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Program Structures and Algorithms

Fall 2021

Assignment No. 3 – WQUPC

- Tasks in the assignment:

Part 1: Height-weighted Quick Union with Path Compression

1. Implemented height weighted quick union where we maintain an array to keep track of the tree height and while union add the shorter tree to the taller tree.
2. Path compression is done using single pass path-halving mechanism where we update root to that of the grandparent. This reduces the time taken by find().
3. Once a site is found we check if they're connected else union is performed.
4. Merge the components which contain site p with site q using strategy mentioned in point 1.

Part 2: Union Find Client

1. Developed a union find client which contains a main() method and a count() method.
2. Takes number of sites from the command line and number of times we wish to double the sites.
3. For each random pair of integer generated, a connection is established if it's not already present. The output is the number of connections till all sites are connected under a single component.
4. I've performed this operation 100 times for each site and taken an average of the connections generated.

Part 3: Conclusions and observations

1. The relationship between the number of objects (*n*) and the number of pairs (*m*) generated till all sites are connected.
2. Observations are mentioned below.

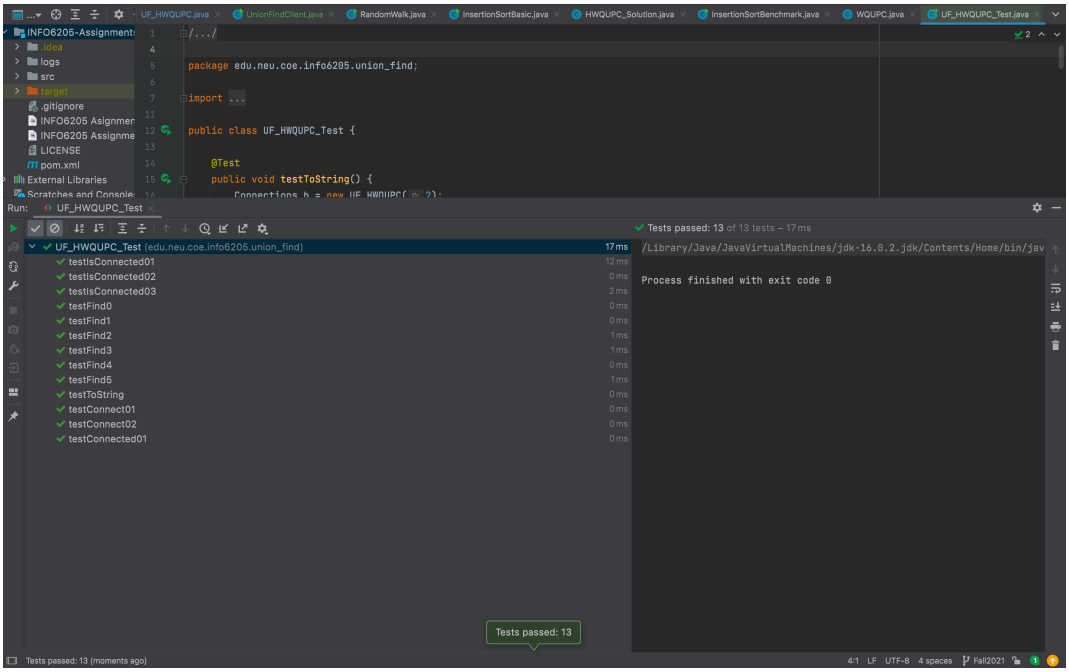
- Relationship Conclusion

Relationship between 'n' ie the number of sites taken and 'm' ie. the count of pairs generated is :

