Team Notebook

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1 Advice

```
Pre-submit:
```

Are time limits close? If so, generate max cases. Is the memory usage fine? Could anything overflow? Make sure to submit the right file.

Wrong answer: Print your solution! Print debug output, as well.

Are you clearing all datastructures between test cases?

Can your algorithm handle the whole range of input?

Read the full problem statement again. Do you handle all corner cases correctly? Have you understood the problem correctly?

Any uninitialized variables? Any

overflows? Confusing N and M, i and j, etc.?

Are you sure

your algorithm works? What special cases have you not thought of?

Are you sure the STL functions you use work as you think? Add some assertions, maybe resubmit Create some testcases to run your algorithm on. Go through the algorithm for a simple case.

Go through this list again. Explain your algorithm to a team $$\operatorname{\mathtt{mate}}$.$

Ask the team mate to look at your code. Go for a small walk, e.g. to the toilet. Is your output format correct?

Rewrite your solution from the start or let a team mate do it.

Runtime error: Have you tested all corner cases locally? Any uninitialized variables? Are you reading or writing outside the range of any vector? Any assertions that might fail? Any possible division by 0? (mod 0 for example)

Any possible infinite recursion? Invalidated pointers or iterators? Are you using too much memory? Debug with resubmits.

Time limit exceeded: Do you have any possible infinite loops $\ref{eq:condition}$

What is the complexity of your algorithm? Are you copying a lot of unnecessary data? (References) How big is the input and output? (consider scanf) Avoid vector, map. (use arrays/unordered_map) What do your team mates think about your algorithm?

Memory limit exceeded: What is the max amount of

memory your algorithm should need? Are you clearing all datastructures between test cases?

Primes - 10001st prime is 1299721, 100001st prime is 15485867 Large primes - 999999937, 1e9+7, 987646789, 987101789 78498 primes less than 10⁶ The number of divisors of n is at most around 100, for n<5e4, 500 for n<1e7, 2000 for n<1e10, 200,000 for n<1e19 7! 5040, 8! 40320, 9! 362880, 10! 362880, 11! 4.0e7, 12! 4.8e8, 15! 1.3e12, 20! 2e18

The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e19

Articulation points and bridges articulation point:- there
exist child : dfslow[child] >= dfsnum[curr] bridge :- tree
 ed:
dfslow[ch] > dfsnum[par];

A connected multigraph has an Euler path but not an Euler circuit if and only if it has exactly two vertices of odd degree

Binomial coefficients - base case ncn and nc0 = 1; recursion is nCk = (n-1)C(k-1)+(n-1)Ck

Catalan numbers - used in valid paranthesis expressions - formula is $Cn = summation\{i=0 \text{ to } n-1\}$ (CiCn-i-1); Another formula is Cn = 2nCn/(n+1). There are Cn binary trees of n nodes and Cn-1 rooted trees of n nodes

```
Derangements - D(n) = (n-1)(D(n-1)+D(n-2))
```

Burnsides Lemma - number of equivalence classes = (summation I(pi))/n : I(pi) are number of fixed points. Usual

formula: [summation {i=0 to n-1} k^gcd(i,n)]/n

Stirling numbers - first kind - permutations of n elements with k disjoint cycles. s(n+1,k) = ns(n,k)+s(n,k-1). s(0,0)

1, s(n,0) = 0 if n>0. Summation $\{k=0 \text{ to } n\}$ s(n,k) = n!

