

CC EXPERIMENT NO. 05

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Aim:

Implement LR(0) parser for Given Grammar.

- (1) $E \rightarrow E \text{ sub } E \text{ sup } E$
- (2) $E \rightarrow E \text{ sub } E$
- (3) $E \rightarrow E \text{ sup } E$
- (4) $E \rightarrow \{ E \}$
- (5) $E \rightarrow c$

Theory:

The LR parser is an efficient bottom up syntax analysis technique that can be used to large class of context-free grammar. This technique is also called LR(0) parsing.

- L stands for left to right scanning
- R stands for rightmost derivation in reverse • 0 stands for no. of input symbols of lookahead.

Augmented grammar:

If P is a grammar with starting symbol S, then G' (augmented grammar for G) is a grammar with a new starting symbol S' and productions $S \rightarrow .S'$. The purpose of this new starting production is to indicate the parser when it should stop parsing. The ' . ' before S indicates the left side of ' . ' has been read by a compiler and the right side of ' . ' is yet to be read by a compiler.

Steps for constructing the LR parsing table :

1. Writing augmented grammar
2. LR(0) collection of items to be found
3. Defining 2 functions: goto(list of terminals) and action(list of non-terminals) in the parsing table.

$I_0 \rightarrow E' \rightarrow \cdot E$
 $E \rightarrow \cdot E \text{ sub } E \text{ sup } E$
 $E \rightarrow \cdot E \text{ sub } E$
 $E \rightarrow \cdot E \text{ sup } E$
 $E \rightarrow \cdot \{E\}$
 $E \rightarrow \cdot C$

$I_1 : E' \rightarrow E \cdot$
 $E \rightarrow E \cdot \text{ sub } E \text{ sup } E$
 $E \rightarrow E \cdot \text{ sub } E$
 $E \rightarrow E \cdot \text{ sup } E$

$I_2 : E \rightarrow \{ \cdot E \}$
 $E \rightarrow \cdot E \text{ sub } E \text{ sup } E$
 $E \rightarrow \cdot E \text{ sub } E$
 $E \rightarrow \cdot E \text{ sup } E$
 $E \rightarrow \cdot \{E\}$
 $E \rightarrow \cdot C$

$I_3 : E \rightarrow C \cdot$

$I_4 : E \rightarrow E \text{ sub } \cdot E \text{ sup } E$
 $E \rightarrow E \text{ sub } \cdot E$
 $E \rightarrow \cdot E \text{ sub } E \text{ sup } E$
 $E \rightarrow \cdot E \text{ sub } E$
 $E \rightarrow \cdot E \text{ sup } E$
 $E \rightarrow \cdot \{E\}$
 $E \rightarrow \cdot C$

$I_5: E \rightarrow E \text{ sup } E$
 $E \rightarrow \cdot E \text{ sup } E \text{ sup } E$
 $E \rightarrow \cdot E \text{ sub } E$
 $E \rightarrow \cdot E \text{ sup } E$
 $E \rightarrow \cdot \{E\}$
 $E \rightarrow \cdot c$

$I_6: E \rightarrow E \cdot \text{sup } E \text{ sup } E$
 $E \rightarrow E \cdot \text{sub } E$
 $E \rightarrow E \cdot \text{sup } E$
 $E \rightarrow \{E \cdot \}$

$I_7: E \rightarrow E \cdot \text{sub } E \text{ sup } E$
 $E \rightarrow E \text{ sub } E \cdot \text{sup } E$
 $E \rightarrow E \cdot \text{sup } E$
 $E \rightarrow E \text{ sub } E \cdot$
 $E \rightarrow E \cdot \text{sup } E$

$I_8: E \rightarrow E \cdot \text{sub } E \text{ sup } E$
 $E \rightarrow E \cdot \text{sub } E$
 $E \rightarrow E \cdot \text{sup } E$
 $E \rightarrow E \text{ sup } E \cdot$