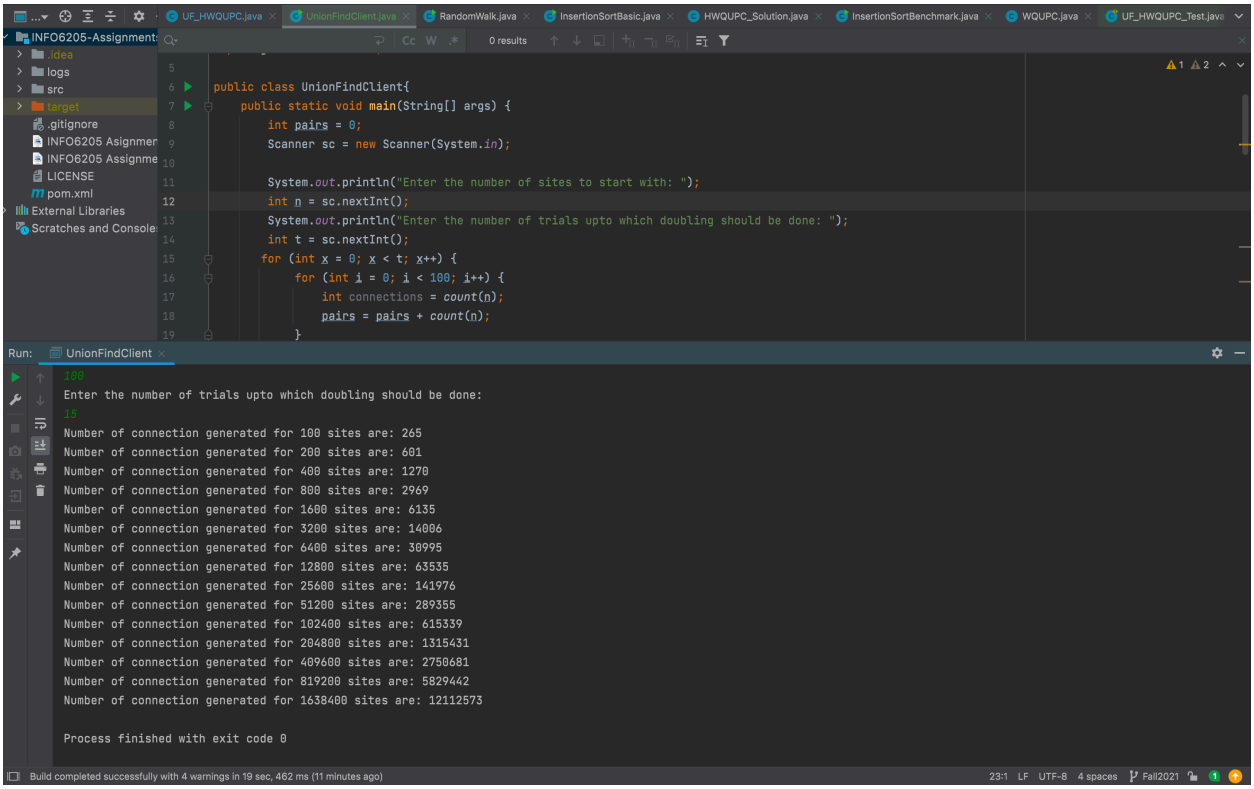
$m=n(\log n)/2$

- Evidence to support conclusions

1. Snapshot of unit test part 1



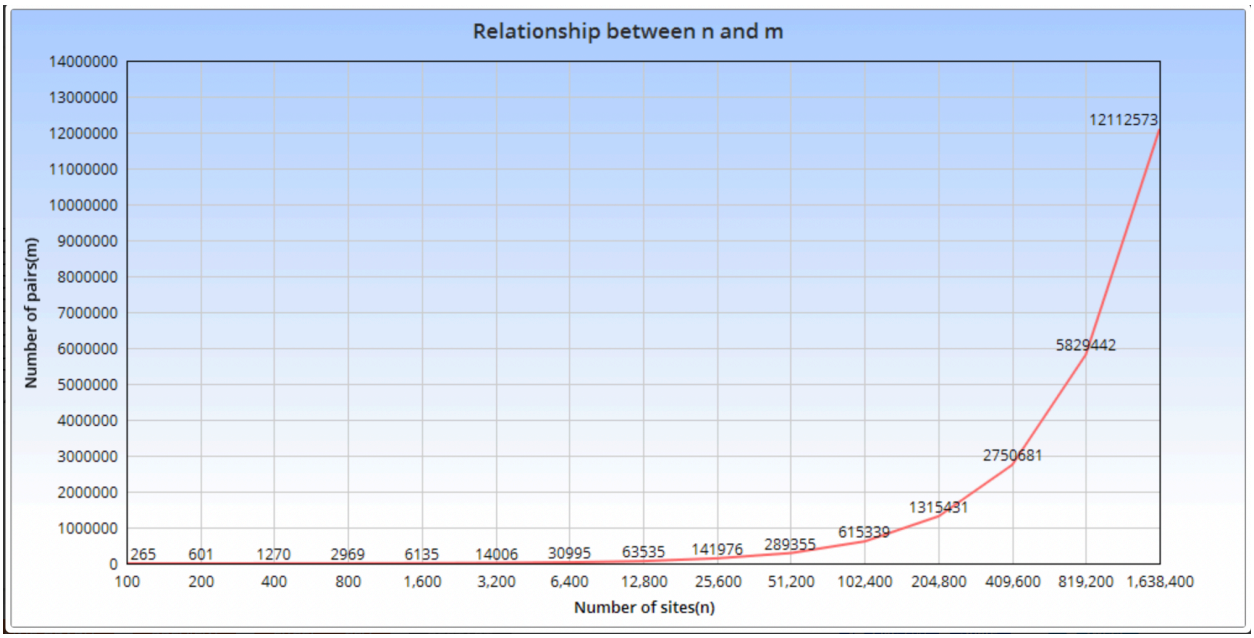
2. Snapshot of output part 2



Observations:

Relationship between n and m

Number of sites(n)	Number of pairs(m) ≈ n(logn)/2	Value of n(logn)
100	265	460
200	601	1,059
400	1,270	2,397
800	2,969	5,348
1,600	6,135	11,804
3,200	14,006	25,827
6,400	30,995	56,090
12,800	63,535	121,052
25,600	141,976	259,849
51,200	289,355	555,187
102,400	615,339	1,181,352
204,800	1,315,431	2,504,661
409,600	2,750,681	5,293,235
819,200	5,829,442	11,154,295
1,638,400	12,112,573	23,444,243



Conclusion:

Weighted quick union makes sure length of the tree is not too long. By storing the height of an array we reduce the the time that find() takes. We also reduce the number of array access thus reducing the number of random pairs generated from MN, where M is the number of connections and N is the number of sites, to M(logN) approximately . In worst case WQCPC will have height of tree to be logN.

From the graph we can infer that the relationship between M and N is logarithmic . The path compression by halving will reduce the time taken by find() operation reducing the height by 0.5

Hence from graph observations and proof we can conclude the relationship between number of sites(n) and number of connections generated(m) is:

$$m= n(\log(n))/2$$

- Unit test results:

All test cases passed successfully .