

## Optimal Approach (Using Binary Search):

Using Binary Search to optimize.

If we observe carefully then our code we see that the ~~the~~ answer space  $[1, n]$  is sorted -

Thus we apply binary search in the answer space.

Edge case: How to eliminate the halves.

- place Low at 1, high at  $m$
- based on ~~mid~~  $mid$  powered to  $n$  we check if we get our answer

- if value is smaller we eliminate left half
- if value is higher we eliminate right half

But - if  $m$  and  $n$  are big

we can't store  $mid^n$  in a variable -  
to resolve this we use a function

$f(n, m, mid)$ :

- first declare variable  $ans$  to store  $mid^n$
- run loop  $n$  times and then multiply with ' $ans$ '
- if at a point  $ans > m$ , return 2
- loop completed  $ans == m$  then return 1
- if value is smaller return 0:



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based on this above function, we check if 'mid' is our possible answer or we eliminate the halves. Thus we can avoid the integer overflow case.

### Algorithm -

1. Place 2 pointer i.e. high and low
2. Calculate mid

$$\text{mid} = (\text{low} + \text{high}) // 2$$

("//" refers to integer division)

3. Eliminate the halves accordingly:  
(1) If  $\text{func}(n, m, \text{mid}) == 1$ :

we can conclude that mid is our answer.

- (2) If  $\text{func}(n, m, \text{mid}) == 0$ :

we can conclude that mid is smaller than answer.

eliminate left half }  
consider right half } (i.e.  $\text{low} = \text{mid} + 1$ ).

- (3) If  $\text{func}(n, m, \text{mid}) == 2$ :

conclude: mid is greater than answer

eliminate right half }  
consider left half } ( $\text{high} = \text{mid} - 1$ )