



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

AY: 2025-26

Class:	T. E.	Semester:	V
Course Code:	CSC501	Course Name:	COMPUTER NETWORKS

Name of Student:	SHRUTI GAUCHANDRA
Roll No. :	18
Assignment No.:	02
Title of Assignment:	ERROR CORRECTION, DETECTION
Date of Submission:	05/08/25
Date of Correction:	05/08/25

Evaluation

Performance Indicator	Max. Marks	Marks Obtained
Completeness	5	05
Demonstrated Knowledge	3	03
Legibility	2	02
Total	10	10

Performance Indicator	Exceed Expectations (EE)	Meet Expectations (ME)	Below Expectations (BE)
Completeness	5	3-4	1-2
Demonstrated Knowledge Legibility	3	2	1
Legibility	2	1	0

Checked by

Name of Faculty : Mrs. SNEHA YADAV

Signature :

Date :

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05/08

SC501.2 Apply communication mechanisms like services, farming, error detection, error correction, multiple access control for flow and error control.

- Q1. A layer pure ALOHA network transmits 200-bit frames on a shared channel of 200 kbps. Find the throughput if the system (all station together) produces
- 1000 frames per second?
 - 500 frames per second?
 - 250 frames per second?

→ (1) Frame transmission time (Average) = t is 200 bit / 200 kbps or 1 ms

(2) The vulnerable time is $2t = 2 \times 1 \text{ ms}$
 $= 2 \text{ ms}$

(3) This means no station should sent 1 ms before or after this station starts transmission.

(4) The throughput for pure ALOHA is,

$$S = Ge^{-2G} \quad \text{(i)}$$

(5) The maximum throughput;
 $S = 0.184 \text{ when } G = 0.5$

- a. If the station produces 1000 frames per second i.e. 1 frame per millisecond.
The load is 1.

Put $G=1$ in equation (i),

$$\therefore S = 1 \times e^{-2 \times 1}$$
$$= 0.135$$

$$\therefore S = 13.5\%$$

Therefore, throughput is $1000 \times 0.135 = 135$ frames.
Only 135 frames out of 1000 would be probably transmitted successfully (survive).

b. If station produces 500 frames per second,
i.e. $1/2$ frame per millisecond.

The load is $\frac{1}{2}$

Put $G = 1/2$ in equation (i),

$$S = \frac{1}{2} (e)^{-1/2 \times 2} = \frac{1}{2} (e)^{-2 \times \frac{1}{2}}$$
$$= 0.184$$

$$\therefore S = 18.4\%$$

Therefore, throughput is $500 \times 0.184 = 92$ frames.
Only 92 frames out of 500 would be probably transmitted successfully (survive).

c. If the station produces 250 frames per second
i.e. $1/4$ frame per millisecond.

The load is $\frac{1}{4}$

Put $G = \frac{1}{4}$ in equation(i),

$$\therefore S = \frac{1}{4} (e)^{-2 \times \frac{1}{4}}$$

$$= 0.152$$

$$\therefore S = 15.2\%$$

Therefore, throughput is $250 \times 0.152 = 38$ frames
Only 38 frames out of 250 would be probably transmitted successfully (survive).

Q2. A bit stream 10011101 is transmitted using the standard CRC method. The generator polynomial is $x^3 + 1$.

1. What is actual bit string transmitted?
2. Suppose the third bit from the left is inverted during transmission, how will receiver detect the error.

→ Given: Data Word (Bit String) : 10011101

Generator Polynomial : $x^3 + 1$

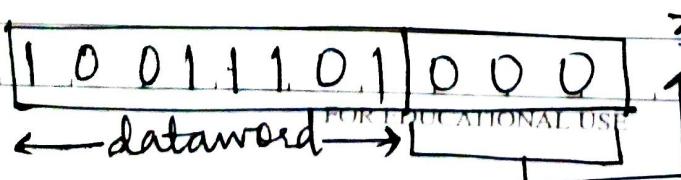
$$: x^3 + 0x^2 + 0x + 1$$

$$: 1001$$

$$= n$$

Step 1: Obtain the Dividend.

Dividend = Data word + $(n-1)$ zeroes.



3 additional zeroes ($n-1$)

Step 2: Carry Out Long division.

Generator Dividend

$$\begin{array}{r}
 \downarrow & 100011 \\
 1001) 10011101000 & \leftarrow \\
 \oplus 1001 & \downarrow \downarrow \downarrow \downarrow \\
 0000\text{:}101 & | \\
 1001 & \downarrow \\
 \oplus 01000 & \downarrow \\
 \oplus 1001 & \downarrow \\
 \hline 000100 & \leftarrow \text{Remainder}
 \end{array}$$

Step 3: Obtain actually transmitted bit stream
Transmitted word is obtained by writing data word followed by remainder.

$$\text{Transmitted word} = \boxed{10011101|0100}$$

\leftarrow Datawork \rightarrow Remainder

~~Error detection:~~

~~Step 4: Write erroneous received word.~~

The received word = 10111010100
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error bit (3rd bit from left is inverted).

At receiver, word is divided by same divider used at transmitter i.e. 1001

$$\begin{array}{r}
 10101000 \\
 1001) \overline{101111010100} \leftarrow \text{received word} \\
 + 1001 \\
 \hline
 001011 \\
 1001 \\
 + 1001 \\
 \hline
 001001 \\
 1001 \\
 + 1001 \\
 \hline
 00000100 \leftarrow \text{remainder}
 \end{array}$$

A non zero remainder indicates that there is an error in received codeword.