

Team Notebook

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1 Advice

Pre-submit:

Are time limits close? If so, generate max cases.

Is the memory usage fine? Could anything overflow? Make sure to submit the right file.

Wrong answer: Print your solution! Print debug output, as well.

Are you clearing all datastructures between test cases?

Can your algorithm handle the whole range of input?

Read the full problem statement again. Do you handle all corner cases correctly? Have you understood the problem correctly?

Any uninitialized variables? Any overflows? Confusing N and M, i and j, etc.?

Are you sure

your algorithm works? What special cases have you not thought of?

Are you sure the STL functions you use work as you think? Add some assertions, maybe resubmit. Create some testcases to run your algorithm on. Go through the algorithm for a simple case.

Go through this list again. Explain your algorithm to a team mate.

Ask the team mate to look at your code. Go for a small walk, e.g. to the toilet. Is your output format correct?

Rewrite your solution from the start or let a team mate do it.

Runtime error: Have you tested all corner cases locally?

Any uninitialized variables? Are you reading or writing outside the range of any vector? Any assertions that might fail? Any possible division by 0? (mod 0 for example)

Any possible infinite recursion? Invalidated pointers or iterators? Are you using too much memory? Debug with resubmits.

Time limit exceeded: Do you have any possible infinite loops?

What is the complexity of your algorithm? Are you copying a lot of unnecessary data? (References) How big is the input and output? (consider scanf) Avoid vector, map. (use arrays/unordered_map) What do your team mates think about your algorithm?

Memory limit exceeded: What is the max amount of

memory your algorithm should need? Are you clearing all datastructures between test cases?

Primes - 10001st prime is 1299721, 100001st prime is 15485867 Large primes - 999999937, $1e9+7$, 987646789, 987101789 78498 primes less than 10^6 The number of divisors of n is at most around 100, for $n < 5e4$, 500 for $n < 1e7$, 2000 for $n < 1e10$, 200,000 for $n < 1e19$ 7! 5040, 8! 40320, 9! 362880, 10! 362880, 11! $4.0e7$, 12! $4.8e8$, 15! $1.3e12$, 20! $2e18$

The number of divisors of n is at most around 100 for $n < 5e4$, 500 for $n < 1e7$, 2000 for $n < 1e10$, 200 000 for $n < 1e19$.

Articulation points and bridges articulation point:- there exist child : $\text{dfslow}[\text{child}] > \text{dfsnum}[\text{curr}]$ bridge :- tree edge:
 $\text{dfslow}[\text{ch}] > \text{dfsnum}[\text{par}]$;

A connected multigraph has an Euler path but not an Euler circuit if and only if it has exactly two vertices of odd degree

Binomial coefficients - base case ncn and $\text{nc0} = 1$;
 recursion is $\text{nCk} = (\text{n-1})\text{C}(\text{k-1}) + (\text{n-1})\text{Ck}$

Catalan numbers - used in valid paranthesis expressions - formula is $C_n = \sum_{i=0}^{n-1} (C_i C_{n-i-1})$; Another formula is $C_n = \frac{2n C_n}{(n+1)}$. There are C_n binary trees of n nodes and C_{n-1} rooted trees of n nodes

Derangements - $D(n) = (n-1)(D(n-1) + D(n-2))$

Burnsides Lemma - number of equivalence classes = $(\sum I(\pi))/n$: $I(\pi)$ are number of fixed points.
 Usual
 formula: $[\sum_{i=0}^{n-1} k^{\text{gcd}(i,n)}]/n$

Stirling numbers - first kind - permutations of n elements with k disjoint cycles. $s(n+1, k) = ns(n, k) + s(n, k-1)$. $s(0, 0) = 1$, $s(n, 0) = 0$ if $n > 0$. $\sum_{k=0}^n s(n, k) = n!$

Stirling numbers - Second kind - partition n objects into k non empty subsets. $S(n+1, k) = kS(n, k) + S(n, k-1)$. $S(0, 0) = 1$, $S(n, 0) = 0$ if $n > 0$. $S(n, k) = (\sum_{j=0}^k [(-1)^{(k-j)} C_{kj} j^n])/k!$

Hermite identity - $\sum_{k=0}^{n-1} \text{floor}[(x+k)/n] =$

$\text{floor}[nx]$

Kirchoff matrix tree theorem - number of spanning trees in a graph is determinant of Laplacian Matrix with one row and column removed, where L = degree matrix - adjacency matrix

Expected value tricks:

1. Linearity of Expectation: $E(X+Y) = E(X) + E(Y)$
2. Contribution to the sum - If we want to find the sum over many ways/possibilities, we should consider every element (maybe a number, or a pair or an edge) and count how many times it will be added to the answer.
3. For independent events - $E(XY) = E(X)E(Y)$
4. Ordered pairs (Super interpretation of square) - The square of the size of a set is equal to the number of ordered pairs of elements in the set. So we iterate over pairs and for each we compute the contribution to the answer. Similarly, the k-th power is equal to the number of sequences (tuples) of length k.
5. Powers technique - If you want to maintain the sum of k-th powers, it might help to also maintain the sum of smaller powers. For example, if the sum of 0-th, 1-th and 2-nd powers is S_0, S_1 and S_2 , and we increase every element by x, the new sums are $S_0, S_1 + S_0x$ and $S_2 + 2S_1x + x^2 S_0$.

2 Aho Corasick

```
struct AhoCorasick{
    enum {alpha=26, first='a'};
    struct Node{
        int back, next[alpha], start = -1, end = -1, nmatches = 0;
        Node(int v){memset(next, v, sizeof(next));}
    };
    vector<Node> N;
    vector<int> backp;
    inline void insert(string &s, int j){
        assert(!s.empty());
        int n=0;
        for(auto &c: s){
            int &m=N[n].next[c-first];
            if(m==-1){n=m=N.size(); N.emplace_back(-1);}
            else n=m;
        }
        if(N[n].end==-1) N[n].start=j;
        backp.push_back(N[n].end);
        N[n].end=j;
    }
};
```

```

    N[n].nmatches++;
}
void clear(){
    N.clear();
    backp.clear();
}
void create(vector<string>& pat){
    N.emplace_back(-1);
    for(int i=0;i<pat.size();++i) insert(pat[i],i);
    N[0].back=N.size();
    N.emplace_back(0);
    queue<int> q;
    for(q.push(0);!q.empty();q.pop()){
        int n=q.front(),prev=N[n].back;
        for(int i=0;i<alpha;++i){
            int &ed=N[n].next[i],y=N[prev].next[i];
            if(ed==-1) ed=y;
            else{
                N[ed].back=y;
                (N[ed].end==-1 ? N[ed].end:backp[N[ed].start])=N[y].end;
                N[ed].nmatches+=N[y].nmatches;
                q.push(ed);
            }
        }
    }
}
ll find(string word){
    int n=0;
    // vector<int> res;
    ll count=0;
    for(auto &c: word){
        n=N[n].next[c-first];
        // res.push_back(N[n].end);
        count+=N[n].nmatches;
    }
    return count;
}
};
struct AhoOnline{
    int sz=0;
    vector<string> v[25];
    AhoCorasick c[25];
    void add(string &p){
        int val=__builtin_ctz(~sz);
        auto &cur=v[val];
        for(int i=0;i<val;++i){
            for(auto &it: v[i]) cur.push_back(it);
            c[i].clear();
            v[i].clear();
        }
    }

```

```

        cur.push_back(p);
        c[val].create(cur);
        ++sz;
    }
    ll query(string &p){
        ll ans=0;
        for(int i=0;i<25;++i){
            if((1<<i)&sz) ans+=c[i].find(p);
            if((1<<i)>=sz) break;
        }
        return ans;
    }
} add,del;

```

3 Centroid Decomposition

```

vector<set<int>> g;
vector<int> par,sub;
int dfs(int u,int p){
    sub[u]=1;
    for(auto &it: g[u]) if(it!=p) sub[u]+=dfs(it,u);
    return sub[u];
}
int find_centroid(int u,int p,int n){
    for(auto &it: g[u]){
        if(it!=p && sub[it]>n/2){
            return find_centroid(it,u,n);
        }
    }
    return u;
}
void decompose(int u,int p=-1){
    int n=dfs(u,p);
    int centroid=find_centroid(u,p,n);
    if(p!=-1) p=centroid;
    // Do stuff here for merges

    // Recurse
    par[centroid]=p;
    for(auto &it: g[centroid]){
        g[it].erase(centroid);
        decompose(it,centroid);
    }
    g[centroid].clear();
}
void reset(int n){
    par.resize(n);
    sub.resize(n);
}

```

```

g.assign(n,set<int>());
}

```

4 Convex Hull and Li Chao tree

```

// Li chao Tree (can be made persistent)
struct Line{
    ll m, c;
    Line(ll mm=0,ll cc=-3e18): m(mm),c(cc){}
    inline ll get(const int &x){return m*x+c;}
    inline ll operator [] (const int &x){return m*x+c;}
};
vector<Line> LN;
struct node{
    node *lt,*rt;
    int Ln;
    node(const int&l): Ln(l),lt(0),rt(0){};
    inline ll operator [] (const int &x){ return LN[Ln].get(x);}
    inline ll get(const int &x){return LN[Ln].get(x);}
};
const static int LX=-(1e9+1),RX=1e9+1;
struct Dynamic_Hull{ /* Max hull */
    node *root=0;
    void add(int l,node* &it,int lx=LX,int rx=RX){
        if(it==0) it=new node(l);
        if(it->get(lx)>LN[l].get(lx) and it->get(rx)>LN[l].get(rx)) return;
        if(it->get(lx)<=LN[l].get(lx) and it->get(rx)<=LN[l].get(rx)){
            it->Ln=l;
            return;
        }
        int mid=(lx+rx)>>1;
        if(it->get(lx)<LN[l].get(lx)) swap(it->Ln,l);
        if(it->get(mid)>=LN[l].get(mid)){
            add(l,it->rt,mid+1,rx);
        }else{
            swap(it->Ln,l);
            add(l,it->lt,lx,mid);
        }
    }
    inline void add(int ind){add(ind,root);}
    inline void add(int m,int c){LN.pb(Line(m,c));add(LN.size()-1,root);}
    ll get(int &x,node* &it,int lx=LX,int rx=RX){
        if(it==0) return -3e18; // Max hull
        ll ret=it->get(x);
        int mid=(lx+rx)>>1;
    }
}

```

```

    if(x<=mid) ret=max(ret,get(x,it->lt,lx,mid));
    else ret=max(ret,get(x,it->rt,mid+1,rx));
    return ret;
}
inline ll get(int x){return get(x,root);}
};

// const static int LX = -(1e9), RX = 1e9;
// struct Dynamic_Hull { /* Max hull */
//     struct Line{
//         ll m, c; // slope, intercept
//         Line(ll mm=0, ll cc=-1e18) { m = mm; c = cc; }
//         ll operator()(const int&x){ return m*x+c; }
//     };
//     struct node {
//         node *lt,*rt; Line Ln;
//         node(const Line &l){lt=rt=nullptr; Ln=l;}
//     };
//     node *root=nullptr;
//     void add(Line l,node*&it,int lx=LX,int rx=RX){
//         if(it==nullptr)it=new node(l);
//         if(it->Ln[lx]>=l[lx] and it->Ln[rx]>=l[rx]) return;
//         if(it->Ln[lx]<=l[lx] and it->Ln[rx]<=l[rx]) {it->Ln=l;
//             return;}
//         int mid = (lx+rx)>>1;
//         if(it->Ln[lx] < l[lx]) swap(it->Ln,l);
//         if(it->Ln[mid] >= l[mid]) add(l,it->rt,mid+1,rx);
//         else { swap(it->Ln,l); add(l,it->lt,lx,mid); }
//     }
//     void add(const ll &m,const ll &c) { add(Line(m,c),root);
//     }
//     ll get(int &x,node*&it,int lx=LX,int rx=RX){
//         if(it==NULL) return -1e18; // Max hull
//         ll ret = it->Ln[x];
//         int mid = (lx+rx)>>1;
//         if(x<=mid) ret = max(ret , get(x,it->lt,lx,mid));
//         else ret = max(ret , get(x,it->rt,mid+1,rx));
//         return ret;
//     }
//     ll get(int x){ return get(x,root); }
// };
struct Hull{
    struct line {
        ll m,c;
        ll eval(ll x){return m*x+c;}
        ld intersectX(line l){return (ld)(c-l.c)/(l.m-m);}
        line(ll m,ll c): m(m),c(c){}
    };
    deque<line> dq;

```

```

    v32 ints;
    Hull(int n){ints.clear(); forn(i,n) ints.pb(i); dq.clear();
    };
    // Dec order of slopes
    void add(line cur){
        while(dq.size()>2 && cur.intersectX(dq[0])>=dq[0].
            intersectX(dq[1]))
            dq.pop_front();
        dq.push_front(cur);
    }
    void add(const ll &m,const ll &c){add(line(m,c));}
    // query sorted dec.
    ll getval(ll x){
        while(dq.size()>2 && dq.back().eval(x)<=dq[dq.size()-2].eval(x))
            dq.pop_back();
        // return dq.back().eval(x);
    }
    // arbitrary query
    ll getval(ll x,deque<line> &dq){
        auto cmp = [&dq](int idx,ll x){return dq[idx].intersectX(
            dq[idx+1])<x;};
        int idx = *lower_bound(ints.begin(),ints.begin()+dq.size()
            ()-1,x,cmp);
        return dq[idx].eval(x);
    }
    ll get(const ll &x){return getval(x,dq);}
};

