

\* Books →

- [1] Chapra → Applied Numerical Methods
- [2] Chapra & Canale → Numerical methods for Eng.

[3] Chapman → Matlab programming

[4] R. Pustap → Getting started with MATLAB

\* Reference →

[1] Devices & Haslett

- A first Course in Computational Physics

[2] Giordino & Nakamishi

- Computational Physics

15

• Mid Term

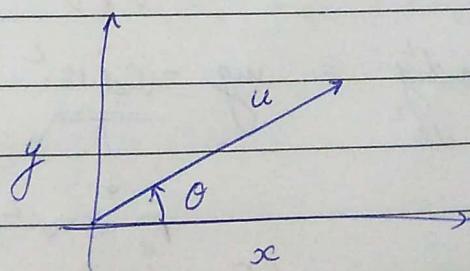
• End Term

• Quiz

→ Lab

Q. 20 Why Computational Physics?

$$\text{Eq. 1} \quad \vec{F} = m \vec{a}$$



$$\frac{d^2y}{dx^2} = -g$$

$$\frac{d^2x}{dt^2} = 0$$

$$\frac{dy}{dt} = -gt + k_1$$

$$y = -\frac{1}{2}gt^2 + k_1 t + k_2$$

$$x = u_x t$$

at  $t=0, x=0$

$$u_x = u \cos \theta$$

5

$y = 0$

$$u_y = u \sin \theta$$

$$x = u_x t$$

$$y = u_y t - \frac{1}{2}gt^2$$

10

\* Restrictions on ideal cond<sup>n</sup>:

- ① Air drag
- ② Wind velocity
- ③ Magnus Effect
- ④ Coriolis Effect

You can see the eg. of Coriolis effect, one is due to earth's rotation the target misses its projectile.

Eg 2.20

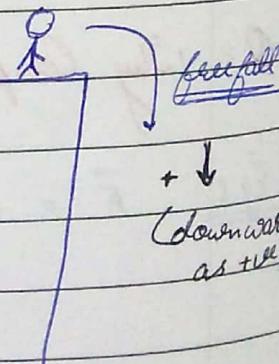
Bungee Jumping  $\Rightarrow$

Cof. of air drag

$$\frac{d^2y}{dt^2} = g - \frac{(C_d)}{m} v^2$$

$$F_A \propto v^2$$

(Force due to air drag)



25

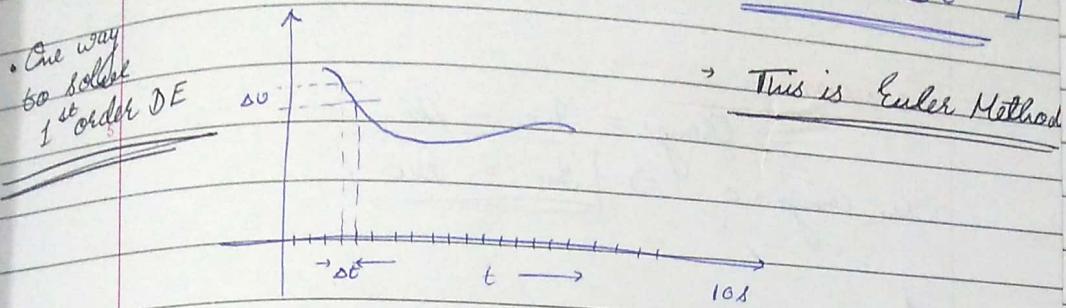
$$\frac{d^2y}{dt^2} = g - \left(\frac{C_d}{m}\right) v^2 \quad [1^{st} \text{ order diff. eqn}]$$

Sol<sup>n</sup>  $\rightarrow$

$$v(t) = \sqrt{\frac{g m}{v_0}} \cdot \tanh \left( \sqrt{\frac{C_d}{m}} t \right)$$

$$\tanh(x) = \frac{e^{ix} - e^{-ix}}{e^{ix} + e^{-ix}}$$

$$\frac{du}{dt} \approx \frac{\Delta u}{\Delta t}$$



t	u	
$t_1 = 0$	$u_1 = 0$	$\frac{u(i+1) - u(i)}{t(i+1) - t(i)} = g - \left(\frac{cd}{m}\right) u^2(i)$
$t_2 = t_1 + \Delta t$	$u_2 = ?$	
$t_3 = t_1 + 2\Delta t$	$u_3 = ?$	$\Rightarrow u(i+1) = u(i) + \Delta u(i)$
$\vdots$	$\vdots$	$\left\{ g - \left(\frac{cd}{m}\right) u^2(i) \right\}$
$t_n$	$u_n$	

For  $i=1 \Rightarrow$

$$u(1) = u(0) + \Delta t \left\{ g - \left(\frac{cd}{m}\right) u^2(0) \right\}$$

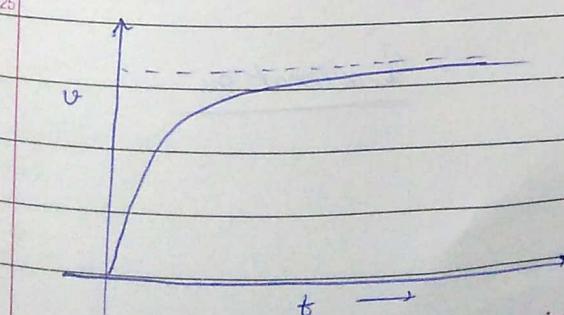
$$\Rightarrow u(1) = u_1$$

$$\Rightarrow u(2) = u_2 \text{ & so on.}$$

$$\Rightarrow u_2 = u_1 + \Delta t \left\{ g - \left(\frac{cd}{m}\right) u_1^2 \right\}$$

Free fall →  
Terminal velocity

$$u = \sqrt{\frac{gm}{cd}}$$



Coding → When you have a problem & its code, run the code for a test case where you already know the sol<sup>n</sup>.

\* Conservation law →

$$5 \quad \text{Change} = \text{Inc.} - \text{Dec.}$$

$$\text{When, } \text{change} = 0 \Rightarrow \text{Inc.} = \text{Dec.}$$

\* All

10       $\begin{array}{l} \text{pi}; \\ \text{format long}; \\ \text{pi}; \end{array} \quad ] \quad \begin{array}{l} \text{More precision upto 15 digits} \\ (\text{Double precision}) \end{array}$

format short;      ]      4 digit precision  
pi;

\* Col "Matrix" →  $\cdot b = [2; 4; 6; 8; 10]$

$$\cdot b = [2 \quad 4]$$

4  
6  
8  
10 ]

$$\cdot b = [2 \quad 4 \quad 6 \quad 8 \quad 10]'$$

20

• who      ]      To know abt the variable.  
whos

25      • MATLAB has no zero index.

x = zeros(1, 11)      % Preallocation

linspace (x1, x2, n) → To generate a row vector

→ logspace() # See it

→ n = 100 # by default.

x = linspace(0, 1, 6)

x = [0, 0.2, 0.4, 0.6, 0.8, 1]

$$h = \frac{x_n - x_1}{n-1}$$

$$h = \frac{x_n - x_1}{n-1} \quad \text{→ Step size (lein width)}$$

logspace(x1, xn, n)

x = logspace(-1, 2, 4) → [10<sup>-1</sup> 10<sup>0</sup> 10<sup>1</sup> 10<sup>2</sup>]

x = [10<sup>x\_1</sup>, 10<sup>x\_2</sup>, 10<sup>x\_3</sup>, ... ]

x<sub>1</sub> = x, # If n is omitted then

x<sub>2</sub> = x<sub>1</sub> + h by default it will be

x<sub>3</sub> = x<sub>2</sub> + 2h 50.

\* Operators → / → Division

\ → left division

<sup>n</sup> → Power

\* System of linear Eq<sup>n</sup> →

$$A = \begin{bmatrix} 3 & 2 & 5 \\ 4 & 8 & 7 \\ 5 & 3 & 8 \end{bmatrix}$$

$$b = \begin{bmatrix} 4 \\ 9 \\ 11 \end{bmatrix}$$

$$\underline{Ax = b} \rightarrow \text{To find } x$$

$$\rightarrow x = A^{-1}b \quad \text{or} \quad \underline{A^{-1}b}$$

[In MATLAB]

$$[A^{-1}Ax = A^{-1}b]$$

$$\rightarrow x = A^{-1}b$$

- \* Matrix Division is undefined, but we can do elementwise division.

- Matrix Multiplication  $\rightarrow B = A * b$

<sup>col<sup>n</sup></sup>  
No. of ~~rows~~ in 1<sup>st</sup> matrix = No. of rows in 2<sup>nd</sup> one.

- \* Element by element operation  $\rightarrow$

- ① Elementwise squaring  $\rightarrow B = A.^2$

- ② Elementwise divide  $\rightarrow C = X ./ Y$

- \* To get help  $\rightarrow$  Command prompt

①  $\gg \text{help log}$

②  $\text{doc log}$

- \* Rounding  $\rightarrow E = [-1.6, -1.5, -1.4, 1.5, 1.6]$

①  $\text{round}(E)$

$$[-2, -2, -1, 2, 2]$$

- ②  $\text{ceil}(F)$  # Push towards  $+\infty$ .

$$[-1, -1, -1, 2, 2]$$

③ floor (E) →

$$\rightarrow [-2, -2, -2, 1, 1]$$

- ① min(E)
  - ② max(E)
  - ③ mean(E)
  - ④ prod(E)
  - ⑤ sum(E)
  - ⑥ sort(E)

Note → Write units in comment.

$$y_{-10} = \frac{g m}{c d} \tanh \left( \frac{\int g c d}{m} t \right)$$

$$g = 9.81 \text{ ; } \% \text{ m s}^{-2}$$

$$m = 68.1 \text{ ; } \% \text{ kg}$$

→ Always use Std SI units while doing a quib. in MATLAB.

$$c_d = 0.25 ; \quad \text{kg/m}$$

$$t = 0 : 2 : 20 ;$$

$$v = \sqrt{[(g * m) / cd]} * \tanh\left(\sqrt{[(g * m) / cd]}\right)$$

20. To continue a ~~one~~ expression in next line ...

plot(v, t);

$\times$  label ("velocity");

yester ('time') ;

25 title (' or os t');

~~grid();~~

Legend ("velocity"):

\* Colour Specifiers →

① Blue → b

Green → g

5 Red → r

Cyan → c

Magenta → m

Yellow → y

Black → k

10 White → w

② Line →

Solid - - -

Dashed - - - : :

Dotted - - - - -

Dashed dot - - - o - - -

③ Symbols → point

circle

o

X-mark

x

15 Plus

+

Star

\*

Square

s

diamond

d

\* 20 Gray scale op. →

① . \* → term by term multiplication

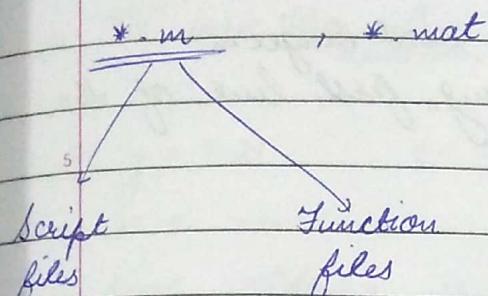
② . / → ——— division

③ . ^ → ——— exponentiation.

$$h = \frac{x_n - x_1}{(n-1)}$$

④ . \ → Element wise inverse multiply

\* MATLAB  $\Rightarrow$



clear variables,  
close all;

\* Multiple plot on different windows  $\rightarrow$

```

figure();
plot(xc, y1);
figure();
plot(x, y2);
  
```

\* multiplot()      o/o works like subplot  
 $\hookrightarrow$  Multiple plots on same windows

\* Function files  $\Rightarrow$

```
function [out1, out2, ...] = funcname(in1, in2, ...)
```

% help comments

Name of the f" file will be  
this

out1 = \_\_\_\_\_

It will tell what the f" will do,

end

what are the i/p's & what are the  
o/p's.

Why Required ??

>> help funcname

$\rightarrow$  It will show what is written as  
comment written after the first line of  
the f". [declaration]

\* Path  $\rightarrow$  Wherever MATLAB searches for a f", it  
searches

or program

H1 line → 1<sup>st</sup> line of the help

→ look for bungee word bungee in

→ search all the f" having first line of the help comment.

\* log. of f" files →

I function v = freefall(t, m, cd)

10 % freefall: bungee velocity with drag

% g =

\* To take value from the user →

→ m = input ('Mass (kg):')

→ This will take as value m in it.

\* Op →

① m % in Command Prompt.

20 ② disp(m) ← % finer way

③ fprintf ('The velocity is %.4f m/s\n',  
 3 digits →                  ↓  
 on left                      4 decimal pt.  
 hand side                    after decimal  
 [before decimal])

\* Format

%d

Description

integer

%e

Scientific format

%F

Decimal format

%f

Compact of %e or %f  
 Newline

\t → Tale.

\* To save the values →

⇒  $g = 9.81$ ;  $t = 4$ ;  $d = 0.25$   
 $v = \underline{\hspace{2cm}}$

save veldrag & d

∴ v & d will be saved in veldrag  
as veldrag.mat

10

. load veldrag ∵ Load the variable

→ If we don't specify v & d then by default all the  
variable will be saved present in workspace.