$I_9: \{E\} \cdot$

$I_{10}: E \rightarrow E \text{ sub } E \text{ sup } \cdot E$
 $E \rightarrow E \text{ sup } \cdot E$
 $E \rightarrow \cdot E \text{ sub } E \text{ sup } E, \cdot E \text{ sub } E, \cdot E \text{ sup } E, \cdot \{E\}$

$I_{11}: E \rightarrow E \cdot \text{sub } E \text{ sup } E$
 $E \rightarrow E \text{ sub } E \text{ sup } \cdot E$
 $E \rightarrow E \cdot \text{sub } E, E \text{ sub } E \cdot, E \cdot \text{sup } E$

Parsing Table:

State	action						goto
	sub	sup	{	}	c	\$	E
0			S2		S3		1
1	S4	S5			✓		
2		S4	S2		S3		6
3	r5	r5		r5		r5	
4			S2		S3		7
5			S2		S3		8
6	S4	S5		S9			
7	S4	S10		r2		r2	
8	S4	S5		r3		r3	
9	r4	r4		r4		r4	
10			S2		S3		11
11	S4	S5		r1		r1	

Implementation Code:

```
import java.util.*;
import java.util.stream.*;

public class Parser {
    private final List<List<String>> production = Arrays.asList(
        Arrays.asList("E", "sub", "E", "sup", "E"),
        Arrays.asList("E", "sub", "E"),
        Arrays.asList("E", "sup", "E"),
        Arrays.asList("E", "{", "E", "}"),
        Arrays.asList("E", "c")
    );

    private final Integer[][] table = new Integer[][]{
        {2, null, null, null, 3, null, 1},
        {null, null, 5, 4, null, -1, null},
        {2, null, null, null, 3, null, 6},
        {null, 5, 5, 5, null, 5, -1},
        {2, null, null, null, 3, null, 7},
    };
}
```

```

        {2, null, null, null, 3, null, 8},
        {null, 9, 5, 4, null, null, -1},
        {null, 2, 4, 10, null, 2, -1},
        {null, 3, 5, 4, null, 3, -1},
        {null, 4, 4, 4, null, 4, -1},
        {2, null, null, null, 3, null, 11},
        {null, 1, 5, 4, null, 1, -1}
    };

    private final Map<String, Integer> terminals = new HashMap<String, Integer>()
    {{
        put("{", 0);
        put("}", 1);
        put("sup", 2);
        put("sub", 3);
        put("c", 4);
        put("$", 5);
    }};

    private final Map<String, Integer> nonTerminals = new HashMap<String,
Integer>() {{
        put("S", 0);
        put("E", 1);
    }};

    public boolean parse(String input) {
        List<String> tokens = Arrays.stream(input.split("
")).collect(Collectors.toList());
        tokens.add("$");

        Deque<Integer> stack = new ArrayDeque<>();
        stack.push(0);
        int ptr = 0;
        boolean error = false;

        while (true) {
            System.out.print("Stack: ");
            System.out.println(stack);

            int sd = stack.peek();
            int bd = terminals.get(tokens.get(ptr));
            Integer ptd = table[sd][bd];

            // Error Occurence
            if (ptd == null) {

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        error = true;
        break;
    }

    // String accepted
    if (ptd == -1) {
        break;
    }

    // REDUCE Case
    if (ptd < 0) {
        List<String> prod = production.get(-ptd - 1);
        int l = prod.get(1).equals("") ? 1 : prod.size();

        if (stack.size() < 2 * l) {
            error = true;
            break;
        }

        for (int i = 0; i < 2 * l; i++) {
            stack.pop();
        }

        sd = stack.peek();
        bd = nonTerminals.get(prod.get(0));

        if (table[sd][bd] == null) {
            error = true;
            break;
        }

        stack.push(bd);
        stack.push(table[sd][bd]);
    }

    // SHIFT Case

    else {
        stack.push(ptd);
        ptr++;
    }
}

if (error) {
    System.out.println("Error: Invalid input");
}