```

5 Dynamic Connectivity

```

int u[LIM],v[LIM],e[LIM],q[LIM];
map<p32,int> ids;
struct dsu{
    int sz;
    v32 par,rk;
    stack<int> st;
    void reset(int n){
        rk.assign(n,1);
        par.resize(n);
        iota(all(par),0);
        sz=n;
    }
    int getpar(int i){
        return (par[i]==i)? i:getpar(par[i]);
    }
    bool con(int i,int j){
        return getpar(i)==getpar(j);
    }
};

```

```

    }
    bool join(int i,int j){
        i=getpar(i),j=getpar(j);
        if(i==j) return 0;
        --sz;
        if(rk[j]>rk[i]) swap(i,j);
        par[j]=i,rk[i]+=rk[j];
        st.push(j);
        return 1;
    }
    int moment(){
        return st.size();
    }
    void revert(int tm){
        while(st.size()>tm){
            auto tp=st.top();
            rk[par[tp]]-=rk[tp];
            par[tp]=tp;
            st.pop();
            ++sz;
        }
    }
} d;
void solve(int l,int r,vp32 &ed){
    if(l>r) return;
    // dbg(ed,l,r,d.sz);
    int mid=(l+r)>>1;
    vp32 low;
    int tm=d.moment();
    forstl(it,ed){
        if(it.se<l or it.fi>r) continue;
        else if(it.fi<=l and it.se>=r) d.join(u[it.fi],v[it.fi]);
        else low.pb(it);
    }
    if(l==r){
        if(q[l]) cout<<(d.con(u[l],v[l])? "YES":"NO")<<ln;
    }else{
        solve(l,mid,low);
        solve(mid+1,r,low);
    }
    d.revert(tm);
}
signed main(){
    fastio;
    cin>>n>>k;
    d.reset(n);
    string t;
    forn(i,k){
        cin>>t;
        cin>>x>>y; --x,--y;
    }
};

```

```

if(x>y) swap(x,y);
u[i]=x,v[i]=y;
if(t[0]=='c'){
    q[i]=1;
}else{
    if(t[0]=='a'){
        ids[mp(x,y)]=i;
        e[i]=k-1;
    }else{
        e[ids[mp(x,y)]] = i;
        e[i]=-1;
    }
}
}
vp32 ed;
for(i,k) if(!q[i] && e[i]!=-1) ed.pb({i,e[i]});
solve(0,k-1,ed);
return 0;
}

```

6 Euler Path

procedure FindEulerPath(V)

1. iterate through all the edges outgoing from vertex V;
remove **this** edge from the graph,
and call FindEulerPath from the second end of **this** edge;
2. add vertex V to the answer.

7 Extended Euclidean GCD

```

int egcd(int a,int b, int* x, int* y){
    if(a==0){
        *x=0;*y=1;
        return b;
    }
    int x1,y1;
    int gcd=egcd(b%a,a,&x1,&y1);
    *x=y1-(b/a)*x1;
    *y=x1;
    return gcd;
}

```

8 Fenwick 2D

```

//BIT<N, M, K> b; N x M x K (3-dimensional) BIT
//b.update(x, y, z, P); // add P to (x,y,z)
//b.query(x1, x2, y1, y2, z1, z2); // query between (x1, y1,
//                                z1) and (x2, y2, z2)
inline int lastbit(int x){
    return x&(-x);
}
template <int N, int... Ns>
struct BIT<N, Ns...> {
    BIT<Ns...> bit[N + 1];
    template<typename... Args>
    void update(int pos, Args... args) {
        for (; pos <= N; bit[pos].update(args...), pos += lastbit
            (pos));
    }
    template<typename... Args>
    int query(int l, int r, Args... args) {
        int ans = 0;
        for (; r >= 1; ans += bit[r].query(args...), r -= lastbit
            (r));
        for (--l; l >= 1; ans -= bit[l].query(args...), l -=
            lastbit(l));
        return ans;
    }
};
// Another implementation
struct FenwickTree2D {
    vector<vector<int>>> bit;
    int n, m;
    // init(...) { ... }
    int sum(int x, int y) {
        int ret = 0;
        for (int i = x; i >= 0; i = (i & (i + 1)) - 1)
            for (int j = y; j >= 0; j = (j & (j + 1)) - 1)
                ret += bit[i][j];
        return ret;
    }
    void add(int x, int y, int delta) {
        for (int i = x; i < n; i = i | (i + 1))
            for (int j = y; j < m; j = j | (j + 1))
                bit[i][j] += delta;
    }
};

```

9 Gaussian Elimination, Base 2

```

struct Gaussbase2{
    int numofbits=20;
    int rk=0;
    v32 Base;
    Gaussbase2() {clear();}
    void clear(){
        rk=0;
        Base.assign(numofbits,0);
    }
    Gaussbase2& operator = (Gaussbase2 &g){
        for(i,numofbits) Base[i]=g.Base[i];
        rk=g.rk;
    }
    bool canbemade(int x){
        rform(i,numofbits-1) x=min(x,x^Base[i]);
        return x==0;
    }
    void Add(int x){
        rform(i,numofbits-1){
            if((x>>i)&1){
                if(!Base[i]){
                    Base[i]=x;
                    rk++;
                    return;
                }else x^=Base[i];
            }
        }
    }
    int maxxor(){
        int ans=0;
        rform(i,numofbits-1){
            if(ans < (ans^Base[i])) ans^=Base[i];
        }
        return ans;
    }
};

```

10 Gaussian Elimination

```

int gauss (vector <vector<double> > a, vector<double> &ans){
    int n = (int) a.size();
    int m = (int) a[0].size()-1;

    vector<int> where(m,-1);
    for(int col=0, row=0;col<m && row<n; ++col){
        int sel = row;
        for(int i=row;i<n;++i){

```

```

        if(abs(a[i][col]) > abs(a[sel][col])){
            sel = i;
        }
    }
    if(abs(a[sel][col])<EPS) continue;
    for(int i=col; i<=m; ++i){
        swap(a[sel][i],a[row][i]);
    }
    where[col] = row;

    for(int i=0;i<n;++i){
        if(i!=row){
            double c = a[i][col]/a[row][col];
            for(int j=col;j<=m;++j){
                a[i][j] -= a[row][j]*c;
            }
        }
    }
    ++row;
}
ans.assign(m,0);
for(int i=0;i<m;++i){
    if(where[i]!=-1){
        ans[i] = a[where[i]][m]/a[where[i]][i];
    }
}
for(int i=0;i<n;++i){
    double sum=0;
    for(int j=0;j<m;++j){
        sum+=ans[j]*a[i][j];
    }
    if(abs(sum-a[i][m])>EPS)
        return 0;
}
for(int i=0;i<m;++i){
    if(where[i]==-1) return MOD;
}
return 1;
}

```

11 General Weighted Matching

```

struct MaxMatchingEdmonds{
    // Assume General Unweighted Directed Graph
    //  $O(V^3)$  edmonds for maximum matching
    vv32 g;
    vv32 match,p,base;
    vector<bool> blossom;

```

```

    int n;
    int lca(int a,int b){
        vector<bool> used(match.size(),0);
        while(1){
            a=base[a];
            used[a]=1;
            if(match[a]==-1) break;
            a=p[match[a]];
        }
        while(1){
            b=base[b];
            if(used[b]) return b;
            b=p[match[b]];
        }
    }
    void markPath(int v,int b,int children) {
        for(;base[v]!=b;v=p[match[v]]){
            blossom[base[v]]=blossom[base[match[v]]]=1;
            p[v]=children;
            children=match[v];
        }
    }
    int findPath(int root) {
        vector<bool> used(n,0);
        p.assign(n,-1);
        base.assign(n,0);
        for(int i=0;i<n;++i) base[i] = i;
        used[root]=1;
        int qh=0;
        int qt=0;
        vv32 q(n,0);
        q[qt++]=root;
        while(qh<qt){
            int v=q[qh++];
            for(int &to:g[v]){
                if(base[v]==base[to] || match[v]==to) continue;
                if(to==root || match[to]!=-1 && p[match[to]]!=-1)
                    {
                        int curbase=lca(v,to);
                        blossom.assign(n,0);
                        markPath(v,curbase,to);
                        markPath(to,curbase,v);
                        for(int i=0;i<n;++i)
                            if(blossom[base[i]]){
                                base[i]=curbase;
                                if(!used[i]){
                                    used[i]=1;
                                    q[qt++]=i;
                                }
                            }
                    }
            }
        }
    }

```

```

        }else if(p[to]==-1){
            p[to]=v;
            if(match[to]==-1) return to;
            to=match[to];
            used[to]=1;
            q[qt++]=to;
        }
    }
    return -1;
}
int maxMatching(vv32 &graph){
    n=graph.size();
    g=graph;
    match.assign(n,-1);
    p.assign(n,0);
    for(int i=0;i<n;++i){
        if(match[i]==-1){
            int v=findPath(i);
            while(v!=-1){
                int pv=p[v];
                int ppv=match[pv];
                match[v]=pv;
                match[pv]=v;
                v=ppv;
            }
        }
    }
    int matches=0;
    for(int i=0;i<n;++i) if(match[i]!=-1) ++matches;
    return matches/2;
}
};

```

12 Geometry

```

const int MAX_SIZE = 1000;
const double PI = 2.0*acos(0.0);
struct PT
{
    double x,y;
    double length() {return sqrt(x*x+y*y);}
    int normalize(){
        // normalize the vector to unit length; return -1 if the
        // vector is 0
        double l = length();
        if(fabs(l)<EPS) return -1;
    }

```

```

    x/=1; y/=1;
    return 0;
}
PT operator-(PT a){
    PT r;
    r.x=x-a.x; r.y=y-a.y;
    return r;
}
PT operator+(PT a){
    PT r;
    r.x=x+a.x; r.y=y+a.y;
    return r;
}
PT operator*(double sc){
    PT r;
    r.x=x*sc; r.y=y*sc;
    return r;
}
};

bool operator<(const PT& a,const PT& b){
    if(fabs(a.x-b.x)<EPS) return a.y<b.y;
    return a.x<b.x;
}
double dist(PT& a, PT& b){
    // the distance between two points
    return sqrt((a.x-b.x)*(a.x-b.x) + (a.y-b.y)*(a.y-b.y));
}
double dot(PT& a, PT& b){
    // the inner product of two vectors
    return(a.x*b.x+a.y*b.y);
}
double cross(PT& a, PT& b){
    return(a.x*b.y-a.y*b.x);
}

// =====
// The Convex Hull
// =====

int sideSign(PT& p1,PT& p2,PT& p3){
    // which side is p3 to the line p1->p2? returns: 1 left, 0
    // on, -1 right
    double sg = (p1.x-p3.x)*(p2.y-p3.y)-(p1.y - p3.y)*(p2.x-p3.
        x);
    if(fabs(sg)<EPS) return 0;
    if(sg>0) return 1;
    return -1;
}

```

```

bool better(PT& p1,PT& p2,PT& p3){
    // used by convex hull: from p3, if p1 is better than p2
    double sg = (p1.y - p3.y)*(p2.x-p3.x)-(p1.x-p3.x)*(p2.y-p3.
        y);
    //watch range of the numbers
    if(fabs(sg)<EPS){
        if(dist(p3,p1)>dist(p3,p2))return true;
        else return false;
    }
    if(sg<0) return true;
    return false;
}

void vex2(vector<PT> vin,vector<PT>& vout){
    // vin is not pass by reference, since we will rotate it
    vout.clear();
    int n=vin.size();
    sort(vin.begin(),vin.end());
    PT stk[MAX_SIZE];
    int pstk, i;
    // hopefully more than 2 points
    stk[0] = vin[0];
    stk[1] = vin[1];
    pstk = 2;
    for(i=2; i<n; i++){
        if(dist(vin[i], vin[i-1])<EPS) continue;
        while(pstk > 1 && better(vin[i], stk[pstk-1], stk[pstk-2]))
            pstk--;
        stk[pstk] = vin[i];
        pstk++;
    }
    for(i=0; i<pstk; i++) vout.push_back(stk[i]);
    // turn 180 degree
    for(i=0; i<n; i++){
        vin[i].y = -vin[i].y;
        vin[i].x = -vin[i].x;
    }
    sort(vin.begin(), vin.end());
    stk[0] = vin[0];
    stk[1] = vin[1];
    pstk = 2;
    for(i=2; i<n; i++){
        if(dist(vin[i], vin[i-1])<EPS) continue;
        while(pstk > 1 && better(vin[i], stk[pstk-1], stk[pstk-2]))
            pstk--;
        stk[pstk] = vin[i];
        pstk++;
    }
}