15

⇒  $A = [5 \ 7 \ 8 \ 3 \ 4 \ 2]$ ,

⇒ save Mymtx.txt  $\xrightarrow{\text{ascii}}$

⇒ clear variables,

⇒ load Mymtx.txt

⇒  $A = \text{load}('Mymtx.txt');$

\* if

if condition  
statements

end

if  
else elseif  
||

Eg:-

$y = \begin{cases} -1 & ; \text{ if } x < 0 \\ 0 & ; \text{ if } x = 0 \\ 1 & ; \text{ if } x > 0 \end{cases}$

function sgn = mysignal(x)

if  $x < 0$

sgn = -1;

elseif  $x == 0$

sgn = 0;

else

sgn = 1;

end

end

10

• for Loop →

for index = start : step : finish

;

15 end

Eg. → for i = 1:2:10  
      disp(i)  
end

for i = 10:-1:0  
      disp(i)  
end

20

\* Vectorization →

i = 0  
for t = 0:0.02:50  
    i = i + 1;  
    y(i) = cos(t.^2);  
end

{ Same Code  
↓  
t = 0:0.02:50;  
y = cos(t.^2);  
↓  
This is Vectorization.

→ Vectorization makes code run faster. So, do vectorization whenever possible.

### \* while loop →

while (1)

statements

if cond " "

break;

end

end

while (condition )

statement

end

5

Eg. →  $x = 19$

while (1)

10

if ( $x < 0$ ) break; end

$x = x - 5$

disp (x)

end

Pre test loop

15

$x = 19$

while (1)

$x = x - 5$

if  $x < 0$ , break;

end

disp (x)

end

Post test

loop

- pause (0.5) % 0.5 ms or sec. ↳ ①

15/01/19

### \*<sup>20</sup> Anonymous f" →

① Script file

② Function file

→ One more way to define f" is Anonymous f".

Gyntax

handle = @ (arglist) expression;

"f1(x,y) =  $x^2 + y^2$ " ↳ It can be written as.

"f1 = @ (x,y)  $x.^2 + y.^2$ ";

22  $a = 4$

$$f(x) = ax^b$$

22  $b = 2$

$$f(x) = a \cdot x^b ;$$

5  $f(3) = 36$

$a = 3$

Run this again, then only value of  $a$  &  $b$  will change.

6  $f(3) = 27$

10

\* Plotting a func.

→  $fplot (func, limit)$ ;

15  $v(t) = \sqrt{\frac{g+m}{d}} \tanh\left(\sqrt{\frac{g+m}{d}} * t\right)$

$g = 9.81$ ;  $m = 68$ ;  $d = 0.25$ ;

"Define"  $v = @v(t) \operatorname{sgt}($  <sup>fixing the values</sup>  $)$

∴ To plot

$fplot (v, [0 12])$ ; ∴ See the step size

Q. Write a script to find the average  $v$  from  $t=0$  to  $t=12$

25  $t = \text{linspace}(0, 12);$

$v = \operatorname{sgt}\left(\frac{g+m}{d}\right) * \tanh\left(\operatorname{sgt}\left(\frac{g+m}{d}\right) * t\right).$

→  $\text{mean} = \text{mean}(v)$

→ We use the expression as anonymous  $f^n$ , if we have to use the expression again & again otherwise not.

5 script file →

$t_{\min} = 0 ; t_{\max} = 12 ; n = 20 ;$

$\text{vel} = @ (t) \sin(t) * \tanh(t)$

$\text{omean} = \text{findmean}(\underline{\text{vel}}, \underline{t_{\min}}, \underline{t_{\max}}, \underline{n})$

$f^n$

data pts.

[If we have to pass the anonymous  $f^n$  to another  $f^n$  we don't have to do  $\text{vel}(t)$   
we pass only  $\text{vel}$ ]

15 Function file →

$\text{function omean} = \text{findmean}(f, a, b, \underline{n})$  \*

$x = \text{linspace}(a, b, n);$

$y = f(x);$

$\text{omean} = \text{mean}(y);$

Saved → findmean.m

end

$t = 0 \text{ to } 12;$

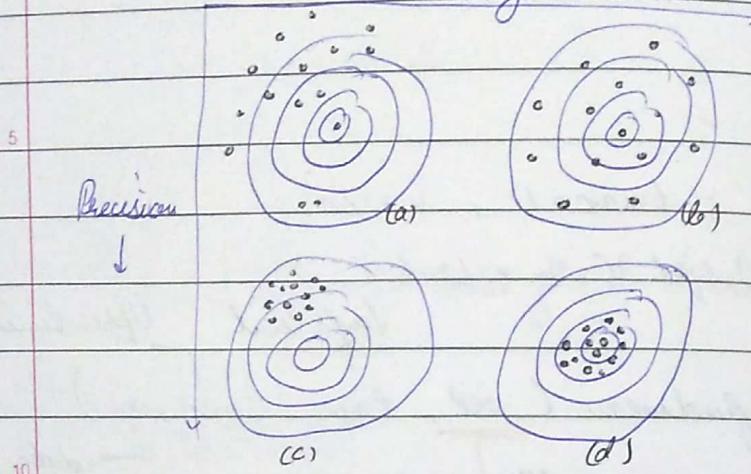
$m) * t);$

\* We will always get approx answer and not exact one.  
So, we calculate accuracy & precision

Accuracy → How closely we computed or measured value agrees with true value.

Precision → How closely individual values agree with one another.

Accuracy →



(a) → Biased & imprecise

(b) Unbiased

(c) Biased & precise

(d) Unbiased & .

15

Error = Biasness + Impreciseness.

\* Error → True value - Approx value.

20 Eg → True value of bridge = 1 km.

Here, both have True chalk = 10 cm.

1 cm. Measured value of bridge = 99,999 cm.  
Error chalk = 9 cm

Prob.

The value that we are measuring should also come into the def. of error.

Fractional Error =  $\frac{\text{Approx Value} - \text{Exp. Value}}{\text{True Value}}$

So, in above eg.  $\rightarrow fe = 10^{-5}$ ,  $fe = 10^{-1}$   
(frac. error of Bridge)

agrees

Further most of the time we don't have true value, so:

$$\text{Fractional error} = \frac{\text{abs}(\text{Present approx} - \text{True approx})}{\text{Present approx}}$$

\* Root finding  $\rightarrow$  Finding the zero of expression.

(1) Graphical Method

(2) Bracket Method  $\rightarrow$  Bisection

(3) Open end method  $\rightarrow$  False Position

(3) Open end method  $\rightarrow$  Newton Raphson

$\rightarrow$  Secant Method

$\rightarrow$  Successive approx.

Q5

. Let's visit Bungee Jumping Prob.  $\rightarrow v = \sqrt{\frac{g}{d}} \tanh\left(\sqrt{\frac{gd}{m}} t\right)$

$$v = 36 \text{ m s}^{-1}, t = 48, d = 0.25 \text{ kg/m}, g = 9.81 \text{ m s}^{-2}$$

$$m = ?$$

25) Mass. mass a person can have so that it didn't cross  
bd.  $36 \text{ m s}^{-1}$  in 48.

\* It is very difficult to put 'm' on left hand side &  
everything on right because  $v \rightarrow$  explicit fn  
 $m \rightarrow$  implicit fn of exp.

Approach  $\rightarrow$

$$f(m) = \sqrt{\frac{g}{d}} \tanh\left(\sqrt{\frac{gd}{m}} \times t\right) - v$$

There will be one value of  $m$  for which exp. becomes zero.

I Graphical Method

$$d = , g = , t , \omega =$$

①

$n = \text{linspace}(40, 25)$

$$f1 = \text{sqr}(\ ) \quad \text{1. Expression}$$

~~Method 1~~ plot(n, f1);

↳ Looking at Graph we see that at  $\approx 150$   
the  $f$  becomes 0.

→ By Expression

②  $g, d, \omega, t$

$$f2 = @(\text{nd}) \text{sqr}(\ ) ;$$

15 → Expression

~~Method 2~~ fplot(f2, [40, 250]);

→ Anonymous  
By func

20 Expression →

$$f1 = \text{sqr}(g * n / d) * \text{tanh}(\text{sqr}(g * d / m) * t)$$

[ $\because$  we are multiplying  
an array with  
another array]

[ $*$  as  
dividing by m  
as array will not  
make sense.]

$$\text{Eg. } f = g / m ;$$

→  $f2$  will be same as  $f1$ .

### ③ Modification of Method 2

g, d, v, t

$$f2 = @ (mid)$$

m = linspace (40, 250);

fval = f2(m); It is a 1D array

plot (m, fval);

plot (m, f2); → Error as f2 is a function  
has no value

10. As a separate fun file →

m, d, v, t

fval = fun4 (m, d, v, t)

plot (m, fval); grid();

11. Function →

function f2 = fun4 (m, d, v, t)

$$free @ g = 9.81;$$

$$f2 = \text{sqrt}(g * m / d) * \tanh(\text{sqrt}(g * d / m) * t)$$

- 19;

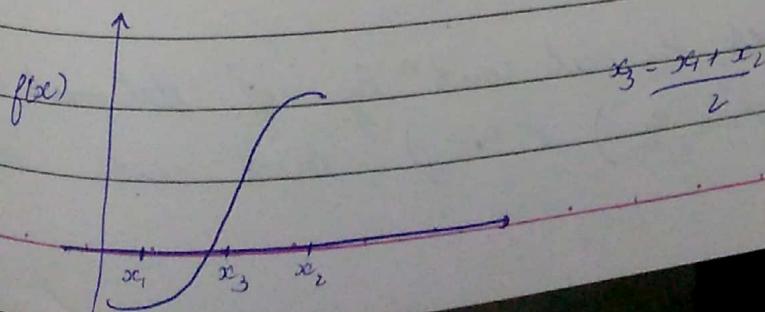
20.

end

### II Bracket Method →

① Bisection method → ② Requires 2 guesses

which should bracket the root



① 2 initial guesses which bracket the root.

$$\textcircled{2} \quad x_3 = \underline{x_1 + x_2}$$

2

③ Check which bin encloses the root.

If  $(x_1, x_3)$  bracket the root then

$$x_1 = \underline{x_3}$$

else if  $(x_3, x_2)$  bracket the root  $\rightarrow$

$$x_2 = \underline{x_3}$$

④ 10 Keep doing above until satisfied.

\* Algorithm →

i/p  $\rightarrow f(x), x_1, x_2, tol$

o/p  $\rightarrow root$

(Write the  
for this) 15

$f''$  Check if  $f(x_1) * f(x_2) < 0$  otherwise exit program.

while (  $\infty$  loop )  $\rightarrow \frac{|x_1 - x_2|}{x_1} \rightarrow tol$

$$x_3 = (x_1 + x_2)/2; \rightarrow \text{some +ve number}$$

if  $x_3 \in f(x_3) * f(x_1) < 0$

$$x_2 = x_3;$$

else

$$x_1 = x_3;$$

end end

25 (Cond "for exit")

$\rightarrow$  while loop ends

You can check the exit cond'n by 'if' statement also:

$y ( \frac{|x_1 - x_2|}{x_1} < tol ) \text{ break; end}$

\* Cond<sup>n</sup> of failure →

①  $f(x_1) * f(x_2) < 0$

→ Odd no. of roots

②  $f(x_1) * f(x_2) > 0$

→ Even no. of roots

③  $f(x_1) * f(x_2) < 0$

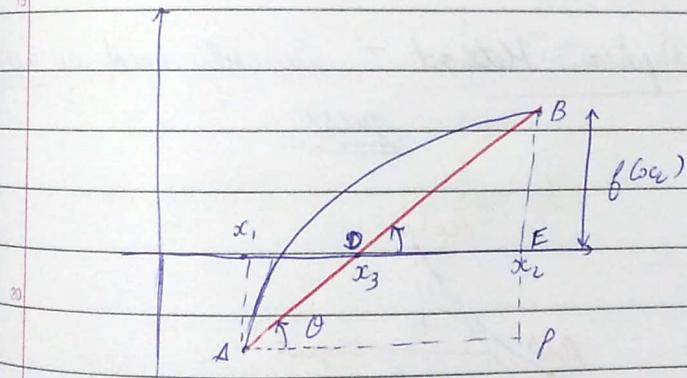
→ Even roots if  $f''$  at root is tangential.

④  $f(x_1) * f(x_2) < 0$

→ Even root if  $f''$  is discontinuous.

⑤ Method of false Position →

15



In  $\triangle BDE \rightarrow \frac{f(x_2)}{(x_2 - x_3)} = \tan \theta$

In  $\triangle BAP \rightarrow \frac{f(x_2) - f(x_1)}{x_2 - x_1} = \tan \theta$

b.  $x_2 - x_3 = \frac{f(x_2) f(x_2 - x_1)}{f(x_2) - f(x_1)}$

$$\Rightarrow x_3 = \frac{x_1 f(x_2) - x_2 f(x_1)}{f(x_2) - f(x_1)}$$

? But here sometime  $f(x_2) - f(x_1)$  will be too small and dividing by a very small no. will lead to large errors.

So,

if  $\text{abs}(f(x_2) - f(x_1)) < \text{tol}$

exit program with proper message.

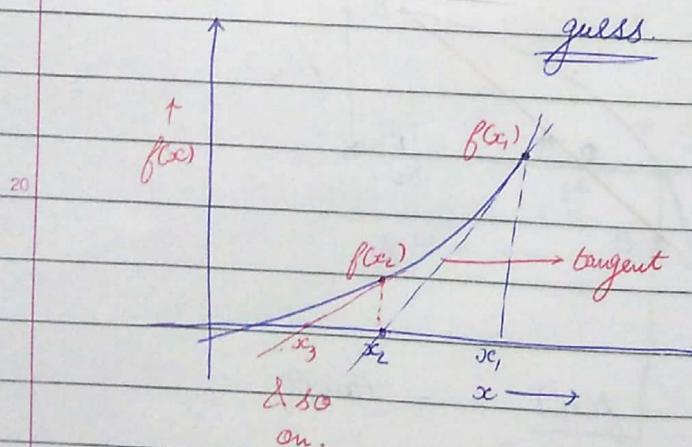
10 // If we have to check this cond "too".

Otherwise algo. wise it will be similar as the above one

### III \* Open End Method $\Rightarrow$

15

\* ① Newton - Raphson Method  $\rightarrow$  We will need an initial guess.

