

Introduction to Artificial Intelligence

Course 16 :198 :440

Recitation 3:

Informed Search

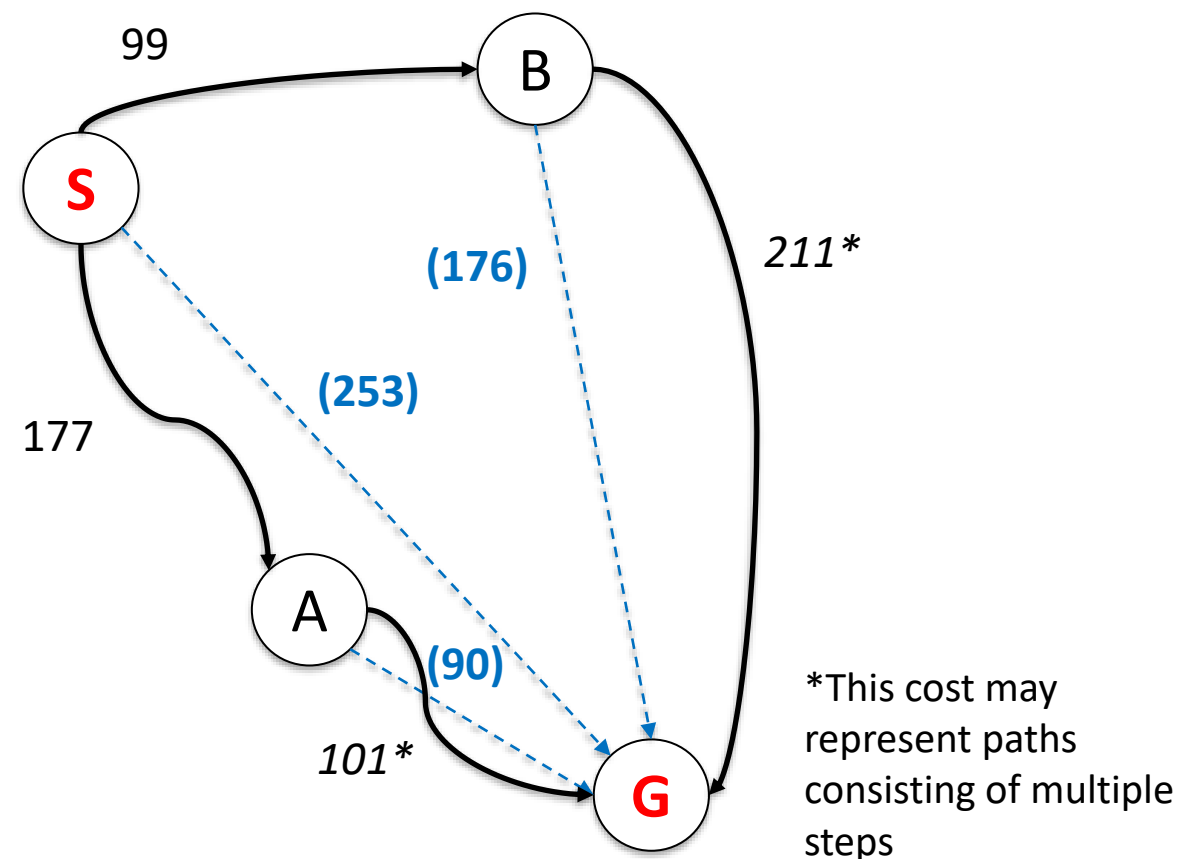


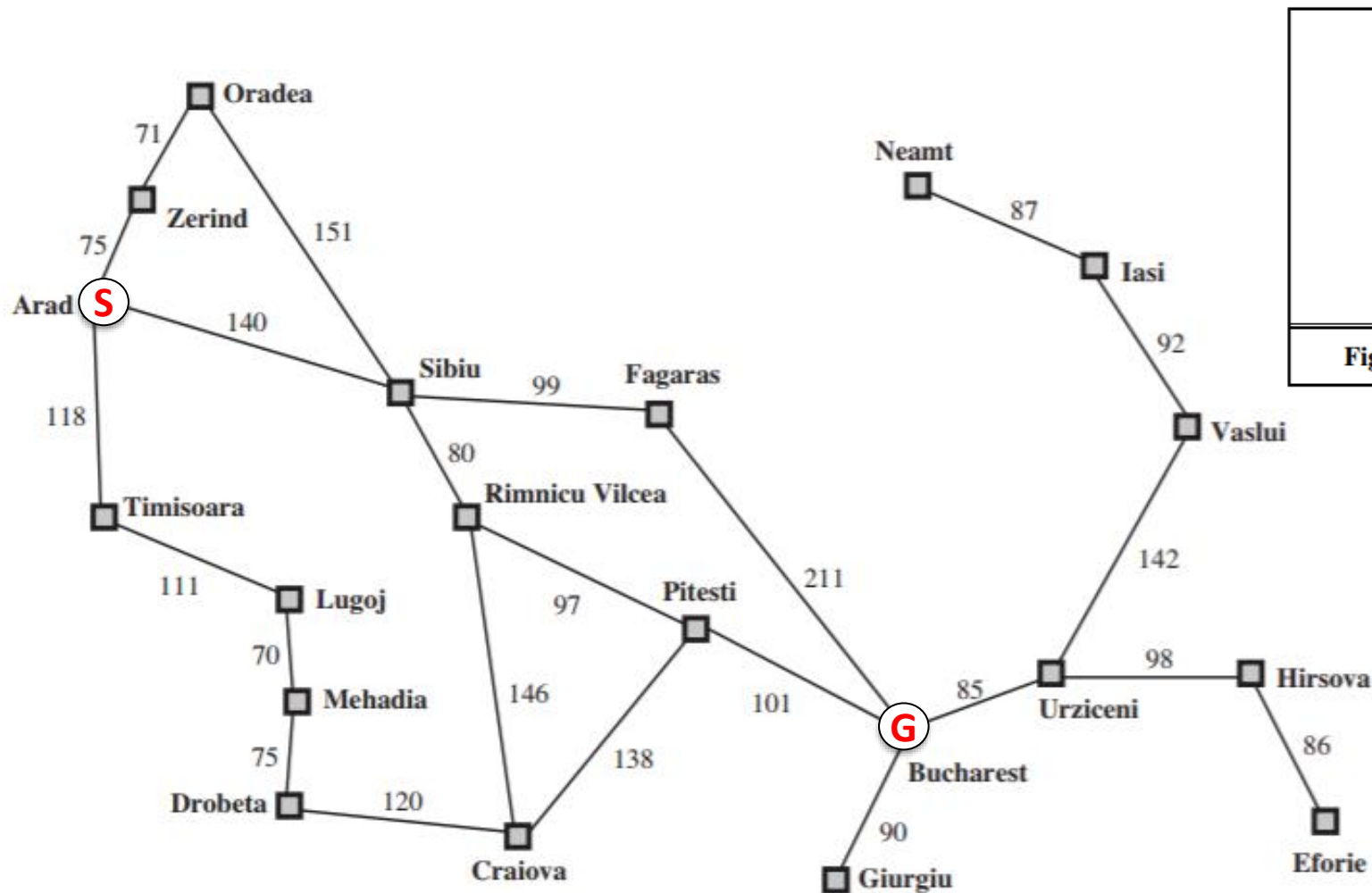
RUTGERS

Consider going from S to G on this piece of a larger graph...

- Geometrically, the 'A' path seems shorter, as we are guided by the spatial proximity of A to G.
- What would uniform-cost expand first, A or B?
- How can an 'intelligent agent' be guided by this logic, such that it will consider A first? This is what brings us to *informed* search – and the A* algorithm.

Additionally, imagine paths that point northwest from S. These should strike you as more likely to be suboptimal. The **Euclidean heuristic** reflects this, and deprioritizes them in the queue.





Arad	366	Mehadia	241
Bucharest	0	Neamt	234
Craiova	160	Oradea	380
Drobeta	242	Pitesti	100
Eforie	161	Rimnicu Vilcea	193
Fagaras	176	Sibiu	253
Giurgiu	77	Timisoara	329
Hirsova	151	Urziceni	80
Iasi	226	Vaslui	199
Lugoj	244	Zerind	374

Figure 3.22 Values of h_{SLD} —straight-line distances to Bucharest.

Execute A*, using the straight line distance heuristic, to find the shortest path from Arad to Bucharest

For Tree Search, admissibility is sufficient to achieve A optimality*

Definition

A heuristic function h is said to be **admissible** if it *never overestimates* the cost to reach the goal, i.e.

$$\forall n : h(n) \leq h^*(n),$$

where $h^*(n)$ is the true cost of the shortest path from n to the goal.

Theorem

If h is admissible then the tree search A is optimal.*

For graph search A^* , we need a stronger requirement on h because repeated states are not allowed.

Definition

A heuristic function h is said to be **consistent** if

$$\forall(n, a, n') : h(n) \leq c(n, a, n') + h(n').$$

where $c(n, a, n')$ is the step cost for going from n to n' using action a .

Theorem

If h is consistent then the graph search A^ is optimal.*