Stirling numbers - Second kind - partition n objects into k non empty subsets. S(n+1,k) = kS(n,k) + S(n,k-1). S(0,0) = 1, S(n,0) = 0 if n>0. $S(n,k) = (summation{j=0 to k} [(-1)^(kj)kCjj^n])/k!$

Hermite identity - summation{k=0 to n-1} floor[(x+k)/n] =

```
floor[nx]
```

Kirchoff matrix tree theorem - number of spanning trees in a graph is determinant of Laplacian Matrix with one row and column removed, where L = degree matrix - adjacency matrix

Expected value tricks:

- 1. Linearity of Expectation: E(X+Y) = E(X)+E(Y)
- 2. Contribution to the sum If we want to find the sum over many ways/possibilities, we should consider every element (maybe a number, or a pair or an edge) and count how many times it will be added to the answer.
- 3. Forcindependent events E(XY) = E(X)E(Y)
- 4. Ordered pairs (Super interpretation of square) The square of the size of a set is equal to the number of ordered pairs of elements in the set. So we iterate over pairs and for each we compute the contribution to the answer. Similarly, the k-th power is equal to the number of sequences (tuples) of length k.

powers, it might help to also maintain the sum of smaller powers. For example, if the sum of 0-th, 1-th and 2-nd powers is S0, S1 and S2, and we increase every element by x, the new sums are S0, S1+S0x and S2 + 2S1x + x^2 S0.

2 Aho Corasick

```
struct AhoCorasick{
enum {alpha=26.first='a'}:
struct Node{
 int back. next[alpha]. start = -1. end = -1. nmatches = 0:
 Node(int v){memset(next.v.sizeof(next)):}
vector<Node> N:
vector<int> backp;
inline void insert(string &s.int i){
 assert(!s.emptv()):
 int n=0:
 for(auto &c: s){
  int &m=N[n].next[c-first];
  if(m==-1){n=m=N.size(); N.emplace_back(-1);}
 if(N[n].end==-1) N[n].start=j;
 backp.push_back(N[n].end);
 N[n].end=j;
```

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```
N[n].nmatches++:
void clear(){
 N.clear():
 backp.clear();
void create(vector<string>& pat){
 N.emplace_back(-1);
 for(int i=0;i<pat.size();++i) insert(pat[i],i);</pre>
 N[0].back=N.size();
 N.emplace_back(0);
 aueue<int> a:
 for(q.push(0);!q.empty();q.pop()){
  int n=q.front(),prev=N[n].back;
  for(int i=0;i<alpha;++i){</pre>
   int &ed=N[n].next[i], y=N[prev].next[i];
   if(ed==-1) ed=v:
   else{
    N[ed].back=v:
    (N[ed].end==-1 ? N[ed].end:backp[N[ed].start])=N[y].end;
    N[ed].nmatches+=N[v].nmatches;
    q.push(ed);
11 find(string word){
 int n=0:
 // vector<int> res;
 11 count=0:
 for(auto &c: word){
  n=N[n].next[c-first];
  // res.push_back(N[n].end);
  count+=N[n].nmatches;
 return count:
};
struct AhoOnline{
int sz=0:
vector<string> v[25]:
AhoCorasick c[25];
void add(string &p){
 int val=__builtin_ctz(~sz);
 auto &cur=v[val];
 for(int i=0:i<val:++i){</pre>
  for(auto &it: v[i]) cur.push_back(it);
  c[i].clear():
  v[i].clear();
 }
```

```
cur.push_back(p);
c[val].create(cur);
++sz;
}
ll query(string &p){
    ll ans=0;
    for(int i=0;i<25;++i){
        if((1<<i)&sz) ans+=c[i].find(p);
        if((1<<i)>=sz) break;
}
return ans;
}
} add,del;
```

3 Centroid Decomposition

```
vector<set<int> > g;
vector<int> par,sub;
int dfs(int u,int p){
sub[u]=1:
for(auto &it: g[u]) if(it!=p) sub[u]+=dfs(it,u);
return sub[u]:
int find_centroid(int u,int p,int n){
for(auto &it: g[u]){
 if(it!=p && sub[it]>n/2){
 return find_centroid(it,u,n);
return u:
void decompose(int u,int p=-1){
int n=dfs(u.p):
int centroid=find_centroid(u,p,n);
if(p==-1) p=centroid;
// Do stuff here for merges
// Recurse
par[centroid]=p;
for(auto &it: g[centroid]){
 g[it].erase(centroid);
 decompose(it,centroid);
g[centroid].clear():
void reset(int n){
par.resize(n):
sub.resize(n);
```

