



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

Name:	SHREYA GANESH WANKHEDE
Roll No:	62
Class/Sem:	SE/IV
Experiment No.:	6
Title:	Prim's Algorithm.
Date of Performance:	
Date of Submission:	
Marks:	
Sign of Faculty:	



Experiment No. 6

Title: Prim's Algorithm.

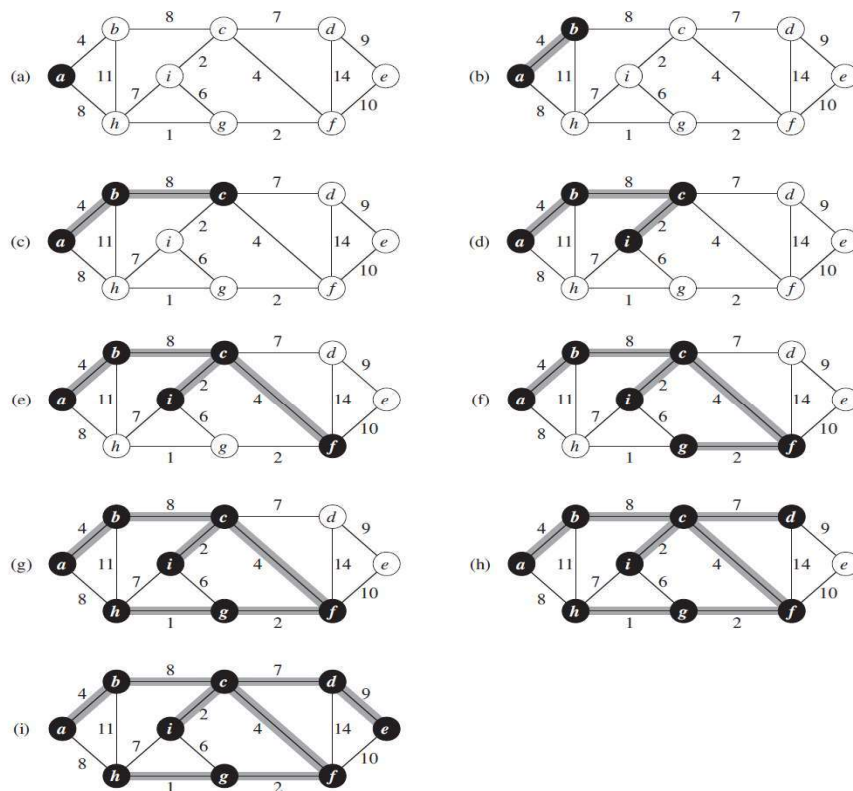
Aim: To study and implement Prim's Minimum Cost Spanning Tree Algorithm.

Objective: To introduce Greedy based algorithms

Theory:

Prim's algorithm is a greedy algorithm that finds a minimum spanning tree for a weighted undirected graph. This means it finds a subset of the edges that forms a tree that includes every vertex, where the total weight of all the edges in the tree is minimized. The algorithm operates by building this tree one vertex at a time, from an arbitrary starting vertex, at each step adding the cheapest possible connection from the tree to another vertex.

Example:





Algorithm and Complexity:

```
1  Algorithm Prim(E, cost, n, t)
2  // E is the set of edges in G. cost[1 : n, 1 : n] is the cost
3  // adjacency matrix of an n vertex graph such that cost[i, j] is
4  // either a positive real number or  $\infty$  if no edge (i, j) exists.
5  // A minimum spanning tree is computed and stored as a set of
6  // edges in the array t[1 : n - 1, 1 : 2]. (t[i, 1], t[i, 2]) is an edge in
7  // the minimum-cost spanning tree. The final cost is returned.
8  {
9      Let (k, l) be an edge of minimum cost in E;
10     mincost := cost[k, l];
11     t[1, 1] := k; t[1, 2] := l;
12     for i := 1 to n do // Initialize near.
13         if (cost[i, l] < cost[i, k]) then near[i] := l;
14         else near[i] := k;
15     near[k] := near[l] := 0;
16     for i := 2 to n - 1 do
17     { // Find n - 2 additional edges for t.
18         Let j be an index such that near[j]  $\neq$  0 and
19         cost[j, near[j]] is minimum;
20         t[i, 1] := j; t[i, 2] := near[j];
21         mincost := mincost + cost[j, near[j]];
22         near[j] := 0;
23         for k := 1 to n do // Update near[ ].
24             if ((near[k]  $\neq$  0) and (cost[k, near[k]] > cost[k, j]))
25                 then near[k] := j;
26     }
27     return mincost;
28 }
```

