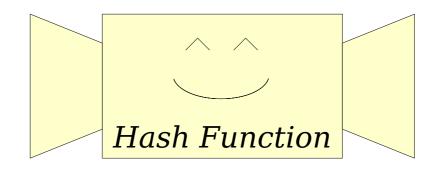
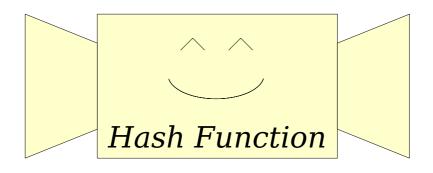
Hashing Part One

Way Back When...

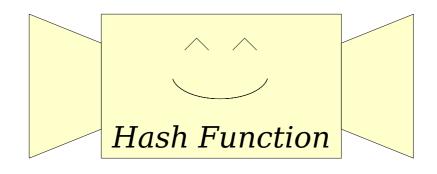
```
int nameHash(string first, string last){
    /* This hashing scheme needs two prime numbers, a large prime and a small
     * prime. These numbers were chosen because their product is less than
     * 2^31 - kLargePrime - 1.
     */
    static const int kLargePrime = 16908799;
    static const int kSmallPrime = 127;
    int hashVal = 0;
    /* Iterate across all the characters in the first name, then the last
     * name, updating the hash at each step.
     */
    for (char ch: first + last) {
        /* Convert the input character to lower case. The numeric values of
         * lower-case letters are always less than 127.
        ch = tolower(ch);
        hashVal = (kSmallPrime * hashVal + ch) % kLargePrime;
    return hashVal:
```



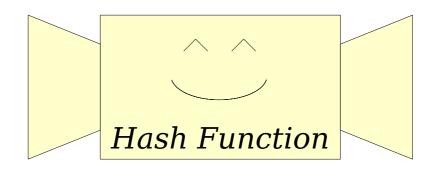
This is a *hash function*. It's a type of function some smart math and CS people came up with.



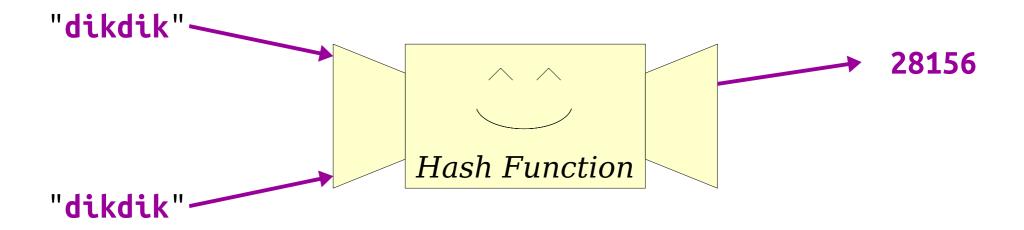
Most hash functions return a number. In CS106B, we'll use the **int** type.



Different hash functions take inputs of different types. In this example, we'll assume it takes string inputs.

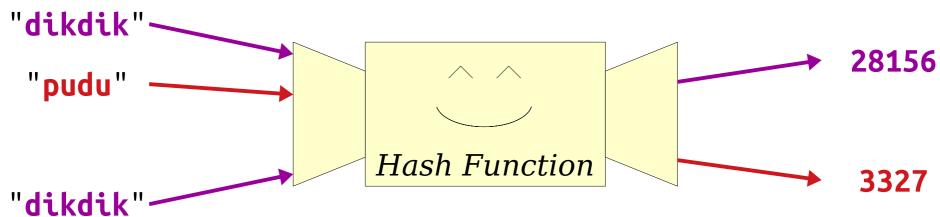


What makes this type of function so special?

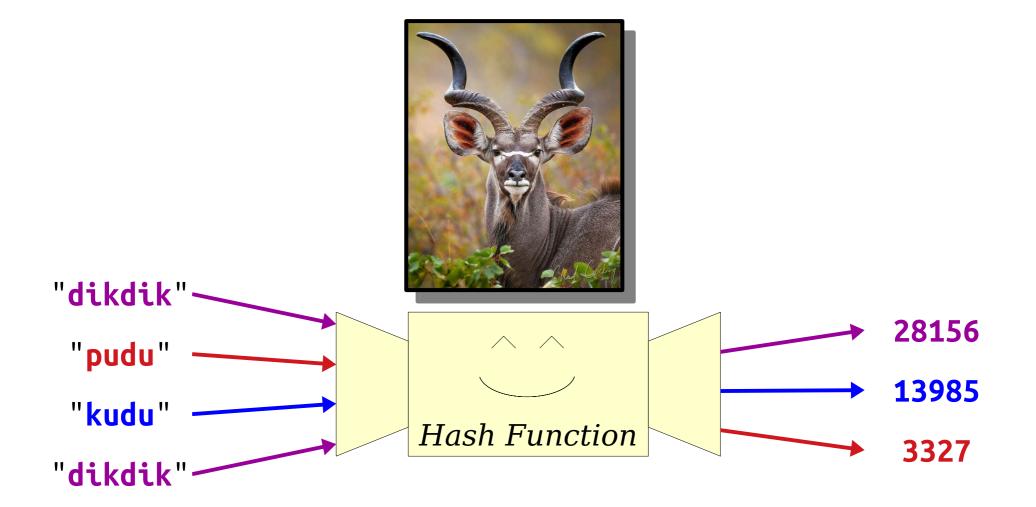


First, if you compute the hash code of the same string many times, you always get the same value.





Second, the hash codes of different inputs are (usually) very different from one another.

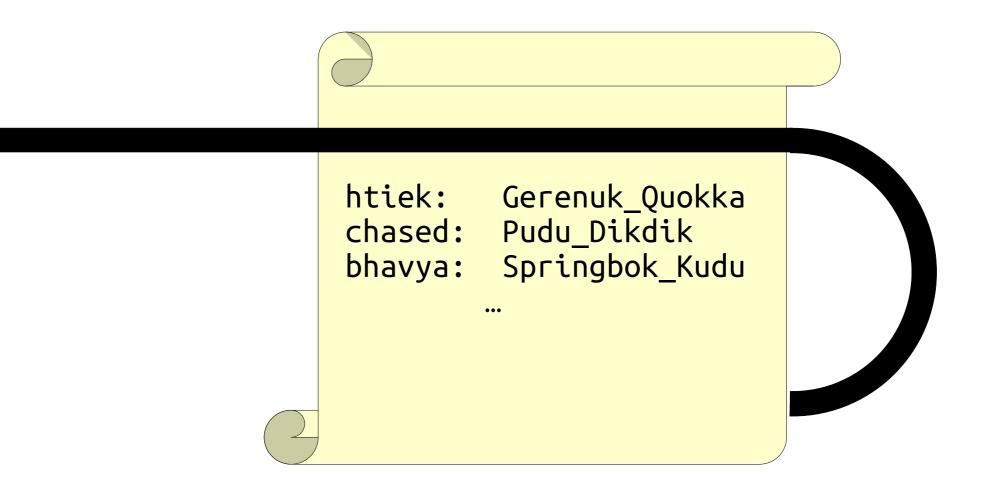


Even very similar inputs give very different outputs!

