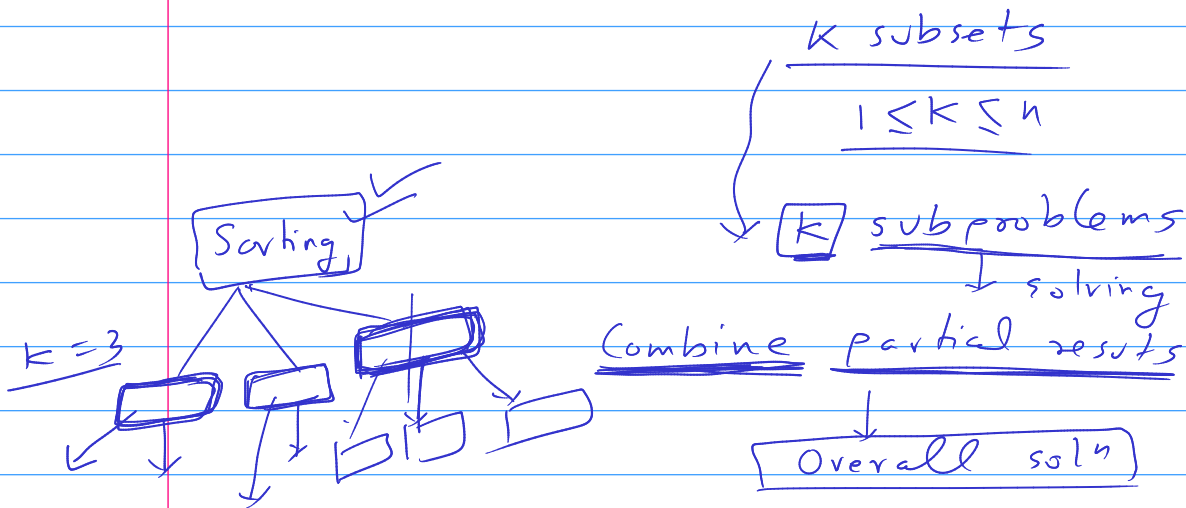
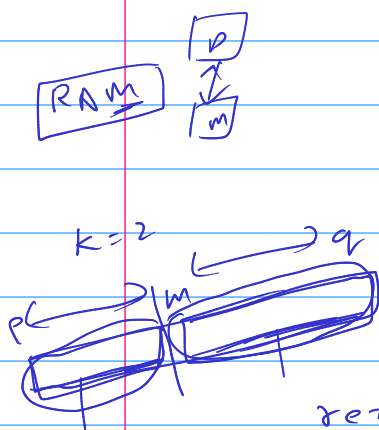


Divide & Conquer

- compute n inputs



- Subproblems are of same type
- Size of a subproblem is smaller than the size of the original problem
- original problem nature with subproblems
- subproblems are solved separately
- ⇒ DANDC can be expressed as recursive procedure



Procedure DANDC(p, q)

global n, A[1, n], int p, q

→ if small(p, q)

then return G(p, q) ← straight forward sol^n

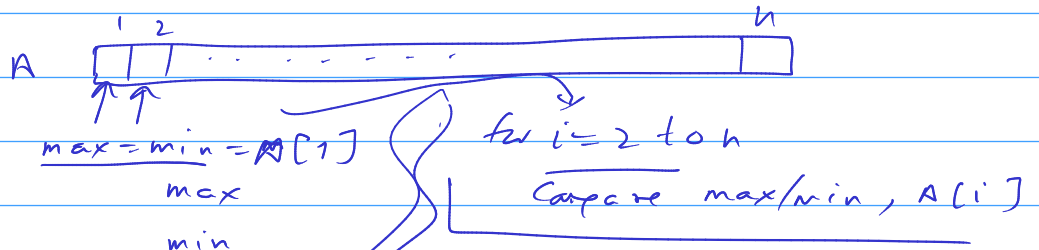
else

→ m ← DIVIDE(p, q) $p \leq m \leq q$

return COMBINE(DANDC(p, m), DANDC(m, q))

terminating step

Finding the maximum and minimum



2(n-1) comparisons

Best case $n-1$

(A)

