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Anagram 🗌
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Given two strings s1 and s2 consisting of lowercase characters. The task is to check whether two given strings are an
anagram of each other or not. An anagram of a string is another string that contains the same characters, only the
order of characters can be different. For example, act and tac are an anagram of each other. Strings s1 and s2 can
only contain lowercase alphabets.
Note: You can assume both the strings s1 & s2 are non-empty.
Examples:
 Input: s1 = "geeks", s2 = "kseeg"
 Output: true
 Explanation: Both the string have same characters with same frequency. So, they are anagrams.
 Input: s1 = "allergy", s2 = "allergic"
 Output: false
 Explanation: Characters in both the strings are not same, so they are not anagrams.
 Input: s1 = "g", s2 = "g"
 Output: true
 Explanation: Character in both the strings are same, so they are anagrams.
Constraints:
1 \le s1.size(), s2.size() \le 10^5
```

```
bool areAnagrams(string& s1, string& s2) {
    if(s1.size()!=s2.size())return false;
    vector<int> arr1(26,0),arr2(26,0);
    for(int i=0;i<s1.size();i++){
        arr1[s1[i]-'a']++;
        arr2[s2[i]-'a']++;
    }
    for(int i=0;i<26;i++){
        if(arr1[i]!=arr2[i])return false;
    }
    return true;
}</pre>
```

TimeComplexity:O(n)
SpaceComplexity:O(26)

```
Longest consecutive subsequence \square
Difficulty: Medium
Given an array arr of non-negative integers. Find the length of the longest sub-sequence such that elements in the
subsequence are consecutive integers, the consecutive numbers can be in any order.
Examples:
 Input: arr[] = [2, 6, 1, 9, 4, 5, 3]
 Output: 6
 Explanation: The consecutive numbers here are 1, 2, 3, 4, 5, 6. These 6 numbers form the longest
  Input: arr[] = [1, 9, 3, 10, 4, 20, 2]
  Output: 4
  Explanation: 1, 2, 3, 4 is the longest consecutive subsequence.
 Input: arr[] = [15, 13, 12, 14, 11, 10, 9]
 Explanation: The longest consecutive subsequence is 9, 10, 11, 12, 13, 14, 15, which has a length of 7.
Constraints:
1 \le \arcsin() \le 10^5
0 <= arr[i] <= 10<sup>5</sup>
```

```
int findLongestConseqSubseq(vector<int>& arr) {
    unordered_set<int> st;
    for(auto a:arr){
        st.insert(a);
    }
    int maxi=1;
    for(int i=0;i<arr.size();i++){
        if(st.find(arr[i]-1)!=st.end()){
            continue;
        }
        int ans=1;
        int num=arr[i];
        while(st.find(num+1)!=st.end()){
            num++;
            ans++;
        }
        maxi=max(maxi,ans);
    }
    return maxi;
}</pre>
```

TimeComplexity:O(n)
SpaceComplexity:O(n)

```
int rowWithMax1s(vector<vector<int> > &arr) {
    for(int i=0;i<arr[0].size();i++){
        for(int j=0;j<arr.size();j++){
            if(arr[j][i]==1){
                return j;
            }
        }
    }
   return -1;
}</pre>
```

TimeComplexity:O(n*m)
SpaceComplexity:O(1)

```
Difficulty: Medium Accuracy: 23.2% Submissions: 306K+ Points: 4

Given a string s, your task is to find the longest palindromic substring within s. A substring is a contiguous sequence of characters within a string, defined as s[i...j] where 0 ≤ i ≤ j < len(s).

A palindrome is a string that reads the same forward and backward. More formally, s is a palindrome if reverse(s) == s.

Note: If there are multiple palindromes with the same length, return the first occurrence of the longest palindromic substring from left to right.

Examples:

Input: s = "aaaaabbaa"
Output: "aabbaa"
Explanation: The longest palindromic substring is "aabbaa".
```

```
string longestPalindrome(string s) {
        vector<vector<int>> arr(s.size(),vector<int>(s.size(),-1));
         int start=0,count=1;
        for(int i=0;i<s.size();i++){</pre>
             arr[i][i]=1;
         for(int i=0;i<s.size()-1;i++){</pre>
             if(s[i]==s[i+1]){
                 arr[i][i+1]=2;
                 start=i;
                 count=2;
         for(int i=3;i<=s.size();i++){</pre>
             for(int j=0;j<s.size()-i+1;j++){</pre>
                 int k=j+i-1;
                 if(arr[j+1][k-1]!=-1 \text{ and } s[k]==s[j]){}
                      arr[j][k]=arr[j+1][k-1]+2;
                      if(arr[j][k]>count){
                          count=arr[j][k];
                          start=j;
                 }
        string ans="";
        for(int i=start;i<start+count;i++){</pre>
             ans+=s[i];
        return ans;
```

TimeComplexity:O(n**2)
SpaceComplexity:O(n**2)

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Difficulty: Medium

Accuracy: 35.75

Submissions: 302K4

Points:

Consider a rat placed at (0,0) in a square matrix **mat** of order n^*n . It has to reach the destination at (n-1,n-1). Find all possible paths that the rat can take to reach from source to destination. The directions in which the rat can move are 'U'(up), 'D'(down), 'L' (left), 'R' (right). Value 0 at a cell in the matrix represents that it is blocked and rat cannot move to it while value 1 at a cell in the matrix represents that rat can be travel through it.

Note: In a path, no cell can be visited more than one time. If the source cell is 0, the rat cannot move to any other cell. In case of no path, return an empty list. The driver will output "-1" automatically.

Examples:

```
oid rec(int r,int c,int n, vector<string> &ans,vector<vector<int>> &visited,string move,vector<vector<int>> &mat){
        if(r==n-1 \text{ and } c==n-1)
            ans.push_back(move);
        if(r+1< n \text{ and } visited[r+1][c]==0 \text{ and } mat[r+1][c]==1){
            visited[r][c]=1;
            rec(r+1,c,n,ans,visited,move+'D',mat);
visited[r][c]=0;
        if(c-1)=0 and visited[r][c-1]==0 and mat[r][c-1]==1){}
            visited[r][c]=1;
            rec(r,c-1,n,ans,visited,move+'L',mat);
            visited[r][c]=0;
        if(c+1< n \text{ and } visited[r][c+1]==0 \text{ and } mat[r][c+1]==1){
            visited[r][c]=1;
            rec(r,c+1,n,ans,visited,move+'R',mat);
visited[r][c]=0;
        if(r-1)=0 and visited[r-1][c]==0 and mat[r-1][c]==1){
            visited[r][c]=1;
            rec(r-1,c,n,ans,visited,move+'U',mat);
visited[r][c]=0;
```

TimeComplexity:O(4**(n*m))
SpaceComplexity:O(n**2)