

EXPERIMENT 1

Aim:

To convert a Regular Expression to optimized DFA.

Theory:

Convert Regular Expression to DFA

- Uses augmented regular expression $r\#$.
- Important states of NFA correspond to positions in regular expression that hold symbols of the alphabet.
- Regular expression is represented as syntax tree where interior nodes correspond to operators representing union, concatenation and closure operations.
- Leaf nodes corresponds to the input symbols
- Construct DFA directly from a regular expression by computing the functions $nullable(n)$, $firstpos(n)$, $lastpos(n)$ and $followpos(i)$ from the syntax tree.
 - **nullable** (n): Is true for $*$ node and node labeled with ϵ . For other nodes it is false.
 - **firstpos** (n): Set of positions at node ti that corresponds to the first symbol of the sub-expression rooted at n .
 - **lastpos** (n): Set of positions at node ti that corresponds to the last symbol of the sub-expression rooted at n .
 - **followpos** (i): Set of positions that follows given position by matching the first or last symbol of a string generated by sub-expression of the given regular expression.

Rules for computing nullable, firstpos and lastpos

Node n	nullable (n)	firstpos (n)	lastpos (n)
A leaf labeled ϵ	True	\emptyset	\emptyset
A leaf with position i	False	$\{i\}$	$\{i\}$
An or node $n = c_1 \mid c_2$	Nullable (c_1) or Nullable (c_2)	firstpos (c_1) U firstpos (c_2)	lastpos (c_1) U lastpos (c_2)

	Nullable (c_2)	firstpos (c_2)	lastpos (c_2)
A cat node $n = c_1c_2$	Nullable (c_1) and Nullable (c_2)	If (Nullable (c_1)) firstpos (c_1) U firstpos (c_2) else firstpos (c_1)	If (Nullable (c_2)) lastpos (c_1) U lastpos (c_2) else lastpos (c_1)
A star node $n = c_1^*$	True	firstpos (c_1)	lastpos (c_1)

Computation of followpos

The position of regular expression can follow another in the following ways:

- If n is a cat node with left child c_1 and right child c_2 , then for every position i in $lastpos(c_1)$, all positions in $firstpos(c_2)$ are in $followpos(i)$.
- For cat node, for each position i in $lastpos$ of its *left child*, the *firstpos* of its *right child* will be in $followpos(i)$.
- If n is a star node and i is a position in $lastpos(n)$, then all positions in $firstpos(n)$ are in $followpos(i)$.
- For star node, the *firstpos* of that node is in *followpos* of all positions in *lastpos* of that node.

Implementation

Code:

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# TE Comps

leaf_no = 0
leaf_array = []
follow_pos = []
print()

# input a regular expression
re = input(" [+] Enter the regular expression: ")

# convert to augmented regular expression
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are = '('
for e in re:
    if are[-1] in [')'] and e not in [')', '|', '*']:
        are = are + '.' + e
    elif are[-1] not in ['(', ')', '|'] and e not in [')', '|', '*']:
        are = are + '.' + e
    else:
        are = are + e
are = are[1:]+".#"
print('\n - Augmented regular expression: ' + are + '\n')

# Class to construct a syntax tree for the ARE
class SyntaxTree():
    content = '.'
    nullable = False
    first_pos = set()
    last_pos = set()
    leaf_number = int()
    left = None
    right = None

    def __init__(self, content, leaf_number, left, right):
        self.content = content
        self.leaf_number = leaf_number
        self.left = left
        self.right = right
        if content in ['*']:
            self.nullable = True

    def __str__(self) -> str:
        return self.content + ' ' + str(self.leaf_number) + ' ' +
str(self.nullable) + ' ' + str(self.first_pos) + ' ' + str(self.last_pos)

    def update_nullable(self):
        if self.content == '|':
            self.nullable = bool(self.right.nullable) or bool(self.left.nullable)
        elif self.content == '.':
            self.nullable = bool(self.right.nullable) and
bool(self.left.nullable)

    def update_first_pos(self):
        if self.content == '*':
            self.first_pos = self.left.first_pos
        elif self.content == '|':
            lfp = self.left.first_pos

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        rfp = self.right.first_pos
        self.first_pos = lfp | rfp
    elif self.content == '.':
        ln = self.left.nullable
        lfp = self.left.first_pos
        rfp = self.right.first_pos
        if ln:
            self.first_pos = lfp | rfp
        else:
            self.first_pos = lfp
    else:
        self.first_pos = {self.leaf_number}

def update_last_pos(self):
    if self.content == '*':
        self.last_pos = self.left.last_pos
    elif self.content == '|':
        llp = self.left.last_pos
        rlp = self.right.last_pos
        self.last_pos = llp | rlp
    elif self.content == '.':
        rn = self.right.nullable
        llp = self.left.last_pos
        rlp = self.right.last_pos
        if rn:
            self.last_pos = llp | rlp
        else:
            self.last_pos = rlp
    else:
        self.last_pos = {self.leaf_number}

def update_nfl(self):
    if self.left:
        self.left.update_nfl()
    if self.right:
        self.right.update_nfl()
    self.update_nullable()
    self.update_first_pos()
    self.update_last_pos()

def print_tree(self):
    if self.left:
        self.left.print_tree()
    print(self)