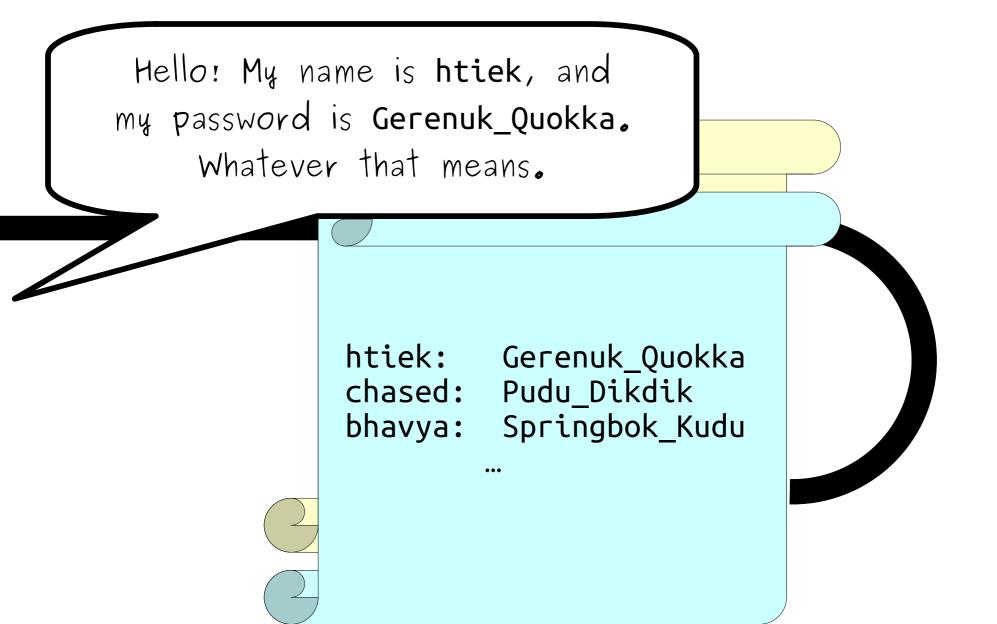
To Recap:

Equal inputs give equal outputs.

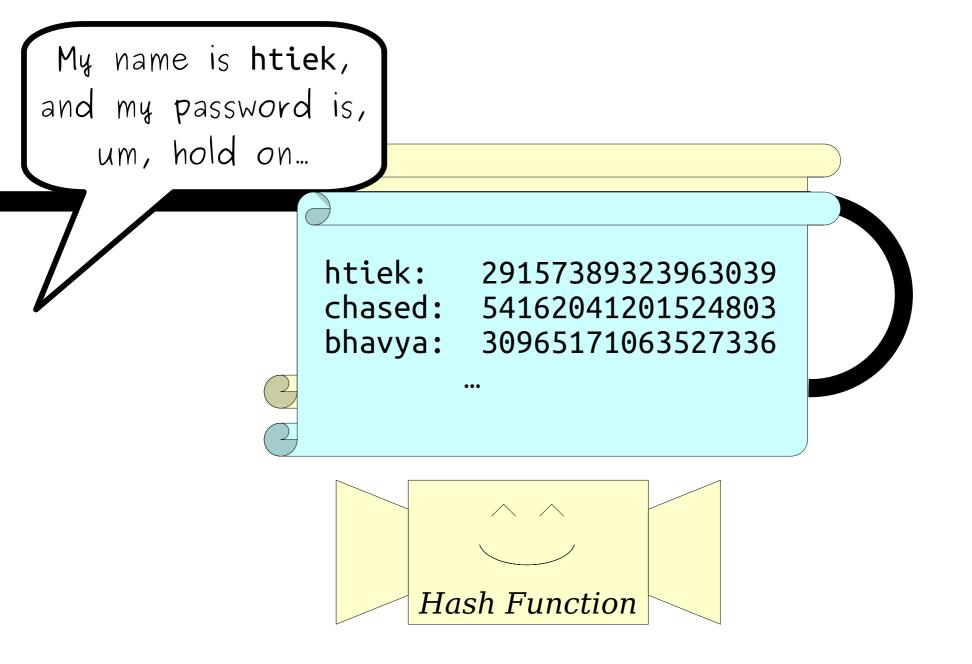
Unequal inputs (usually) give very different outputs.



How do servers store passwords?



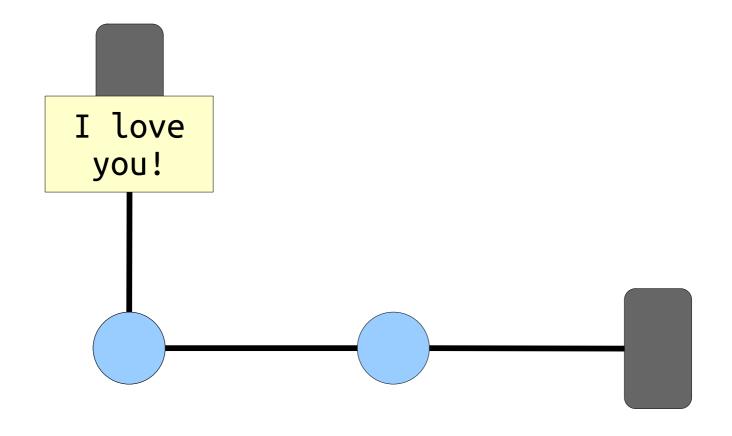
How do servers store passwords?

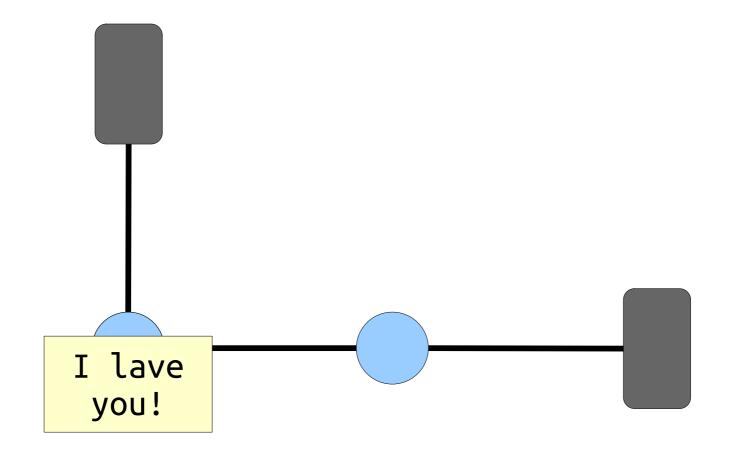


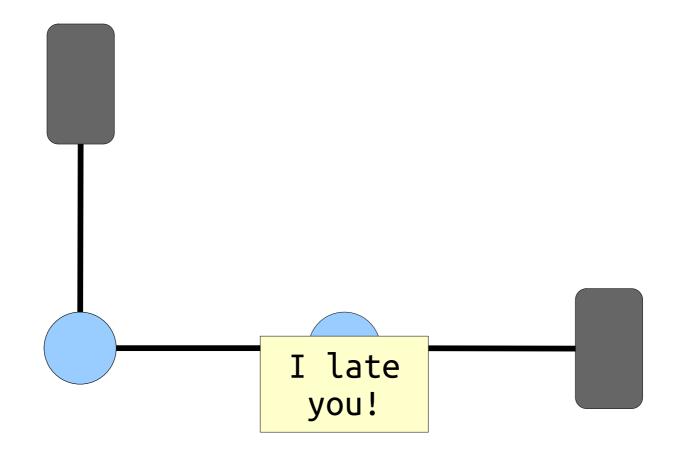
How do servers store passwords?

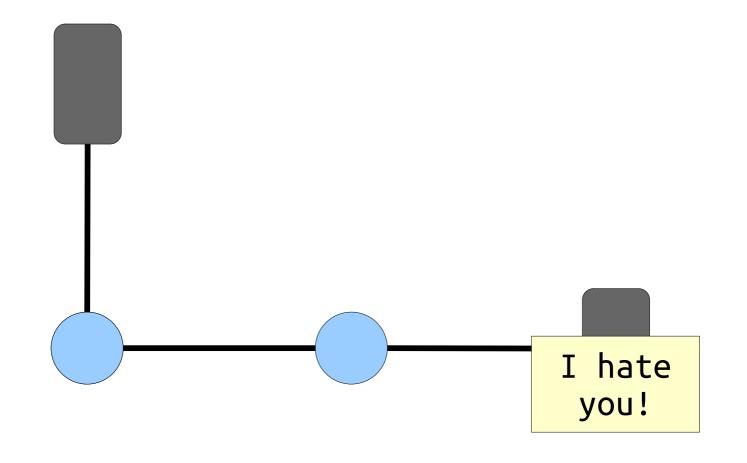
This is how passwords are typically stored. Look up *salting and hashing* for details!

