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SUBJECT	Design and Analysis of Algorithms
EXPERIMENT NO:	2(B)
AIM:	Understanding more concepts regarding quick sort algorithm
Algorithm:	<p>Quick Sort Algorithm</p> <pre> partition (arr[], low, high) { // pivot (Element to be placed at right position) pivot = arr[high]; i = (low - 1) // Index of smaller element and indicates the // right position of pivot found so far for (j = low; j <= high- 1; j++){ // If current element is smaller than the pivot if (arr[j] < pivot){ i++; // increment index of smaller element swap arr[i] and arr[j] } } swap arr[i + 1] and arr[high]) </pre>

	<pre> return (i + 1) } quickSort(arr[], low, high) { if (low < high) { /* pi is partitioning index, arr[pi] is now at right place */ pi = partition(arr, low, high); quickSort(arr, low, pi - 1); // Before pi quickSort(arr, pi + 1, high); // After pi } </pre>
Code:	<pre> #include <stdio.h> #include <stdlib.h> #include <time.h> long SWAP = 0; void merge(int arr[], int p, int q, int r) { int i, j, k; int n1 = q - p + 1; int n2 = r - q; int L[n1], R[n2]; for (i = 0; i < n1; i++) L[i] = arr[p + i]; for (j = 0; j < n2; j++) R[j] = arr[q + 1 + j]; i = 0; j = 0; k = p; while (i < n1 && j < n2) { </pre>

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        if (L[i] <= R[j]) {
            arr[k] = L[i];
            i++;
        } else {
            arr[k] = R[j];
            j++;
        }
        k++;
    }
    while (i < n1) {
        arr[k] = L[i];
        i++;
        k++;
    }
    while (j < n2) {
        arr[k] = R[j];
        j++;
        k++;
    }
}

void mergeSort(int arr[], int l, int r)
{
    if (l < r) {
        int m = l + (r - l) / 2;
        mergeSort(arr, l, m);
        mergeSort(arr, m + 1, r);
        merge(arr, l, m, r);
    }
}

int quicksort(int a[], int start, int end) {
    int pivot = a[end];
    //int pivot = a[start];
    //int random = start + rand() % (end - start);
    //int pivot = a[random];

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//int mid = start + (end - start)/2;
//int pivot = a[mid];
int i = (start - 1);
for (int j = start; j <= end - 1; j++) {
    if (a[j] < pivot) {
        i++;
        int t = a[i];
        a[i] = a[j];
        a[j] = t;
        SWAP++;
    }
}
int t = a[i + 1];
a[i + 1] = a[end];
a[end] = t;
SWAP++;
return (i + 1);
}

double quick(int a[], int start, int end) {
    if (start < end) {
        int p = quicksort(a, start, end);
        quick(a, start, p - 1);
        quick(a, p + 1, end);
    }
}