```

```

for(i=1; i<pstk-1; i++){
    stk[i].x= -stk[i].x; // dont forget rotate 180 d back.
    stk[i].y= -stk[i].y;
    vout.push_back(stk[i]);
}
}

int isConvex(vector<PT>& v){
    // test whether a simple polygon is convex
    // return 0 if not convex, 1 if strictly convex,
    // 2 if convex but there are points unnecessary
    // this function does not work if the polygon is self
    // intersecting
    // in that case, compute the convex hull of v, and see if
    // both have the same area
    int i,j,k;
    int c1=0; int c2=0; int c0=0;
    int n=v.size();
    for(i=0;i<n;i++){
        j=(i+1)%n;
        k=(j+1)%n;
        int s=sideSign(v[i], v[j], v[k]);
        if(s==0) c0++;
        if(s>0) c1++;
        if(s<0) c2++;
    }
    if(c1 && c2) return 0;
    if(c0) return 2;
    return 1;
}

// =====
// Areas
// =====
double trap(PT a, PT b){
    // Used in various area functions
    return (0.5*(b.x - a.x)*(b.y + a.y));
}

double area(vector<PT> &vin){
    // Area of a simple polygon, not neccessary convex
    int n = vin.size();
    double ret = 0.0;
    for(int i = 0; i < n; i++) ret += trap(vin[i], vin[(i+1)%n
        ]);
    return fabs(ret);
}

double peri(vector<PT> &vin){
    // Perimeter of a simple polygon, not neccessary convex
    int n = vin.size();
    double ret = 0.0;

```

```

for(int i = 0; i < n; i++) ret += dist(vin[i], vin[(i+1)%n
]);
return ret;
}

double triarea(PT a, PT b, PT c){
return fabs(trap(a,b)+trap(b,c)+trap(c,a));
}

double height(PT a, PT b, PT c){
// height from a to the line bc
double s3 = dist(c, b);
double ar=triarea(a,b,c);
return(2.0*ar/s3);
}

// =====
// Points and Lines
// =====
int intersection( PT p1, PT p2, PT p3, PT p4, PT &r ) {
// two lines given by p1->p2, p3->p4 r is the intersection
point
// return -1 if two lines are parallel
double d = (p4.y - p3.y)*(p2.x-p1.x) - (p4.x - p3.x)*(p2.y
- p1.y);
if( fabs( d ) < EPS ) return -1;
// might need to do something special!!!
double ua, ub;
ua = (p4.x - p3.x)*(p1.y-p3.y) - (p4.y-p3.y)*(p1.x-p3.x);
ua /= d;
// ub = (p2.x - p1.x)*(p1.y-p3.y) - (p2.y-p1.y)*(p1.x-p3.x)
;
//ub /= d;
r = p1 + (p2-p1)*ua;
return 0;
}

void closestpt( PT p1, PT p2, PT p3, PT &r ){
// the closest point on the line p1->p2 to p3
if( fabs( triarea( p1, p2, p3 ) ) < EPS ) { r = p3; return;
}
PT v = p2-p1;
v.normalize();
double pr; // inner product
pr = (p3.y-p1.y)*v.y + (p3.x-p1.x)*v.x;
r = p1+v*pr;
}

int hcenter( PT p1, PT p2, PT p3, PT& r ){
// point generated by altitudes

```

```

if( triarea( p1, p2, p3 ) < EPS ) return -1;
PT a1, a2;
closestpt( p2, p3, p1, a1 );
closestpt( p1, p3, p2, a2 );
intersection( p1, a1, p2, a2, r );
return 0;
}

int center( PT p1, PT p2, PT p3, PT& r ){
// point generated by circumscribed circle
if( triarea( p1, p2, p3 ) < EPS ) return -1;
PT a1, a2, b1, b2;
a1 = (p2+p3)*0.5;
a2 = (p1+p3)*0.5;
b1.x = a1.x - (p3.y-p2.y);
b1.y = a1.y + (p3.x-p2.x);
b2.x = a2.x - (p3.y-p1.y);
b2.y = a2.y + (p3.x-p1.x);
intersection( a1, b1, a2, b2, r );
return 0;
}

int bcenter( PT p1, PT p2, PT p3, PT& r ){
// angle bisection
if( triarea( p1, p2, p3 ) < EPS ) return -1;
double s1, s2, s3;
s1 = dist( p2, p3 );
s2 = dist( p1, p3 );
s3 = dist( p1, p2 );
double rt = s2/(s2+s3);
PT a1,a2;
a1 = p2*rt+p3*(1.0-rt);
rt = s1/(s1+s3);
a2 = p1*rt+p3*(1.0-rt);
intersection( a1,p1, a2,p2, r );
return 0;
}

// =====
// Angles
// =====
double angle(PT& p1, PT& p2, PT& p3){
// angle from p1->p2 to p1->p3, returns -PI to PI
PT va = p2-p1;
va.normalize();
PT vb; vb.x=-va.y; vb.y=va.x;
PT v = p3-p1;
double x,y;
x=dot(v, va);
y=dot(v, vb);
return(atan2(y,x));
}

```

```

}

double angle(double a, double b, double c){
// in a triangle with sides a,b,c, the angle between b and
c
// we do not check if a,b,c is a triangle here
double cs=(b*b+c*c-a*a)/(2.0*b*c);
return(acos(cs));
}

void rotate(PT p0, PT p1, double a, PT& r){
// rotate p1 around p0 clockwise, by angle a
// dont pass by reference for p1, so r and p1 can be the
same
p1 = p1-p0;
r.x = cos(a)*p1.x-sin(a)*p1.y;
r.y = sin(a)*p1.x+cos(a)*p1.y;
r = r+p0;
}

void reflect(PT& p1, PT& p2, PT p3, PT& r){
// p1->p2 line, reflect p3 to get r.
if(dist(p1, p3)<EPS) {r=p3; return;}
double a=angle(p1, p2, p3);
r=p3;
rotate(p1, r, -2.0*a, r);
}

// =====
// points, lines, and circles
// =====

int pAndSeg(PT& p1, PT& p2, PT& p){
// the relation of the point p and the segment p1->p2.
// 1 if point is on the segment; 0 if not on the line; -1
if on the line but not on the segment
double s=triarea(p, p1, p2);
if(s>EPS) return(0);
double sg=(p.x-p1.x)*(p.x-p2.x);
if(sg>EPS) return(-1);
sg=(p.y-p1.y)*(p.y-p2.y);
if(sg>EPS) return(-1);
return(1);
}

int lineAndCircle(PT& oo, double r, PT& p1, PT& p2, PT& r1,
PT& r2){
// returns -1 if there is no intersection
// returns 1 if there is only one intersection
PT m;

```



```

closestpt(p1,p2,oo,m);
PT v = p2-p1;
v.normalize();
double r0=dist(oo, m);
if(r0>r+EPS) return -1;
if(fabs(r0-r)<EPS){
    r1=r2=m;
    return 1;
}
double dd = sqrt(r*r-r0*r0);
r1 = m-v*dd; r2 = m+v*dd;
return 0;
}

int CAndC(PT o1, double r1, PT o2, double r2, PT &q1, PT& q2)
){

    // intersection of two circles
    // -1 if no intersection or infinite intersection
    // 1 if only one point

    double r=dist(o1,o2);
    if(r1<r2) { swap(o1,o2); swap(r1,r2); }
    if(r<EPS) return(-1);
    if(r>r1+r2+EPS) return(-1);
    if(r<r1-r2-EPS) return(-1);
    PT v = o2-o1; v.normalize();
    q1 = o1+v*r1;
    if(fabs(r-r1-r2)<EPS || fabs(r+r2-r1)<EPS)
    { q2=q1; return(1); }
    double a=angle(r2, r, r1);
    q2=q1;
    rotate(o1, q1, a, q1);
    rotate(o1, q2, -a, q2);
    return 0;
}

int pAndPoly(vector<PT> pv, PT p){
    // the relation of the point and the simple polygon
    // 1 if p is in pv; 0 outside; -1 on the polygon
    int i, j;
    int n=pv.size();
    pv.push_back(pv[0]);
    for(i=0;i<n;i++) if(pAndSeg(pv[i], pv[i+1], p)==1) return
        (-1);
    for(i=0;i<n;i++) pv[i] = pv[i]-p;
    p.x=p.y=0.0;
    double a, y;
    while(1){
        a=(double)rand()/10000.00;

```

```

j=0;
for(i=0;i<n;i++){
    rotate(p, pv[i], a, pv[i]);
    if(fabs(pv[i].x)<EPS) j=1;
}
if(j==0){
    pv[n]=pv[0];
    j=0;
    for(i=0;i<n;i++) if(pv[i].x*pv[i+1].x < -EPS){
        y=pv[i+1].y-pv[i+1].x*(pv[i].y-pv[i+1].y)/(pv[i].x-pv[i+1].x);
        if(y>0) j++;
    }
    return(j%2);
}
return 1;
}

```

13 Giant Step Baby Step

```

// Giant Step - Baby Step for discrete log
// find x with a^x = b mod MOD
// Find one soln can be changed to find all
// 0(root(MOD)*log(MOD)) can be reduced with unordered map
// or array
ll solve(ll a,ll b,ll MOD){
    int n=(int)sqrt(MOD+.0)+1;
    ll an=1,cur;
    forn(i,n) an=(an*a)%MOD;
    cur=an;
    vector<pair<ll,int> > vals;
    forsn(i,1,n+1){
        vals.pb(mp(cur,i));
        cur=(cur*an)%MOD;
    }
    cur=b;
    sort(all(vals));
    forn(i,n+1){
        auto in=lower_bound(all(vals),mp(cur,-1))-vals.begin();
        if(in!=vals.size() && vals[in].fi==cur){
            ll ans=n*(ll)vals[in].se-i;
            if(ans<MOD) return ans;
        }
        cur=(cur*a)%MOD;
    }
    return -1;
}

```

14 Hashtable

```

struct hashtable{
    v64 hash1,hash2,inv1,inv2;
    ll MOD1=MOD,MOD2=MOD+2;
    ll pr1=31,pr2=37;
    void create(string &p){
        int len=p.size();
        hash1.resize(len);hash2.resize(len);
        inv1.resize(len);inv2.resize(len);
        ll p1=1,p2=1;
        int i=0;
        while(p[i]){
            hash1[i]= (i==0)? 0:hash1[i-1];
            hash2[i]= (i==0)? 0:hash2[i-1];
            hash1[i]= (hash1[i]+p[i]*p1)%MOD1;
            hash2[i]= (hash2[i]+p[i]*p2)%MOD2;
            p1=p1*pr1%MOD1;
            p2=p2*pr2%MOD2;
            i++;
        }
        ll iv1=inv(pr1,MOD1),iv2=inv(pr2,MOD2);
        inv1[0]=1,inv2[0]=1;
        forsn(i,1,len){
            inv1[i]=inv1[i-1]*iv1%MOD1;
            inv2[i]=inv2[i-1]*iv2%MOD2;
        }
    }
    p64 gethash(int l,int r){
        ll ans1=hash1[r-1];
        if(l!=0) ans1+=MOD1-hash1[l-1];
        ll ans2=hash2[r-1];
        if(l!=0) ans2+=MOD2-hash2[l-1];
        ans1=ans1*inv1[l]%MOD1;
        ans2=ans2*inv2[l]%MOD2;
        return mp(ans1,ans2);
    }
};

```

15 Heavy Light Decomposition

```

struct SegTree{
    v32 T,lazy;
    int N,MX;
    void clear(int n,int mx){
        N=n,MX=mx;
        T.assign(4*N,0);
    }
};

```

```

    lazy.assign(4*N,0);
}
void build(int a[],int v,int tl,int tr){
    if(tl==tr){
        T[v]=a[tl];
    }else{
        int tm=(tl+tr)>>1,lf=v<<1,rt=lf^1;;
        build(a,lf,tl,tm);
        build(a,rt,tm+1,tr);
        T[v]=min(T[lf],T[rt]);
    }
}
void push(int v){
    int lf=v<<1,rt=lf^1;
    T[lf]=(T[lf]+lazy[v]);
    lazy[lf]=(lazy[lf]+lazy[v]);
    T[rt]=(T[rt]+lazy[v]);
    lazy[rt]=(lazy[rt]+lazy[v]);
    lazy[v]=0;
}
void update(int v,int tl,int tr,int l,int r,int val){
    if(l>r or tl>r or tr<l) return;
    if(l<=tl && tr<=r){
        T[v]=T[v]+val;
        lazy[v]=(lazy[v]+val);
    }else{
        if(tl==tr) return;
        push(v);
        int tm=(tl+tr)>>1,lf=v<<1,rt=lf^1;;
        update(lf,tl,tm,l,r,val);
        update(rt,tm+1,tr,l,r,val);
        T[v]=max(T[lf],T[rt]);
    }
}
int query(int v,int tl,int tr,int l,int r){
    if(l>r) return MX;
    if(l<=tl && tr<=r) return T[v];
    push(v);
    int tm=(tl+tr)>>1,lf=v<<1,rt=lf^1;
    return max(query(lf,tl,tm,l,min(r,tm)),query(rt,tm+1,tr,
        max(l,tm+1),r));
}
int q(int l,int r){
    return query(1,0,N-1,l,r);
}
void u(int l,int r,int val){
    update(1,0,N-1,l,r,val);
}
} st;
struct hld{