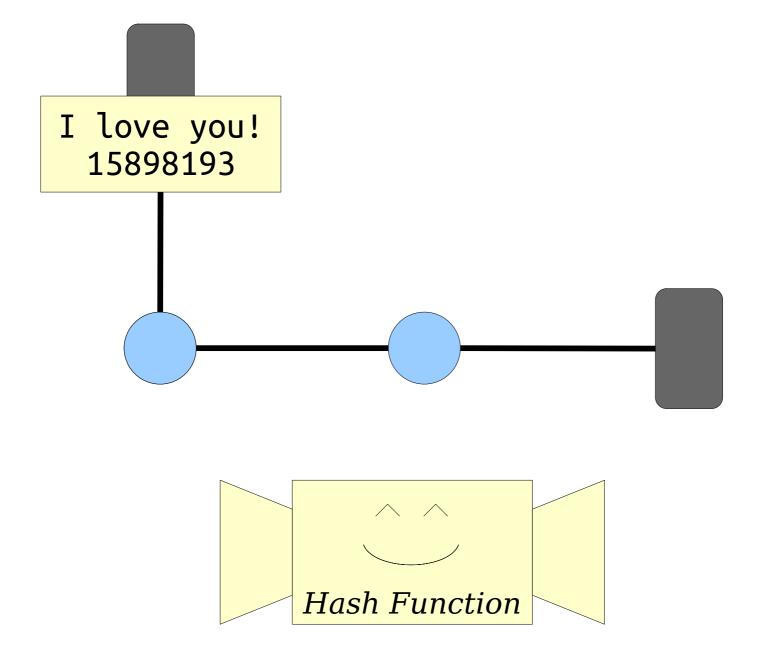
And look up *commitment schemes* if you want to see some even cooler things!