```
g.assign(n,set<int>());
```

4 Convex Hull and Li Chao tree

```
// Li chao Tree (can be made persistent)
struct Line{
11 m. c:
Line(ll mm=0.11 cc=-3e18): m(mm).c(cc){}
inline 11 get(const int &x){return m*x+c;}
inline 11 operator [](const int &x){return m*x+c;}
vector<Line> LN:
struct node{
node *lt,*rt;
int Ln:
node(const int&1): Ln(1),lt(0),rt(0){};
inline 11 operator[](const int &x){ return LN[Ln].get(x);}
inline 11 get(const int &x){return LN[Ln].get(x);}
const static int LX=-(1e9+1),RX=1e9+1;
struct Dynamic_Hull{ /* Max hull */
node *root=0:
void add(int 1,node* &it,int lx=LX,int rx=RX){
 if(it==0) it=new node(1):
 if(it->get(lx)>=LN[l].get(lx) and it->get(rx)>=LN[l].get(
 if(it->get(lx)<=LN[1].get(lx) and it->get(rx)<=LN[1].get(</pre>
      rx)){
  it->Ln=1:
  return:
 int mid=(lx+rx)>>1:
 if(it->get(lx)<LN[1][lx]) swap(it->Ln,1);
 if(it->get(mid)>=LN[1][mid]){
  add(l.it->rt.mid+1.rx):
 }else{
   swap(it->Ln.1):
   add(l.it->lt.lx.mid):
 inline void add(int ind){add(ind,root);}
 inline void add(int m,int c){LN.pb(Line(m,c));add(LN.size
      ()-1.root):}
 11 get(int &x,node* &it,int lx=LX,int rx=RX){
   if(it==0) return -3e18; // Max hull
   ll ret=it->get(x);
   int mid=(lx+rx)>>1;
```

```
if(x<=mid) ret=max(ret,get(x,it->lt,lx,mid));
   else ret=max(ret,get(x,it->rt,mid+1,rx));
   return ret;
 inline 11 get(int x){return get(x,root);}
}:
// const static int LX = -(1e9), RX = 1e9;
// struct Dynamic_Hull { /* Max hull */
// struct Line{
      11 m. c: // slope, intercept
      Line(11 mm=0, 11 cc=-1e18) { m = mm; c = cc; }
      11 operator[](const int&x){ return m*x+c; }
// struct node {
      node *lt.*rt: Line Ln:
      node(const Line &1){lt=rt=nullptr; Ln=1;}
// node *root=nullptr:
// void add(Line l,node*&it,int lx=LX,int rx=RX){
      if(it==nullptr)it=new node(1);
      if(it->Ln[lx]>=l[lx] and it->Ln[rx]>=l[rx]) return;
      if(it->Ln[lx]<=l[lx] and it->Ln[rx]<=l[rx]) {it->Ln=l;
      return:}
      int mid = (lx+rx)>>1;
      if(it->Ln[lx] < l[lx]) swap(it->Ln,1);
      if(it->Ln[mid] >= l[mid]) add(l.it->rt.mid+1.rx):
      else { swap(it->Ln,1); add(1,it->lt,lx,mid); }
// }
// void add(const ll &m,const ll &c) { add(Line(m,c),root);
// ll get(int &x,node*&it,int lx=LX,int rx=RX){
      if(it==NULL) return -1e18; // Max hull
      11 \text{ ret} = it - \sum [x]:
      int mid = (1x+rx)>>1:
      if(x<=mid) ret = max(ret , get(x,it->lt,lx,mid));
      else ret = max(ret , get(x,it->rt,mid+1,rx));
      return ret;
// ll get(int x){ return get(x,root); }
struct Hull{
 struct line {
   11 m,c;
   11 eval(11 x){return m*x+c:}
   ld intersectX(line 1){return (ld)(c-1.c)/(1.m-m);}
   line(ll m.ll c): m(m).c(c){}
 deque<line> dq;
```