```

```

    int n,t;
    v32 sz,in,out,root,par,depth;
    vv32 g;
    SegTree tree;
    void dfs_sz(int v=0,int p=0){
        sz[v]=1;
        for(auto &u: g[v]){
            if(u==p) continue;
            dfs_sz(u,v);
            sz[v]+=sz[u];
            if(sz[u]>sz[g[v][0]]) swap(u, g[v][0]);
        }
    }
    void dfs_hld(int v=0,int p=0){
        in[v]=t++;
        par[v]=p;
        depth[v]=depth[p]+1;
        for(auto u: g[v]){
            if(u==p) continue;
            root[u]= (u==g[v][0] ? root[v]:u);
            dfs_hld(u,v);
        }
        out[v]=t;
    }
}
void pre(vv32 &v){
    g=v;n=v.size();t=0;
    sz.assign(n,0);in.assign(n,0);out.assign(n,0);
    root.assign(n,0);par.assign(n,0);depth.assign(n,0);
    depth[0]=-1;
    dfs_sz();dfs_hld();
    tree.clear(n,-MOD);
}
template <class BinaryOperation>
void processPath(int u,int v,BinaryOperation op){
    for(;root[u]!=root[v];v=par[root[v]]){
        if(depth[root[u]] > depth[root[v]]) swap(u,v);
        op(in[root[v]],in[v]);
    }
    if(depth[u]>depth[v]) swap(u,v);
    op(in[u],in[v]);
}
void modifyPath(int u,int v,const int &value){
    processPath(u,v,[this,&value](int l,int r){tree.u(l,r,
        value);}); // [l,r]
}
void modifySubtree(int u,const int &value){
    tree.u(in[u],out[u]-1,value);
}
int queryPath(int u,int v){
    int res=-MOD;

```

```

        auto add=[](int &a,const int &b){a=max(a,b);};
        processPath(u,v,[this,&res,&add](int l,int r){add(res,
            tree.q(l,r));});
        return res;
    }
    int querySubtree(int u){
        return tree.q(in[u],out[u]-1);
    }
};

```

16 Hopcraft Karp

```

// Max matching
//1 indexed Hopcroft-Karp Matching in O(E sqrtV)
struct Hopcroft_Karp{
    static const int inf = 1e9;
    int n;
    vector<int> matchL, matchR, dist;
    vector<vector<int>> > g;
    Hopcroft_Karp(int n):n(n),matchL(n+1),matchR(n+1),dist(n+1),
        g(n+1){}
    void addEdge(int u, int v){
        g[u].pb(v);
    }
    bool bfs(){
        queue<int> q;
        for(int u=1;u<=n;u++){
            if(!matchL[u]){
                dist[u]=0;
                q.push(u);
            }else dist[u]=inf;
        }
        dist[0]=inf;
        while(!q.empty()){
            int u=q.front();
            q.pop();
            for(auto v:g[u]){
                if(dist[matchR[v]] == inf){
                    dist[matchR[v]] = dist[u] + 1;
                    q.push(matchR[v]);
                }
            }
        }
        return (dist[0]!=inf);
    }
    bool dfs(int u){
        if(!u) return true;
        for(auto v:g[u]){

```

```

    if(dist[matchR[v]] == dist[u]+1 &&dfs(matchR[v])){
        matchL[u]=v;
        matchR[v]=u;
        return true;
    }
}
dist[u]=inf;
return false;
}
int max_matching(){
    int matching=0;
    while(bfs()){
        for(int u=1;u<=n;u++){
            if(!matchL[u])
                if(dfs(u)) matching++;
        }
    }
    return matching;
}
};

```

17 Hungarian Algorithm

```

struct Hungarian{
    //Important: cost matrix a[1..n][1..m]>=0,n<=m (works with
        negative costs)
    // 0(V^3) Use p to find matching of 1..m
    vv64 a;
    v64 u,v;
    v32 p,way;
    int n,m;
    Hungarian(int n,int m): n(n),m(m),u(n+1,0),v(m+1,0),p(m
        +1,0),way(m+1,0),a(n+1,v64(m+1,0)){}
    void addEdge(int u,int v,ll val){
        a[u][v]=val;
    }
    ll solveAssignmentProblem(){
        for(int i=1;i<=n;++i){
            p[0]=i;
            int j0=0;
            v64 minv(m+1,2e17+10);
            vector<bool> used(m+1,0);
            do{
                used[j0] = true;
                int i0=p[j0];
                ll delta=2e17+10;
                int j1=0;
                for(int j=1;j<=m;++j){

```

```

                    if(!used[j]){
                        ll cur=a[i0][j]-u[i0]-v[j];
                        if(cur<minv[j]){
                            minv[j]=cur;
                            way[j]=j0;
                        }
                    }
                    if(minv[j]<delta){
                        delta=minv[j];
                        j1=j;
                    }
                }
            }
            for(int j=0;j<=m;++j){
                if(used[j]){
                    u[p[j]]+=delta;
                    v[j]-=delta;
                }else minv[j]-=delta;
            }
            j0=j1;
        }while(p[j0]!=0);
        do{
            int j1=way[j0];
            p[j0]=p[j1];
            j0=j1;
        }while(j0!=0);
    }
    return -v[0];
}
};

```

18 Linear Sieve

```

int mu[LIM],is_com[LIM];
v32 pr;
void sieve(){
    mu[1]=1;
    forsn(i,2,LIM){
        if(!is_com[i]) pr.pb(i),mu[i]=-1;
        forstl(it,pr){
            if(it*i>=LIM) break;
            is_com[i*it]=1;
            if(i%it==0){
                mu[i*it]=0;
                break;
            }else{
                mu[i*it]=mu[i]*mu[it];
            }
        }
    }
}

```

```

    }
}

```

19 Lowest Common Ancestor

```

vv32 v;
v32 tin,tout,dist;
vv32 up;
int l;
void dfs(int i,int par,int lvl){
    tin[i]= ++t;
    dist[i]= lvl;
    up[i][0] = par;
    forsn(j,1,l+1) up[i][j]= up[up[i][j-1]][j-1];
    forstl(it,v[i]) if(it!=par) dfs(it,i,lvl+1);
    tout[i] = ++t;
}
bool is_ancetor(int u, int v){
    return tin[u]<=tin[v] && tout[u]>=tout[v];
}
int lca(int u, int v){
    if (is_ancetor(u, v)) return u;
    if (is_ancetor(v, u)) return v;
    rfor(n,i,1) if(!is_ancetor(up[u][i], v)) u=up[u][i];
    return up[u][0];
}
int get_dis(int u,int v){
    int lcauv=lca(u,v);
    return dist[u]+dist[v]-2*dist[lcauv];
}
void preprocess(int root){
    tin.resize(n);
    tout.resize(n);
    dist.resize(n);
    t=0;
    l=ceil(log2((double)n));
    up.assign(n,v32(l+1));
    dfs(root,root,0);
}

```

20 Lucas Theorem

```

//Lucas Theorem: Find (n Choose m) mod p for prime p and
    large n,m. in O(log(m*n))
// nCm mod p by lucas theorem for large n,m >=0

```

```
// p prime, require fact(factorial) & invfact(inverse
    factorial)
v32 fact,invfact;
ll lucas(ll n,ll m,int p){
    ll res=1;
    while(n || m) {
        ll a=n%p,b=m%p;
        if(a<b) return 0;
        res=((res*fact[a]%p)*(invfact[b]%p)*(invfact[a-b]%p)%p;
        n/=p; m/=p;
    }
    return res;
}
```

21 Manacher

Manacher

```
// Given a string s of length N, finds all palindromes as
    its substrings.
// p[0][i] = half length of longest even palindrome around
    pos i
// p[1][i] = longest odd at i (half rounded down i.e len 2*x
    +1).
//Time: O(N)
void manacher(const string& s){
    int n=s.size();
    v32 p[2]={v32(n+1),v32(n)};
    forn(z,2) for(int i=0,l=0,r=0;i<n;++i){
        int t=r-i+!z;
        if(i<r) p[z][i]=min(t,p[z][l+t]);
        int L=i-p[z][i],R=i+p[z][i]-!z;
        while(L>=1 && R+1<n && s[L]==s[R+1]) p[z][i]++,L--,R++;
        if(R>r) l=L,r=R;}}}
```

22 Matrix

```
int MOD1=MOD;
inline ll add(ll a,ll b){
    return (a+b)%MOD1;
}
inline ll mult(ll a,ll b){
    return a*b%MOD1;
}
struct matrix{
    int arr[105][105]={0};
```

```
int SZ;
void reset(int sz){
    SZ=sz;
    //memset(arr,0,sizeof(arr));
}
void makeiden(int sz){
    reset(sz);
    for(int i=0;i<SZ;i++){
        arr[i][i]=1;
    }
}
matrix operator +(const matrix &o)const{
    matrix res;
    res.reset(SZ);
    for(int i=0;i<SZ;i++){
        for(int j=0;j<SZ;j++){
            res.arr[i][j]=add(arr[i][j],o.arr[i][j]);
        }
    }
    return res;
}
matrix operator *(const matrix &o)const{
    matrix res;
    res.reset(SZ);
    for(int i=0;i<SZ;i++){
        for(int j=0;j<SZ;j++){
            res.arr[i][j]=0;
            for(int k=0;k<SZ;k++){
                res.arr[i][j]=add(res.arr[i][j],mult(arr[i][k],o.arr[k][j]));
            }
        }
    }
    return res;
}
matrix mpower(matrix a,int sz,ll b){
    matrix res;
    res.makeiden(sz);
    while(b){
        if(b&1){
            res=res*a;
        }
        a=a*a;
        b>>=1;
    }
    return res;
}
```

23 Merge Sort Tree

```
// Merge sort Tree
const int MAXN=1e5+5;
v32 T[4*MAXN]; // nlogn memory
void build(int a[],int v,int tl,int tr){
    if(tl==tr){
        T[v]=v32(1,a[tl]);
    }else{
        int tm=(tl+tr)>>1;
        build(a,v<<1,tl,tm);
        build(a,(v<<1)^1,tm+1,tr);
        merge(all(T[v<<1]),all(T[(v<<1)^1]),back_inserter(T[v]));
        // built in combine in sorted order (2pointer)
    }
}
// number of numbers <=x in [l,r]
int query(int v,int tl,int tr,int l,int r,int x){
    if(l>r) return 0;
    if(l<=tl && tr<=r){
        return upper_bound(all(T[v]),x)-T[v].begin();
    }
    int tm=(tl+tr)>>1;
    return query(v<<1,tl,tm,l,min(r,tm),x)+query((v<<1)^1,tm+1,
        tr,max(l,tm+1),r,x);
}
// Number of distinct integers in [l,r]
int b[MAXN];
void convert(int a[],int n){ // b store next occ index
    m32 m; // Can be replaced by vv32 in small numbers
    rforn(i,n-1){
        auto it=m.find(a[i]);
        if(it==m.end()) b[i]=MOD;
        else b[i]=it->se;
        m[a[i]]=i;
    }
    build(b,1,0,n-1);
}
inline int q(int l,int r){ // no. of val in [l,r] with nxt
    ind > r
    return (r-l+1)-query(1,0,n-1,l,r,r);
}
```

24 Miller Rabin

```
using u64 = uint64_t;
using u128 = __uint128_t;
```

```

u64 binpower(u64 base, u64 e, u64 mod) {
    u64 result = 1;
    base %= mod;
    while (e) {
        if (e & 1)
            result = (u128)result * base % mod;
        base = (u128)base * base % mod;
        e >>= 1;
    }
    return result;
}

bool check_composite(u64 n, u64 a, u64 d, int s) {
    u64 x = binpower(a, d, n);
    if (x == 1 || x == n - 1)
        return false;
    for (int r = 1; r < s; r++) {
        x = (u128)x * x % n;
        if (x == n - 1)
            return false;
    }
    return true;
};

bool MillerRabin(u64 n) { // returns true if n is prime,
    else returns false.
    if (n < 2)
        return false;

    int r = 0;
    u64 d = n - 1;
    while ((d & 1) == 0) {
        d >>= 1;
        r++;
    }

    for (int a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}) {
        if (n == a)
            return true;
        if (check_composite(n, a, d, r))
            return false;
    }
    return true;
}

```

25 Min Cost Max Flow

```

// Mincost Maxflow : O(E^2)
// [Hell-Johnson MinCostMaxFlow using Dijkstra with potential
//   & Fibonnaci Heap]
// Negative cost cycles are not supported.
struct MCMF{
    struct Edge{
        int u,v,rind;
        FLOW cap,flow;
        COST cost;
    };
    int N;
    vector<COST> pot,dist;
    vector<vector<Edge>> > v;
    vector<pair<int,int>> > par;
    MCMF(int n): N(n),dist(n),v(n),par(n){}
    void AddEdge(int to,int from,int cap,int cost){
        if(to==from){
            assert(cost>=0);
            return;
        }
        int i1=v[to].size(),i2=v[from].size();
        v[to].push_back({to,from,i2,cap,0,cost});
        v[from].push_back({from,to,i1,0,0,-cost});
    }
    void setpi(int s){
        pot.assign(N,CINF);
        pot[s]=0;
        int ch=1,ite=N;
        COST cur,nw;
        while(ch-- && ite--){
            for(int i=0;i<N;++i){
                if(pot[i]!=CINF){
                    cur=pot[i];
                    for(auto &e: v[i]){
                        if(e.cap>0 && (nw=cur+e.cost)<pot[e.v]){
                            pot[e.v]=nw; ch=1;
                        }
                    }
                }
            }
        }
        assert(ite>=0); // Else negative cycle
    }
    bool path(int s,int t){
        fill(dist.begin(),dist.end(),CINF);
        dist[s]=0;
        __gnu_pbds::priority_queue<pair<COST,int>> > pq;
        vector<decltype(pq)::point_iterator> its(N);