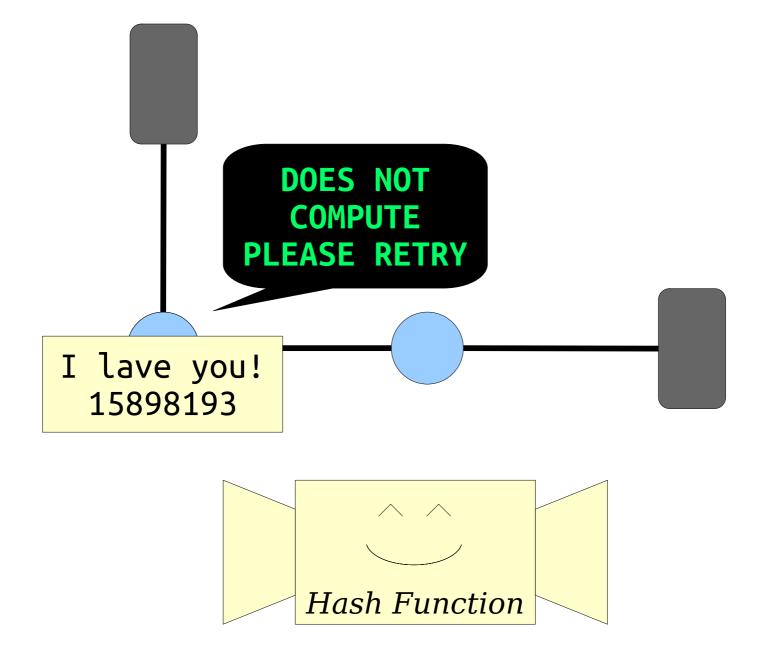


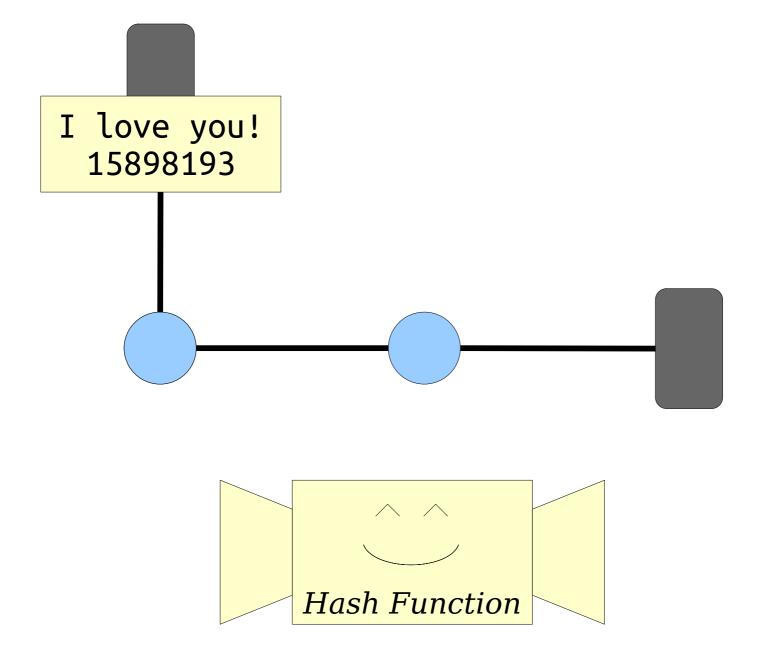


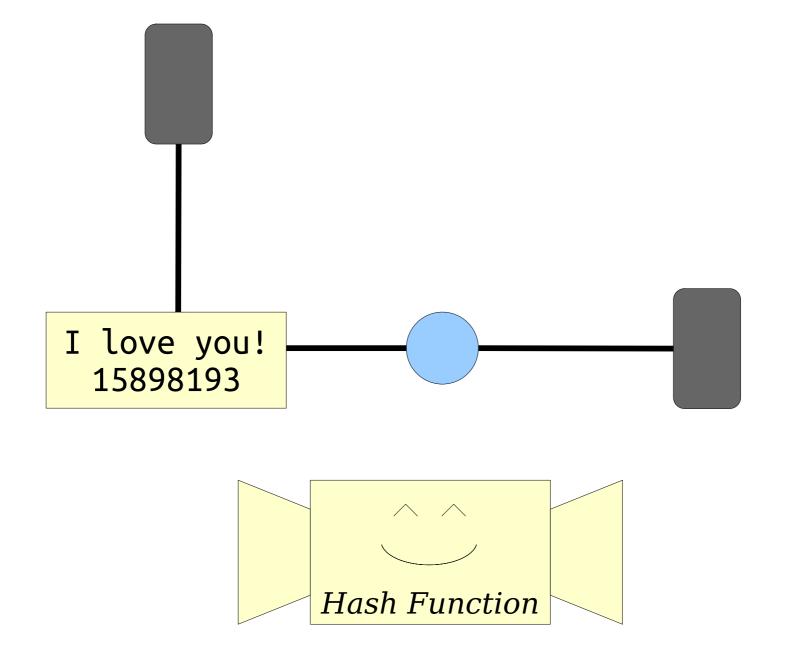


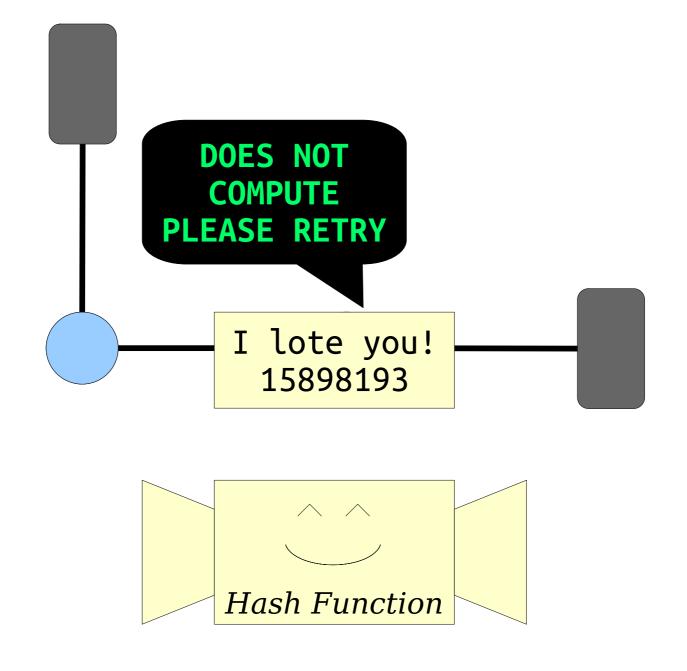


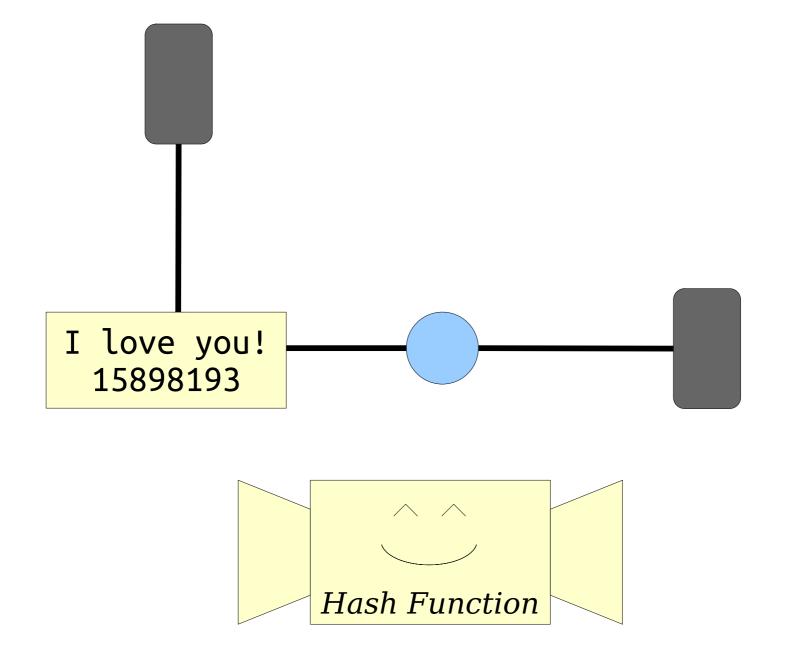


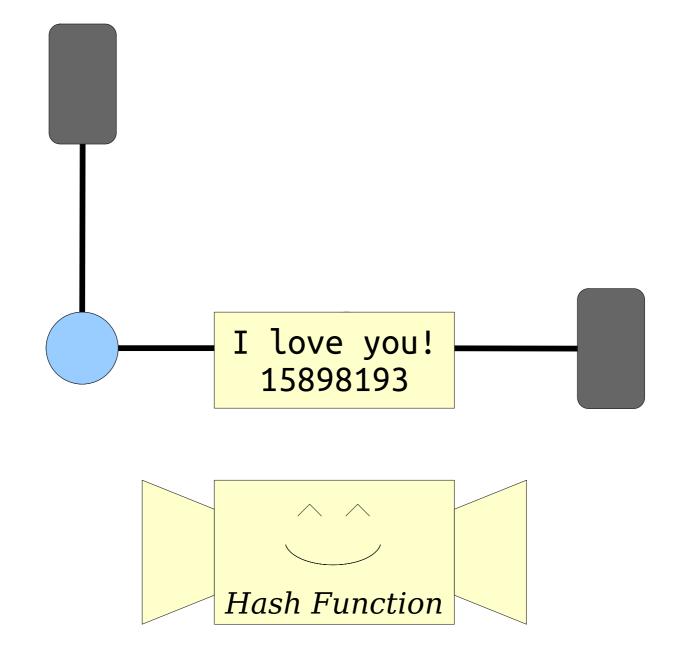


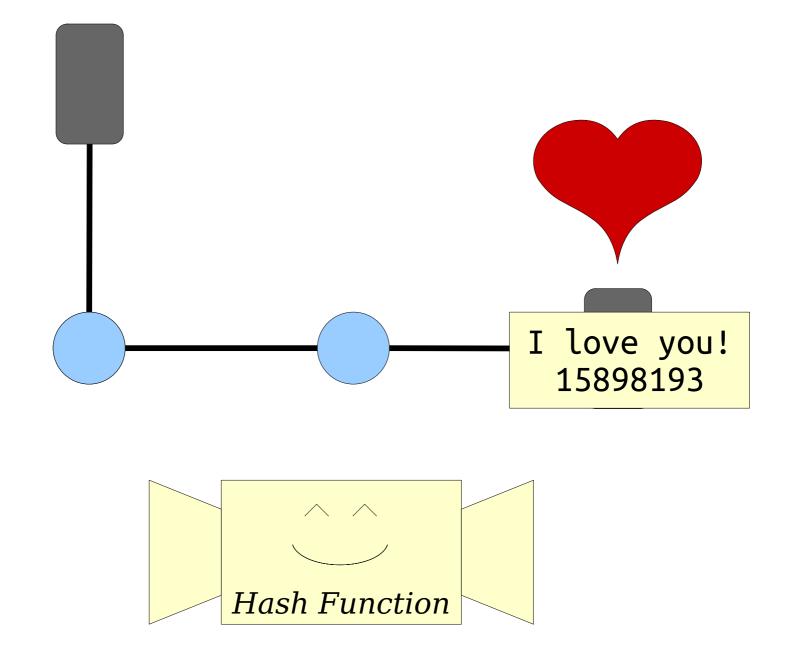










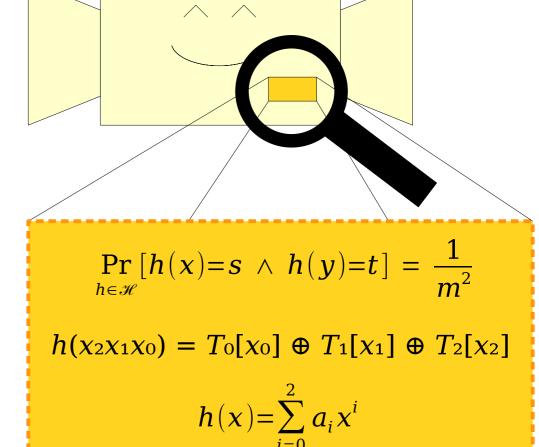


This is done in practice!

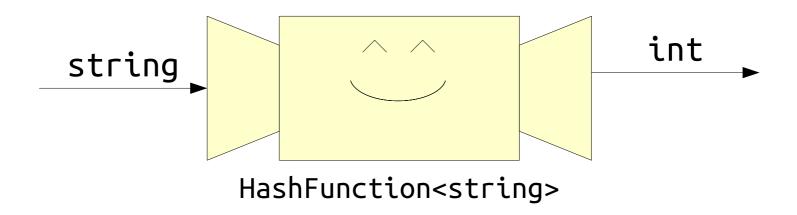
Look up *SHA-256*, the *Luhn algorithm*, and *CRC32* for some examples!

Designing Hash Functions

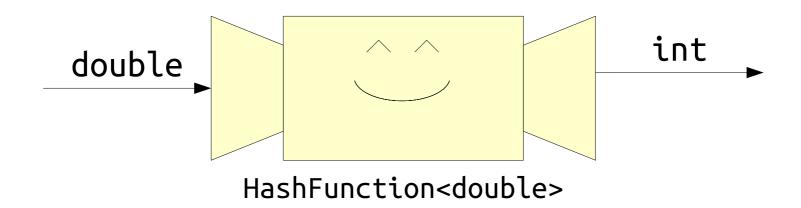
- Designing good hash functions is challenging, and it's beyond the scope of what we'll explore in CS106B.
- Interested in things like independent random variables, finite fields, and the like? Come talk to me after class and I'll give the rundown.