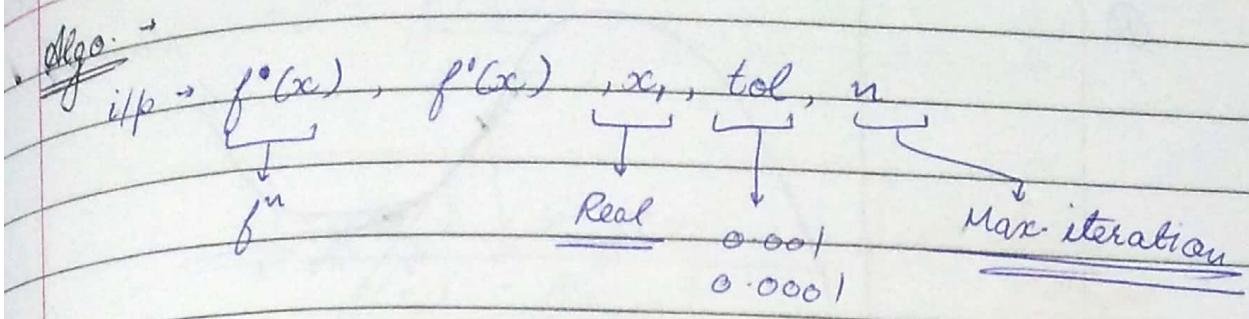


25  $\rightarrow$  The advantage here is it will converge too fast.  
ACD  $\Rightarrow$

$$\frac{f(x_1)}{(x_1 - x_2)} = \tan \theta = f'(x_1)$$

derivative of f<sub>i</sub> at x<sub>i</sub>

$$\rightarrow x_2 = x_1 - \frac{f(x_1)}{f'(x_1)}$$



→ There can be case when this method will diverge,  
so we have to limit the no. of iteration (n).

ifp → root.

Algo. →

```

while (n > 0)
    if (abs( $f^*(x_n)$ ) > tol)
        print msg & exit;
    end
     $x_{n+1} = x_n - (f(x_n) / f'(x_n))$ ;
  
```

① if ( $f(x_{n+1}) < tol$ )  
 then  
 break; end

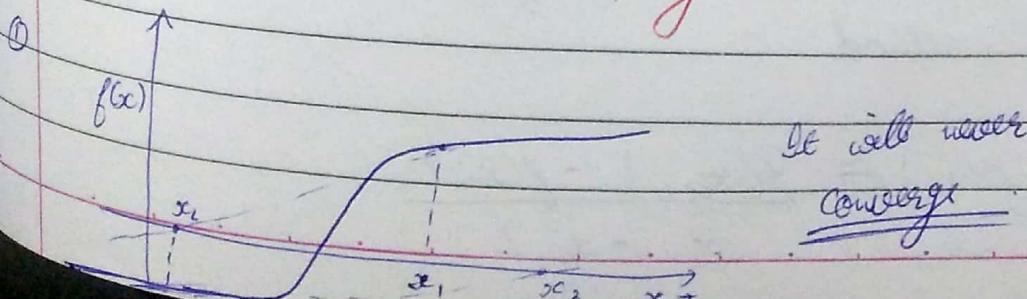
\*\*  ~~$x_n = x_{n+1}$~~ ; // jump line.  
 $n = n - 1$ ;

end

root =  $x_{n+1}$ ;

→ We need to find the derivative of the  $f^n$  on paper  
 before if we are using this method.

\* Cond' when it will not converge →



(2)

5

Never C.

10

(3)

15

Never C.

(4)

20

 $f(x)$  $x_1$ 

(2) Secant Method →

→ We don't want to find the derivatives, use this method.

$$f'(x_i) \approx \frac{f(x_{i+1}) - f(x_i)}{(x_{i+1} - x_i)}$$

$$x_{i+1} = x_i - \frac{f(x_i)}{f'(x_i)}$$

$$x_{i+1} = \frac{x_{i-1} f(x_i) - x_i f(x_{i-1})}{f(x_i) - f(x_{i-1})}$$

Secant Method

→ This is similar to false position method but the two initial guess need not to bracket the root.

10. Algo. →

inp →  $f(x)$ ,  $x_1$ ,  $x_2$ , tol, n

opp → Root

for  $i = 1 : 1 : n$

15 if ( $\text{abs}(f(x_1) - f(x_2)) < \text{tol}$ ) Msg; break; end

$$x_3 = (x_1 * f(x_1) - x_2 * f(x_2)) / (f(x_2) - f(x_1));$$

if ( $\text{abs}((x_1 - x_2) / x_1) < \text{tol}$ )

Msg; break;

end

~~temp =  $x_2$~~

$x_2 = x_3$  ;

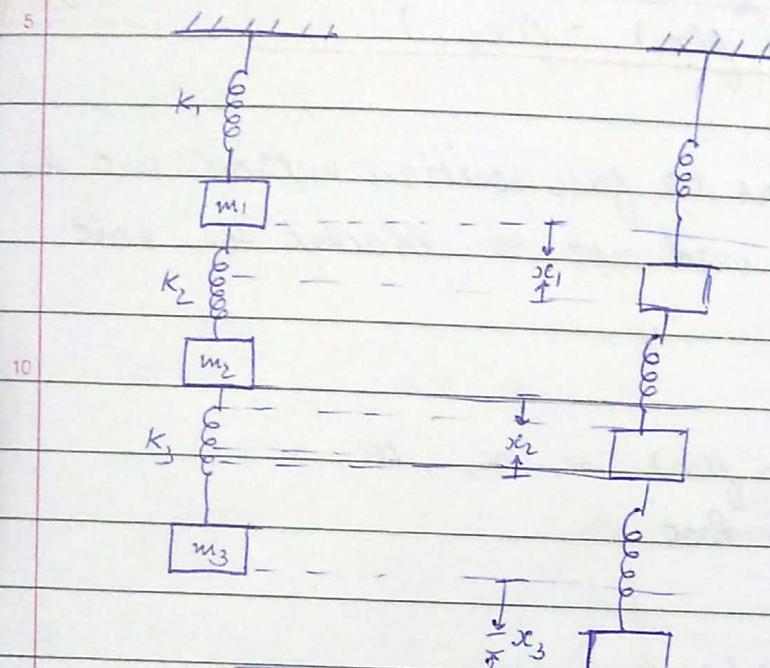
$x_3 = x_2$  ;

end

25 root =  $x_3$  ;

use this

22/01/19

Sys. of linear Eq<sup>n</sup> =>Problem =>Initial

$$m_1 \frac{d^2x_1}{dt^2} = m_1 g + k_2(x_2 - x_1) - k_1 x_1$$

$$m_2 \frac{d^2x_2}{dt^2} = m_2 g + k_3(x_3 - x_2) - k_2(x_2 - x_1)$$

$$m_3 \frac{d^2x_3}{dt^2} = m_3 g - k_3(x_3 - x_2)$$

Steady State Solution (Eqf. sol<sup>n</sup>) =>

$$(k_1 + k_2)x_1 - k_2 x_2 = m_1 g$$

$$-k_2 x_1 + (k_2 + k_3)x_2 - k_3 x_3 = m_2 g$$

$$-k_3 x_2 + k_3 x_3 = m_3 g$$

Mtx sol<sup>n</sup>  $\Rightarrow$

$$C = \begin{bmatrix} k_1+k_2 & -k_2 & 0 \\ -k_2 & k_2+k_3 & -k_3 \\ 0 & -k_3 & k_3 \end{bmatrix}$$

5.  $b = \begin{bmatrix} m_1 g \\ m_2 g \\ m_3 g \end{bmatrix}$

$$x = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix}$$

$$Cx = b$$

10.  $x = b/C$  or,  $C \setminus b$

. Symmetric Mtx.  $\Rightarrow a_{ij} = a_{ji}$

15.  $A = \begin{bmatrix} 5 & 1 & 2 \\ 1 & 3 & 7 \\ 2 & 7 & 8 \end{bmatrix}$

20. Diagonal Matrix  $\Rightarrow \begin{bmatrix} a_{11} & 0 & 0 \\ 0 & a_{22} & 0 \\ 0 & 0 & a_{33} \end{bmatrix}$

25. Upper triangular Mtx.  $\Rightarrow \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ 0 & a_{22} & a_{23} \\ 0 & 0 & a_{33} \end{bmatrix}$

Lower triangular Mtx.  $\Rightarrow \begin{bmatrix} a_{11} & 0 & 0 \\ a_{21} & a_{22} & 0 \\ a_{31} & a_{32} & a_{33} \end{bmatrix}$

Banded Mtx.  $\Rightarrow$

$$\begin{bmatrix} a_{11} & a_{12} & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 \\ 0 & a_{32} & a_{33} & a_{34} \\ 0 & 0 & a_{43} & a_{44} \end{bmatrix}$$

Elements above & below the diagonal Mtx are non-zero, and the rest are zero.

Methods to solve linear eq<sup>n</sup>  $\Rightarrow$

① Elimination of unknown  $\Rightarrow$

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1 \quad \textcircled{1}$$

$$a_{21}x_1 + a_{22}x_2 + a_{23}x_3 = b_2 \quad \textcircled{2}$$

$$\textcircled{1} \times a_{21} \quad \textcircled{2} \times a_{11}$$

$$\Rightarrow x_2 = \frac{a_{11}b_2 - a_{21}b_1}{a_{11}a_{22} - a_{21}a_{12}}$$

$$\textcircled{2} \quad a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1 \quad \textcircled{1}$$

$$a_{21}x_1 + a_{22}x_2 + a_{23}x_3 = b_2 \quad \textcircled{2}$$

$$a_{31}x_1 + a_{32}x_2 + a_{33}x_3 = b_3 \quad \textcircled{3}$$

$$C = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix}$$

$$b = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}$$

$$A = \left[ \begin{array}{ccc|c} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \end{array} \right]$$

$\hookrightarrow$  Augmented Mtx.

$$\textcircled{2} - \textcircled{1} \quad (a_{21} - a_{11})$$

$$(a_{21} - a_{11})$$

$$\textcircled{3} - \textcircled{1} \quad (a_{31} - a_{11})$$

$$(a_{31} - a_{11})$$

Now,  $x_1$  is

$$a_{11}x_1 +$$

\* Pseudo Code

for  $i = 2$

$$u =$$

$$a_{ii} =$$

$$\textcircled{2} - \textcircled{1} \left( \frac{a_{21}}{a_{11}} \right)$$

$$\left( a_{21} - \frac{a_{12} \times a_{11}}{a_{11}} \right) x_1 + \left( a_{22} - \frac{a_{12}}{a_{11}} a_{12} \right) x_2 + \dots$$

$$\left( a_{23} - \frac{a_{12}}{a_{11}} a_{13} \right) x_3 =$$

$$\left( a_{24} - \frac{a_{12}}{a_{11}} a_{14} \right)$$

$$\textcircled{3} - \textcircled{1} \left( \frac{a_{31}}{a_{11}} \right) \Rightarrow \textcircled{3} - \textcircled{1} u_{30} u_{31}$$

$u_{30}$

$$(a_{31} - u_{30} a_{11}) x_1 + (a_{32} - u_{30} a_{12}) x_2 + (a_{33} - u_{30} a_{13}) x_3 = (a_{34} - u_{30} a_{14})$$

Now,  $x_1$  is eliminated from  $\textcircled{2}$  &  $\textcircled{3}$  eq<sup>n</sup>  $\Rightarrow$

$$a_{11} x_1 + a_{12} x_2 + a_{13} x_3 = b_1 \quad \textcircled{1}$$

$$a_{21}' x_2 + a_{23}' x_3 = b_2' a_{24}' \quad \textcircled{6}$$

$$a_{32}' x_2 + a_{33}' x_3 = b_3' a_{34}' \quad \textcircled{7}$$

\* Pseudo Code  $\Rightarrow$

for  $i = 2$  to  $3$  in step of  $1$ .

$$u = a_{ii} / a_{11}$$

$$a_{ii} = 0$$

for  $j = 2 \text{ to } 4$  in step of 1

$$a_{ij} = a_{ij} - u a_{ij}$$

end

end

5

→ Eliminate  $x_1$  from all eq<sup>n</sup> below eq<sup>n</sup> ①.

Now, we will eliminate  $x_2$  below 2<sup>nd</sup> eq<sup>n</sup>.

Then,  $a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1$ ,

$$a_{21}'x_1 + a_{22}'x_2 = b_2'$$

$$a_{23}'x_3 = b_2''$$

→ In general we will eliminate  $x_k$  from all eq<sup>n</sup> below the  $k^{\text{th}}$  eq<sup>n</sup>.

15

for  $i=2 \text{ to } n$

$$u = a_{ii} / a_{11}$$

$$a_{ii} = 0$$

for  $j = 2 \text{ to } (n+1)$

20

$$a_{ij} = a_{ij} - u a_{1j}$$

end

end

25

Now we have to eliminate  $x_k$  from all eq<sup>n</sup> below the  $k^{\text{th}}$  eq<sup>n</sup> →

for  $k = 1 \text{ to } n-1$

for  $i = (k+1) \text{ to } n$

$$u = a_{ik} / a_{kk}$$

$$a_{ik} = 0$$

The al

In this

end

5 end

→ This is

Now, tc

10 do back

$x_3 =$

$x_2 =$

$x_1 =$

15

→  $x_n =$

algo.  
\*\*

for

20

end

$x_n =$

for

25

end.

The above method is Naive - Gauss Elimination

Process.

In this method the diagonal element can't be zero.

for  $j = (k+1)$  to  $(n+1)$

$$a_{ij} = a_{ij} - u_{ik} \cdot a_{kj}$$

end

end

end

→ This is forward Elimination.