```

```

        pq.push({0,s});
        COST curr,val;
        int node,cnt;
        bool ok=0;
        while(!pq.empty()){
            tie(curr,node)=pq.top();
            pq.pop();
            curr=-curr;
            if(curr!=dist[node]) continue;
            curr+=pot[node];
            if(node==t) ok=1;
            cnt=0;
            for(auto &e: v[node]){
                if(e.cap>e.flow && (val=curr+e.cost-pot[e.v])<dist[e.v]){
                    dist[e.v]=val;
                    par[e.v]=make_pair(node,cnt);
                    if(its[e.v]==pq.end()) its[e.v]=pq.push({-val,e.v});
                }
                else pq.modify(its[e.v],{-val,e.v});
            }
            ++cnt;
        }
    }
    for(int i=0;i<N;++i){
        pot[i]=min(pot[i]+dist[i],FINF);
    }
    return ok;
}

pair<FLOW,COST> SolveMCMF(int s,int t,FLOW need=FINF,bool
    neg=0){
    FLOW tot=0,cflow=0; COST tcost=0;
    if(s==t) return {tot,tcost};
    if(!neg) pot.assign(N,0);
    else setpi(s);
    int cntr=0;
    while(path(s,t) && need>0){
        cflow=need;
        for(int node=t,u,ind;node!=s;node=u){
            u=par[node].first;
            ind=par[node].second;
            cflow=min(cflow,v[u][ind].cap-v[u][ind].flow);
        }
        tot+=cflow; need-=cflow;
        for(int node=t,u,ind,rind;node!=s;node=u){
            u=par[node].first;
            ind=par[node].second;
            rind=v[u][ind].rind;
            v[u][ind].flow+=cflow;
            v[node][rind].flow-=cflow;

```

```

    }
}
return {tot,tcost};
}
};

```

26 Nearest Pair of Points

```

vector<pt> t;

void rec(int l, int r) {
    if (r - l <= 3) {
        for (int i = l; i < r; ++i) {
            for (int j = i + 1; j < r; ++j) {
                upd_ans(a[i], a[j]);
            }
        }
        sort(a.begin() + l, a.begin() + r, cmp_y());
        return;
    }

    int m = (l + r) >> 1;
    int midx = a[m].x;
    rec(l, m);
    rec(m, r);

    merge(a.begin() + l, a.begin() + m, a.begin() + m, a.
        begin() + r, t.begin(), cmp_y());
    copy(t.begin(), t.begin() + r - l, a.begin() + l);

    int tsz = 0;
    for (int i = l; i < r; ++i) {
        if (abs(a[i].x - midx) < mindist) {
            for (int j = tsz - 1; j >= 0 && a[i].y - t[j].y <
                mindist; --j)
                upd_ans(a[i], t[j]);
            t[tsz++] = a[i];
        }
    }

    // In main, call as:

    t.resize(n);
    sort(a.begin(), a.end(), cmp_x());
    mindist = 1E20;
    rec(0, n);

```

27 Number Theoretic Transform

```

const int mod=998244353;
// 998244353=1+7*17*2^23 : g=3
// 1004535809=1+479*2^21 : g=3
// 469762049=1+7*2^26 : g=3
// 7340033=1+7*2^20 : g=3
// For below change mult as overflow:
// 10000093151233=1+3^3*5519*2^26 : g=5
// 1000000523862017=1+10853*1373*2^26 : g=3
// 100000000949747713=1+2^29*3*73*8505229 : g=2
// For rest find primitive root using Shoup's generator
// algorithm
// root_pw: power of 2 >= maxn, Mod-1=k*root_pw => w =
// primitive^k
template<long long Mod,long long root_pw,long long primitive
>
struct NTT{
    inline long long powm(long long x,long long pw){
        x%=Mod;
        if(abs(pw)>Mod-1) pw%=(Mod-1);
        if(pw<0) pw+=Mod-1;
        ll res=1;
        while(pw){
            if(pw&1LL) res=(res*x)%Mod;
            pw>>=1;
            x=(x*x)%Mod;
        }
        return res;
    }
    inline ll inv(ll x){
        return powm(x,Mod-2);
    }
    ll root,root_1;
    NTT(){
        root=powm(primitive,(Mod-1)/root_pw);
        root_1=inv(root);
    }
    void ntt(vector<long long> &a,bool invert){
        int n=a.size();
        for(long long i=1,j=0;i<n;i++){
            long long bit=n>>1;
            for(;j&bit;bit>>=1) j^=bit;
            j^=bit;
            if(i<j) swap(a[i],a[j]);
        }
        for(long long len=2;len<=n;len<=1){
            long long wlen= invert ? root_1:root;
            for(long long i=len;i<root_pw;i<=1) wlen=wlen*wlen%Mod;
            for(long long i=0;i<n;i+=len){

```

```

                long long w=1;
                for(long long j=0;j<len/2;j++){
                    long long u=a[i+j],v=a[i+j+len/2]*w%Mod;
                    a[i+j]= u+v<Mod ? u+v:u+v-Mod;
                    a[i+j+len/2]= u-v>=0 ? u-v:u-v+Mod;
                    w=w*wlen%Mod;
                }
            }
        }
        if(invert){
            ll n_1=inv(n);
            for(long long &x: a) x=x*n_1%Mod;
        }
    }
    vector<long long> multiply(vector<long long> const& a,
        vector<ll> const& b){
        vector<long long> fa(a.begin(),a.end()),fb(b.begin(),b.end
            ());
        int n=1;
        while(n<a.size()+b.size()) n<=1;
        point(fa,1,n);
        point(fb,1,n);
        for(int i=0;i<n;++i) fa[i]=fa[i]*fb[i]%Mod;
        coef(fa);
        return fa;
    }
    void point(vector<long long> &A,bool not_pow=1,int atleast
        ==-1){
        if(not_pow){
            if(atleast===-1){
                atleast=1;
                while(atleast<A.size()) atleast<=1;
            }
            A.resize(atleast,0);
        }
        ntt(A,0);
    }
    void coef(vector<long long> &A,bool reduce=1){
        ntt(A,1);
        if(reduce) while(A.size() and A.back()==0) A.pop_back();
    }
    void point_power(vector<long long> &A,long long k){
        for(long long &x: A) x=powm(x,k);
    }
    void coef_power(vector<long long> &A,int k){
        while(A.size() and A.back()==0) A.pop_back();
        int n=1;
        while(n<k*A.size()) n<=1;
        point(A,1,n);
        point_power(A,k);
    }

```

```

    coef(A);
}
vector<long long> power(vector<long long> a, ll p){
    while(a.size() and a.back()==0) a.pop_back();
    vector<long long> res;
    res.pb(1);
    while(p){
        if(p&1) res=multiply(res,a);
        a=multiply(a,a);
        p/=2;
    }
    return res;
}
};
NTT<mod,1<<20,3> ntt;

```

28 Ordered Set

```

// Set/Map using Leftist Trees
// * To get a map, change {null_type to some value}.
#include <bits/extc++.h> /** keep-include */
using namespace __gnu_pbds;
template<class T>
using Tree = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
void example() {
    Tree<int> t, t2; t.insert(8);
    auto it = t.insert(10).first;
    assert(it == t.lower_bound(9));
    assert(t.order_of_key(10) == 1);
    assert(t.order_of_key(11) == 2);
    assert(*t.find_by_order(0) == 8);
    t.join(t2); // assuming T < T2 or T > T2, merge t2 into t
}

```

29 Polynomial

```

namespace algebra{
    const int inf = 1e9;
    const int magic = 500; // threshold for sizes to run the
        naive algo
    namespace fft {
        const int maxn = 1 << 18;
        typedef double ftype;
        typedef complex<ftype> point;

```

```

        point w[maxn];
        const ftype pi = acos(-1);
        bool initiated = 0;
        void init() {
            if(!initiated) {
                for(int i = 1; i < maxn; i *= 2) {
                    for(int j = 0; j < i; j++) {
                        w[i + j] = polar(ftype(1), pi * j / i);
                    }
                }
                initiated = 1;
            }
        }
        template<typename T>
        void fft(T *in, point *out, int n, int k = 1) {
            if(n == 1) {
                *out = *in;
            } else {
                n /= 2;
                fft(in, out, n, 2 * k);
                fft(in + k, out + n, n, 2 * k);
                for(int i = 0; i < n; i++) {
                    auto t = out[i + n] * w[i + n];
                    out[i + n] = out[i] - t;
                    out[i] += t;
                }
            }
        }
        template<typename T>
        void mul_slow(vector<T> &a, const vector<T> &b) {
            vector<T> res(a.size() + b.size() - 1);
            for(size_t i = 0; i < a.size(); i++) {
                for(size_t j = 0; j < b.size(); j++) {
                    res[i + j] += a[i] * b[j];
                }
            }
            a = res;
        }
        template<typename T>
        void mul(vector<T> &a, const vector<T> &b) {
            if(min(a.size(), b.size()) < magic) {
                mul_slow(a, b);
                return;
            }
            init();
            static const int shift = 15, mask = (1 << shift) - 1;
            size_t n = a.size() + b.size() - 1;
            while(__builtin_popcount(n) != 1) n++;
            a.resize(n);
            static point A[maxn], B[maxn];
            static point C[maxn], D[maxn];
            for(size_t i = 0; i < n; i++) {
                A[i] = point(a[i] & mask, a[i] >> shift);
                if(i < b.size()) B[i] = point(b[i] & mask, b[i] >> shift);
                ;
                else B[i] = 0;
            }
        }

```

```

        fft(A, C, n); fft(B, D, n);
        for(size_t i = 0; i < n; i++) {
            point c0 = C[i] + conj(C[(n - i) % n]);
            point c1 = C[i] - conj(C[(n - i) % n]);
            point d0 = D[i] + conj(D[(n - i) % n]);
            point d1 = D[i] - conj(D[(n - i) % n]);
            A[i] = c0 * d0 - point(0, 1) * c1 * d1;
            B[i] = c0 * d1 + d0 * c1;
        }
        fft(A, C, n); fft(B, D, n);
        reverse(C + 1, C + n);
        reverse(D + 1, D + n);
        int t = 4 * n;
        for(size_t i = 0; i < n; i++) {
            int64_t A0 = llround(real(C[i]) / t);
            T A1 = llround(imag(D[i]) / t);
            T A2 = llround(imag(C[i]) / t);
            a[i] = A0 + (A1 << shift) + (A2 << 2 * shift);
        }
        return;
    }
    template<typename T>
    T bpow(T x, size_t n) {
        return n ? n % 2 ? x * bpow(x, n - 1) : bpow(x * x, n / 2)
            : T(1);
    }
    template<typename T>
    T bpow(T x, size_t n, T m) {
        return n ? n % 2 ? x * bpow(x, n - 1, m) % m : bpow(x * x
            % m, n / 2, m) : T(1);
    }
    template<int m>
    struct modular{
        int64_t r;
        modular() : r(0) {}
        modular(int64_t rr) : r(rr) {if(abs(r) >= m) r %= m; if(r
            < 0) r += m;}
        modular inv() const {return bpow(*this, m - 2);}
        modular pow(int64_t k) const {return bpow(*this, k);}
        modular operator * (const modular &t) const {return (r * t
            .r) % m;}
        modular operator / (const modular &t) const {return *this
            * t.inv();}
        modular operator += (const modular &t) {r += t.r; if(r >=
            m) r -= m; return *this;}
        modular operator -= (const modular &t) {r -= t.r; if(r <
            0) r += m; return *this;}
        modular operator + (const modular &t) const {return
            modular(*this) += t;}
        modular operator - (const modular &t) const {return
            modular(*this) -= t;}
    }

```



```

modular operator *= (const modular &t) {return *this = *
    this * t;}
modular operator /= (const modular &t) {return *this = *
    this / t;}
bool operator == (const modular &t) const {return r == t.r
    ;}
bool operator != (const modular &t) const {return r != t.r
    ;}
operator int64_t() const {return r;}
};
template<int T>
istream& operator >> (istream &in, modular<T> &x) {
    return in >> x.r;
}
template<typename T>
struct poly {
    vector<T> a;
    void normalize() { // get rid of leading zeroes
        while(!a.empty() && a.back() == T(0)) a.pop_back();}
    poly(){}
    poly(T a0) : a{a0}{normalize();}
    poly(vector<T> t) : a(t){normalize();}
    poly operator += (const poly &t) {
        a.resize(max(a.size(), t.a.size()));
        for(size_t i = 0; i < t.a.size(); i++) {
            a[i] += t.a[i];
        }
        normalize();
        return *this;
    }
    poly operator -= (const poly &t) {
        a.resize(max(a.size(), t.a.size()));
        for(size_t i = 0; i < t.a.size(); i++) {
            a[i] -= t.a[i];
        }
        normalize();
        return *this;
    }
    poly operator + (const poly &t) const {return poly(*this)
        += t;}
    poly operator - (const poly &t) const {return poly(*this)
        -= t;}
    poly mod_xk(size_t k) const { // get same polynomial mod x
        ^k
        k = min(k, a.size());
        return vector<T>(begin(a), begin(a) + k);
    }
    poly mul_xk(size_t k) const { // multiply by x^k
        poly res(*this);
        res.a.insert(begin(res.a), k, 0);