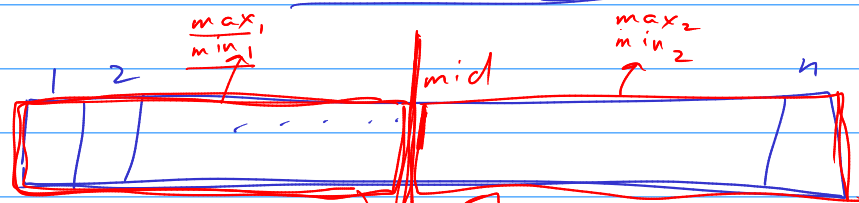
$A[i] > \text{max}$
 $\text{max} \leftarrow A[i]$
 $A[i] \leq \text{min}$
 $\text{min} \leftarrow A[i]$
 $\Rightarrow O(n)$

Re-written

(B)

$\Rightarrow O(n)$

- 1) Split
- 2) Solving subproblems
- 3) Combine



$\rightarrow \text{mid} = \lfloor \frac{lb+ub}{2} \rfloor$

$k=2$

$O(1)$

$\frac{n}{2} \text{ max} \leftarrow \text{MAX}(\text{max}_1, \text{max}_2)$
 $\frac{n}{2} \text{ min} \leftarrow \text{MIN}(\text{min}_1, \text{min}_2)$

$n=2^k$ true integer

$n=8=2^3$



1st subproblem

$n=4$

$\text{min}_1 = 1$ $\text{min}_2 = 5$
 $\text{max}_1 = 4$ $\text{max}_2 = 6$

$\text{max}_1 = 6$
 $\text{min}_1 = 1$

2nd subproblem

$n=4$

$\text{max}_2 = 10$
 $\text{min}_2 = 2$

$\text{max} = 10$
 $\text{min} = 1$

straight forward

$n=1$

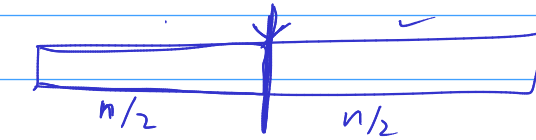
$\text{max} = \text{min} = A[i]$

$n=2$

Time complexity of D&C method for MZNM MAX

Recurrence relation

$$n = 2^k$$



$$T_n = C_2 + T_{n/2} + T_{n/2} + C_1$$

$$\Downarrow \boxed{T_n = 2T_{n/2} + C} \quad T_1 = 1$$

$$T_n = 2T_{n/2} + C$$

$$= 2(2T_{n/4} + C) + C$$

(substitution method)

$$=$$

$$T_{n/8}$$



$$T_{n/2^k} \rightarrow T_1$$

$$=$$

$$= \boxed{nT_1 + (n-1)C} = O(n)$$

$$= n + nC - C$$

$$n = 2^k$$



Next Lab experiment: D&C implementation

0111

Multiplication of 2 n-bit numbers

$$n = 2^k$$

$$\begin{array}{r} x \\ y \end{array} \times \quad Z = x * y$$

$$\begin{array}{ccccccc} x_n & x_{n-1} & & & & & x_1 \\ * & y_n & y_{n-1} & & & & y_1 \\ \hline & & & & & & \downarrow \downarrow \end{array}$$

$O(n^2)$ multiplications

$O(n^2)$ addition

$O(n^2)$

$$\begin{array}{r} 1101 \\ \times 0101 \\ \hline 1101 \\ 0000 \\ 1101 \\ 0000 \\ \hline \end{array}$$

addition

$$X \quad \boxed{A \quad B}$$

$$M \times 2^{n/2} \quad \underline{M000 \dots 0}$$

$$Y \quad \boxed{C \quad D}$$

$$+ \quad \begin{array}{r} A \quad \overset{n/2}{00 \dots 0} \\ \hline B \end{array}$$

$$X = A * 2^{n/2} + B$$

$$Y = C * 2^{n/2} + D$$

$$Z = X * Y$$

$$Z = X * Y = (A * 2^{n/2} + B) (C * 2^{n/2} + D)$$

$$Z = \underline{AC * 2^n} + \underline{(AD + BC) * 2^{n/2}} + \underline{BD}$$

(1) \downarrow (2) \downarrow (3) \downarrow (4) \downarrow
 $n/2$ bit multiplications (4 Nos.)

(II) Shift opⁿ ($O(n)$)

(III) Binary addition ($O(n)$)

$$\begin{aligned} T_n &= 4T_{n/2} + O(n) \\ &= \underline{O(n^2)} \end{aligned}$$