Now, to find the value of  $x_1$  to  $x_n$  we will  
do back - substitution.

$$x_3 = a_{3n} / a_{33}$$

$$x_2 = (a_{2n} - a_{23} x_3) / a_{22}$$

$$x_1 = (a_{1n} - a_{12} x_2 - a_{13} x_3) / a_{11}$$

→  $x_n = a_{nn} / a_{nn}$

for  $i = (n-1)$  to 1 in step of (-1)

$$x_i = (a_{ii} - \sum_{j=i+1}^n a_{ij} x_j) / a_{ii}$$

end

$$x_n = a_{nn} / a_{nn}$$

for  $i = (n-1) : (-1) : 1$

$$\text{temp} = 0;$$

for  $j = i+1 : 1 : n$

$$\text{temp} = \text{temp} + (a(i)(j) \cdot x(j))$$

end

$$x(i) = (a(i)(n+1) - \text{temp}) / a(i)(i);$$

end.

$$Q. A = \begin{bmatrix} k_1+k_2 & -k_2 & 0 \\ -k_2 & k_2+k_3 & -k_3 \\ 0 & -k_3 & k_3 \end{bmatrix};$$

(in laptop)

$$5. b = g \begin{bmatrix} m_1 \\ m_2 \\ m_3 \end{bmatrix}; \quad k_1 = 10^3 N/m \quad m_1 = 0.07 \\ k_2 = 15^2 N/m \quad m_2 = 0.065 \\ k_3 = 20^5 N/m \quad m_3 = 0.075$$

2 files  $\Rightarrow$  script file, f' file

$$g = 9.8 m s^{-2}$$

Ans.

$$x_1 = 0.6860$$

$$x_2 = 1.3720$$

$$x_3 = 1.5190$$

24/01/19

In Forward Eqs.  $\rightarrow$ 

- The loop can be run in a single line using vectorization.

$$20. a(i, (k+1):(n+1)) = a(i, (k+1):(n+1)) - u * a(k, (k+1):(n+1));$$

for  $j = (k+1):(n+1)$

$$25. a(i,j) = a(i,j) - u * a(k,j);$$

end

Backward Eqs.

$$x(n) = a(n, n+1) / a(n, n);$$

$$\text{for } i = (n-1):(-1):1$$

$$x(i) = (a(i, n+1) - a(i, (i+1):n) * x(i+1)) / a(i, i);$$

\*  $\rightarrow$  not used  
because we use:

OR,

$\rightarrow \text{temp} = 0$  // It is also not needed.

$\text{temp} = \cancel{\text{temp}} + a(i, i+1:n) * x(i+1:n);$

5  $x(i) = (a(i, n+1) - \text{temp}) / a(i, i)$

Original loop  $\Rightarrow$

$\text{temp} = 0$

for  $j = (i+1):n$

10  $\text{temp} = \text{temp} + a(i, j) * x(j)$

end

P.S.  $\rightarrow$

15 In above method diagonal element can't be zero or some small digit.

So, we will see the modification of the above method  $\Rightarrow$

### \* Pivot Gauss Elimination Process.

$$2x_2 + 3x_3 = 8$$

$$4x_1 + 6x_2 + 7x_3 = -3$$

$$2x_1 + 3x_2 + 6x_3 = 5$$

$$A = \begin{bmatrix} 0 & 2 & 3 \\ 4 & 6 & 7 \\ 2 & -3 & 6 \end{bmatrix} \xrightarrow{\text{Swap}} B = \begin{bmatrix} 8 \\ -3 \\ 5 \end{bmatrix}$$

which ever is highest element  
in 1<sup>st</sup> col swap that  
row with 1<sup>st</sup> row (here)

Note → we are choosing the highest value as we need not to divide by small no.

$$I \quad a = \left[ \begin{array}{ccccc} a_{11} & a_{12} & a_{13} & a_{14} & | & a_{15} \\ a_{21} & a_{22} & a_{23} & a_{24} & | & a_{25} \\ a_{31} & a_{32} & a_{33} & a_{34} & | & a_{35} \\ a_{41} & a_{42} & a_{43} & a_{44} & | & a_{45} \end{array} \right] \quad \begin{matrix} \text{Augmented} \\ \text{Matrix} \\ \text{4 Eqn} \end{matrix}$$

5  
 $k = 1$

→ Consider the sub col<sup>n</sup>

→ Find the highest absolute value (let  $a_{21}$  here).

→ Interchange the 1<sup>st</sup> & 2<sup>nd</sup> row (here).

10 → Proceed to use forward elimination for  $k=1$ .

$$II \quad a = \left( \begin{array}{ccccc} a_{11} & a_{12} & a_{13} & a_{14} & a_{15} \\ 0 & a_{22} & a_{23} & a_{24} & a_{25} \\ 0 & a_{32} & a_{33} & a_{34} & a_{35} \\ 0 & a_{42} & a_{43} & a_{44} & a_{45} \end{array} \right)$$

15  
 $k = 2$

Start here and proceed as above.

20  
III  $k = 3$

IV  $k = 4$

### 1. Partial pivoting

① Take the  $k^{\text{th}}$  column and find the highest value in the  $k^{\text{th}}$  col<sup>n</sup> starting from  $k^{\text{th}}$  row upto the  $n^{\text{th}}$  row.

② Let the biggest value have index  $p$ .

③ If  $p$  is not equal to  $k$

swap the two rows ( $p \leftrightarrow k$ ).

for  $i = (k+1) \text{ to } n$  & so on... (previous code).

$$a = \begin{pmatrix} 3 & 4 & 6 & 6 & 7 \\ 0 & 3 & 9 & 5 & 2 \\ 0 & -8 & 6 & 10 & 1 \\ 0 & 2 & 5 & 9 & 8 \end{pmatrix}$$

$n=4, k=2$

$$a(2:4, 2) \leftrightarrow a(k:n, k)$$



→ We have to choose max of abs. values of col?

$$\max(\text{abs}(a(k:n, k))) = 8$$

|| to get index of 8

$$[\text{big value}, p] = \max(\text{abs}(a(k:n, k))),$$

Index → 2 ← Because we consider the <sup>sub</sup> matrix as red dotted.

$$q = p + (k-1); \rightarrow \text{Index of the original Mtx.}$$

• Partial Pivoting Code ⇒

$$\text{for } k = 1 : (n-1)$$

$$[\text{big value}, p] = \max(\text{abs}(a(k:n, k))),$$

$$q = p + (k-1);$$

Swapping if ( $q \neq k$ )

$$a([k, q], :) = a([q, k], :);$$

end

(Rest similar to the code before)

H.W.

$$\text{Q. } C = \begin{bmatrix} 0 & 2 & 3 \\ 4 & 6 & 7 \\ 2 & -3 & 6 \end{bmatrix}, \quad b = \begin{bmatrix} 8 \\ -3 \\ 5 \end{bmatrix}$$

→ Apply Pivotal Gauss Eli. method to do this

\* Tridiagonal Mtx. →

$$\begin{bmatrix} f_1 & g_1 & 0 & 0 \\ e_2 & f_2 & g_2 & 0 \\ 0 & e_3 & f_3 & g_3 \\ 0 & 0 & e_4 & f_4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} z_1 \\ z_2 \\ z_3 \\ z_4 \end{bmatrix}$$



Band Mtx.

or Tridiagonal Mtx. prob.

Putting this in upper triangular Mtx format

$$\begin{bmatrix} f_1 & g_1 & & \\ f_2 & g_2 & & \\ f_3 & g_3 & & \\ f_4 & g_4 & & \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} z_1 \\ z_2 \\ z_3 \\ z_4 \end{bmatrix}$$

Now, do this we will proceed in this way:

$$\textcircled{2} - \textcircled{1} (e_1/f_1)$$

$$e_2 = e_2 - f_1 \times (e_1/f_1) = 0$$

$$f_2 = f_2 - g_1 \left( \frac{e_1}{f_1} \right)$$

$$g_2 = b_2 - a_{21}(e_1/f_1)$$

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→ Forward Elimination

$$\textcircled{3} - \textcircled{2} \times (e_3/f_2)$$

$$f_3 = b_3 - g_2(e_3/f_2)$$

$$g_4 = b_4 - a_{43}(e_3/f_2)$$

$$\textcircled{4} - \textcircled{3} \times (e_4/f_3)$$

$$f_4 = b_4 - g_3(e_4/f_3)$$

$$g_5 = b_5 - a_{54}(e_4/f_3)$$

∴ So here we get the upper slar matx sys.  $\Rightarrow$

$$i = 2 \text{ to } n \text{ in step of 1.}$$

$$f_i = f_{i-1} - g_{i-1}(e_i/f_{i-1})$$

$$g_i = a_i - a_{i-1}(e_i/f_{i-1})$$

} In Code

Now, we will do Backward Substitution here  $\Rightarrow$

$$f_1 x_1 + g_1 x_2 = a_1$$

$$f_2 x_2 + g_2 x_3 = a_2$$

$$f_3 x_3 + g_3 x_4 = a_3$$

$$f_4 x_4 = a_4$$

$$x_4 = a_4/f_4$$

$$x_3 = \frac{a_3 - g_3 x_4}{f_3}$$

$$x_2 = \frac{a_2 - g_2 x_3}{f_2}$$

$$x_1 = \frac{a_1 - g_1 x_2}{f_1}$$

We have,

$$x_n = a_n/f_n;$$

$$x_i = \left( \frac{a_i - g_i * x_{i+1}}{f_i} \right) . \quad i = (n-1) \text{ to } 1 \text{ in steps of } (-1).$$

⇒ Only next (Tuesday)  $\approx$  30 min. Root finding

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H.W.

$$\text{Q.} \quad \left[ \begin{array}{cccc|c} 2.04 & -1 & 0 & 0 & x_1 \\ -1 & 2.04 & -1 & 0 & x_2 \\ 0 & -1 & 2.04 & -1 & x_3 \\ 0 & 0 & -1 & 2.04 & x_4 \end{array} \right] = \left[ \begin{array}{c} 40.8 \\ 0.8 \\ 0.8 \\ 200.8 \end{array} \right]$$

5

(By code).

Ans.

$$x_1 = 20.287$$

$$x_2 = 0.5886$$

$$x_3 = 0.3971$$

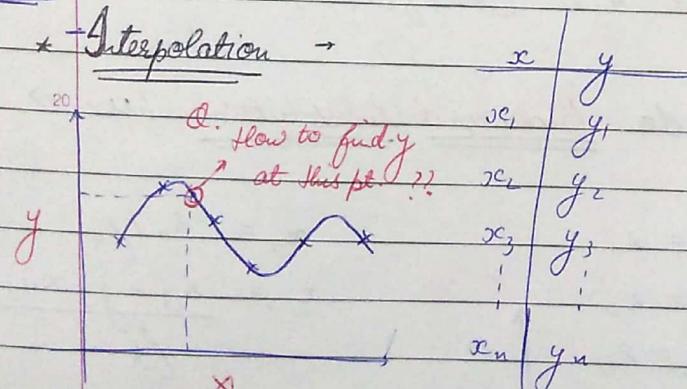
$$x_4 = 0.0102$$

Note → Don't create an augmented M+x. here.

$$\left\{ \begin{array}{l} e = [0 \ -1 \ -1 \ -1] \\ f = [2.04 \ 2.04 \ 2.04 \ 2.04] \\ g = [-1 \ -1 \ -1 \ 0] \\ h = [40.8 \ 0.8 \ 0.8 \ 200.8] \end{array} \right.$$

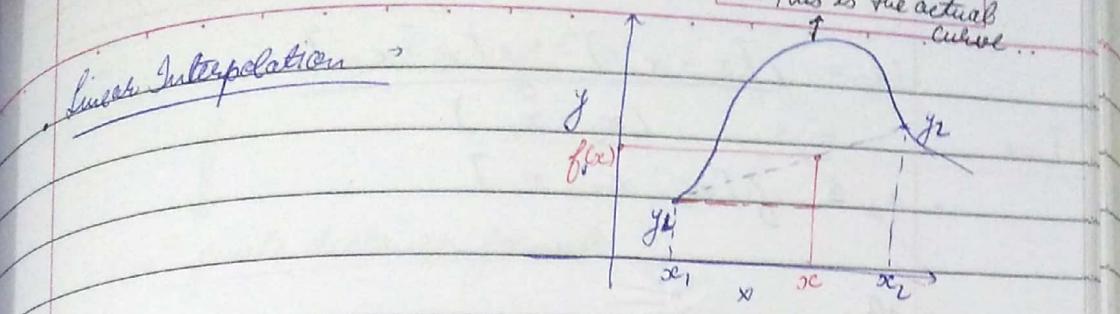
29/01/13

\* Interpolation →



→ This can be done by Interpolation.

I Newton's Interpolating Polynomial →



$x$	$y$
$x_1$	$y_1 = f(x_1)$
$x_2$	$y_2 = f(x_2)$

$$\frac{f(x) - f(x_1)}{(x - x_1)} = \frac{f(x_2) - f(x_1)}{(x_2 - x_1)}$$

$$\Rightarrow f(x) = f(x_1) + \frac{f(x_2) - f(x_1)}{(x_2 - x_1)} (x - x_1)$$

↳ 1<sup>st</sup> order polynomial fitting.