int main() {
    double qust, mest;
    srand(time(0));
    FILE * fp, * file;
    fp = fopen("random.txt", "w");
    for (int i = 0; i < 100000; i++) {
        fprintf(fp, "%d\n", rand() % 900001 + 100000);
    }
    int upper_limit = 100;
    fclose(fp);
}

```

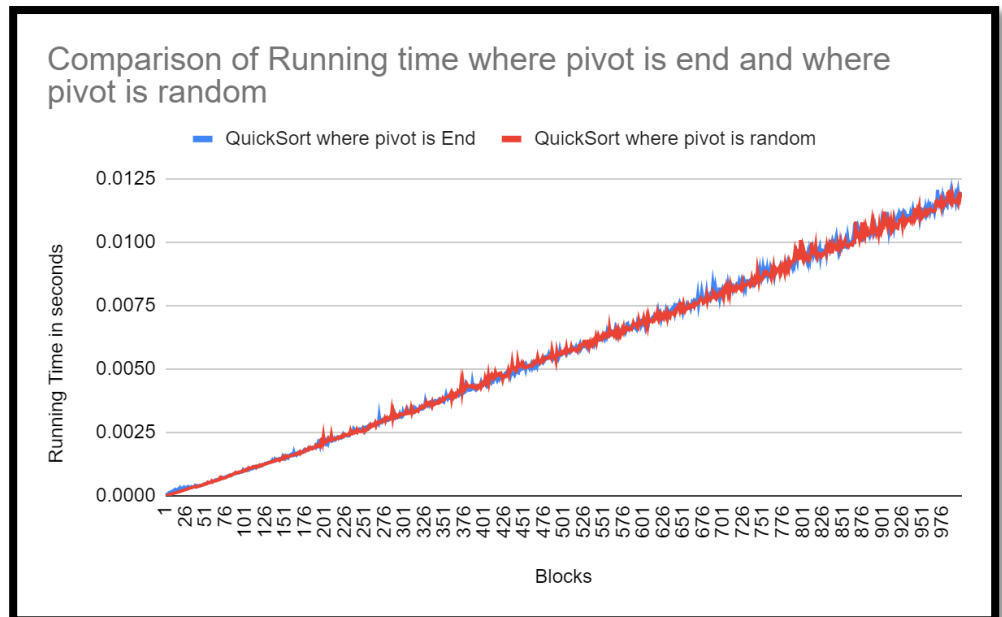
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file = fopen("outputEnd.txt", "w");
fprintf(file,
"Block\tMerSort\tQuickSort\tSwaps\n");
for (int i = 0; i < 1000; i++) {
    fp = fopen("random.txt", "r");
    int arr1[upper_limit], arr2[upper_limit],
temp_num;
    for (int j = 0; j < upper_limit; j++) {
        fscanf(fp, "%d", & temp_num);
        arr1[j] = temp_num;
        arr2[j] = temp_num;
    }
    fclose(fp);
    clock_t t;
    t = clock();
    mergeSort(arr2, 0, upper_limit - 1);
    t = clock() - t;
    mest = ((double) t) / CLOCKS_PER_SEC;
    clock_t t1;
    t1 = clock();
    qust = quick(arr1, 0, upper_limit - 1);
    t1 = clock() - t1;
    qust = ((double) t1) / CLOCKS_PER_SEC;
    fprintf(file, "%d\t%lf\t%lf\t%ld\n", i + 1,
mest, qust, SWAP);
    fflush(stdout);
    upper_limit += 100;
}
return 0;
}

```

Graphs and Observation:

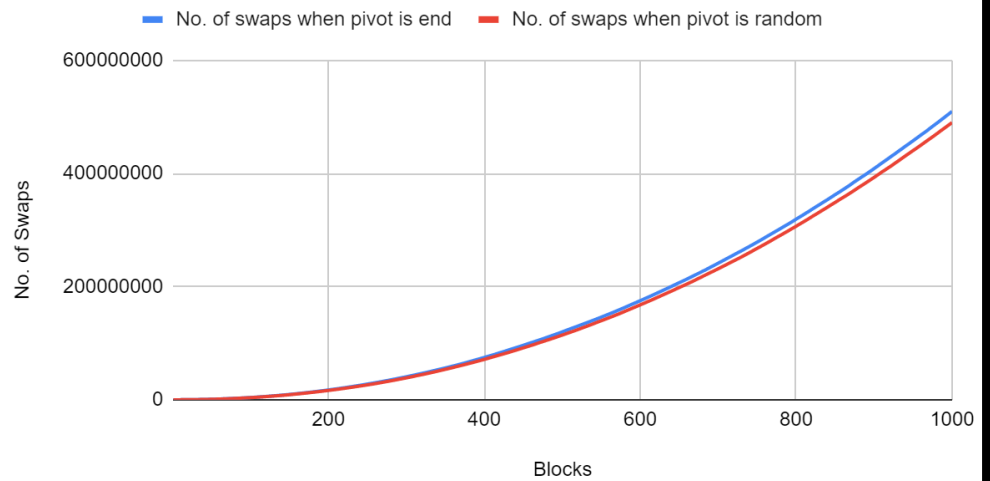
Running time comparison for Different Pivot Positions



- Here, we can see that the time complexity of quick sort is nearly the same even when varied pivot points are taken into account.
- We can see that a quick sort when the pivot is at a random position takes longer than when the pivot is in an end position at the conclusion of execution.
- Despite the fact that both executions are finished in 0.1 seconds.

Number of swaps considering different pivot positions

Comparison of No. of swaps required when pivot is end and where pivot is random



- We can see from this that fewer swaps are needed for a quick sort when the pivot is in a random position as opposed to when it is at the end.
- The average number of swaps is over 500,000,000.
- Throughout the entire execution, the number of swaps keeps rising.

Conclusion:

Thus, we have provided observations for different pivots for quick sort algorithms.