

Tutorial-6 (DAA)

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Ques 1 What do you mean by Min Spanning Tree? What are the applications of MST?

Ans Minimum spanning tree is a subset of edges of a connected edge-weighted graph that connects all the vertices together without any cycles & with the minimum possible edge weight.

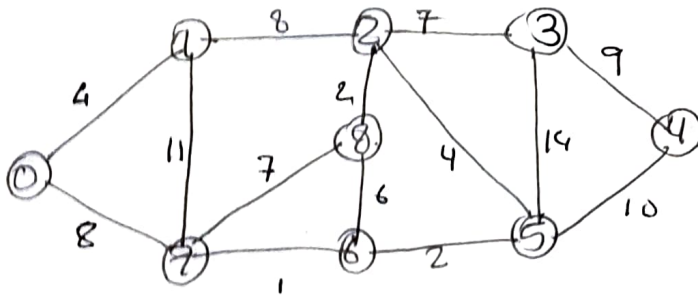
Applications:

1. Consider n stations are to be linked using a communication network and using of communication link between any two stations involved a cost. The ideal condition is to extract a subgraph termed as min cost spanning tree.
2. Designing LAN
3. Laying pipeline connecting offshore drilling sites, refineries & consumer markets

Ques 2 Analyze time & space complexity of different algorithms

<u>Ans</u>	Algorithm	Time Comp	Space Comp
	Prim's Algorithm	$O(E \log V)$	$O(V)$
	Kruskal's Algorithm	$O(E \log E)$	$O(V)$
	Dijkstra's Algorithm	$O(V^2)$	$O(V^2)$
	Bellman Ford's Algorithm	$O(VE)$	$O(E)$

Ans-3

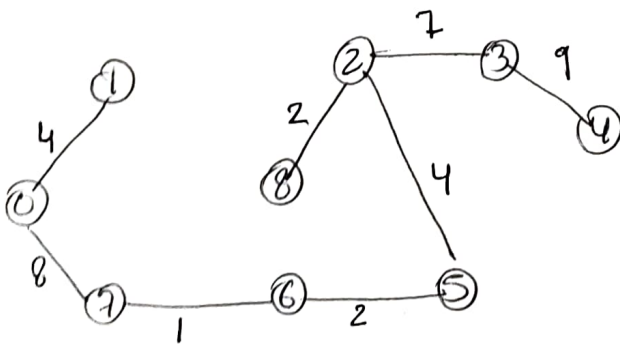


Kruskal's Algorithm :-

0	V	W	
6	7	1	✓
5	6	2	✓
2	8	2	✓
0	1	4	✓
2	5	4	✓
6	8	6	X
2	3	7	✓
7	8	7	X
0	7	8	✓
1	2	8	X
4	3	9	✓
4	5	10	X
1	7	11	X
3	5	14	X

Prim's Algorithm

$$\text{weight} = 4 + 8 + 2 + 4 + 2 + 7 + 9 + 3 = 37$$

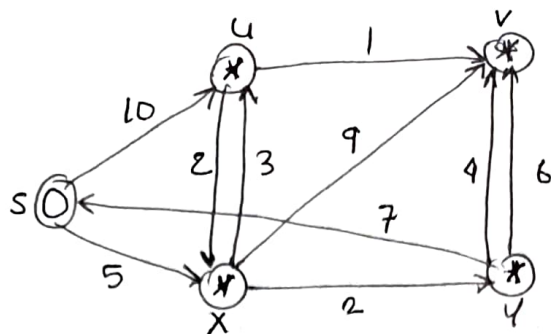


$$\text{weight} = 1 + 2 + 2 + 4 + 4 + 7 + 8 + 9 = 37$$

Ans-4

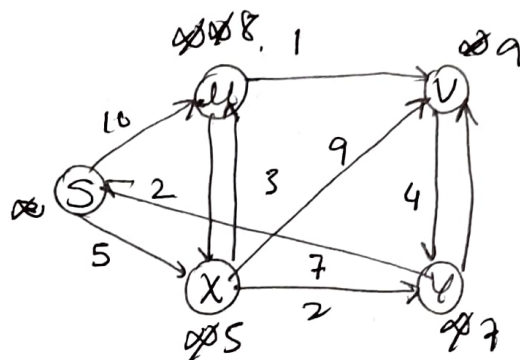
- (i) The shortest path may change. The reason is that there may be different no. of edges in different path from S to t . For eg let the shortest path of weight 15 and has edges 5. Let there be another path with 2 edge and total weight 25. The weight of shortest path is measured by 5×10 and becomes 15. $15 + 50$ weight of other path is increased by 2×10 & becomes 25. So the shortest path changes to other path with weight as 45.
- (ii) If we multiply all edges weight by 10, the shortest path do not change. The reason is that weights of all path from S to t gets multiplied by same unit the number of edges or path doesn't matter.

Ans-5



Dijkstra's Algorithm:

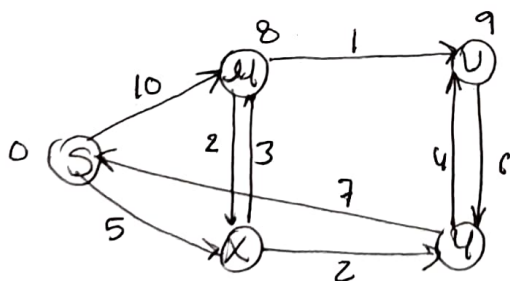
Node	Shortest dist from source node
u	8
x	5
v	9
y	7



Bellman Ford Algorithm:

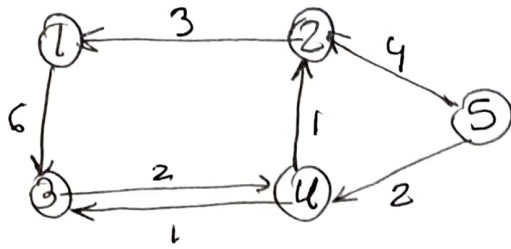
1 st →	S	U	V	X	Y
2 nd →	S	10	11	5	10
3 rd →	S	8	9	5	7
4 th →	S	8	9	5	7

graph does not have -ve cycle



Final Graph

Ans-6



	1	2	3	4	5
1	0	∞	6	3	∞
2	2	0	∞	∞	∞
3	∞	∞	0	2	∞
4	∞	1	1	0	∞
5	∞	4	∞	2	0

	1	2	3	4	5
1	0	∞	6	3	∞
2	2	0	∞	5	∞
3	∞	∞	0	2	∞
4	∞	1	1	0	∞
5	∞	∞	∞	2	0

	1	2	3	4	5
1	0	∞	6	3	∞
2	2	0	8	5	∞
3	∞	∞	0	2	∞
4	3	1	1	0	∞
5	6	4	12	2	0

	1	2	3	4	5
1	0	∞	6	3	∞
2	2	0	8	5	∞
3	∞	∞	0	2	∞
4	3	1	1	0	∞
5	6	4	12	2	0

Ans

Time comp $\rightarrow O(V^3)$

Space comp $\rightarrow O(V^2)$