - 2(4)^2 + 4 \\&= 64 - 32 + 4 \\&= 32 + 4 \\&= 36\end{aligned}$$

$$\begin{aligned}g(4) &= (-4)^2 \\&= 16\end{aligned}$$

$n=5$

$$\begin{aligned}f(5) &= 5^3 - 2(5)^2 + 5 \\&= 125 - 50 + 5 \\&= 85\end{aligned}$$

$$\begin{aligned}g(5) &= (-5)^2 \\&= 25\end{aligned}$$

$$f(n) \geq g(n)$$

So it is best case according to asymptotic notations

$$f(n) = \Omega(g(n))$$

- ⑨ Determine whether $h(n) = n \log n + n$ is in $\Theta(n \log n)$
prove a rigorous proof for your conclusion

1. Upper Bound (O notation)

we need to find c_1 and n_0 such that

$$h(n) \leq c_1 \cdot n \log n \text{ for all } n \geq n_0$$

$$h(n) = n \log n + n$$

$$\leq n \log n + n \log n \text{ (since } \log n \text{ is increasing)}$$

$$= 2n \log n$$

now, let $c_1 = 2$, then $h(n) \leq 2n \log n$ for all $n \geq 1$

So, $h(n)$ is $O(n \log n)$

2. Lower Bound (Ω -notation)

We need to find c_2 and n_0 such that

$$h(n) \geq c_2 \cdot n \log n \text{ for all } n \geq n_0$$

$$h(n) = n \log n + n$$

$$\geq \frac{1}{2} \cdot n \log n \text{ (for } n \geq 2\text{)}$$

- ⑪ Given an array of $[4, -2, 5, 3, 10, -5, 2, 8, -3, 6, 7, 4, 18, 9, -1, 0, -6, -8, 11, -9]$ integers. Find the maximum and minimum product that can be obtained by multiplying 2 integers.

Sort

$$a: [-9, -8, -6, -5, -4, -3, -2, -1, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11]$$

Maximum product

$$\Rightarrow 10 \times 11 = 110$$

$$\Rightarrow -9 \times -8 = 72$$

Minimum product

$$\Rightarrow -9 \times -8 = 72$$

$$\Rightarrow -9 \times 11 = -99$$

$$\Rightarrow -8 \times 11 = -88$$

Maximum product = 110

Minimum product = -99