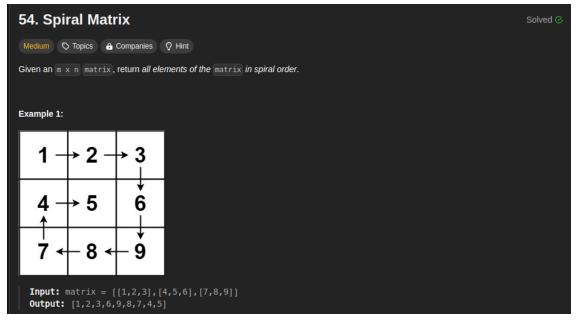


```
Int rec(vector<vector<int>> &dp,vector<vector<int>> &visited,vector<vector<int>> & grid,int i,int j){
    //cout<<ir/>
    //cout<<ir/>
    if(i==grid.size()-1 and j==grid[0].size()-1){
        //cout<<ir/>
        //cout<<ir/>
        //cout<<ir/>
        return grid[i][j];
    }
    if(dp[i][j]!=-1)return dp[i][j];
    int down=INT_MAX,right=INT_MAX;
    if(j+1<grid[0].size() and visited[i][j+1]==-1){
        visited[i][j+1]=0;
        right=rec(dp,visited,grid,i,j+1);
        visited[i][j+1]=-1;
        //cout<<"right" <<ri>
        //cout<<"right" <</ri>
        if(i+1<grid.size() and visited[i+1][j]==-1){
            visited[i+1][j]=0;
            down=rec(dp,visited,grid,i+1,j);
            visited[i+1][j]=-1;
            //cout<<"down" <<down<<endl;
        }
        return dp[i][j]=grid[i][j]+min(right,down);
}
int minPathSum(vector<vector<int>>& grid) {
        vector<vector<int>> visited(grid.size(),vector<int> (grid[0].size(),-1));
        vector<vector<int>> bq(grid.size(),vector<int> (grid[0].size(),-1));
        return rec(dp,visited,grid,0,0);
}
```

TimeComplexity:O(2\*\*n)
SpaceComplexity:O(m\*n)



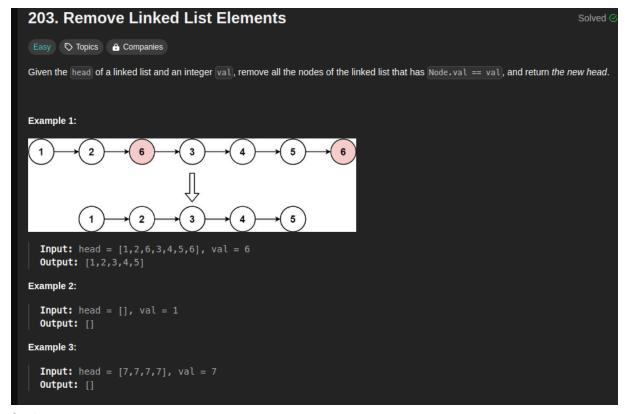
```
vector<int> spiralOrder(vector<vector<int>>& matrix) {
        int left = 0, right = matrix[0].size() - 1, top = 0,
        bottom = matrix.size() - 1;
        vector<int> ans;
        while (left <= right and top <= bottom) {</pre>
            for (int i = left; i <= right; i++) {</pre>
                 ans.push_back(matrix[top][i]);
            top++;
            for (int i = top; i <= bottom; i++) {</pre>
                ans.push_back(matrix[i][right]);
            right--;
            if (top <= bottom) {</pre>
                for (int i = right; i >= left; i--) {
                     ans.push back(matrix[bottom][i]);
                bottom--;
            if (left <= right) {</pre>
                 for (int i = bottom; i >= top; i--) {
                     ans.push_back(matrix[i][left]);
                 left++:
            }
        return ans;
```

TimeComplexity:O(m\*n)
SpaceComplexity:O(m\*n)

```
3. Longest Substring Without Repeating Characters
                                                                                                     Solved ©
Medium ♥ Topics ♠ Companies ♥ Hint
Given a string s, find the length of the longest substring without repeating characters.
Example 1:
  Input: s = "abcabcbb"
  Output: 3
  Explanation: The answer is "abc", with the length of 3.
  Input: s = "bbbbb"
  Output: 1
  Explanation: The answer is "b", with the length of 1.
Example 3:
  Input: s = "pwwkew"
  Output: 3
  Explanation: The answer is "wke", with the length of 3.
 Notice that the answer must be a substring, "pwke" is a subsequence and not a substring.
```

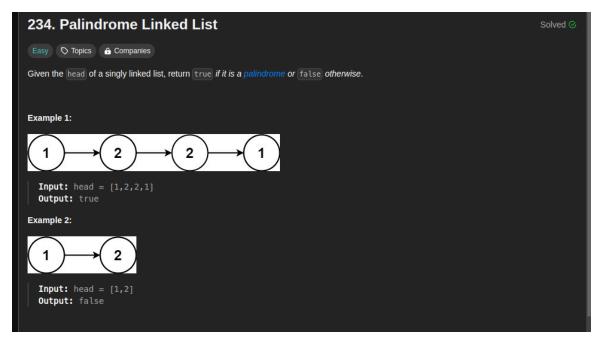
```
int lengthOfLongestSubstring(string s) {
    int maxi=0;
    unordered_map<char,int> hash;
    int left=0;
    for(int i=0;i<s.size();i++){
        hash[s[i]]++;
        while(hash[s[i]]>1){
            hash[s[left]]--;
            left++;
        }
        maxi=max(maxi,i-left+1);
    }
    return maxi;
}
```