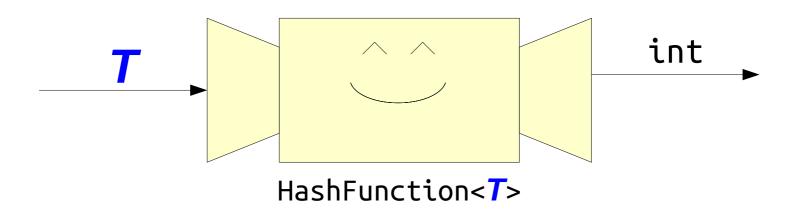
- Every programming language has a different way for programmers to work with hash functions.
- In CS106B, we'll represent hash functions using the type HashFunction<T>.



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- In CS106B, we'll represent hash functions using the type HashFunction<T>.



- Sometimes, you want a hash function that outputs values in a wide range.
 - For example, when storing hashes of passwords. (Why?)
- Sometimes, you want a hash function that outputs values in a small range.
 - For example, assigning tasks to volunteers.
- Our HashFunction<T> returns a value in the range 0, 1, 2, ..., n 1, where n is some number you provide to the constructor.

An Application:

Map and Set

```
class OurSet {
public:
    OurSet();
    void add(const std::string& str);
    bool contains(const std::string& str) const;
    int size() const;
    bool isEmpty() const;
private:
    /* What goes here? */
};
```

In header files, we refer to the string type as std::string. It's an Endearing C++ Quirk. Feel free to ask me about this after class if you're curious why.

```
class OurSet {
public:
    OurSet();
    void add(const std::string& str);
    bool contains(const std::string& str) const;
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private:
    /* What goes here? */
```

An Example: Clothes





For Large Values of *n*



Our Strategy

- Maintain a large number of small collections called buckets (think drawers).
- Find a *rule* that lets us tell where each object should go (think knowing which drawer is which).
- To find something, only look in the bucket assigned to it (think looking for socks).

Our Strategy

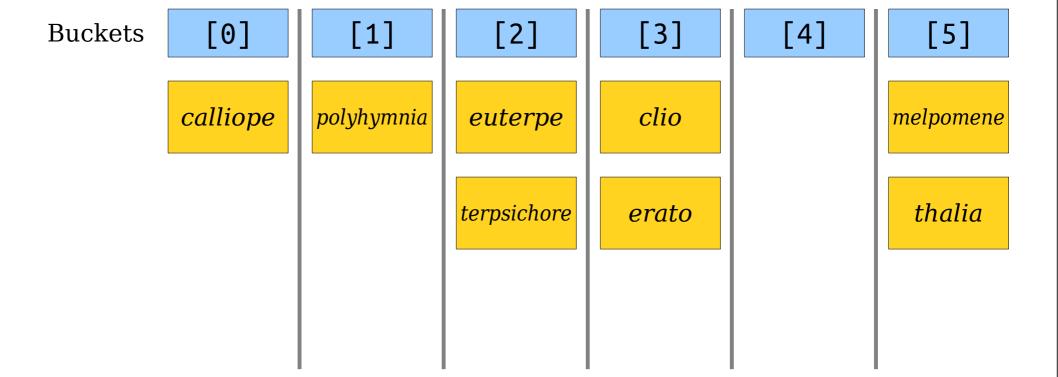
Maintain a large number of small collections called *buckets* (think drawers).

• Find a *rule* that lets us tell where each object should go (think knowing which drawer is which).

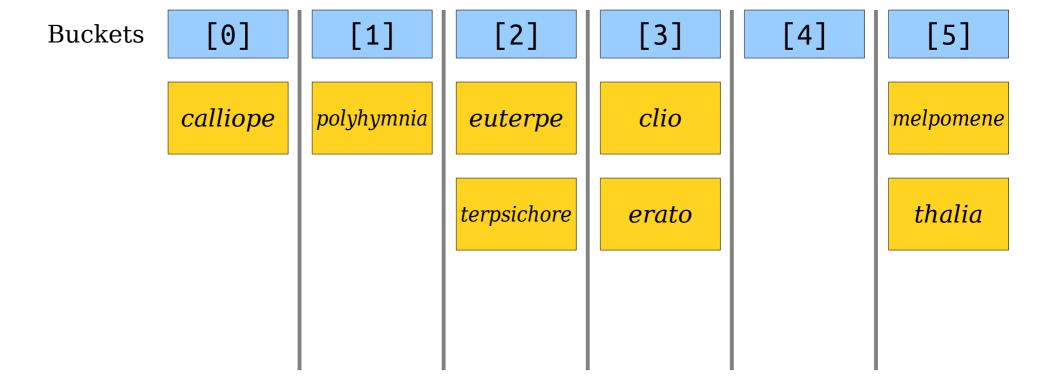
Use a hash

To find something, on function! e bucket assigned to it (think looking for socks).

Buckets	[0]	[1]	[2]	[3]	[4]	[5]
	calliope	polyhymnia	euterpe	clio		melpomene
			terpsichore	erato		thalia



erato



```
Buckets [0] [1] [2] [3] [4] [5]

calliope
polyhymnia
euterpe
clio
melpomene

terpsichore
erato
```

```
bool OurSet::contains(const string& value) const {
   int bucket = hashFn(value);
}
erato
```

Buckets [0] [1] [2] [3] [4] [5]

calliope polyhymnia euterpe clio melpomene

terpsichore erato

```
bool OurSet::contains(const string& value) const {
   int bucket = hashFn(value);
}
erato
(bucket 3)
```

Buckets [0] [1] [2] [3] [4] [5]

calliope polyhymnia euterpe clio melpomene

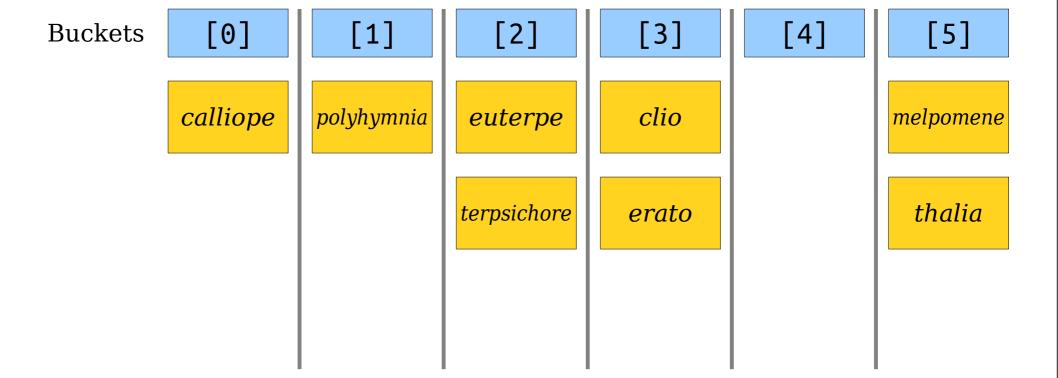
terpsichore erato

```
bool OurSet::contains(const string& value) const {
   int bucket = hashFn(value);
   return buckets[bucket].contains(value);
}

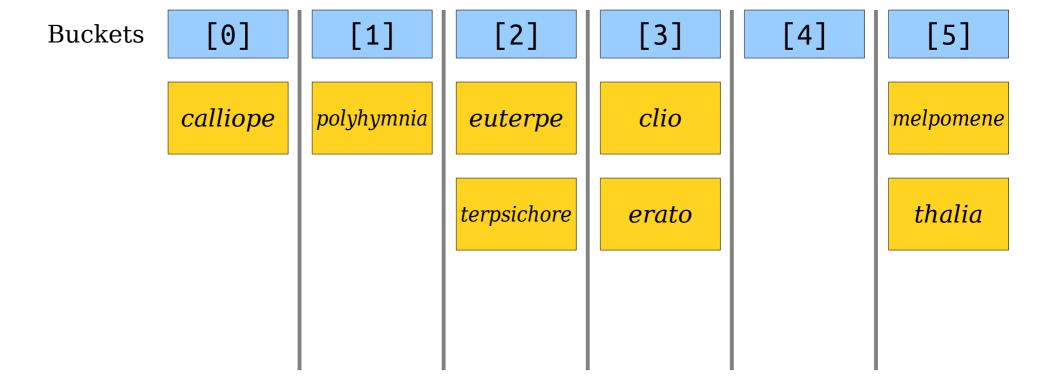
e
(bucket)
```