```
v32 ints:
Hull(int n){ints.clear(); forn(i,n) ints.pb(i); dq.clear()
// Dec order of slopes
void add(line cur){
  while(dq.size()>=2 && cur.intersectX(dq[0])>=dq[0].
      intersectX(dq[1]))
    dq.pop_front();
  dq.push_front(cur);
void add(const 11 &m.const 11 &c){add(line(m.c));}
// guery sorted dec.
// 11 getval(11 x){
// while(dq.size()>=2 && dq.back().eval(x)<=dq[dq.size()</pre>
     -21.eval(x)
      dq.pop_back();
// return dq.back().eval(x);
// }
// arbitary guery
11 getval(ll x,deque<line> &dq){
  auto cmp = [&dq](int idx,ll x){return dq[idx].intersectX(
       da[idx+1])<x:}:
  int idx = *lower_bound(ints.begin(),ints.begin()+dq.size
       ()-1,x,cmp);
  return dq[idx].eval(x);
11 get(const 11 &x){return getval(x,dq);}
```

5 Dynamic Connectivity

```
int u[LIM],v[LIM],e[LIM],q[LIM];
map<p32,int> ids;
struct dsu{
  int sz;
    v32 par,rk;
    stack<int> st;
    void reset(int n){
        rk.assign(n,1);
        par.resize(n);
        iota(all(par),0);
        sz=n;
    }
  int getpar(int i){
        return (par[i]==i)? i:getpar(par[i]);
    }
    bool con(int i,int j){
        return getpar(i)==getpar(j);
}
```

```
bool join(int i,int j){
       i=getpar(i), j=getpar(j);
       if(i==j) return 0;
       --sz;
       if(rk[i]>rk[i]) swap(i,i):
       par[j]=i,rk[i]+=rk[j];
       st.push(j);
       return 1;
   int moment(){
    return st.size():
   }
   void revert(int tm){
    while(st.size()>tm){
     auto tp=st.top();
     rk[par[tp]]-=rk[tp];
     par[tp]=tp;
     st.pop();
     ++sz:
   }
void solve(int 1,int r,vp32 &ed){
if(l>r) return:
// dbg(ed,1,r,d.sz);
int mid=(l+r)>>1;
vp32 low:
int tm=d.moment();
forstl(it.ed){
 if(it.se<l or it.fi>r) continue;
 else if(it.fi<=l and it.se>=r) d.join(u[it.fi],v[it.fi]);
 else low.pb(it);
if(l==r){
 if(q[1]) cout<<(d.con(u[1],v[1])? "YES":"NO")<<ln:
 solve(1,mid,low);
 solve(mid+1,r,low);
d.revert(tm):
signed main(){
fastio:
cin>>n>>k;
d.reset(n):
string t;
forn(i,k){
 cin>>t:
 cin>>x>>y; --x,--y;
```

```
if(x>y) swap(x,y);
u[i]=x,v[i]=v;
if(t[0]=='c'){
  q[i]=1;
}else{
 if(t[0]=='a'){
  ids[mp(x,y)]=i;
  e[i]=k-1;
 }else{
  e[ids[mp(x,y)]]=i;
  e[i]=-1:
}
vp32 ed:
forn(i,k) if(!q[i] && e[i]!=-1) ed.pb({i,e[i]});
solve(0.k-1.ed):
return 0;
```

6 Euler Path

```
procedure FindEulerPath(V)

1. iterate through all the edges outgoing from vertex V;
remove this edge from the graph,
and call FindEulerPath from the second end of this
edge;
2. add vertex V to the answer.
```

7 Extended Euclidean GCD

```
int egcd(int a,int b, int* x, int* y){
    if(a==0) {
        *x=0;*y=1;
        return b;
    }
    int x1,y1;
    int gcd=egcd(b%a,a,&x1,&y1);
    *x=y1-(b/a)*x1;
    *y=x1;
    return gcd;
}
```

8 Fenwick 2D