```

```

        return false;
    } else {
        System.out.println("Success: Valid input");
        return true;
    }
}

public static void main(String[] args) {
    Parser parser = new Parser();

    // Test input strings
    String input1 = "c";
    String input2 = "{ c }";
    String input3 = "{ sub c }";
    String input4 = "{ c sup c }";
    String input5 = "{ sub c sup c }";
    String input6 = "{ sub c } sup c";
    String input7 = "{ sub { c sup c } }";
    String input8 = "{ sub { c sub c } sup { c sub c } }";

    parser.parse(input1);
    parser.parse(input2);
    parser.parse(input3);
    parser.parse(input4);
    parser.parse(input5);
    parser.parse(input6);
    parser.parse(input7);
    parser.parse(input8);
}
}

```

Result:

The Parsing Table:

	{	}	sup	sub	c	\$	E
0	(S, 2)	None	None	None	(S, 3)	None	1
1	None	None	(S, 5)	(S, 4)	None	Accept	-1
2	(S, 2)	None	None	None	(S, 3)	None	6
3	None	(r, 5)	(r, 5)	(r, 5)	None	(r, 5)	-1
4	(S, 2)	None	None	None	(S, 3)	None	7
5	(S, 2)	None	None	None	(S, 3)	None	8
6	None	(S, 9)	(S, 5)	(S, 4)	None	None	-1
7	None	(r, 2)	(S, 4)	(S, 10)	None	(r, 2)	-1
8	None	(r, 3)	(S, 5)	(S, 4)	None	(r, 3)	-1
9	None	(r, 4)	(r, 4)	(r, 4)	None	(r, 4)	-1
10	(S, 2)	None	None	None	(S, 3)	None	11
11	None	(r, 1)	(S, 5)	(S, 4)	None	(r, 1)	-1

Enter your input string: c sub { c }

Your input: ['c', 'sub', '{', 'c', '}', '\$']

Stack: 0
 Stack: 0 c 3
 Stack: 0 E 1
 Stack: 0 E 1 sub 4
 Stack: 0 E 1 sub 4 { 2
 Stack: 0 E 1 sub 4 { 2 c 3
 Stack: 0 E 1 sub 4 { 2 E 6
 Stack: 0 E 1 sub 4 { 2 E 6 } 9
 Stack: 0 E 1 sub 4 E 7
 Stack: 0 E 1

[+] The given input is VALID!

The Parsing Table:

	{	}	sup	sub	c	\$	E
0	(S, 2)	None	None	None	(S, 3)	None	1
1	None	None	(S, 5)	(S, 4)	None	Accept	-1
2	(S, 2)	None	None	None	(S, 3)	None	6
3	None	(r, 5)	(r, 5)	(r, 5)	None	(r, 5)	-1
4	(S, 2)	None	None	None	(S, 3)	None	7
5	(S, 2)	None	None	None	(S, 3)	None	8
6	None	(S, 9)	(S, 5)	(S, 4)	None	None	-1
7	None	(r, 2)	(S, 4)	(S, 10)	None	(r, 2)	-1
8	None	(r, 3)	(S, 5)	(S, 4)	None	(r, 3)	-1
9	None	(r, 4)	(r, 4)	(r, 4)	None	(r, 4)	-1
10	(S, 2)	None	None	None	(S, 3)	None	11
11	None	(r, 1)	(S, 5)	(S, 4)	None	(r, 1)	-1

Enter your input string: { c sup c

Your input: ['{', 'c', 'sup', 'c', '\$']

Stack: 0
 Stack: 0 { 2
 Stack: 0 { 2 c 3
 Stack: 0 { 2 E 6
 Stack: 0 { 2 E 6 sup 5
 Stack: 0 { 2 E 6 sup 5 c 3
 Stack: 0 { 2 E 6 sup 5 E 8
 Stack: 0 { 2 E 6

[+] The given input is INVALID..

Conclusion:

From the above experiment, I was able to implement code and programmatically execute and verify the working of LR(0) parser by manually finding the parsing table for a given grammar.