erato

(bucket 3)



urania



```
void OurSet::add(const string& value) {
    urania
}
```

```
Buckets [0] [1] [2] [3] [4] [5]

calliope polyhymnia euterpe clio melpomene

terpsichore erato
```

```
void OurSet::add(const string& value) {
   int bucket = hashFn(value);
   urania
}
```

```
Buckets [0] [1] [2] [3] [4] [5]

calliope polyhymnia euterpe clio melpomene

terpsichore erato
```

```
void OurSet::add(const string& value) {
   int bucket = hashFn(value);
   urania
}
```

```
Buckets [0] [1] [2] [3] [4] [5]

calliope polyhymnia euterpe clio melpomene

terpsichore erato
```

```
Buckets [0] [1] [2] [3] [4] [5]

calliope polyhymnia euterpe clio melpomene

terpsichore urania
```

Buckets [0] [1] [2] [3] [4] [5]

calliope polyhymnia euterpe clio melpomene

terpsichore urania

```
void OurSet::add(const string& value) {
   int bucket = hashFn(value);
   buckets[bucket] += value;
   numElems++;
}

(bucket 2)
```

Buckets [0] [1] [2] [3] [4] [5]

calliope polyhymnia euterpe clio melpomene

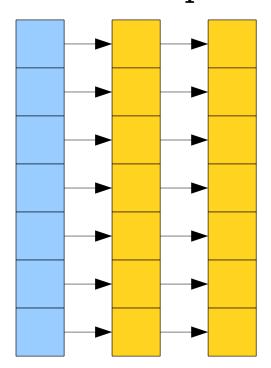
terpsichore urania

```
void OurSet::add(const string& value) {
   if (contains(value)) return;
   int bucket = hashFn(value);
   buckets[bucket] += value;
   numElems++;
}
```

How efficient is this?

Efficiency Concerns

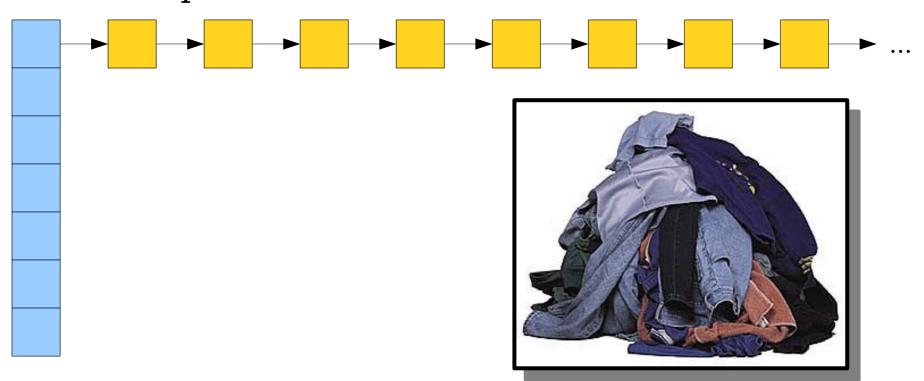
- Each hash table operation
 - chooses a bucket and jumps there, then
 - potentially scans everything in the bucket.
- *Claim:* The efficiency of our hash table depends on how well-spread the elements are.





Efficiency Concerns

- Each hash table operation
 - chooses a bucket and jumps there, then
 - potentially scans everything in the bucket.
- *Claim:* The efficiency of our hash table depends on how well-spread the elements are.



Efficiency Concerns

- For a hash table to be fast, we need a hash function that spreads things around nicely.
- We'll assume our HashFunction<T> type distributes elements more or less randomly.
- Writing good hash functions or quantifying how good they are – is the domain of courses like CS161, CS166, and CS265. Come talk to me after class if you're curious!

Analyzing our Efficiency

- Let's suppose we have a "strong" hash function that distributes elements fairly evenly.
- Imagine we have b buckets and n elements in our table.
- On average, how many elements will be in a bucket?

Answer: n / b

• The *expected* cost of an insertion, deletion, or lookup is therefore

$$O(1 + n / b)$$
.

Load Factors

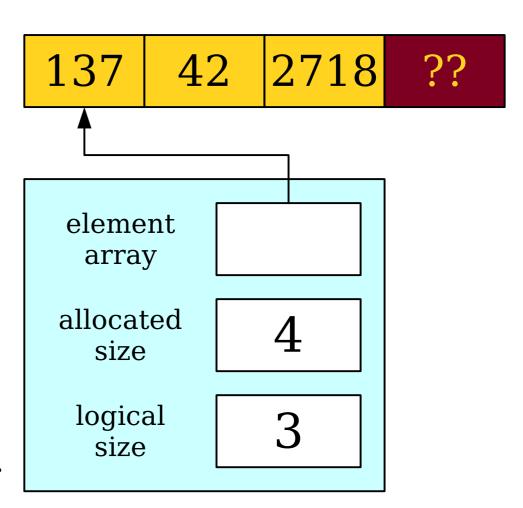
• The *load factor* of a hash table with n elements and b buckets is denoted α and given by the expression

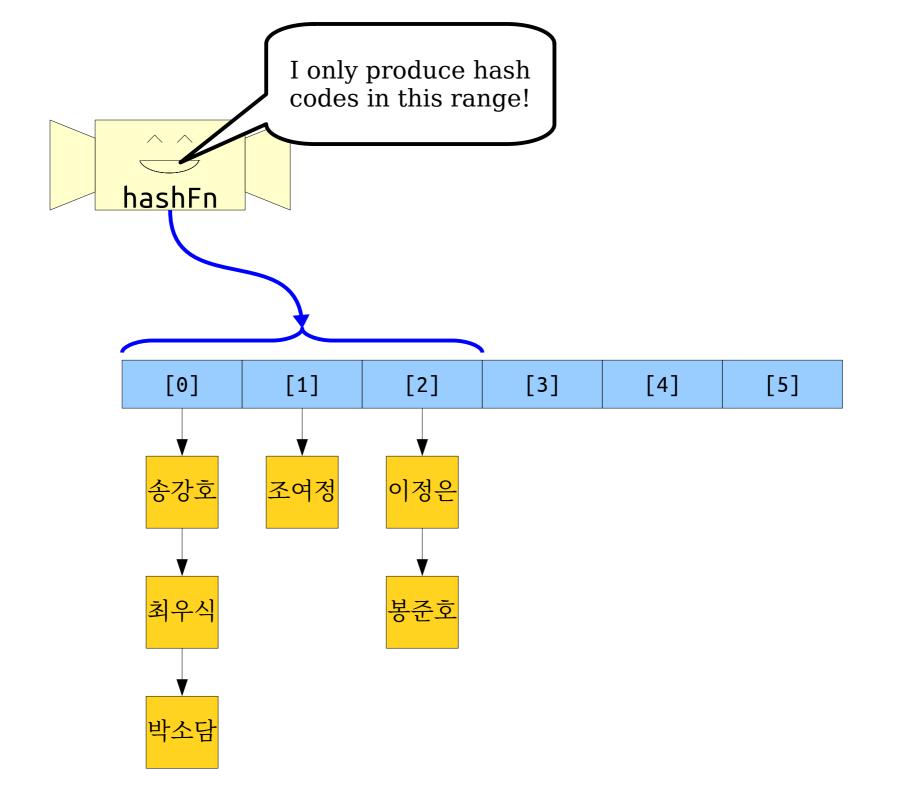
$$\alpha = n / b$$
.