```
//BIT<N, M, K> b; N x M x K (3-dimensional) BIT
//b.update(x, y, z, P); // add P to (x,y,z)
//b.query(x1, x2, y1, y2, z1, z2); // query between (x1, y1,
     z1) and (x2, y2, z2)
inline int lastbit(int x){
 return x&(-x):
template <int N. int... Ns>
struct BIT<N, Ns...> {
 BIT<Ns...> bit[N + 1]:
 template<typename... Args>
 void update(int pos, Args... args) {
   for (; pos <= N; bit[pos].update(args...), pos += lastbit</pre>
        (pos));
 template<tvpename... Args>
 int query(int 1, int r, Args... args) {
   for (; r \ge 1; ans += bit[r].query(args...), r -= lastbit
   for (--1: 1 >= 1: ans -= bit[1].guerv(args...), 1 -=
        lastbit(1));
   return ans:
// Another implementation
struct FenwickTree2D {
   vector<vector<int>> bit:
   int n. m:
   // init(...) { ... }
   int sum(int x, int y) {
       int ret = 0:
       for (int i = x; i \ge 0; i = (i & (i + 1)) - 1)
           for (int j = y; j \ge 0; j = (j & (j + 1)) - 1)
              ret += bit[i][j];
       return ret;
   void add(int x, int y, int delta) {
       for (int i = x; i < n; i = i | (i + 1))</pre>
           for (int j = y; j < m; j = j | (j + 1))
              bit[i][i] += delta:
};
```

9 Gaussian Elimination, Base 2

```
struct Gaussbase2{
   int numofbits=20;
   int rk=0:
   v32 Base;
   Gaussbase2() {clear():}
   void clear(){
       rk=0;
       Base.assign(numofbits,0);
   Gaussbase2& operator = (Gaussbase2 &g){
       forn(i.numofbits) Base[i]=g.Base[i]:
       rk=g.rk;
   bool canbemade(int x){
       rforn(i,numofbits-1) x=min(x,x^Base[i]);
       return x==0:
   void Add(int x){
       rforn(i.numofbits-1){
          if((x>>i)&1){
              if(!Base[i]){
                  Base[i]=x:
                  rk++;
                  return:
              }else x^=Base[i];
      }
   int maxxor(){
       int ans=0:
       rforn(i,numofbits-1){
           if(ans < (ans^Base[i])) ans^=Base[i];</pre>
       return ans;
};
```

10 Gaussian Elimination

```
int gauss (vector <vector<double> > a, vector<double> &ans){
   int n = (int) a.size();
   int m = (int) a[0].size()-1;

   vector<int> where(m,-1);
   for(int col=0, row=0;col<m && row<n; ++col){
      int sel = row;
      for(int i=row;i<n;++i){</pre>
```

```
if(abs(a[i][col]) > abs(a[sel][col])){
           sel = i;
       }
   if(abs(a[sel][col]) < EPS) continue;</pre>
   for(int i=col; i<=m; ++i){</pre>
       swap(a[sel][i],a[row][i]);
   where[col] = row;
   for(int i=0:i<n:++i){</pre>
       if(i!=row){
           double c = a[i][col]/a[row][col];
           for(int j=col; j<=m;++j){</pre>
               a[i][i] -= a[row][i]*c;
       }
   }
    ++row:
ans.assign(m,0);
for(int i=0:i<m:++i){</pre>
   if(where[i]!=-1){
       ans[i] = a[where[i]][m]/a[where[i]][i];
   }
for(int i=0:i<n:++i){</pre>
   double sum=0:
   for(int j=0;j<m;++j){</pre>
       sum+=ans[i]*a[i][i]:
   if(abs(sum-a[i][m])>EPS)
       return 0:
for(int i=0:i<m:++i){</pre>
   if(where[i]==-1) return MOD:
return 1;
```

11 General Weighted Matching