### Quadratic Interpolation →

$$f(x) = c_0 + c_1 x + c_2 x^2$$

we need atleast 3 datapoints to fit find  $c_0, c_1, c_2$

so, by generalization to find  $n^{th}$  order polynomial  
we need atleast  $(n+1)$  data points.

$$f(x) = a_0 + a_1 (x - x_1) + a_2 (x - x_1)(x - x_2)$$

$$f(x_1) = a_0$$

$$f(x_2) = a_0 + a_1 (x_2 - x_1) = f(x_1) + a_1 (x_2 - x_1)$$

$$f(x) =$$

$$\Rightarrow a_2 = \frac{f(x_2) - f(x_1)}{x_2 - x_1} = f[x_1, x_2]$$

$$a_3 = \frac{f[x_3 - x_1] - f[x_2 - x_1]}{(x_3 - x_1)}$$

$$\therefore a_3 \equiv f[x_3, x_2, x_1]$$

Now,

$$f(x) = f(x_1) + f[x_2, x_1](x - x_1) + \\ f[x_3, x_2, x_1](x - x_1)(x - x_2)$$

\* General form of Newton's Interpolating Poly.

$$f_{n-1}(x) = a_0 + a_1(x - x_1) + a_2(x - x_1)(x - x_2) + \\ \dots + a_n(x - x_1)(x - x_2) \dots (x - x_{n-1})$$

$$\underline{x < x_1} \quad \text{or} \quad \underline{x > x_n}$$

$$\Rightarrow \boxed{x_1 < x < x_n}$$

*If we know the limits then  
only we can find the point  
betw them only.*

$$a_0 = f(x_1)$$

$$a_1 = f[x_2, x_1]$$

!

$$a_n = f[x_n, x_{n-1}, \dots, x_1]$$

where,

$$f[x_i, x_j] = \frac{f(x_i) - f(x_j)}{(x_i - x_j)}$$

$$f[x_i, x_j, x_k] = \frac{f[x_i, x_j] - f[x_j, x_k]}{(x_i - x_k)}$$

$$f[x_0, x_1, \dots, x_n] = f[x_0, x_1, \dots, x_n] - \frac{f[x_{n-1}, \dots, x_n]}{(x_n - x_1)}$$

(n<sup>th</sup> finite difference formula)

Eg. (for n=4)  $\Rightarrow$  no. of data points

$$f[x_0, x_1, x_2, x_3] = f[x_0, x_1, x_2] - \frac{f[x_1, x_2, x_3]}{(x_3 - x_1)}$$

$$f(x) = f(x_0) + f[x_0, x_1](x - x_0) + \\ f[x_0, x_1, x_2](x - x_0)(x - x_1) + \\ f[x_0, x_1, x_2, x_3](x - x_0)(x - x_1)(x - x_2)$$

To code this we write it in table format  $\Rightarrow$

$x_i$	$y_i = f(x_i)$	$j=1$	$j=2$	$j=3$	$j=4$
$x_0$	$f(x_0) = b_{01}$		$f[x_0, x_1] = b_{12}$	$f[x_0, x_1, x_2] = b_{13}$	$f[x_0, x_1, x_2, x_3] = b_{14}$
$x_1$	$f(x_1) = b_{21}$		$f[x_1, x_2] = b_{22}$	$f[x_1, x_2, x_3] = b_{23}$	$f[x_1, x_2, x_3, x_4] = b_{24}$
$x_2$	$f(x_2) = b_{31}$			$b_{32} = 0$	$b_{33} = 0$
$x_3$	$f(x_3) = b_{41}$			$b_{42} = 0$	$b_{43} = 0$
$x_4$	$f(x_4) = b_{41}$				

So, we need to determine this B matrix  $\Rightarrow$

$$\underline{\text{2nd col}}$$

$$b_{12} = \frac{b_{21} - b_{11}}{(x_2 - x_1)}$$

$$\underline{\text{3rd col}}$$

$$b_{13} = \frac{b_{31} - b_{21}}{(x_3 - x_1)}$$

$$b_{22} = \frac{b_{31} - b_{21}}{(x_3 - x_2)}$$

$$b_{23} = \frac{b_{32} - b_{22}}{(x_4 - x_2)}$$

$$b_{24} = \frac{b_{41} - b_{31}}{(x_4 - x_3)}$$

$4^{\text{th}} \text{ col}^4 \Rightarrow$

$$b_{14} = \frac{b_{13} - b_{12}}{(x_4 - x_1)}$$

for  $j = 2 \text{ to } n$        $\overbrace{*}^{i_{\max}}$

for  $i = 1 \text{ to } \underline{n-j+1}$

$$b(i, j) = \frac{b(i+1, j-1) - b(i, j-1)}{x(i+j-1) - x(i)}$$

10

end

end

$*p$

↓

For  $p$

Proposed one

15

$j=2$	$i=1$	$p=2$	$i+j=3$	$i+j-1=2$
	2	3	4	3
	3	4	5	4
$j=3$	$i=1$	3	4	3
	2	4	5	4
$j=4$	$i=1$	4	5	4

20

Correct one

$$y_f = f(x) = f_0(x) + f[x_1, x_2](x_f - x_1) + \\ f[x_2, x_3, x_1](x_f - x_2)(x_f - x_1) + \\ f[x_3, x_4, x_2, x_1](x_f - x_3)(x_f - x_2)(x_f - x_1)$$

$$y_f = b_{11} + b_{12}(x_f - x_1) + b_{13}(x_f - x_2)(x_f - x_1) + \\ b_{14}(x_f - x_3)(x_f - x_2)(x_f - x_1)$$

# MATLAB : polyfit & polyval

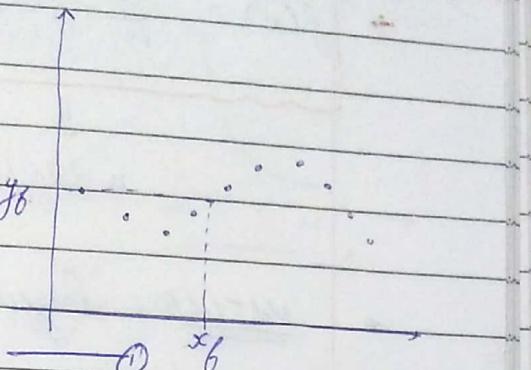
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Lagrange's Interpolation →

→ No error in the datapoints.

$n=3$

$$f(x) = c_1 (x-x_2)(x-x_3) + c_2 (x-x_1)(x-x_3) + c_3 (x-x_1)(x-x_2)$$



$x_1$	$y_1 = f(x_1)$
$x_2$	$y_2 = f(x_2)$
$x_3$	$y_3 = f(x_3)$

$$c_1 = \frac{f(x_1)}{(x_1 - x_2)(x_1 - x_3)}$$

$$c_2 = \frac{f(x_2)}{(x_2 - x_1)(x_2 - x_3)}$$

$$c_3 = \frac{f(x_3)}{(x_3 - x_1)(x_3 - x_2)}$$

Replace  $c_1, c_2, c_3$  in eq "①"

$$f(x) = \frac{f(x_1)}{(x_1 - x_2)(x_1 - x_3)} (x-x_2)(x-x_3) + \frac{f(x_2)}{(x_2 - x_1)(x_2 - x_3)} (x-x_1)(x-x_3)$$

$$f(x_2) \frac{(x-x_1)(x-x_3)}{(x_2 - x_1)(x_2 - x_3)} +$$

$$f(x_3) \frac{(x-x_1)(x-x_2)}{(x_3 - x_1)(x_3 - x_2)}$$

Now, for  $n$  data points ⇒

$$\rightarrow f(x) = \sum_{i=1}^n f(x_i) \prod_{\substack{j=1 \\ i \neq j}}^n \frac{(x - x_j)}{(x_i - x_j)}$$

↓  
n data points

\* MATLAB : Polyfit →

Maths →  $f(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n$

10 MATLAB →

$$f(x) = p_1 x^n + p_2 x^{n-1} + \dots + p_{n-2} x^2 + p_{n-1} x$$

Eg →  $T = [300, 400, 500]$

$$\text{density} = [0.616, 0.525, 0.427]$$

$$T_f = 350;$$

$$p = \text{polyfit}(T, \text{density}, 2)$$

↓  
Order of poly.

20 Array of Coefficient

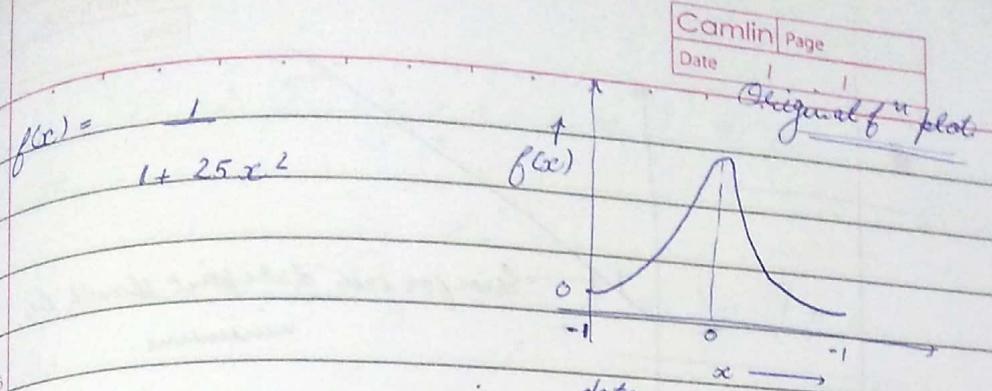
$$p = [p_1, p_2, p_3]$$

$$df = \text{polyval}(p, T_f);$$

↓  
density

↓  
It gives the value of polynomial p at T\_f

\* Dangers of Higher Order Polynomial →



Let we have  $n = 12 \rightarrow$  no. of data points

$x_f = \text{linspace}(-1, 1)$

$p = \text{polyfit}(x_c, y, 11)$

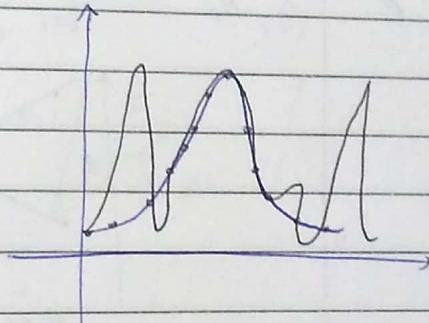
$y_f = \text{polyval}(p, x_f)$

$\text{plot}(x_c, y);$

$\text{hold('on')}$ :

$\text{plot}(x_f, y_f);$  →

$\text{hold('off')}$ :



\* The plot with the help of `polyfit` & `polyval` will not fit at the end perfectly.

\* Also, plotting higher order polynomial is can be erroneous. So, if you have 20 dp then you don't need to make a 19 order polynomial.

\* Linear Regression →

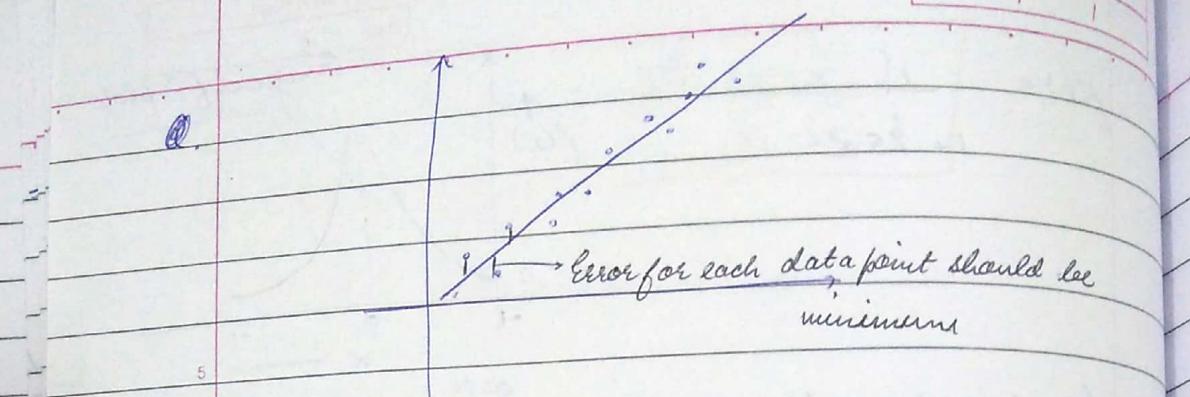
\* When linear data has some error & we need to find the best fit line.

$$y = a_1 + a_2 x + e$$

$$y_1 = a_1 + a_2 x_1 + e_1$$

$$y_2 = a_1 + a_2 x_2 + e_2$$

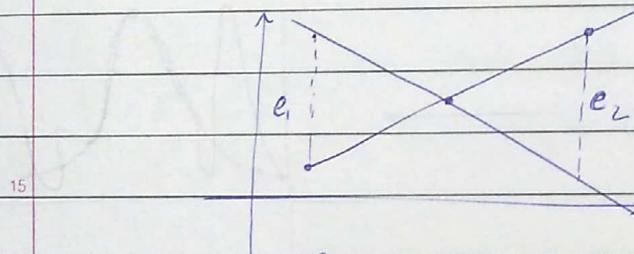
Q. Which one is the best fit line??



$$e = y - a_1 - a_2 x$$

. Criteria for best fit  $\Rightarrow$

$$\text{I } e = \sum_{i=1}^n e_i = \sum_{i=1}^n (y_i - a_1 - a_2 x_i)$$



Both line can be best fit as the +ve & -ve error cancel out each other.

So, it is not a good criteria.

$$\text{II } e = \sum_{i=1}^n |e_i|$$

$\Rightarrow$  It is also not a very good criteria.

$$\text{III } S_e = \sum_{i=1}^n e_i^2 = \sum_{i=1}^n (y_i - a_1 - a_2 x_i)^2$$

$\Rightarrow$  We need to minimize  $S_e$  (summation of square of error).

and find  $a_1$  &  $a_2$ .

$$\frac{\partial S_a}{\partial a_1} = -2 \sum (y_i - a_1 - a_2 x_i) = 0 \quad (i)$$

$$\frac{\partial S_a}{\partial a_2} = -2 \sum [(y_i - a_1 - a_2 x_i) x_i] = 0 \quad (ii)$$

→ Taking (i) & (ii) →

$$\Sigma y_i - n a_1 = n a_2 + (\sum x_i) a_2$$

$$(\sum x_i) a_1 + (\sum x_i^2) a_2 = \sum x_i y_i$$

So, we have →

$$\begin{bmatrix} n \sum x_i \\ \sum x_i \sum x_i^2 \end{bmatrix} \begin{bmatrix} a_1 \\ a_2 \end{bmatrix} = \begin{bmatrix} \sum y_i \\ \sum x_i y_i \end{bmatrix}$$

If we have two datasets →

x	y
10	25
20	70
30	380
40	550
50	610
60	1220
70	830
80	1450

→ Find the best fit line.