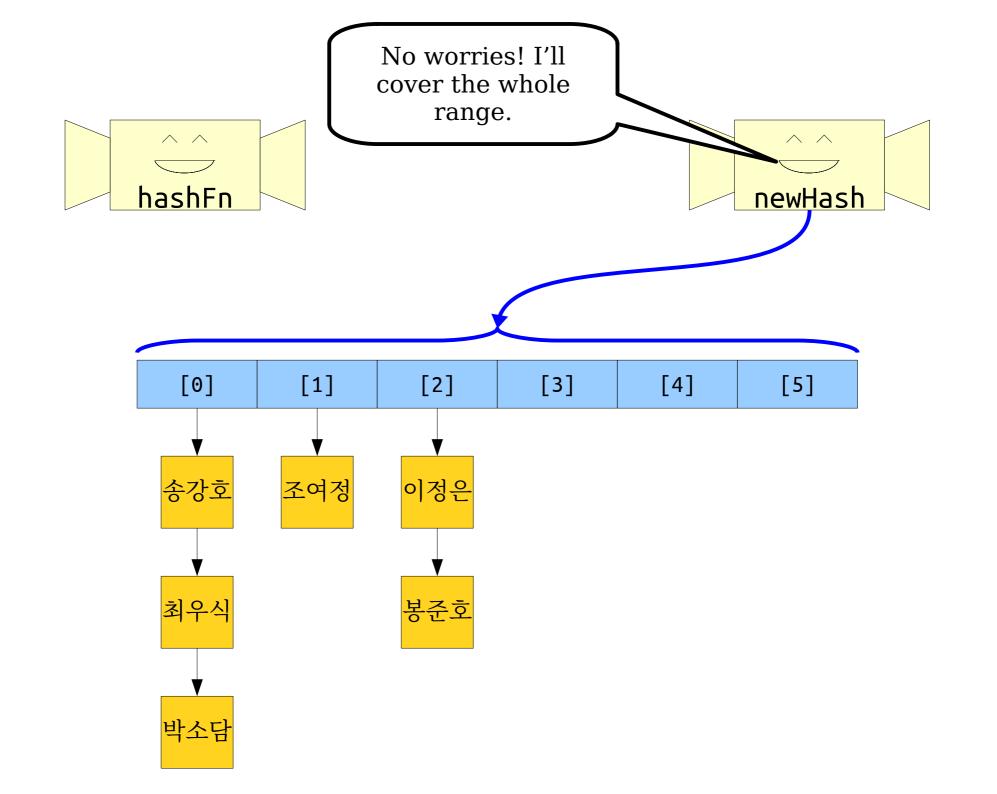
- The expected cost of a lookup in a hash table is $O(1 + n / b) = O(1 + \alpha)$.
 - If α gets too big, the hash table will be too slow.
 - If α gets too low, the hash table will waste too much space.
- How do we balance things?

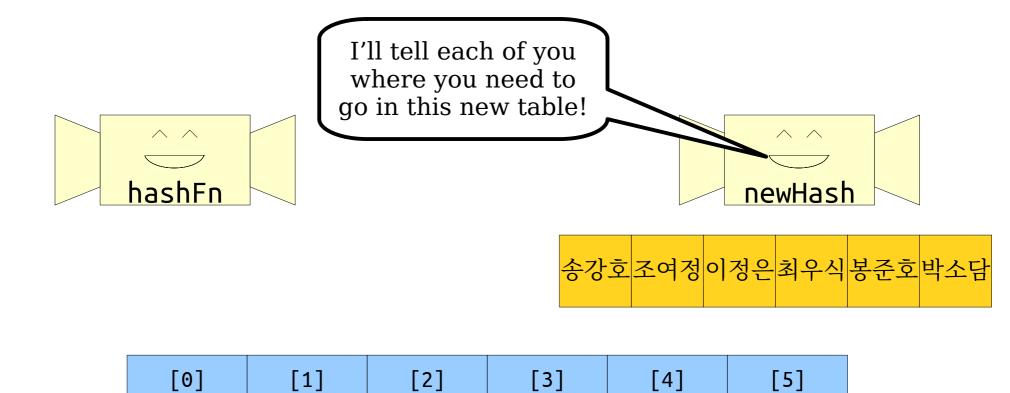
Remember When?

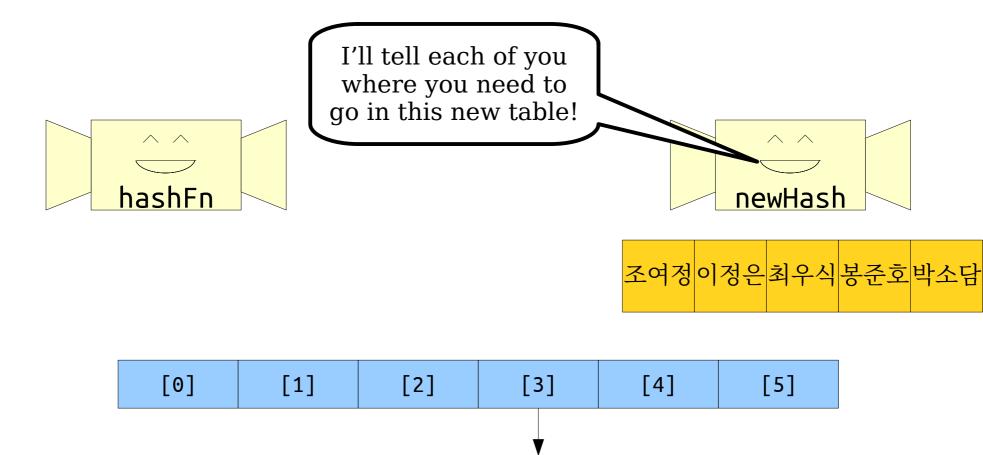
- Think back to how we implemented the Stack.
- Initially, we had a fixed number of slots.
- Once we ran out of space, we doubled the number of slots and transferred things over.
- Can we do that here?
- *Idea*: Double the table size whenever $n / b \ge 2$.











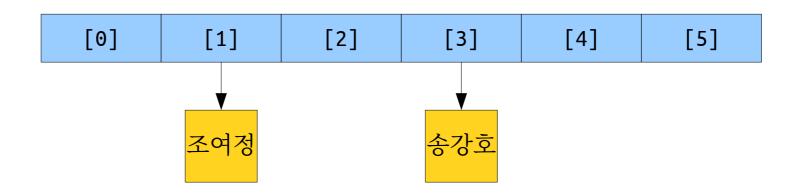
송강호



I'll tell each of you where you need to go in this new table!

newHash

이정은 최우식 봉준호 박소담

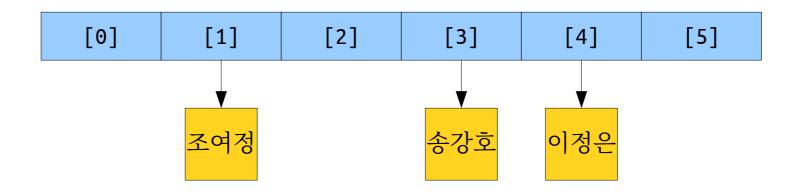


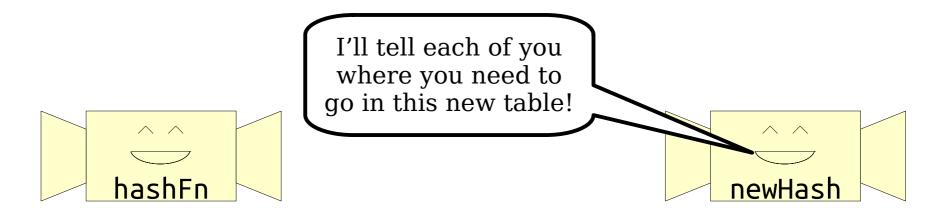


I'll tell each of you where you need to go in this new table!