```
struct MaxMatchingEdmonds{
    // Assume General Unweighted Directed Graph
    // O(V^3) edmonds for maximum matching
    vv32 g;
    v32 match,p,base;
    vector<bool> blossom;
```

```
int n:
int lca(int a,int b){
   vector<bool> used(match.size(),0);
   while(1){
     a=base[a];
     used[a]=1:
     if(match[a] == -1) break;
     a=p[match[a]];
   while(1){
     b=base[b]:
     if(used[b]) return b:
     b=p[match[b]];
}
void markPath(int v,int b,int children) {
   for(:base[v]!=b:v=p[match[v]]){
     blossom[base[v]]=blossom[base[match[v]]]=1;
     p[v]=children:
     children=match[v]:
int findPath(int root) {
   vector<bool> used(n,0);
   p.assign(n,-1);
   base.assign(n,0);
   for(int i=0:i<n:++i) base[i] = i:</pre>
   used[root]=1:
   int qh=0;
   int qt=0;
   v32 q(n,0);
   q[qt++]=root;
   while(qh<qt){</pre>
     int v=q[qh++];
     for(int &to:g[v]){
       if(base[v]==base[to] || match[v]==to) continue;
       if(to==root || match[to]!=-1 && p[match[to]]!=-1)
         int curbase=lca(v,to);
         blossom.assign(n,0);
         markPath(v.curbase.to):
         markPath(to,curbase,v);
         for(int i=0:i<n:++i)</pre>
           if(blossom[base[i]]){
             base[i]=curbase;
             if(!used[i]){
              used[i]=1;
              q[qt++]=i;
          }
```

```
}else if(p[to]==-1){
             p[to]=v:
             if(match[to] == -1) return to;
             to=match[to]:
             used[to]=1;
             q[qt++]=to;
       return -1;
   int maxMatching(vv32 &graph){
       n=graph.size();
       g=graph;
       match.assign(n,-1);
       p.assign(n,0);
       for(int i=0:i<n:++i){</pre>
         if(match[i]==-1){
           int v=findPath(i):
           while (v!=-1) {
             int pv=p[v];
             int ppv=match[pv];
             match[v]=pv;
             match[pv]=v;
             v=ppv;
       int matches=0;
       for(int i=0:i<n:++i) if(match[i]!=-1) ++matches:</pre>
       return matches/2;
};
```

12 Geometry

```
const int MAX_SIZE = 1000;
const double PI = 2.0*acos(0.0);
struct PT
{
    double x,y;
    double length() {return sqrt(x*x+y*y);}
    int normalize(){
        // normalize the vector to unit length; return -1 if the
            vector is 0
    double l = length();
    if(fabs(1)<EPS) return -1;</pre>
```

```
x/=1: v/=1:
 return 0:
 PT operator-(PT a){
 PT r;
 r.x=x-a.x; r.y=y-a.y;
 return r;
 PT operator+(PT a){
 PT r;
 r.x=x+a.x; r.y=y+a.y;
 return r:
 PT operator*(double sc){
 r.x=x*sc; r.y=y*sc;
 return r:
}
}:
bool operator<(const PT& a,const PT& b){
if(fabs(a.x-b.x) < EPS) return a.y < b.y;</pre>
return a.x<b.x;</pre>
double dist(PT& a. PT& b){
// the distance between two points
return sqrt((a.x-b.x)*(a.x-b.x) + (a.v-b.v)*(a.v-b.v)):
double dot(PT& a, PT& b){
// the inner product of two vectors
return(a.x*b.x+a.y*b.y);
double cross(PT& a, PT& b){
return(a.x*b.y-a.y*b.x);
// =========
// The Convex Hull
// =========
int sideSign(PT& p1.PT& p2.PT& p3){
// which side is p3 to the line p1->p2? returns: 1 left, 0
    on. -1 right
 double sg = (p1.x-p3.x)*(p2.y-p3.y)-(p1.y - p3.y)*(p2.x-p3.y)
 if(fabs(sg)<EPS) return 0:</pre>
if(sg>0) return 1;
return -1:
```