```

```

        return res;
    }
    poly div_xk(size_t k) const { // divide by x^k, dropping
        coefficients
        k = min(k, a.size());
        return vector<T>(begin(a) + k, end(a));
    }
    poly substr(size_t l, size_t r) const { // return mod_xk(r
        ).div_xk(l)
        l = min(l, a.size());
        r = min(r, a.size());
        return vector<T>(begin(a) + l, begin(a) + r);
    }
    poly inv(size_t n) const { // get inverse series mod x^n
        poly ans = a[0].inv();
        size_t a = 1;
        while(a < n) {
            poly C = (ans * mod_xk(2 * a)).substr(a, 2 * a);
            ans -= (ans * C).mod_xk(a).mul_xk(a);
            a *= 2;
        }
        return ans.mod_xk(n);
    }
    poly operator *= (const poly &t) {fft::mul(a, t.a);
        normalize(); return *this;}
    poly operator * (const poly &t) const {return poly(*this)
        *= t;}
    poly reverse(size_t n, bool rev = 0) const { // reverses
        and leaves only n terms
        poly res(*this);
        if(rev) { // If rev = 1 then tail goes to head
            res.a.resize(max(n, res.a.size()));
        }
        std::reverse(res.a.begin(), res.a.end());
        return res.mod_xk(n);
    }
    pair<poly, poly> divmod_slow(const poly &b) const { //
        when divisor or quotient is small
        vector<T> A(a);
        vector<T> res;
        while(A.size() >= b.a.size()) {
            res.push_back(A.back() / b.a.back());
            if(res.back() != T(0)) {
                for(size_t i = 0; i < b.a.size(); i++) {
                    A[A.size() - i - 1] -= res.back() * b.a[b.a.size() - i
                        - 1];
                }
            }
            A.pop_back();
        }

```

```

        std::reverse(begin(res), end(res));
        return {res, A};
    }
    pair<poly, poly> divmod(const poly &b) const { // returns
        quotient and remainder of a mod b
        if(a.size() < b.a.size()) {
            return {poly{0}, *this};
        }
        int d = a.size() - b.a.size();
        if(min(d, b.a.size()) < magic) {
            return divmod_slow(b);
        }
        poly D = (reverse(d + 1) * b.reverse(d + 1).inv(d + 1)).
            mod_xk(d + 1).reverse(d + 1, 1);
        return {D, *this - D * b};
    }
    poly operator / (const poly &t) const {return divmod(t).
        first;}
    poly operator % (const poly &t) const {return divmod(t).
        second;}
    poly operator /= (const poly &t) {return *this = divmod(t)
        .first;}
    poly operator %= (const poly &t) {return *this = divmod(t)
        .second;}
    poly operator *= (const T &x) {
        for(auto &it: a) {
            it *= x;
        }
        normalize();
        return *this;
    }
    poly operator /= (const T &x) {
        for(auto &it: a) {
            it /= x;
        }
        normalize();
        return *this;
    }
    poly operator * (const T &x) const {return poly(*this) *=
        x;}
    poly operator / (const T &x) const {return poly(*this) /=
        x;}
    T operator [](int idx) const {
        return idx >= (int)a.size() || idx < 0 ? T(0) : a[idx];
    }
    T& coef(size_t idx) { // mutable reference at coefficient
        return a[idx];
    }
    bool operator == (const poly &t) const {return a == t.a;}
    bool operator != (const poly &t) const {return a != t.a;}

```



```

poly deriv() { // calculate derivative
    vector<T> res;
    for(int i = 1; i <= a.size(); i++) {
        res.push_back(T(i) * a[i]);
    }
    return res;
}

poly integr() { // calculate integral with C = 0
    vector<T> res = {0};
    for(int i = 0; i <= a.size(); i++) {
        res.push_back(a[i] / T(i + 1));
    }
    return res;
}

size_t leading_xk() const { // Let  $p(x) = x^k * t(x)$ ,
    return k;
    int res = 0;
    while(a[res] == T(0)) {
        res++;
    }
    return res;
}

poly log(size_t n) { // calculate  $\log p(x) \bmod x^n$ 
    assert(a[0] == T(1));
    return (deriv().mod_xk(n) * inv(n)).integr().mod_xk(n);
}

poly exp(size_t n) { // calculate  $\exp p(x) \bmod x^n$ 
    if(a.empty()) return T(1);
    assert(a[0] == T(0));
    poly ans = T(1);
    size_t a = 1;
    while(a < n) {
        poly C = ans.log(2 * a).div_xk(a) - substr(a, 2 * a);
        ans -= (ans * C).mod_xk(a).mul_xk(a);
        a *= 2;
    }
    return ans.mod_xk(n);
}

poly pow_slow(size_t k, size_t n) { // if k is small
    return k % 2 ? (*this * pow_slow(k - 1, n)).mod_xk(n)
        : (*this * *this).mod_xk(n).pow_slow(k / 2, n) : T(1);
}

poly pow(size_t k, size_t n) { // calculate  $p^k(x) \bmod x^n$ 
    if(a.empty()) return *this;
    if(k < magic) return pow_slow(k, n);
    int i = leading_xk();
    T j = a[i];
    poly t = div_xk(i) / j;

```

```

        return bpow(j, k) * (t.log(n) * T(k)).exp(n).mul_xk(i * k)
            .mod_xk(n);
    }
};

template<typename T>
poly<T> operator * (const T& a, const poly<T>& b) {
    return b * a;
}

using namespace algebra;
const int mod = 998244353;
const int lim = 2e6+5;
typedef modular<mod> base;
typedef poly<base> polyn;
base fact[lim], invfact[lim];
void pre(){
    fact[0]=1;
    invfact[lim-1]=367642781;
    for(int i=1; i<lim; ++i) fact[i]=fact[i-1]*base(i);
    for(int i=lim-1; i>0; --i) invfact[i-1]=invfact[i]*base(i);
}

void example(){
    int t,n,k;
    cin>>t;
    while(t--){
        cin>>n>>k;
        if(k==1){
            base nk=1,N=n,Ans;
            for(int i=n-2; i>=0; --i){
                base ans=nk*invfact[i]*invfact[n-i]*invfact[n-2-i];
                Ans+=ans;
                if(i) nk*=N;
            }
            Ans=Ans*fact[n-2]*fact[n]/nk;
            cout<<Ans<<endl;
        }else{
            vector<base> a(n+1);
            for(int i=0; i<=n; ++i){
                a[i]=base(i+1).pow(k)*invfact[i];
            }
            polyn A(a), B(A.pow(n,n));
            base Ans=(B[n-2]/base(n).pow(n-2))*fact[n-2];
            cout<<Ans<<endl;
        }
    }
}

```

30 Primitive Root

```

// Primitive root Exist for  $n=1,2,4,(\text{odd prime power}), 2*(\text{odd prime power})$ 
//  $0(\text{Ans.log}(p).\text{logp} + \text{sqrt}(\text{phi})) \leq 0((\text{log } p)^8 + \text{root}(p))$ 
// Change phi when not prime
// Include powm (inverse)
ll phi_cal(ll n){
    ll result=n;
    for(ll i=2; i*i<=n; ++i){
        if(n%i==0){
            while(n%i==0) n/=i;
            result-=result/i;
        }
    }
    if(n>1) result-=result/n;
    return result;
}

ll generator(ll p){
    v64 fact;
    ll phi=p-1; // Call phi_cal if not prime
    ll n=phi;
    for(ll i=2; i*i<=n; ++i){
        if(n%i==0){
            fact.push_back(i);
            while(n%i==0) n/=i;
        }
    }
    if(n>1) fact.push_back(n);
    for(ll res=2; res<=p; ++res){
        bool ok=true;
        for(size_t i=0; i<fact.size() && ok; ++i)
            ok&=(powm(res, phi/fact[i], p)!=1);
        if(ok) return res;
    }
    return -1;
}

```

31 Push Relabel

```

//Push-Relabel Algorithm for Flows - Gap Heuristic,
//Complexity:  $O(V^3)$ 
//To obtain the actual flow values, look at all edges with
//capacity > 0
//Zero capacity edges are residual edges
struct edge{
    int from, to, cap, flow, index;
}

```

```

edge(int from, int to, int cap, int flow, int index):
    from(from), to(to), cap(cap), flow(flow), index(index) {}
};
struct PushRelabel{
    int n;
    vector<vector<edge>> > g;
    vector<long long> excess;
    vector<int> height, active, count;
    queue<int> Q;
    PushRelabel(int n): n(n), g(n), excess(n), height(n), active(n)
        , count(2*n) {}
    void addEdge(int from, int to, int cap){
        g[from].push_back(edge(from, to, cap, 0, g[to].size()));
        if(from==to) g[from].back().index++;
        g[to].push_back(edge(to, from, 0, 0, g[from].size()-1));
    }
    void enqueue(int v){
        if(!active[v] && excess[v]>0){
            active[v]=true;
            Q.push(v);
        }
    }
    void push(edge &e){
        int amt=(int)min(excess[e.from], (long long)e.cap - e.flow)
            ;
        if(height[e.from]<=height[e.to] || amt==0) return;
        e.flow += amt;
        g[e.to][e.index].flow -= amt;
        excess[e.to] += amt;
        excess[e.from] -= amt;
        enqueue(e.to);
    }
    void relabel(int v){
        count[height[v]]--;
        int d=2*n;
        for(auto &it:g[v]){
            if(it.cap-it.flow>0) d=min(d, height[it.to]+1);
        }
        height[v]=d;
        count[height[v]]++;
        enqueue(v);
    }
    void gap(int k){
        for(int v=0; v<n; v++){
            if(height[v]<k) continue;
            count[height[v]]--;
            height[v]=max(height[v], n+1);
            count[height[v]]++;
            enqueue(v);
        }
    }
};

```

```

}
void discharge(int v){
    for(int i=0; excess[v]>0 && i<g[v].size(); i++) push(g[v][i]);
    if(excess[v]>0){
        if(count[height[v]]==1) gap(height[v]);
        else relabel(v);
    }
}
long long max_flow(int source, int dest){
    count[0] = n-1;
    count[n] = 1;
    height[source] = n;
    active[source] = active[dest] = 1;
    for(auto &it:g[source]){
        excess[source]+=it.cap;
        push(it);
    }
    while(!Q.empty()){
        int v=Q.front();
        Q.pop();
        active[v]=false;
        discharge(v);
    }
    long long max_flow=0;
    for(auto &e:g[source]) max_flow+=e.flow;
    return max_flow;
}
};

```

32 Simplex

```

// Two-phase simplex algorithm for solving linear programs
// of the form
//
//      maximize    c^T x
//      subject to  Ax <= b
//                  x >= 0
//
// INPUT: A -- an m x n matrix
//         b -- an m-dimensional vector
//         c -- an n-dimensional vector
//         x -- a vector where the optimal solution will be
//              stored
//
// OUTPUT: value of the optimal solution (infinity if
//         unbounded
//         above, nan if infeasible)

```

```

//
// To use this code, create an LPSolver object with A, b,
// and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>

using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;

struct LPSolver {
    int m, n;
    VI B, N;
    VVD D;

    LPSolver(const VVD &A, const VD &b, const VD &c) :
        m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n +
            2)) {
        for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D
            [i][j] = A[i][j];
        for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1;
            D[i][n + 1] = b[i]; }
        for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j];
            }
        N[n] = -1; D[m + 1][n] = 1;
    }

    void Pivot(int r, int s) {
        for (int i = 0; i < m + 2; i++) if (i != r)
            for (int j = 0; j < n + 2; j++) if (j != s)
                D[i][j] -= D[r][j] * D[i][s] / D[r][s];
        for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] /= D
            [r][s];
        for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] /= -D
            [r][s];
        D[r][s] = 1.0 / D[r][s];
        swap(B[r], N[s]);
    }

    bool Simplex(int phase) {
        int x = phase == 1 ? m + 1 : m;
        while (true) {

```

```

int s = -1;
for (int j = 0; j <= n; j++) {
    if (phase == 2 && N[j] == -1) continue;
    if (s == -1 || D[x][j] < D[x][s] || D[x][j] == D[x][s]
        && N[j] < N[s]) s = j;
}
if (D[x][s] > -EPS) return true;
int r = -1;
for (int i = 0; i < m; i++) {
    if (D[i][s] < EPS) continue;
    if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] /
        D[r][s] ||
        (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s])
        && B[i] < B[r]) r = i;
}
if (r == -1) return false;
Pivot(r, s);
}
}

DOUBLE Solve(VD &x) {
    int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
    if (D[r][n + 1] < -EPS) {
        Pivot(r, n);
        if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -
            numeric_limits<DOUBLE>::infinity();
        for (int i = 0; i < m; i++) if (B[i] == -1) {
            int s = -1;
            for (int j = 0; j <= n; j++)
                if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i][s]
                    && N[j] < N[s]) s = j;
            Pivot(i, s);
        }
    }
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return D[m][n + 1];
}

int main() {
    const int m = 4;
    const int n = 3;
    DOUBLE _A[m][n] = {