TimeComplexity:O(n)
SpaceComplexity:O(n)



```
ListNode* removeElements(ListNode* head, int val) {
    if(head==NULL)return NULL;
    while(head!=NULL and head->val==val){
        head=head->next;
    }
    ListNode *ptr=head,*preptr=head;
    while(ptr){
        if(ptr->val==val){
            preptr->next;
            ptr=preptr->next;
        }
     else{
        preptr=ptr;
        ptr=ptr->next;
    }
}
return head;
}
```

TimeComplexity:O(n)
SpaceComplexity:O(1)



```
ListNode* reverse(ListNode* head){
        ListNode *ptr=head,*preptr=NULL;
        while(ptr!=NULL){
            ListNode *forward=ptr->next;
            ptr->next=preptr;
            preptr=ptr;
            ptr=forward;
        return preptr;
bool isPalindrome(ListNode* head) {
        if(head==NULL or head->next==NULL)return true;
        ListNode *slow=head,*fast=head;
        while(fast->next!=NULL && fast->next!=NULL){
            slow=slow->next;
            fast=fast->next->next;
        slow=reverse(slow);
        fast=head;
        while(slow and fast){
            if(fast->val!=slow->val)return false;
            slow=slow->next;
            fast=fast->next;
        return true;
```

TimeComplexity:O(n)
SpaceComplexity:O(1)

```
31. Next Permutation
                                                                                                                                   Solved ©
Medium ♥ Topics ♠ Companies
A permutation of an array of integers is an arrangement of its members into a sequence or linear order.
• For example, for arr = [1,2,3], the following are all the permutations of arr: [1,2,3], [1,3,2], [2, 1, 3], [2, 3, 1], [3,1,2],
The next permutation of an array of integers is the next lexicographically greater permutation of its integer. More formally, if all the permutations of
the array are sorted in one container according to their lexicographical order, then the next permutation of that array is the permutation that follows
it in the sorted container. If such arrangement is not possible, the array must be rearranged as the lowest possible order (i.e., sorted in ascending
• For example, the next permutation of arr = [1,2,3] is [1,3,2].
• Similarly, the next permutation of arr = [2,3,1] is [3,1,2].
• While the next permutation of arr = [3,2,1] is [1,2,3] because [3,2,1] does not have a lexicographical larger rearrangement.
Given an array of integers nums, find the next permutation of nums.
The replacement must be in place and use only constant extra memory.
Example 1:
  Input: nums = [1,2,3]
  Output: [1,3,2]
```

```
void nextPermutation(vector<int>& nums) {
int index=-1;
        for(int i=nums.size()-1;i>0;i--){
            if(nums[i]>nums[i-1]){
                index=i-1;
                break:
            }
        if(index==-1){
            reverse(nums.begin(),nums.end());
            return ;
        for(int i=nums.size()-1;i>=index;i--){
            if(nums[index]<nums[i]){</pre>
                swap(nums[index],nums[i]);
                break;
            }
        reverse(nums.begin()+index+1,nums.end());
```

TimeComplexity:O(n)
SpaceComplexity:O(1)

```
127. Word Ladder

Hard ○ Topics ♠ Companies

A transformation sequence from word beginWord to word endword using a dictionary wordList is a sequence of words beginWord → s₁ → s₂ → ... → sk such that:

• Every adjacent pair of words differs by a single letter.

• Every si for 1 ← i ← k is in wordList. Note that beginWord does not need to be in wordList.

• sk = endWord

Given two words, beginWord and endword, and a dictionary wordList, return the number of words in the shortest transformation sequence from beginWord to endword, or 0 if no such sequence exists.

Example 1:

Input: beginWord = "hit", endWord = "cog", wordList = ["hot", "dot", "dog", "lot", "log", "cog"] Output: 5

Explanation: One shortest transformation sequence is "hit" -> "hot" -> "dot" -> "dog" -> cog", which is 5 words long.
```

```
.nt wordLadderLength(string startWord, string targetWord,vector<string> &wordList)
       queue<pair<string, int>> q;
       q.push({startWord, 1});
       unordered_set<string> set(wordList.begin(), wordList.end());
       set.erase(startWord);
       while (!q.empty())
           string temp = q.front().first;
           int count = q.front().second;
           q.pop();
           if (temp == targetWord)
               return count;
           for (int i = 0; i < temp.size(); i++)</pre>
               char original = temp[i];
               for (char ch = 'a'; ch <= 'z'; ch++)</pre>
                    temp[i] = ch;
                   if (set.find(temp) != set.end())
                       set.erase(temp);
                        q.push({temp, count + 1});
               temp[i] = original;
       return 0;
```

TimeComplexity:O(n\*\*2)
SpaceComplexity:O(n\*\*2)