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        if self.right:
            self.right.print_tree()

# nullable, firstpos, lastpos, followpos
# re -> dfa

# function to create a syntax tree and return its root node
def create_syntax_tree(are):
    print(are)
    global leaf_no, leaf_array, follow_pos
    if len(are) == 1:
        leaf_no += 1
        head = SyntaxTree(are, leaf_no, None, None)
        leaf_array.append(head)
        follow_pos.append(set())
        return head

    stack = 0
    flag = True
    for e in are:
        if e == '(':
            stack += 1
        if e == ')':
            stack -= 1

        if (e == '.' or e == '|') and stack == 0:
            flag = False

    if flag:
        re = are
        if re[-1] == '*':
            if re[0] == '(':
                left = create_syntax_tree(re[1:-2])
            else:
                left = create_syntax_tree(re[:-1])

            head = SyntaxTree('*', -1, left, None)
            return head
        if re[0] == '(':
            return create_syntax_tree(re[1:-1])

    stack = 0
    temp = ''
    left = None

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right = None
prev = None
root = None
for e in are+'.':
    if e == '(':
        stack += 1
    if e == ')':
        stack -= 1

    if (e == '.' or e == '|') and stack == 0:
        if left == None:
            left = create_syntax_tree(temp)
        elif right == None:
            right = create_syntax_tree(temp)
            root = SyntaxTree(prev, -1, left, right)
        else:
            left = root
            right = create_syntax_tree(temp)
            root = SyntaxTree(prev, -1, left, right)
        prev = e
        temp = ''
    else:
        temp = temp + e
return root

# calculate follow pos of the syntax tree
def caluculate_follow_pos(head):
    if head:
        global follow_pos
        caluculate_follow_pos(head.left)

        if head.content == '*':
            for i in head.last_pos:
                follow_pos[i-1] = follow_pos[i-1] | head.first_pos

        if head.content == '.':
            for i in head.left.last_pos:
                follow_pos[i-1] = follow_pos[i-1] | head.right.first_pos

        caluculate_follow_pos(head.right)

head = create_syntax_tree(are)

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head.update_nfl()
print("The tree is:")
head.print_tree()
print("-----")
caluculate_follow_pos(head)
print(" FOLLOW-POS TABLE ")
for i, leaf in enumerate(leaf_array):
    print(leaf.content , '\t' , leaf.leaf_number , '\t' , follow_pos[i])
print()

# to get the unique terminals
terminals = []
for i in leaf_array:
    terminal = i.content
    if terminal == '#':
        continue
    if terminal not in terminals:
        terminals.append(terminal)

# Making of the DFA table
states = [head.first_pos]
table = []
ptr = 0

while ptr < len(states):
    sub_table = []
    for terminal in terminals:
        cur_state = set()
        for i in states[ptr]:
            if leaf_array[i-1].content == terminal:
                cur_state = cur_state.union(follow_pos[i-1])
            if cur_state not in states:
                states.append(cur_state)
        sub_table.append(states.index(cur_state))
    table.append(sub_table)
    ptr += 1

# Printing the final DFA table
A = ord('A')

print("\n Minimized DFA TABLE ")
for i in [''] + terminals:
    print(i, end='\t')

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print("\n-----"+"-----"*len(terminals))

for id, row in enumerate(table):
    print(chr(A+id), end='\t')
    for column in row:
        print(chr(A+column), end='\t')
    print()
print()

```

Output:

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[+] Enter the regular expression: (a|b)*ab
- Augmented regular expression: (a|b)*.a.b.#

```

FOLLOW-POS TABLE

a	1	{1, 2, 3}
b	2	{1, 2, 3}
a	3	{4}
b	4	{5}
#	5	set()

Minimized DFA TABLE

	a	b
A	B	A
B	B	C
C	B	A

```

[+] Enter the regular expression: (a|b)ab(a|b)*

```

```

- Augmented regular expression: (a|b).a.b.(a|b)*. #

```

FOLLOW-POS TABLE

a	1	{3}
b	2	{3}
a	3	{4}
b	4	{5, 6, 7}
a	5	{5, 6, 7}
b	6	{5, 6, 7}
#	7	set()

Minimized DFA TABLE

	a	b
A	B	B
B	C	D
C	D	E
D	D	D
E	E	E

Conclusion:

From the above experiment, I was able to implement code and programatically convert a Regular Expression to optimized DFA

Ref.:

<https://ecomputernotes.com/compiler-design/convert-regular-expression-to-dfa>

<https://www.youtube.com/watch?v=rGRSiPSmhwE>