(MATLAB)

square of

\* Linearization of non-linear data →

$$x = [ ] ; y = [ ] ;$$

$$\textcircled{1} \quad y = \alpha_1 x^{\beta_1}$$

→ Fitting non-linear data <sup>in</sup> linear line.

$$\begin{matrix} 5 & \ln y = \ln \alpha_1 + \beta_1 x \\ \downarrow & \downarrow \\ y' & a_1 \end{matrix}$$

$$\downarrow \quad \downarrow \quad \downarrow$$

$$y' \quad a_1 \quad a_2$$

$$\textcircled{2} \quad y = \alpha_2 x^{\beta_2}$$

$$\begin{matrix} 10 & \log y = \log \alpha_2 + \beta_2 \log x \\ \downarrow & \downarrow \quad \downarrow \\ y' & a_1 \quad a_2 \end{matrix}$$

$$\downarrow \quad \downarrow \quad \downarrow$$

$$y' \quad a_1 \quad a_2 \quad x'$$

$x'$	$y'$

$$\textcircled{3.19} \quad y = \frac{x_3}{\beta_3 + x}$$

$$\Rightarrow \beta_3 y + xy = x_3 x$$

$$\Rightarrow x = \left( \frac{\beta_3}{x_3} \right) y + \frac{xy}{x_3}$$

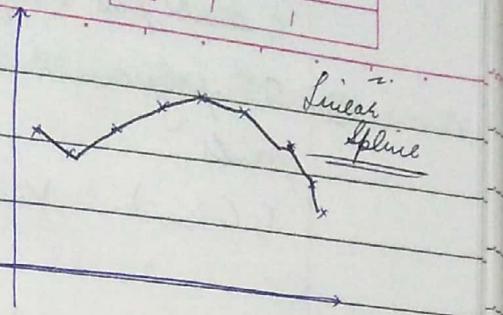
$$\Rightarrow \left( \frac{x}{y} \right) = \left( \frac{\beta_3}{x_3} \right) + \left( \frac{x}{x_3} \right)$$

$$\begin{matrix} 25 & 1 = \left( \frac{\beta_3}{x_3} \right) x + \left( \frac{1}{x_3} \right) \\ \downarrow & \downarrow \quad \downarrow \\ y' & a_2 \quad x' \quad a_1 \end{matrix}$$

- These all are the ways of linearization of data

## \* Spline Interpolation →

### I Linear Spline →



→ For  $n$  data points, fits a polynomial between a pair of data points  $\rightarrow (n-1)$

$x$	$y$	}
$x_1$	$y_1$	
$x_2$	$y_2$	
$\vdots$	$\vdots$	
$x_n$	$y_n$	

→ Linear Spline

$$S_i = a_i + b_i(x - x_i)$$

$$\text{No. of unknowns} = 2(n-1)$$

(Steps)  $\Rightarrow$

- ① Find the coefficients  $a_i, b_i, \dots$
- ② Check  $x_c < x_f < x_n$ .
- ③ Find the  $k^{\text{th}}$  bin where  $x_f$  lies.
- ④  $y_f = \sum_{i=1}^k b_i(x_f) = \dots$

- ⑤ The  $S_i^{\text{th}}$  spline / polynomial should pass through  $i^{\text{th}}$  point.

$$y_i = S_i(x_i) = a_i + b_i(x_i - x_i)$$

$$\Rightarrow y_i = a_i \quad ; \quad i = 1, 2, \dots, (n-1)$$

of data

$s_1$  should pass through pt. 2,  $s_2$  pass through pt. 3 & so on.

②  $s_i^{th}$  polynomial should pass through  $(x_{i+1})^{th}$  data points

$$s_i(x_{i+1}) = Y_{i+1} = a_i + b_i(x_{i+1} - x_i)$$

$$b_i = \left( \frac{Y_{i+1} - Y_i}{x_{i+1} - x_i} \right); i = 1, 2, \dots, (n-1)$$

Find the  $\frac{k}{th}$  term  $\rightarrow$

1 Sequential search  $\Rightarrow$

if ( $x_f < x(1) \& x_f > x(n)$ )  
exit with suitable msg.

else.

$$bin = 0;$$

for  $i = 1 : n$

if ( $x(i) < x_f$ ) ~~Observe  $x_f < x(i)$~~

bin++;

else

break;

end

end.

end.

bin's  $\Rightarrow$  ①  $n = \text{length}(x)$

② Check whether  $x_f$  lies in  $x(1), x(2), \dots, x(n)$  or not.

③  $k = 1$

while(1)

if  $x_f$  less than  $x(k+1)$

break loop

endk = k+1;

But this is not a good method if no. of data points are very large.

So, we will use II Binary Search.

i/p  $\rightarrow x_c, x_f$

o/p  $\rightarrow k$  [bin no.]

for while(~~around~~)

index =  $((1+n)/2)$ ; i2 = 1;

while (1)

if ( $x_f < x(index)$ )  $\&$   $x_f > x(index-1)$ )  
 $k = index - 1$ ; break;

else

if ( $x_f < x(index)$ )

~~index = i2~~ <sup>round</sup>

index =  $((i2 + index)/2)$

i

upto point @ same

③ iL = 1; iU = n;

while (1)

if ( $iU - iL \leq 1$ )

~~k = iL~~, break;

end

iM =  $\text{floor}((iL + iU)/2)$

index if ( $x(iM) < x_f$ )

$xL = iM$ ;

else

$xU = xM$ ;

end

end

$k = iL$ ;

HW MATLAB)  $x = [-40, 0, 20, 50, 100, 150, 100, 250, 300,$   
 $400, 500, 450, 7]$

$x_f = 350$ ; find the  $k^{\text{bin}}$ .

II Quadratic Spline

$$S_i(x) = a_i + b_i(x - x_i) + c_i(x - x_i)^2$$

5 Unknowns  $\rightarrow 3(n-1)$

$i = 1, 2, 3, \dots, (n-1)$

①  $S_i$  must pass through  $i^{th}$  data points.

$$10 \quad S_i(x_i) = y_i = a_i$$

$$\Rightarrow a_i = y_i \quad \rightarrow i = 1, 2, \dots, (n-1)$$

$$\text{Unknowns} = 3(n-1) - (n-1) = 2(n-1)$$

②  $S_i$  must pass through  $(i+1)^{th}$  point.

$$S_i(x_{i+1}) = y_{i+1} = y_i + b_i(x_{i+1} - x_i) + c_i(x_{i+1} - x_i)^2$$

$$20 \quad = y_i + b_i h_i + c_i h_i^2$$

$$\text{; where } h_i = (x_{i+1} - x_i)$$

$$\text{Unknowns} \rightarrow 2(n+1) - (n-1) = (n-1)$$

25 [As we have  $(n-1)$  cond<sup>n</sup> to find out the unknown from above]

③ Visually smooth  $\rightarrow$

$\rightarrow 1^{st}$  order derivative at the nodes must be equal

$$S_i'(x) = b_i + 2c_i(x - x_i)$$

$\Rightarrow$  To make smooth  $\rightarrow$

$$S_i'(x_{i+1}) = S_{i+1}'(x_{i+1})$$

$$\therefore i = 2, 3, \dots, (n-1)$$

$\Rightarrow (n-2)$  cond'n

$$\Rightarrow (n-1) - (n-2) = 1$$

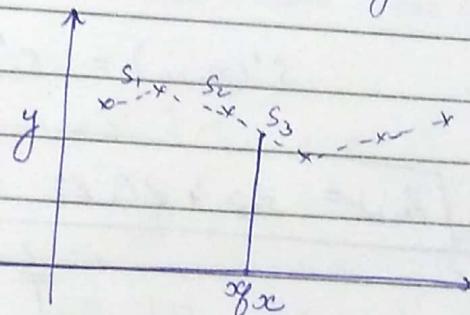
### (i) Natural Spline

$$S_i''(x) = 2c_i$$

$$c_i = 0$$

1<sup>st</sup> spline is a straight line

### \* Cubic Spline $\rightarrow$



$$S_i(x) = a_i + b_i(x - x_i) + c_i(x - x_i)^2 + d_i(x - x_i)^3$$

$i^{\text{th}}$  interval (line)

$$\text{Unknowns} = 4(n-1)$$

① So must pass through  $(x_i, y_i)$

$$S_i(x_i) = y_i = a_i \Rightarrow a_i = y_i$$

$$(i = 1, 2, \dots, (n-1))$$

$$\text{No. of reqd. const.} = 4(n-1) - (n-1) = 3(n-1)$$

Linear Spline

1<sup>st</sup> order derivative of line 1-2

- ②  $S_i(x)$  must pass through  $(x_{i+1}, y_{i+1})$

$$y_{i+1} = S_i(x_{i+1}) = a_i + b_i h_i + c_i (h_i)^2 + d_i h_i$$

$\uparrow$   
 $\underbrace{y_i}_{\text{---}}$

$$i = 1, 2, \dots, (n-1)$$

$$h_i = (x_{i+1} - x_i) ; i = 1, 2, \dots, n-1$$

$$\text{No. of regd. constraints} = 3(n-1) - (n-1) = 2(n-1)$$

- ③ 1<sup>st</sup> order derivative for the interior nodes must be equal.

$$S'_i(x) = b_i + 2c_i(x - x_i) + 3d_i(x - x_i)^2$$

$$S'_i(x_{i+1}) = S'_{i+1}(x_{i+1})$$

$$\boxed{b_{i+1} = b_i + 2c_i h_i + 3d_i h_i^2}$$

$i = 2, 3, \dots, (n-1)$

No. of cons. regd.  $\rightarrow$   
 $2(n-1) - (n-2) = \underline{n}$

As const. is for interior nodes

- ④ 2<sup>nd</sup> order derivative at the interior nodes must be equal.

$$S''_i(x) = 2c_i + 6d_i(x - x_i)$$

$$S''_i(x_{i+1}) = S''_{i+1}(x_{i+1})$$

$$\Rightarrow \boxed{d_i = \frac{c_{i+1} - c_i}{3h_i}}$$

$$\text{No. of const. segd.} = n(n-2) = \frac{2}{2} \\ i = 2, 3, \dots, (n-1)$$

⑤ Natural Spline Cond<sup>n</sup> - (to get 2 more cond<sup>n</sup>)  
2nd order derivative of end nodes are zero

$$S_1''(x_1) = c_1 + 6d_1(x_1 - x_0) = 0 \\ \Rightarrow \boxed{c_1 = 0}$$

$$10. S_{n-1}''(x_n) = 2c_{n-1} + 6d_{n-1}(x_n - x_{n-1}) \\ = c_{n-1} + 3d_{n-1}h_{n-1} = 0 \\ \downarrow \\ c_n = 0$$

 $S_n(x)$ 

→ Last 2 cond<sup>n</sup> are free.

$$h_i, c_{i-1} + 2(h_{i-1} + h_i)c_i + h_i c_{i+1} = \\ 3(y[x_{i+1}, x_i] - y[x_i, x_{i-1}])$$

~~Algebra~~ To find  $c_i$   $i = 2, 3, \dots, (n-1)$

$$y[x_{i+1}, x_j] = y_{ij} = \frac{y_i - y_j}{x_i - x_j} = \frac{y_i - y_j}{h_i}$$

$$25. b_{i+1} = b_i + 2c_i h_i + 3d_i h_i^2 \quad \textcircled{2}$$

$$b_i = \frac{y_{i+1} - y_i}{h_i} - \frac{h_i}{3} (2c_i + c_{i+1})$$

$$d_i = \frac{c_{i+1} - c_i}{3h_i} \quad i = 2, 3, \dots, (n-1) \quad \boxed{a_i = y_0}$$

$$h_i = (x_{i+1} - x_i) \quad ; \quad i = 1, 2, \dots, n$$

i=1:

$$C_1 = a.$$

i=2:

$$h_1 C_1 + 2(h_1 + h_2) C_2 + h_2 C_3 = \\ 3(y[x_3, x_2] - y[x_2, x_1])$$

i=3:

$$h_2 C_2 + 2(h_2 + h_3) C_3 + h_3 C_4 = \\ 3(y[x_4, x_3] - y[x_3, x_2])$$

i=(n-1)

$$h_{n-2} C_{n-2} + 2(h_{n-2} + h_{n-1}) C_{n-1} + h_{n-1} C_n = \\ 3(y[x_n, x_{n-1}] - y[x_{n-1}, x_n])$$

i=n:

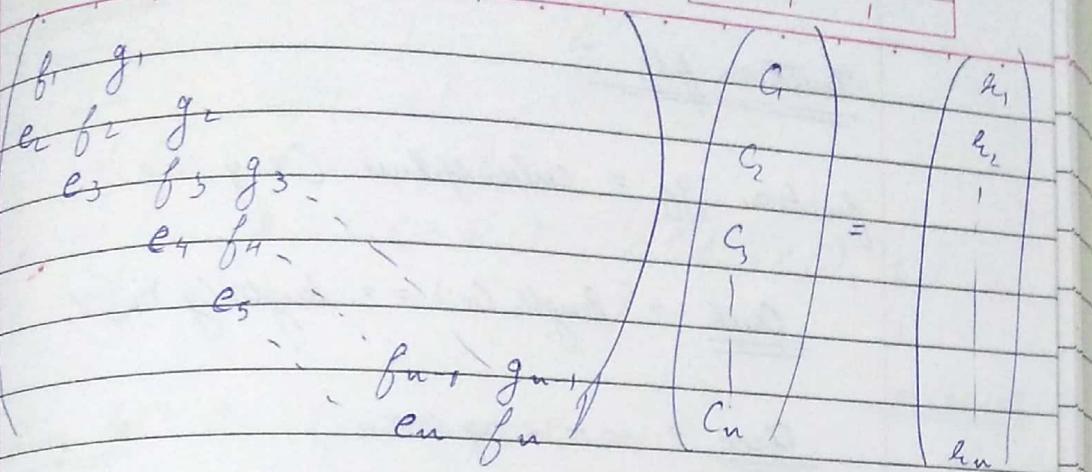
$$C_n = a$$

+ In Mtx. form the above eq<sup>n</sup> will form:  
Tridiagonal Mtx.

$$\begin{pmatrix} 1 & 0 & & & \\ h_1 & 2(h_1 + h_2) & h_2 & & \\ & h_2 & 2(h_2 + h_3) & h_3 & \\ & & h_3 & \ddots & \\ & & & h_{n-2} & 2(h_{n-2} + h_{n-1}) \\ & & & & h_{n-1} \end{pmatrix} \begin{pmatrix} C_1 \\ C_2 \\ \vdots \\ C_n \end{pmatrix} = \begin{pmatrix} a \\ 3(y_{3,2} - y_{2,1}) \\ 3(y_{4,3} - y_{3,2}) \\ \vdots \\ 3(y_{n,n-1} + y_{n-1,n-2}) \end{pmatrix}$$

$a, b, d \rightarrow$  array  $(n-1)$  size  
 $c \rightarrow$  array

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$$c = [a, h_1, h_2, \dots, h_{n-1}, 0] = [a, h(1:(n-1)), 0]$$

$$g = [0, h_2, h_3, \dots, h_n, 0] = [0, h(2:(n-1)), 0]$$

$$f_1 = 1, f_n = 1$$

$$f_i = 2(h_{i-1} + h_i)$$

$$h_1 = 0$$

$$h_n = 0$$

$$a_i = \frac{3(g_{i+1} - g_i)}{h_i} - \frac{3(g_i - g_{i-1})}{h_{i-1}}$$

$$; i = 2, 3, \dots, (n-1)$$

$$h_i = x_{i+1} - x_i ; i = 1, 2, \dots, (n-1)$$

\* Algo  $\Rightarrow$

Script file  $\Rightarrow$

$$x = [ ]$$

$$y = [ ]$$

$$x_f = \underline{\quad}$$

$y_f = \text{cubic spline}(x, y, x_f)$

plot(x, y, 'b-', x\_f, y\_f, 'ro');

Function file  $\Rightarrow$

function  $y_f = \text{cubic spline}(x, y, x_f)$

5

Check  $\rightarrow \text{length}(x) == \text{length}(y)$

10

Check  $\rightarrow x_c < x_f < x_n$

find  $n$

find  $h \rightarrow (n-1)$

find  $e, f, g, a$   $\rightarrow n$

$n \rightarrow c = \text{tridiag}(e, f, g, a)$

15

find  $a, b, d$   $\rightarrow (n-1)$

(Integer)  $k = \text{look up bin}(x_c, x_f)$

[Done by Binary search]

①  $y_f$   
 $(k+1) \downarrow$   
 giving the apt ans.  
 in place of  $k$ .

$$y_f = a(k) + b(k)(x_f - x(k)) + c(k)(x_f - x(k))^2 + d(k)(x_f - x(k))^3$$

end  
 It is  $k$ .

$\rightarrow$   $f^n$  file, 1 script file.

25

Q.  $x = [3, 4.5, 2, 3]$

$y = [2.5, 1, 2.5, 0.5]$

$x_f = 5$

$y_f = ?$

$\rightarrow$  ① this is not apt. in some cases

→ Ans. is coming right when you are using (K+H)  
in place of K in the eqn.  
It is K.

h → 3

e → [0, h(1), 0]

g → [0, h(2), 0]

f → [1, 2/3(h<sub>1</sub> + h<sub>2</sub>), 2(h<sub>2</sub> + h<sub>3</sub>), 1]

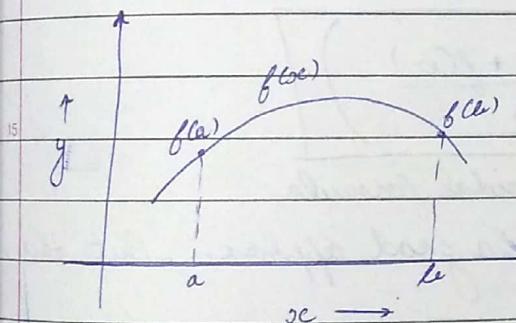
h → [0, \_\_\_\_\_, \_\_\_\_\_] = 3

\* Only Sol<sup>n</sup> → (you know the ans.)

12/02/19

Integration →

$$I = \int_a^b f(x) dx$$



There are two types of integration:

① f(x) → difficult f<sup>n</sup> to integrate analytically.

② (x, y) → set of data points given.

$$I = \int_a^b f(x) dx = \int_a^b f_{n,p}(x) dx$$

↑  
Order of polynomial

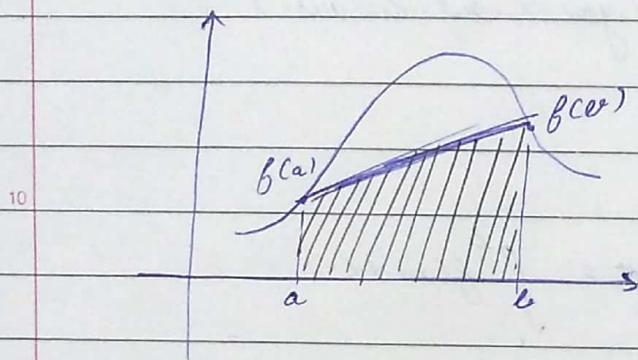
*I Newton-Cotes formula.*

Substitute f(x) → complicated f<sup>n</sup> with a proper poly. f<sub>n,p</sub>(x).

•  $n=1 \rightarrow$  Linear Polynomial  
 $\downarrow$   
 (LM)

Name: Trapezoidal Method (for n=1)

$$I = \int_a^b f(x) dx = \int_a^b \left\{ f(a) + \frac{(f(b)-f(a))}{b-a} (x-a) \right\} dx$$



$$I = (b-a) \left( \frac{f(a) + f(b)}{2} \right)$$

Trapezoidal formula

→ It will not give us a good approx., but this is the beginning.

• Error in Trapezoidal Method →

① It will give good results if a & b are very close to each other otherwise not.

→ Expand  $f(x)$  in the interval  $[a, b]$  in Taylor series abt zero.

$$I = \int_a^b f(x) dx$$

$$\rightarrow b - a$$

$h = \text{diff. bet } n \text{ a \& b}$  is very small.

(Since Post Rule)  $\rightarrow$  Google it Camlin Page  
Date if confuse

$$f(x) = f(a) + xf'(a) + \frac{x^2}{2!} f''(a) + \frac{x^3}{3!} f'''(a) + \dots$$

$$I_1 = \int_a^b f(x) dx = \int_0^h [ \text{Taylor Series} ] dx$$

$$I_1 = h f(a) + \frac{h^2}{2} f'(a) + \frac{h^3}{3!} f''(a) + \dots$$

$$I_2 = \frac{h}{2} [ f'(a) + f(a) ] \rightarrow \text{Trapezoidal TM}$$

• Expand  $f(x)$  about 0.

$$= \frac{h}{2} \left[ f(a) + \left\{ f(a) + \frac{hf'(a)}{2!} + \frac{h^2 f''(a)}{3!} + \dots \right\} \right]$$

$$I_2^{\text{trap.}} = h f(a) + \frac{h^2}{2} f'(a) + \frac{h^3}{4} f''(a) + \dots$$

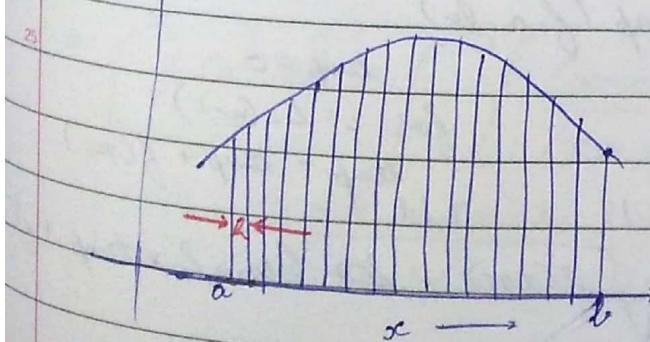
$$\Rightarrow e_t^{\theta} = -\frac{h^3}{12} f'''(a) - \frac{h^4}{24} f^{(4)}(a)$$

$$e_t = I_1^{\text{exact}} - I_2^{\text{trap.}}$$

\* Composite ~~the~~ TM  $\rightarrow$

$$e_t \approx -\frac{h^3}{12} f''(c)$$

Error due to TM



$$h = \frac{(b-a)}{n-1}$$

$$I = \int_a^b f(x) dx$$

$$= \int_a^{x_1} f(x) dx + \int_{x_1}^{x_2} f(x) dx + \dots + \int_{x_{n-1}}^{x_n} f(x) dx$$

Imp. one

(This is derived from the pic. TM)

$$I = \frac{h}{2} \left[ \{f(x_1) + f(x_2)\} + \{f(x_2) + f(x_3)\} + \dots + \{f(x_{n-1}) + f(x_n)\} \right]$$

$$\Rightarrow I = \frac{h}{2} \left[ f(x_1) + 2 \sum_{i=2,3}^{n-1} f(x_i) + f(x_n) \right]$$

Code  $\Rightarrow$ 

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clear variables;

script file

```

a = 0 ;
b = 3 * pi;
f = @(x) sin(x^2);
res = trap(f,a,b);

```

function res = trap(f,a,b)

f file

```

h = 0.01;
temp = 0
x = a:h:b;
for i = 2:(n-1)
    temp = temp + f(xi)
end
temp = temp + f(xn)
res = h/2 * {f(xa) + 2 * temp + f(xn)}

```

\* Simpson's  $\frac{1}{3}$  rule  $\Rightarrow$

$f_2(x)$  = quadratic polynomial

$$I = \int_a^b f(x) dx \approx \int_a^b f_2(x) dx$$

$$f_2(x) = C_1(x - x_1)(x - x_2) + C_2(x - x_2)(x - x_3) + C_3(x - x_1)(x - x_3)$$

$$\underline{x_1 = a}, \quad \underline{x_2 = x_3 = b}$$

$$C_1 = \frac{f(x_1)}{(x_2 - x_1)(x_3 - x_1)}$$

$$C_2 = \frac{f(x_2)}{(x_3 - x_2)(x_2 - x_1)}$$

$$C_3 = \frac{f(x_3)}{(x_3 - x_1)(x_3 - x_2)}$$

$$I = \int_{x_1}^{x_3} f_2(x) dx$$

$$= \int_{x_1}^{x_3} f_2(x) dx = \int_{x_1}^{x_3} \left[ \frac{(x - x_1)(x - x_2)}{(x_1 - x_2)(x_2 - x_3)} f(x_1) + \frac{(x - x_2)(x - x_3)}{(x_1 - x_2)(x_2 - x_3)} f(x_2) + \frac{(x - x_1)(x - x_3)}{(x_2 - x_1)(x_2 - x_3)} f(x_3) \right] dx.$$

$$I = \frac{4}{3} [f(x_1) + 4f(x_2) + f(x_3)]$$

Simpson's  $\frac{1}{3}$  rule

Require 3  
data points

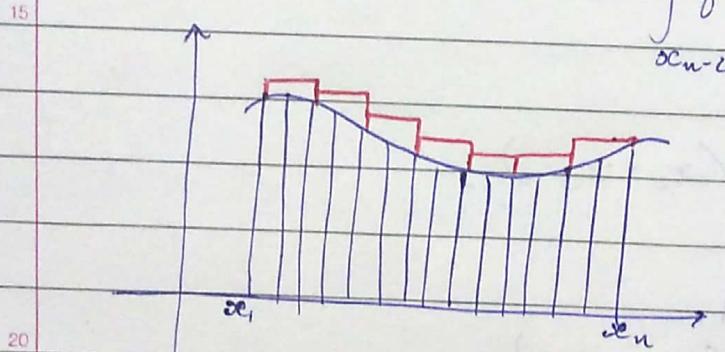
$$e_{S\frac{1}{3}} = -\frac{h^5}{90} f'''(c)$$

$$h = (x_2 - x_1) = (x_3 - x_2)$$

\*  $n$  should be odd for using this.

\* Composite Simpson's  $\frac{1}{3}$  formula rule →

$$I = \int_{x_1}^{x_3} f(x) dx + \int_{x_3}^{x_5} f(x) dx + \dots + \int_{x_n}^{x_{n+2}} f(x) dx$$



$n$  should be odd.

$$I = \frac{h}{3} \left[ \{f(x_1) + 4f(x_2) + f(x_3)\} + \right.$$

$$\left. \{f(x_3) + 4f(x_4) + f(x_5)\} + \dots + \right]$$

$$\left. \{f(x_{n-2}) + 4f(x_{n-1}) + f(x_n)\} \right]$$

$h =$

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$$I = \frac{h}{3} \left[ f(x_1) + 4 \sum_{i=2,4,6,\dots}^{n-1} f(x_i) + 2 \sum_{i=3,5,7,\dots}^{n-2} f(x_i) + f(x_n) \right]$$

\* If  $n$  is even, then  $\Rightarrow$

Use TM for the last interval & for the rest use the SM.