```
bool better(PT& p1.PT& p2.PT& p3){
// used by convec hull: from p3, if p1 is better than p2
double sg = (p1.y - p3.y)*(p2.x-p3.x)-(p1.x-p3.x)*(p2.y-p3.x)
//watch range of the numbers
if(fabs(sg)<EPS){</pre>
 if(dist(p3,p1)>dist(p3,p2))return true;
 else return false;
if(sg<0) return true;</pre>
return false:
void vex2(vector<PT> vin.vector<PT>& vout){
// vin is not pass by reference, since we will rotate it
vout.clear();
int n=vin.size():
sort(vin.begin(),vin.end());
PT stk[MAX_SIZE];
int pstk. i:
// hopefully more than 2 points
stk[0] = vin[0]:
stk[1] = vin[1];
pstk = 2:
for(i=2: i<n: i++){</pre>
 if(dist(vin[i], vin[i-1]) < EPS) continue;</pre>
 while(pstk > 1 && better(vin[i], stk[pstk-1], stk[pstk-2])
 pstk--;
 stk[pstk] = vin[i]:
 pstk++;
for(i=0; i<pstk; i++) vout.push_back(stk[i]);</pre>
// turn 180 degree
for(i=0: i<n: i++){</pre>
 vin[i].v = -vin[i].v:
 vin[i].x = -vin[i].x:
sort(vin.begin(), vin.end());
stk[0] = vin[0]:
stk[1] = vin[1]:
pstk = 2;
for(i=2: i<n: i++){</pre>
 if(dist(vin[i], vin[i-1]) < EPS) continue;</pre>
 while(pstk > 1 && better(vin[i], stk[pstk-1], stk[pstk-2])
     )
 pstk--;
 stk[pstk] = vin[i]:
 pstk++:
```

```
for(i=1: i<pstk-1: i++){</pre>
 stk[i].x= -stk[i].x; // dont forget rotate 180 d back.
 stk[i].y= -stk[i].y;
 vout.push back(stk[i]):
}
int isConvex(vector<PT>& v){
// test whether a simple polygon is convex
// return 0 if not convex, 1 if strictly convex,
// 2 if convex but there are points unnecesary
// this function does not work if the polycon is self
     intersecting
// in that case, compute the convex hull of v, and see if
     both have the same area
int i,j,k;
int c1=0: int c2=0: int c0=0:
int n=v.size():
for(i=0:i<n:i++){</pre>
 i=(i+1)%n:
 k=(j+1)%n;
 int s=sideSign(v[i], v[j], v[k]);
 if(s==0) c0++:
 if(s>0) c1++:
 if(s<0) c2++;
if(c1 && c2) return 0;
if(c0) return 2:
return 1;
// =========
// Areas
// =======
double trap(PT a, PT b){
// Used in various area functions
return (0.5*(b.x - a.x)*(b.y + a.y));
double area(vector<PT> &vin){
// Area of a simple polygon, not neccessary convex
int n = vin.size();
double ret = 0.0;
for(int i = 0; i < n; i++) ret += trap(vin[i], vin[(i+1)%n]
     1):
return fabs(ret);
double peri(vector<PT> &vin){
// Perimeter of a simple polygon, not neccessary convex
int n = vin.size();
double ret = 0.0:
```