```

```

        { 6, -1, 0 },
        { -1, -5, 0 },
        { 1, 5, 1 },
        { -1, -5, -1 }
    };
    DOUBLE _b[m] = { 10, -4, 5, -5 };
    DOUBLE _c[n] = { 1, -1, 0 };

    VVD A(m);
    VD b(_b, _b + m);
    VD c(_c, _c + n);
    for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);

    LPSolver solver(A, b, c);
    VD x;
    DOUBLE value = solver.Solve(x);

    cerr << "VALUE: " << value << endl; // VALUE: 1.29032
    cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
    for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
    cerr << endl;
    return 0;
}

```

33 Stock Span

```

void stockspan(v32 &d, int n) {
    // smallest index j such that j>i and d[j]>d[i]
    nxt.assign(n, n);
    // largest index j such that j<i and d[j]>=d[i]
    pre.assign(n, -1);
    int stk[n+10], ptr=0;
    forn(i, n) {
        while(ptr && d[i]>d[stk[ptr-1]]) nxt[stk[ptr-1]]=i;
        pre[i] = (ptr? stk[ptr-1]:-1);
        stk[ptr++]=i;
    }
    v32 nxt, pr; // nxt index with val < cur and prev index with
        val < cur
    void stockspan(v32 &d, int n) {
        int stk[n+10], ptr=0;
        pr.assign(n, -1);
        forn(i, n) {
            while(ptr && d[i]<=d[stk[ptr-1]]) --ptr;
            pr[i] = (ptr? stk[ptr-1]:-1);
            stk[ptr++]=i;
        }
    }
}

```

```

ptr=0;
nxt.assign(n, n);
forn(i, n) {
    while(ptr && d[i]<d[stk[ptr-1]]) nxt[stk[ptr-1]]=i;
    stk[ptr++]=i;
}
}

```

34 Suffix Array

```

vector<int> sort_cyclic_shifts(string const& s) {
    int n = s.size();
    const int alphabet = 256;
    vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
    for (int i = 0; i < n; i++)
        cnt[s[i]]++;
    for (int i = 1; i < alphabet; i++)
        cnt[i] += cnt[i-1];
    for (int i = 0; i < n; i++)
        p[--cnt[s[i]]] = i;
    c[p[0]] = 0;
    int classes = 1;
    for (int i = 1; i < n; i++) {
        if (s[p[i]] != s[p[i-1]])
            classes++;
        c[p[i]] = classes - 1;
    }
    vector<int> pn(n), cn(n);
    for (int h = 0; (1 << h) < n; ++h) {
        for (int i = 0; i < n; i++) {
            pn[i] = p[i] - (1 << h);
            if (pn[i] < 0)
                pn[i] += n;
        }
        fill(cnt.begin(), cnt.begin() + classes, 0);
        for (int i = 0; i < n; i++)
            cnt[c[pn[i]]]++;
        for (int i = 1; i < classes; i++)
            cnt[i] += cnt[i-1];
        for (int i = n-1; i >= 0; i--)
            p[--cnt[c[pn[i]]]] = pn[i];
        cn[p[0]] = 0;
        classes = 1;
        for (int i = 1; i < n; i++) {
            pair<int, int> cur = {c[p[i]], c[(p[i] + (1 << h)) % n]};
            pair<int, int> prev = {c[p[i-1]], c[(p[i-1] + (1 << h)) % n]};

```

```

        if (cur != prev)
            ++classes;
        cn[p[i]] = classes - 1;
    }
    c.swap(cn);
}
return p;
}
vector<int> suffix_array_construction(string s) {
    s += "$";
    vector<int> sorted_shifts = sort_cyclic_shifts(s);
    sorted_shifts.erase(sorted_shifts.begin());
    return sorted_shifts;
}

```

35 Suffix Automaton

```

struct state {
    int len, link;
    map<char, int> next;
};

const int MAXLEN = 100000;
state st[MAXLEN * 2];
int sz, last;

void sa_init() {
    st[0].len = 0;
    st[0].link = -1;
    sz++;
    last = 0;
}

void sa_extend(char c) {
    int cur = sz++;
    st[cur].len = st[last].len + 1;
    int p = last;
    while (p != -1 && !st[p].next.count(c)) {
        st[p].next[c] = cur;
        p = st[p].link;
    }
    if (p == -1) {
        st[cur].link = 0;
    } else {
        int q = st[p].next[c];
        if (st[p].len + 1 == st[q].len) {
            st[cur].link = q;
        } else {

```

```

            int clone = sz++;
            st[clone].len = st[p].len + 1;
            st[clone].next = st[q].next;
            st[clone].link = st[q].link;
            while (p != -1 && st[p].next[c] == q) {
                st[p].next[c] = clone;
                p = st[p].link;
            }
            st[q].link = st[cur].link = clone;
        }
    }
    last = cur;
}

```

36 Suffix Tree

```

string s;
int n;

struct node {
    int l, r, par, link;
    map<char, int> next;

    node (int l=0, int r=0, int par=-1)
        : l(l), r(r), par(par), link(-1) {}
    int len() { return r - l; }
    int &get (char c) {
        if (!next.count(c)) next[c] = -1;
        return next[c];
    }
};

node t[MAXN];
int sz;

struct state {
    int v, pos;
    state (int v, int pos) : v(v), pos(pos) {}
};

state ptr (0, 0);

state go (state st, int l, int r) {
    while (l < r)
        if (st.pos == t[st.v].len()) {
            st = state (t[st.v].get( s[l] ), 0);
            if (st.v == -1) return st;
        }
        else {
            if (s[ t[st.v].l + st.pos ] != s[l])

```

```

            return state (-1, -1);
        if (r-l < t[st.v].len() - st.pos)
            return state (st.v, st.pos + r-l);
        l += t[st.v].len() - st.pos;
        st.pos = t[st.v].len();
    }
    return st;
}

int split (state st) {
    if (st.pos == t[st.v].len())
        return st.v;
    if (st.pos == 0)
        return t[st.v].par;
    node v = t[st.v];
    int id = sz++;
    t[id] = node (v.l, v.l+st.pos, v.par);
    t[v.par].get( s[v.l] ) = id;
    t[id].get( s[v.l+st.pos] ) = st.v;
    t[st.v].par = id;
    t[st.v].l += st.pos;
    return id;
}

int get_link (int v) {
    if (t[v].link != -1) return t[v].link;
    if (t[v].par == -1) return 0;
    int to = get_link (t[v].par);
    return t[v].link = split (go (state(to, t[to].len()), t[v]
        ].l + (t[v].par==0), t[v].r));
}

void tree_extend (int pos) {
    for(;;) {
        state nptr = go (ptr, pos, pos+1);
        if (nptr.v != -1) {
            ptr = nptr;
            return;
        }

        int mid = split (ptr);
        int leaf = sz++;
        t[leaf] = node (pos, n, mid);
        t[mid].get( s[pos] ) = leaf;

        ptr.v = get_link (mid);
        ptr.pos = t[ptr.v].len();
        if (!mid) break;
    }
}

```

```
void build_tree() {
    sz = 1;
    for (int i=0; i<n; ++i)
        tree_extend(i);
}
```

37 Template

```
#pragma GCC optimize ("-O2")
#pragma GCC optimize("Ofast")
// ~ #pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,abm
,mmx,avx,tune=native")
// ~ #pragma GCC optimize("unroll-loops")
#include <bits/stdc++.h>
using namespace std;
#define fastio ios_base::sync_with_stdio(0);cin.tie(0);cout.
tie(0)
#define pb push_back
#define mp make_pair
#define fi first
#define se second
#define all(x) x.begin(),x.end()
#define memreset(a) memset(a,0,sizeof(a))
#define testcase(t) int t;cin>>t;while(t-->0)
#define forstl(i,v) for(auto &i: v)
#define forn(i,e) for(int i=0;i<e;++i)
#define forsn(i,s,e) for(int i=s;i<e;++i)
#define rforn(i,s) for(int i=s;i>0;--i)
#define rforfn(i,s,e) for(int i=s;i>e;--i)
#define bitcount(a) __builtin_popcount(a) // set bits (add
ll)
#define ln '\n'
#define getcurrtime() cerr<<"Time = "<<((double)clock()/
CLOCKS_PER_SEC)<<endl
#define dbgarr(v,s,e) cerr<<#v<<" = "; forsn(i,s,e) cerr<<v[
i]<<" "; cerr<<endl
#define inputfile freopen("input.txt", "r", stdin)
#define outputfile freopen("output.txt", "w", stdout)
#define dbg(args...) { string _s = #args; replace(_s.begin()
, _s.end(), ',', ' '); \
stringstream _ss(_s); istream_iterator<string> _it(_ss); err
(_it, args); }
void err(istream_iterator<string> it) { cerr<<endl; }
template<typename T, typename... Args>
void err(istream_iterator<string> it, T a, Args... args) {
    cerr << *it << " = " << a << "\t"; err(++it, args...);
}
```

```
template<typename T1,typename T2>
ostream& operator <<(ostream& c,pair<T1,T2> &v){
    c<<"("<<v.fi<<" "<<v.se<<" "; return c;
}
template <template <class...> class TT, class ...T>
ostream& operator<<(ostream& out,TT<T...>& c){
    out<<"{" <<" ";
    forstl(x,c) out<<x<<" ";
    out<<"}"; return out;
}
typedef long long ll;
typedef unsigned long long ull;
typedef long double ld;
typedef pair<ll,ll> p64;
typedef pair<int,int> p32;
typedef pair<int,p32> p96;
typedef vector<ll> v64;
typedef vector<int> v32;
typedef vector<v32> vv32;
typedef vector<v64> vv64;
typedef vector<p32> vp32;
typedef vector<p64> vp64;
typedef vector<vp32> vvp32;
typedef map<int,int> m32;
const int LIM=1e5+5,MOD=1e9+7;
const ld EPS = 1e-9;
mt19937 rng(chrono::steady_clock::now().time_since_epoch().
count());
```

38 Topological Sort

```
int n; // number of vertices
vector<int> adj[LIM]; // adjacency list of graph
vector<bool> visited;
vector<int> ans;

void dfs(int v) {
    visited[v] = true;
    for (int u : adj[v]) {
        if (!visited[u])
            dfs(u);
    }
    ans.push_back(v);
}

void topological_sort() {
    visited.assign(n, false);
    ans.clear();
}
```

```
for (int i = 0; i < n; ++i) {
    if (!visited[i])
        dfs(i);
}
reverse(ans.begin(), ans.end());
}
```

39 Treap

```
mt19937 rng(chrono::steady_clock::now().time_since_epoch().
count());

int getRand(int l, int r){
    uniform_int_distribution<int> uid(l, r);
    return uid(rng);
}

struct Treap{
    struct data{
        int priority, val, cnt;
        data *l, *r;

        data(){
            val = 0, cnt = 0, l = NULL, r = NULL;
        }
        data (int _val){
            val = _val, cnt = 1;
            l = NULL, r = NULL;
            priority = getRand(1, 2e9);
        }
    };
    typedef struct data* node;
    node head;

    Treap(): head(0) {}

    int cnt(node cur){
        return cur ? cur->cnt : 0;
    }

    void updateCnt(node cur)
    {
        if(cur)
            cur->cnt = cnt(cur->l) + cnt(cur->r) + 1;
    }

    void push(node cur) //Lazy Propagation
    {
}
```

```

    };
}

void combine(node &cur, node l, node r)
{
    if(!l)
    {
        cur = r;
        return;
    }
    if(!r)
    {
        cur = l;
        return;
    }
    //Merge Operations like in segment tree
}

void reset(node &cur) //To reset other fields of cur
                        except value and cnt
{
    if(!cur)
        return;
}

void operation(node &cur)
{
    if(!cur)
        return;
    reset(cur);
    combine(cur, cur->l, cur);
    combine(cur, cur, cur->r);
}

void splitPos(node cur, node &l, node &r, int k, int add
              = 0)
{
    if(!cur)
        return void(l = r = 0);
    push(cur);
    int idx = add + cnt(cur->l);
    if(idx <= k)
        splitPos(cur->r, cur->r, r, k, idx + 1), l = cur;
    else
        splitPos(cur->l, l, cur->l, k, add), r = cur;
    updateCnt(cur);
    operation(cur);
}

void split(node cur, node &l, node &r, int k)

```

```

{
    if(!cur)
        return void(l = r = 0);
    push(cur);
    int idx = cur->val;
    if(idx <= k)
        split(cur->r, cur->r, r, k), l = cur;
    else
        split(cur->l, l, cur->l, k), r = cur;
    updateCnt(cur);
    operation(cur);
}

void merge(node &cur, node l, node r)
{
    push(l);
    push(r);
    if(!l || !r)
        cur = l ? l : r;
    else if(l->priority > r->priority)
        merge(l->r, l->r, r), cur = l;
    else
        merge(r->l, l, r->l), cur = r;
    updateCnt(cur);
    operation(cur);
}

void insert(int val)
{
    if(!head)
    {
        head = new data(val);
        return;
    }
    node l, r, mid, mid2, rr;
    mid = new data(val);
    split(head, l, r, val - 1);
    merge(l, l, mid);
    split(r, mid2, rr, val);
    merge(head, l, rr);
}

void erase(int val)
{
    node l, r, mid;
    split(head, l, r, val - 1);
    split(r, mid, r, val);
    merge(head, l, r);
}