11/02/19  
(One class missing)

### \* System of Linear Equations:

\* LU decomposition without pivoting  $\rightarrow$  (Recap)

$$\textcircled{1} [A] = [L][U] \quad \text{--- } \textcircled{1}$$

$$\textcircled{2} [L][d] = [b] \quad \text{--- } \textcircled{2}$$

$$\textcircled{3} [U][x] = [d] \quad \text{--- } \textcircled{3}$$

(Forward Eli. 1)

1) Decompose  $[A]$  into  $[L]$  and  $[U]$ . [FE1]

2) Use eq<sup>n</sup>  $\textcircled{2}$  to find  $[d]$ . [FE2]

3) Use eq<sup>n</sup>  $\textcircled{3}$  to find  $[x]$ . [BS].

(Back Substitution)

\* LU decomposition with pivoting  $\rightarrow$

We will introduce a permutation matrix  $[P]$  which will keep a track of the swapping of rows.

Initially,  $[P] = \text{identity matrix}$ .

$$[L]([U][x] - [d]) = [P]([A][x] - [b]).$$

Eg.  $[A] = \begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix}$ ,  $[x] = \begin{pmatrix} x_1 \\ x_2 \end{pmatrix}$

$[b] = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$ ,  $[P] = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$

Let's assume  $a_{11} \ll a_{21}$ .

Pivoting.

$[A] = \begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix} = \begin{pmatrix} a'_{11} & a'_{12} \\ a'_{21} & a'_{22} \end{pmatrix}$

$[P] = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$

After Pivoting we proceed to FE1.

I. Proceed for FE1  $\Rightarrow$

$$[L][d] = [P][b]$$

# Before FE2 pre-multiply the  $[b]$  mtx with  $[P]$ .

Earlier we created augmented mtx & swapped  $b$  too along with  $L$  but here the Permutation mtx. is taking track of the swapping &  $b$  is not swapped.

$[P][b] = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} b_1 \\ b_2 \end{pmatrix} = \begin{pmatrix} b_2 \\ b_1 \end{pmatrix} = \begin{pmatrix} b'_1 \\ b'_2 \end{pmatrix}$

$\Rightarrow$  Let's see the relation:

Naive Gauss

LU decompose without pivot

Pivot Gauss

LU decompose with pivot

Code

Pivot Gauss

→ Here there will two checks:

I) FET  $P = \text{eye}(n)$

for  $k = 1 : (n-1)$

①  $A$  is a square m $\times$ n.

② No. of col<sup>n</sup> of  $A =$  No. of rows of  $A$ .

Pivoting }  $[ \text{leigvalue}, p ] = \max(\text{abs}(\text{aug}(k:n, k)))$

$q = k+p-1;$

if  $q \approx k$

$\text{aug}(k, q, :) = \text{aug}(q, k, :)$

end  $P(k, q, :) = P(q, k, :)$

for  $i = (k+1) : n$

factor =  $\text{aug}(i, k) / \text{aug}(k, k)$

$\text{aug}(i, k) = \text{factor}$

$\text{aug}(i, (k+1) : (n+1)) = -$

end

end

Now, after for LU decompose without pivoting:

✓ 1. No augmented matrix. [Replace aug by A]

✓ 2. Pivot matrix → Swap rows of [P] m $\times$ n when pivoting.

✓ 3. Store factor.

II FE II →

$$[L][d] = [P][b]$$

$$b = P \times b$$

$$d = zeros(n, 1)$$

for  $i = 1:n$

|  
|  
|

end

\* 10 Backward Substitution →

similar to what done in LU decomposition

$$x = zeros(n, 1)$$

$$x[n] =$$

for  $i = (n-1):-1:1$

$$x(i) = \underline{\hspace{2cm}}$$

end

Q.  $A = \begin{bmatrix} 0.1 & -0.1 & -0.2 \\ 3 & -0.2 & -0.3 \\ 0.3 & 7.0 & 10 \end{bmatrix}$

$b = \begin{bmatrix} 7.85 \\ -19.3 \\ 71.4 \end{bmatrix}$ , find  $x$ .

Ans  $x = \begin{bmatrix} -6.53 \\ 249.26 \\ -167.14 \end{bmatrix}$

## Matrix decomposition without pivoting $\rightarrow$

iff  $\rightarrow A \rightarrow$  square matrix  
 off  $\rightarrow A_{m \times n}$

① Check whether  $A$  is a square matrix or not.

② Decompose  $A$  into  $L$  &  $U$  with pivoting  $P = [1, 1]$

$\rightarrow$  FE1  $\rightarrow$  we have generate  $P$  matrix as before

③ Create  $b$  matrix.  $\rightarrow$

$$1^{\text{st}} \text{ iteration: } b = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$$

Do FE2 & backward substitution to find  $x$ .

$$A_{m \times n}(:, 1) = x(:, 1)$$

$$\text{Create } b \text{ matrix again: } b = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}$$

& so on.

$$Q. A = \begin{bmatrix} 3 & -0.1 & -0.2 \\ 0.1 & 7 & 0.3 \\ 0.3 & 0.2 & 10 \end{bmatrix}$$

$$A_{m \times n} = \text{inv}(A).$$

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$$Ax = b$$

$$x = A^{-1}b$$

$$\begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} a_{11}' & a_{12}' & a_{13}' \\ a_{21}' & a_{22}' & a_{23}' \\ a_{31}' & a_{32}' & a_{33}' \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} a_{11}' \\ a_{21}' \\ a_{31}' \end{pmatrix}$$

# Computational Physics

## Assignment: Miscellaneous

1. Given the equations

$$\begin{aligned} 2x_1 - 6x_2 - x_3 &= -38 \\ -3x_1 - x_2 + 7x_3 &= -34 \\ -8x_1 + x_2 - 2x_3 &= -20 \end{aligned}$$

Solve by Gauss elimination with pivoting. Check your answer by substituting your results into the original equations.

2. It is known that the tensile strength of a plastic increases as a function of the time it is heat treated. The following data are collected:

Time (min)	10	15	20	25	40	50	55	60	75
Tensile Strength (N/m <sup>2</sup> )	5	20	18	40	33	54	70	60	78

Fit a straight line to these data and use the equation to determine the tensile strength ( $s_f$ ) at a time ( $t_f$ ) of 32 min. Plot the data, the fitted curve and the value ( $t_f, s_f$ ).

3. The population of a city in a census taken once in ten years is given below. Estimate the population in the year 1975 with the help of cubic spline interpolation technique without using any of the internal matlab commands. Plot the data along with the interpolated value.

Year	1921	1931	1941	1951	1961	1971	1981
Population	35	42	58	84	120	165	220

4. The bending moment ( $M$ ) for a beam of length  $L = 11$  m is given by

$$M = \int_0^L f(x) dx, \quad \text{where,} \quad f(x) = 5 + 0.25x^2.$$

Here,  $f(x)$  is the shear force. Calculate  $M$  using (a) trapezoidal method and (b) Simpson's 1/3 rule. Choose suitable  $h$  and  $n$ .

5. Integration provides a means to compute how much mass ( $M$ ) enters or leaves a reactor over a specified time period as a product of flow ( $Q$ ) and concentration ( $c$ ) as

$$M = \int_{t_0}^{t_1} Q \times c dt.$$

Use Simpson's 1/3 rule (along with trapezoidal method) to evaluate this equation for the data listed below:

$t$ (min)	0	10	20	30	35	40	45	50
$Q$ ( $m^3/min$ )	4	4.8	5.2	5.0	4.6	4.3	4.3	5.0
$c$ ( $mg/m^3$ )	10	35	55	52	40	37	32	34

$$\text{diskat tetlik} \quad n = \text{floor} \left( \frac{(t(\text{end}) - t(\text{beg})) + 1}{\frac{31}{6}} \right) + 1$$

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Computational Physics

H.

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### Assignment: Matlab Fundamentals

1. Use the linspace function to create vectors identical to the following created with colon notation:

$$(b) x = -4:2;$$

$$\left( \frac{35-4}{6} \right)$$

Correction

- ① (a)  $t = 4:6:35$  ;  $x_n = ux_k + x_1$ ,

$$(b) x = \text{linspace}(8, 4, 5, 8);$$

2. Use colon notation to create vectors identical to the following created with the linspace function:

$$(a) v = \text{linspace}(-2, 1.5, 8);$$

*Note: No need to apply floor or round here*

3. If a force  $F(N)$  is applied to compress a spring, its displacement  $x$  (m) can often be modeled by Hooke's law:  $F = kx$ , where  $k$  = spring constant (N/m). The potential energy stored in the spring  $U$  (J) can then be computed as  $U = (1/2)kx^2$ . Five springs are tested and the following data compiled:

$F$ , (N)	14	18	8	9	13
$x$ , (m)	0.013	0.020	0.009	0.010	0.012

Use MATLAB to store  $F$  and  $x$  as vectors and then compute vectors of the spring constants and the potential energies. Use the max function to determine the maximum potential energy.

4. It is general practice in engineering and science that equations be plotted as lines and discrete data as symbols. Here are some data for concentration ( $c$ ) versus time ( $t$ ) for the photo-degradation of aqueous bromine:

$t$ , (min)	10	20	30	40	50	60
$c$ , (ppm)	3.4	2.6	1.6	1.3	1.0	0.5

These data can be described by the following function:

$$c = 4.84 \exp(-0.034 t)$$

Use MATLAB to create a plot displaying both the data (using red diamond-shape) and the function (using a green, dashed line). Plot the function for  $t = 0$  to  $70$  min.

5. The butterfly curve is given by the following parametric equations:

$$x = \sin(t) \left[ e^{\cos(t)} - 2 * \cos(4 * t) - \sin^5 \left( \frac{t}{12} \right) \right]$$
$$y = \cos(t) \left[ e^{\cos(t)} - 2 * \cos(4 * t) - \sin^5 \left( \frac{t}{12} \right) \right]$$

Generate the values of  $x$  and  $y$  for values of  $t$  from  $0$  to  $100$  with  $\Delta t = 1/16$ . Construct the plot of  $y$  vs  $x$ .

6. The butterfly curve can also be represented in polar coordinates as

$$r = \exp(\sin(\theta)) - 2 \cos(4\theta) - \sin \left( \frac{2\theta - \pi}{24} \right)^5$$

Generate values of  $r$  for values of  $\theta$  from  $0$  to  $8\pi$  with  $\Delta = \pi/32$ . Use the MATLAB function polar or polarplot to generate the polar plot of the butterfly curve with a dashed red line.

Note → polarplot is more recommended than polar

polarplot(θ, r),

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Comput.

Akshat Jethlia

## Computational Physics

### Assignment: Matlab Programming

1. The sine function can be evaluated by the following infinite series:

$$\sin(x) = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots$$

$$\sin(x) = \sum_{i=1}^n \frac{(-1)^{i-1} x^{(2i-1)}}{(2i-1)!}$$

Create an M-file to implement the summation formula so that it computes and displays the values of  $\sin(x)$  as each term in the series is added. In other words, compute and display in sequence the values for

$$\sin(x) = x$$

$$\sin(x) = x - \frac{x^3}{3!}$$

⋮

up to  $n = 5$  order, i.e. up to the term  $x^9/9!$  for  $x = 0.9$ . For each iteration compute the fractional error given by the formula

$$\text{fractional\_error} = \text{abs}\left(\frac{\text{true\_value} - \text{series\_approx}}{\text{true\_value}}\right)$$

**Hint:** Use the `factorial()` command, but not the `sum()` command.

2. Develop a vectorized version of the following code:

```
tstart=0; tend=20; ni=8;
t(1)=tstart; y(1)= 10 + 5 * cos(2*pi*t(1)/(tend-tstart));
for i=2:ni+1
    t(i)=t(i-1)+(tend-tstart)/ni;
    y(i)=10 + 5*cos(2*pi*t(i)/(tend-tstart));
end
```

3. Develop an M-file to determine the polar coordinates  $(r, \theta)$  for the values of  $(x, y)$  given by

```
x = [2 2 0 -3 -2 -1 0 0 2] ;
y = [0 1 3 1 0 -2 0 -2 2] ;
```

The radius  $r$  can be computed by the formula

$$r = \sqrt{x^2 + y^2}$$

whereas  $\theta$  can be computed with the following table

$x$	$y$	$\theta$
$> 0$	All values	$\tan^{-1}(y/x)$
$< 0$	$> 0$	$\tan^{-1}(y/x) + \pi$
$< 0$	$< 0$	$\tan^{-1}(y/x) - \pi$
$= 0$	$= 0$	$\pi$
$= 0$	$> 0$	$\pi/2$
$= 0$	$< 0$	$-\pi/2$
$= 0$	$= 0$	$0$

4. A function  $f(x)$  is given by the formula

$$f(x) = \frac{\sin(x)}{x^a + b}$$

where  $a = 2$ ,  $b = 1$  are parameters. Write a function file `findmean1.m` to find the average of  $f(x)$ . The first couple of lines of the `findmean1.m` file is given below:

```
function avgval = findmean1(func,x)
% Input:
%   func: Input function whose average has to be calculated.
%   x: Range of x over which the function has to be evaluated.
% Output:
%   avgval: Function average value.
% ...
end
```

Hint: `yavg = findmean1(func,x);`

5. Repeat the problem of Q 4, to create another function file `findmean2.m`. In this case pass the parameters  $a, b$  as anonymous variables to the function file. The first couple of lines of the `findmean2.m` file is given below:

```
function [avgval] = findmean2(f,x,varargin)
% Find the mean of a function.
% Input:
%   func: Input function whose average has to be calculated.
%   x: Range of x over which the function has to be evaluated.
%   varargin: Variable number of arguments
% Output:
%   avgval: Function average value.
% ...
end
```

Hint: `yavg = findmean2(func,x,a,b);`