Time Complexity is $O(n^2)$, Where, *n* = number of vertices **Theory:**

Implementation:

```
// A C program for Prim's Minimum
// Spanning Tree (MST) algorithm. The program is
// for adjacency matrix representation of the graph
#include <limits.h>
#include <stdbool.h>
#include <stdio.h>
// Number of vertices in the graph
#define V 5
// A utility function to find the vertex with
// minimum key value, from the set of vertices
// not yet included in MST
int minKey(int key[], bool mstSet[])
```



```
{
// Initialize min value
int min = INT_MAX, min_index;

for (int v = 0; v < V; v++)
    if (mstSet[v] == false && key[v] < min)
        min = key[v], min_index = v;
return min_index;
}

// A utility function to print the
// constructed MST stored in parent[]
int printMST(int parent[], int graph[V][V])
{
    printf("Edge \tWeight\n");
    for (int i = 1; i < V; i++)
        printf("%d - %d \t%d \n", parent[i], i,
            graph[i][parent[i]]);
}

// Function to construct and print MST for
// a graph represented using adjacency
// matrix representation
void primMST(int graph[V][V])
{
    // Array to store constructed MST
    int parent[V];
    // Key values used to pick minimum weight edge in cut
    int key[V];
    // To represent set of vertices included in MST
    bool mstSet[V];

    // Initialize all keys as INFINITE
    for (int i = 0; i < V; i++)
        key[i] = INT_MAX, mstSet[i] = false;

    // Always include first 1st vertex in MST.
    // Make key 0 so that this vertex is picked as first
    // vertex.
    key[0] = 0;
    // First node is always root of MST
    parent[0] = -1;
    // The MST will have V vertices
    for (int count = 0; count < V - 1; count++) {
```



```
// Pick the minimum key vertex from the
// set of vertices not yet included in MST
int u = minKey(key, mstSet);
// Add the picked vertex to the MST Set
mstSet[u] = true;
// Update key value and parent index of
// the adjacent vertices of the picked vertex.
// Consider only those vertices which are not
// yet included in MST
for (int v = 0; v < V; v++)
    // graph[u][v] is non zero only for adjacent
    // vertices of m mstSet[v] is false for vertices
    // not yet included in MST Update the key only
    // if graph[u][v] is smaller than key[v]
    if (graph[u][v] && mstSet[v] == false
        && graph[u][v] < key[v])
        parent[v] = u, key[v] = graph[u][v];
}
// print the constructed MST
printMST(parent, graph);
}
// Driver's code
int main()
{
    int graph[V][V] = { { 0, 2, 0, 6, 0 },
                        { 2, 0, 3, 8, 5 },
                        { 0, 3, 0, 0, 7 },
                        { 6, 8, 0, 0, 9 },
                        { 0, 5, 7, 9, 0 } };

    // Print the solution
    primMST(graph);

    return 0;
}
```



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

Output:

Edge	Weight
0 - 1	2
1 - 2	3
0 - 3	6
1 - 4	5

Conclusion: Implementing Prim's algorithm has proven to be effective in generating minimum spanning trees, efficiently connecting all nodes in a graph while minimizing total edge weight. This experiment underscores the algorithm's practical applicability in optimizing network connectivity, demonstrating its importance in various real-world scenarios.