Advanced Algorithms

LAB -1

Language: C / C++ / Python

Brute-Force Algorithm

Initially, P is aligned with T at the first index position. P is then compared with T from **left-to-right**. If a mismatch occurs, "slide" P to *right* by 1 position, and start the comparison again.

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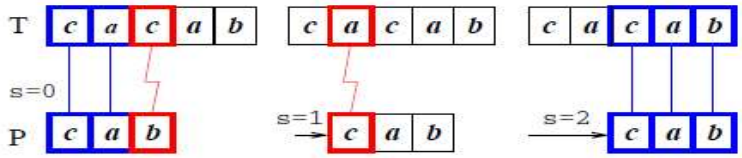

Complexity : Θ (m(n-m+1)) is equivalent to Θ(mn)

Brute-Force Algorithm

```
BF_StringMatcher(T, P) {
 n = length(T);
 m = length(P);

 // s increments by 1 in each iteration
 // => slide P to right by 1
 for (s=0; s<=n-m; s++) {
  // starts the comparison of P and T again
  i=1; j=1;
  while (j<=m && T[s+i]==P[j]) {
   // corresponds to compare P and T from
   // left-to-right
   i++; j++;
  }
  if (j==m+1)
   print "Pattern occurs with shift=", s
 }
}
```

Horspool's Algorithm

Algorithm 2.11: Horspool
Input: text $T = T[0 \ldots n)$, pattern $P = P[0 \ldots m)$
Output: position of the first occurrence of P in T
Preprocess:
(1) for $c \in \Sigma$ do $shift[c] \leftarrow m$
(2) for $i \leftarrow 0$ to $m-2$ do $shift[P[i]] \leftarrow m-1-i$

Horspool's Algorithm

Text: JIMY_HAILED_THE_LEADER_TO_STOP
Pattern: LEADER

```
JIMY_RAN_AND_HAILED_THE_LEADER_TO_STOP
     ||       |  |       |       ||   |
LEADER        |  |       |       ||   |
```

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                            LEADER
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The worst case cost is $\Theta(nm)$, but for random text is $\Theta(n)$.

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(1) for $c \in \Sigma$ do $shift[c] \leftarrow m$
(2) for $i \leftarrow 0$ to $m-2$ do $shift[P[i]] \leftarrow m-1-i$
Search:
(3) $j \leftarrow 0$
(4) while $j + m \leq n$ do
(5) if $P[m-1] = T[j+m-1]$ then
(6) $i \leftarrow m-2$
(7) while $i \geq 0$ and $P[i] = T[j+i]$ do $i \leftarrow i-1$
(8) if $i = -1$ then return j
(9) $j \leftarrow j + shift[T[j+m-1]]$
(10) return n

Try Yourself

Text: `JIM_SAW_ME_IN_A_BARBER_SHOP`

Pattern: `BARBER`

Try Yourself (Solution)