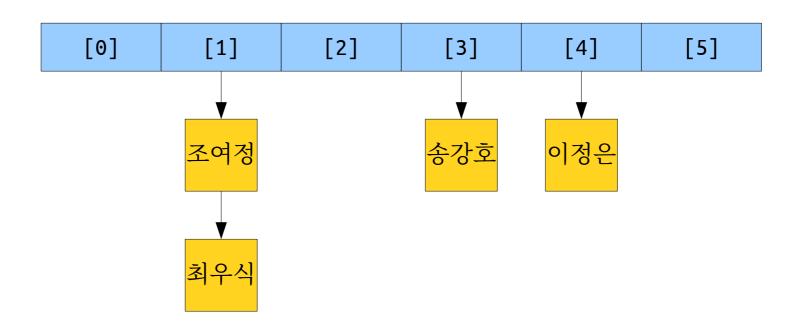
최우식 봉준호 박소담

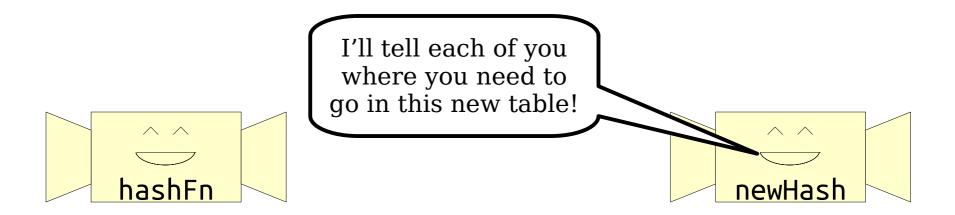
newHash



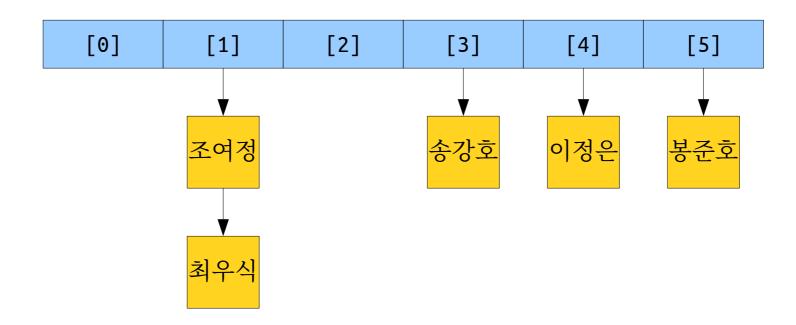


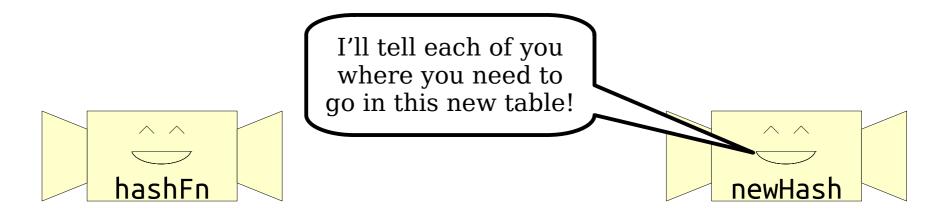
<mark>봉준호</mark>박소담

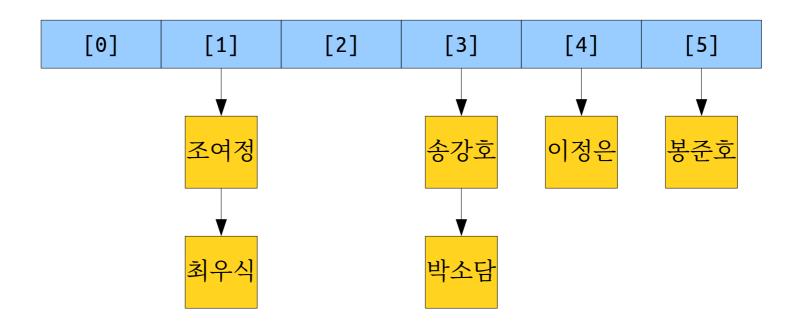




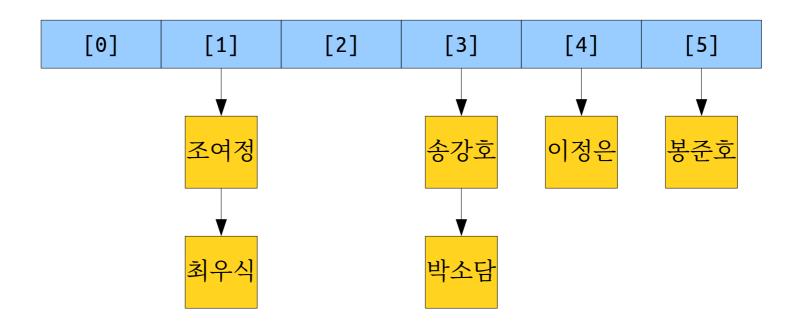
박소담



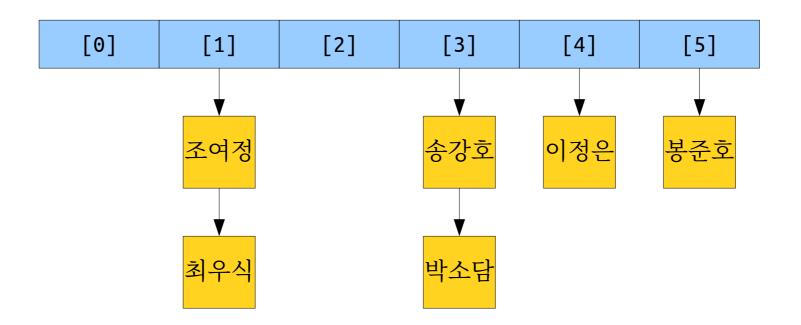




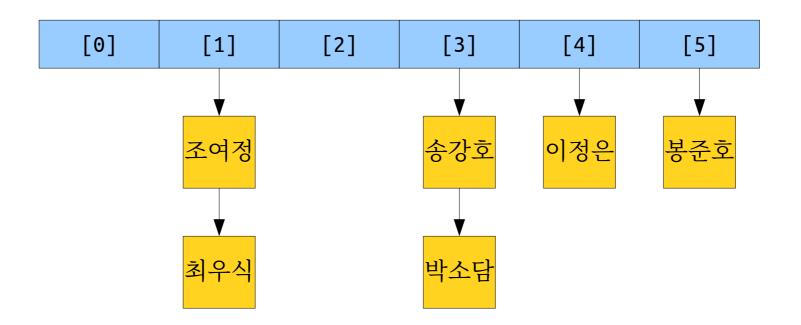












Rehashing

- To perform a rehash, do the following:
 - Get a new list of buckets, twice as big as before.
 - Get a new hash function that distributes elements across the wider range.
 - Redistribute the elements from the old buckets into the new ones, using the new hash function.
 - Use the new buckets and hash functions going forward.
- Time required is O(n). However, this happens so rarely that the extra work averages out to O(1) per insert.

The Final Scorecard

- Assuming we cap the load factor α at some constant (say, 2), then $1 + \alpha = O(1)$.
 - That is, $1 + \alpha$ doesn't grow as a function of n, the number of elements in the hash table.
- The expected cost of a lookup is O(1).
- The expected cost of an insertion is O(1).
 - (It's actually *expected amortized* O(1), since we do some work to copy things over, but only very infrequently.)
- This is about as good as it gets!

Your Action Items

- Work on Assignment 6
 - Aim to complete the HeapPQueue implementation by Wednesday.
 - Need help or support? Come talk to us at LaIR, in office hours, or over EdStem!

Next Time

- Linear Probing
 - A different strategy for building hash tables.
- Robin Hood Hashing
 - A clever and fast hashing strategy.