```

```

void inorder(node cur)
{
    if(!cur)
        return;
    push(cur);
    inorder(cur->l);
    cout<<cur->val<<" ";
    inorder(cur->r);
}

void inorder()
{
    inorder(head);
    cout<<endl;
}

void clear(node cur)
{
    if(!cur)
        return;
    clear(cur->l), clear(cur->r);
    delete cur;
}

void clear()
{
    clear(head);
}

int find_by_order(int x) //1 indexed
{
    if(!x)
        return -1;
    x--;
    node l, r, mid;
    splitPos(head, l, r, x - 1);
    splitPos(r, mid, r, 0);
    int ans = -1;
    if(cnt(mid) == 1)
        ans = mid->val;
    merge(r, mid, r);
    merge(head, l, r);
    return ans;
}

int order_of_key(int val) //1 indexed
{
    node l, r, mid;
    split(head, l, r, val - 1);
    split(r, mid, r, val);
}

```

```

    int ans = -1;
    if(cnt(mid) == 1)
        ans = 1 + cnt(1);
    merge(r, mid, r);
    merge(head, 1, r);
    return ans;
}
};

```

40 Tree Bridge

```

struct dsu{
    v32 par,rk;
    dsu(){
        dsu(int n) {reset(n);}
    void reset(int n){
        rk.assign(n,0);
        par.resize(n);
        iota(all(par),0);
    }
    int getpar(int i){
        return (par[i]==i)? i:(par[i]=getpar(par[i]));
    }
    bool con(int i,int j){
        return getpar(i)==getpar(j);
    }
    bool join(int i,int j){
        i=getpar(i),j=getpar(j);
        if(i==j) return 0;
        if(rk[i]>rk[j]) par[j]=i;
        else{
            par[i]=j;
            if(rk[i]==rk[j]) rk[j]++;
        }
        return 1;
    }
};

vector<vector<int>> > getBridgeTree(vector<vector<int>> > &v,
    int n){
    int in[n]={0},low[n],ctm=0;
    vector<pair<int,int>> > bridge;
    dsu d(n);
    function<void(int,int)> dfs=[&](int u,int p){
        in[u]=low[u]=++ctm;
        for(auto &it: v[u]){
            if(it==p) continue;
            if(in[it]){
                if(low[it]<low[u]) low[u]=low[it];

```

```

            }else{
                dfs(it,u);
                if(low[it]<low[u]) low[u]=low[it];
                if(low[it]>in[u]) bridge.push_back({u,it});
                else d.join(u,it);
            }
        }
    };
    for(int i=0;i<n;++i) if(!in[i]) dfs(i,i);
    int par[n],id[n];
    for(int i=0;i<n;++i) par[i]=d.getpar(i),id[i]=-1;
    vector<vector<int>> > g;
    for(auto &it: bridge){
        it.first=par[it.first],it.second=par[it.second];
        if(id[it.first]==-1){
            id[it.first]=g.size();
            g.push_back(vector<int>(0));
        }
        if(id[it.second]==-1){
            id[it.second]=g.size();
            g.push_back(vector<int>(0));
        }
        g[id[it.first]].push_back(id[it.second]);
        g[id[it.second]].push_back(id[it.first]);
    }
    return g;
}
// dfs find all bridgees and g contain bridge tree rest is
// diameter calculation

```

41 Z Algorithm

```

// Z Algorithm
// Z[i] is the length of the longest substring starting from
// S[i]
// which is also a prefix of S
// O(n)
void z_func(v32 &s,v32 &z){
    int L=0,R=0;
    int sz=s.size();
    z.assign(sz,0);
    forsn(i,1,sz){
        if(i>R){
            L=R=i;
            while(R<sz && s[R-L]==s[R]) R++;
            z[i]=R-L; R--;
        }else{
            int k=i-L;

```

```

            if(z[k]<R-i+1) z[i]=z[k];
            else{
                L=i;
                while(R<sz && s[R-L]==s[R]) R++;
                z[i]=R-L; R--;
            }
        }
    }
}

```

42 Z Ideas

Gray codes:

Applications:

1. Gray code of n bits forms a Hamiltonian cycle on a hypercube, where each bit corresponds to one dimension.
2. Gray code can be used to solve the Towers of Hanoi problem. Let n denote number of disks. Start with Gray code of length n which consists of all zeroes ($G(0)$) and move between consecutive Gray codes (from $G(i)$ to $G(i+1)$). Let i -th bit of current Gray code represent n -th disk (the least significant bit corresponds to the smallest disk and the most significant bit to the biggest disk). Since exactly one bit changes on each step, we can treat changing i -th bit as moving i -th disk. Notice that there is exactly one move option for each disk (except the smallest one) on each step (except start and finish positions). There are always two move options for the smallest disk but there is a strategy which will always lead to answer: if n is odd then sequence of the smallest disk moves looks like $ftrftr \dots$ where f is the initial rod, t is the terminal rod and r is the remaining rod), and if n is even: $frtfrt \dots$.

```

int gray (int n) {
    return n ^ (n >> 1);
}

int rev_g (int g) {
    int n = 0;
    for (; g; g >>= 1)
        n ^= g;
    return n;
}

```

Enumerating all submasks of a bitmask:

```
for (int s=m; ; s=(s-1)&m) {
    ... you can use s ...
    if (s==0) break;
}
```

Sparse Table:

```
int st[MAXN][K];
```

```
for (int i = 0; i < N; i++)
    st[i][0] = array[i];
```

```
for (int j = 1; j <= K; j++)
    for (int i = 0; i + (1 << j) <= N; i++)
        st[i][j] = min(st[i][j-1], st[i + (1 << (j - 1))][j - 1]);
```

To get minimum of a range:

```
int j = log[R - L + 1];
int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);
```

Divide and Conquer DP:

Some dynamic programming problems have a recurrence of [this](#) form:

```
dp(i,j)=mink {dp(i1,k)+C(k,j)}
where C(k,j) is some cost function.
```

Say 1 in and 1 jm , and evaluating C takes $O(1)$ time. Straightforward evaluation of the above recurrence is $O(nm^2)$.

There are nm states, and m transitions for each state.

Let $opt(i,j)$ be the value of k that minimizes the above expression.

If $opt(i,j)$ $opt(i,j+1)$ for all i,j , then we can apply divide-and-conquer DP. This known as the monotonicity condition.

The optimal "splitting point" for a fixed i increases as j increases.

```
int n;
long long C(int i, int j);
vector<long long> dp_before(n), dp_cur(n);
```

```
// compute dp_cur[l], ... dp_cur[r] (inclusive)
void compute(int l, int r, int optl, int opttr)
```

```
{
    if (l > r)
        return;
    int mid = (l + r) >> 1;
    pair<long long, int> best = {INF, -1};

    for (int k = optl; k <= min(mid, opttr); k++) {
        best = min(best, {dp_before[k] + C(k, mid), k});
    }

    dp_cur[mid] = best.first;
    int opt = best.second;

    compute(l, mid - 1, optl, opt);
    compute(mid + 1, r, opt, opttr);
}

Knuth Optimization:
dp[i][j] = mini < k < j {dp[i][k] + dp[k][j]} + C[i][j]
monotonicity : C[b][c] <= C[a][d]
quadrangle inequality: C[a][c]+C[b][d] <= C[a][d]+C[b][c]

Lyndon factorization: We can get the minimum cyclic shift.
Factorize the string as s = w1w2w3...wn

string min_cyclic_string(string s) {
    s += s;
    int n = s.size();
    int i = 0, ans = 0;
    while (i < n / 2) {
        ans = i;
        int j = i + 1, k = i;
        while (j < n && s[k] <= s[j]) {
            if (s[k] < s[j])
                k = i;
            else
                k++;
            j++;
        }
        while (i <= k)
            i += j - k;
    }
    return s.substr(ans, n / 2);
}

Rank of a matrix:
const double EPS = 1E-9;
```

```
int compute_rank(vector<vector<double>> A) {
    int n = A.size();
    int m = A[0].size();

    int rank = 0;
    vector<bool> row_selected(n, false);
    for (int i = 0; i < m; ++i) {
        int j;
        for (j = 0; j < n; ++j) {
            if (!row_selected[j] && abs(A[j][i]) > EPS)
                break;
        }

        if (j != n) {
            ++rank;
            row_selected[j] = true;
            for (int p = i + 1; p < m; ++p)
                A[j][p] /= A[j][i];
            for (int k = 0; k < n; ++k) {
                if (k != j && abs(A[k][i]) > EPS) {
                    for (int p = i + 1; p < m; ++p)
                        A[k][p] -= A[j][p] * A[k][i];
                }
            }
        }
    }
    return rank;
}
```

Determinant of a matrix:

```
const double EPS = 1E-9;
int n;
vector<vector<double>> a(n, vector<double>(n));

double det = 1;
for (int i=0; i<n; ++i) {
    int k = i;
    for (int j=i+1; j<n; ++j)
        if (abs(a[j][i]) > abs(a[k][i]))
            k = j;
    if (abs(a[k][i]) < EPS) {
        det = 0;
        break;
    }
    swap(a[i], a[k]);
    if (i != k)
        det = -det;
    det *= a[i][i];
}
```



```

    for (int j=i+1; j<n; ++j)
        a[i][j] /= a[i][i];
    for (int j=0; j<n; ++j)
        if (j != i && abs (a[j][i]) > EPS)
            for (int k=i+1; k<n; ++k)
                a[j][k] -= a[i][k] * a[j][i];
}

cout << det;

Generating all k-subsets:

vector<int> ans;

void gen(int n, int k, int idx, bool rev) {
    if (k > n || k < 0)
        return;

    if (!n) {
        for (int i = 0; i < idx; ++i) {
            if (ans[i])
                cout << i + 1;
        }
        cout << "\n";
        return;
    }

    ans[idx] = rev;
    gen(n - 1, k - rev, idx + 1, false);
    ans[idx] = !rev;
    gen(n - 1, k - !rev, idx + 1, true);
}

void all_combinations(int n, int k) {
    ans.resize(n);
    gen(n, k, 0, false);
}

Simpsons formula for integration:

const int N = 1000 * 1000; // number of steps (already
                             multiplied by 2)

double simpson_integration(double a, double b){
    double h = (b - a) / N;
    double s = f(a) + f(b); // a = x_0 and b = x_2n
    for (int i = 1; i <= N - 1; ++i) { // Refer to final
        Simpson's formula
        double x = a + h * i;

```

```

        s += f(x) * ((i & 1) ? 4 : 2);
    }
    s *= h / 3;
    return s;
}

Picks theorem:

Given a certain lattice polygon with non-zero area.

We denote its area by S, the number of points with integer
coordinates lying strictly inside the polygon by I and the
number of points lying on polygon sides by B.

Then, the Pick formula states:  $S = I + B/2$ 
In particular, if the values of I and B for a polygon are
given,
the area can be calculated in  $O(1)$  without even knowing the
vertices.

Strongly Connected component and Condensation Graph:

vector < vector<int> > g, gr;
vector<bool> used;
vector<int> order, component;

void dfs1 (int v) {
    used[v] = true;
    for (size_t i=0; i<g[v].size(); ++i)
        if (!used[ g[v][i] ])
            dfs1 (g[v][i]);
    order.push_back (v);
}

void dfs2 (int v) {
    used[v] = true;
    component.push_back (v);
    for (size_t i=0; i<gr[v].size(); ++i)
        if (!used[ gr[v][i] ])
            dfs2 (gr[v][i]);
}

int main() {
    int n;
    ... reading n ...
    for (;;) {
        int a, b;
        ... reading next edge (a,b) ...
        g[a].push_back (b);

```

```

        gr[b].push_back (a);
    }

    used.assign (n, false);
    for (int i=0; i<n; ++i)
        if (!used[i])
            dfs1 (i);
    used.assign (n, false);
    for (int i=0; i<n; ++i) {
        int v = order[n-1-i];
        if (!used[v]) {
            dfs2 (v);
            ... printing next component ...
            component.clear();
        }
    }
}

FFT Matrices:

XOR FFT: 1 1 / 1 -1, AND FFT: 0 1/ 1 1, OR FFT: 1 1/ 1 0

Harmonic lemma:

for (int i = 1, la; i <= n; i = la + 1) {
    la = n / (n / i);
    v.pb(mp(n/i, la-i+1));
    //la-i+1 denotes the number of occurrences of n/i
    //n / x yields the same value for i <= x <= la.
}

Mobius inversion theory:

if f and g are multiplicative, then their dirichlet
convolution,
i.e  $\sum_{d|x} f(d)g(x/d)$  is also multiplicative. eg. choose
    g = 1

Properties:

1. If  $g(n) = \sum_{d|n} f(d)$ , then  $f(n) = \sum_{d|n} g(d)u(n/d)$ .
2.  $\sum_{d|n} u(d) = [n=1]$ 

Standard question: Number of co-prime integers in range 1,n
Answer:  $f(n) = \sum_{d=1}^n u(d) \text{floor}(n/d)^2$ 

Euler totient:  $\phi(\text{totient fn}) = u * n$  (dirichlet convolution)

```