```
for(int i = 0; i < n; i++) ret += dist(vin[i], vin[(i+1)%n]
return ret;
double triarea(PT a, PT b, PT c){
return fabs(trap(a,b)+trap(b,c)+trap(c,a));
double height(PT a, PT b, PT c){
// height from a to the line bc
double s3 = dist(c, b):
double ar=triarea(a,b,c);
return(2.0*ar/s3):
// ==========
// Points and Lines
int intersection( PT p1, PT p2, PT p3, PT p4, PT &r ) {
// two lines given by p1->p2, p3->p4 r is the intersection
// return -1 if two lines are parallel
double d = (p4.v - p3.v)*(p2.x-p1.x) - (p4.x - p3.x)*(p2.v
if( fabs( d ) < EPS ) return -1:
// might need to do something special!!!
double ua. ub:
ua = (p4.x - p3.x)*(p1.y-p3.y) - (p4.y-p3.y)*(p1.x-p3.x);
// ub = (p2.x - p1.x)*(p1.y-p3.y) - (p2.y-p1.y)*(p1.x-p3.x)
//ub /= d:
r = p1 + (p2-p1)*ua;
return 0:
void closestpt( PT p1, PT p2, PT p3, PT &r ){
// the closest point on the line p1->p2 to p3
if( fabs( triarea( p1, p2, p3 ) ) < EPS ) { r = p3; return;</pre>
     }
PT v = p2-p1;
v.normalize():
double pr; // inner product
pr = (p3.y-p1.y)*v.y + (p3.x-p1.x)*v.x;
r = p1+v*pr:
int hcenter( PT p1, PT p2, PT p3, PT& r ){
// point generated by altitudes
```

```
if( triarea( p1, p2, p3 ) < EPS ) return -1;</pre>
PT a1. a2:
closestpt( p2, p3, p1, a1 );
closestpt( p1, p3, p2, a2 );
intersection( p1, a1, p2, a2, r );
return 0:
int center( PT p1, PT p2, PT p3, PT& r ){
// point generated by circumscribed circle
if( triarea( p1, p2, p3 ) < EPS ) return -1;</pre>
PT a1. a2. b1. b2:
a1 = (p2+p3)*0.5:
a2 = (p1+p3)*0.5;
b1.x = a1.x - (p3.y-p2.y);
b1.v = a1.v + (p3.x-p2.x);
b2.x = a2.x - (p3.y-p1.y);
b2.v = a2.v + (p3.x-p1.x):
intersection(a1, b1, a2, b2, r);
int bcenter( PT p1, PT p2, PT p3, PT& r ){
// angle bisection
if( triarea( p1, p2, p3 ) < EPS ) return -1;</pre>
double s1, s2, s3:
s1 = dist(p2, p3);
s2 = dist( p1, p3 );
s3 = dist(p1, p2):
double rt = s2/(s2+s3);
PT a1.a2:
a1 = p2*rt+p3*(1.0-rt);
rt = s1/(s1+s3);
a2 = p1*rt+p3*(1.0-rt):
intersection( a1,p1, a2,p2, r );
return 0:
// ==========
// ===========
double angle(PT& p1, PT& p2, PT& p3){
// angle from p1->p2 to p1->p3, returns -PI to PI
PT va = p2-p1:
va.normalize():
PT vb; vb.x=-va.y; vb.y=va.x;
PT v = p3-p1:
double x,y;
x=dot(v, va):
v=dot(v, vb):
return(atan2(v.x)):
```

```
double angle(double a, double b, double c){
// in a triangle with sides a.b.c. the angle between b and
// we do not check if a,b,c is a triangle here
double cs=(b*b+c*c-a*a)/(2.0*b*c);
return(acos(cs)):
void rotate(PT p0, PT p1, double a, PT& r){
// rotate p1 around p0 clockwise, by angle a
// dont pass by reference for p1, so r and p1 can be the
p1 = p1-p0:
r.x = cos(a)*p1.x-sin(a)*p1.y;
r.y = \sin(a)*p1.x+\cos(a)*p1.y;
r = r + p0;
void reflect(PT% p1, PT% p2, PT p3, PT% r){
// p1->p2 line, reflect p3 to get r.
if(dist(p1, p3) < EPS) {r=p3; return;}</pre>
double a=angle(p1, p2, p3);
rotate(p1, r, -2.0*a, r);
// points, lines, and circles
// ===========
int pAndSeg(PT& p1, PT& p2, PT& p){
// the relation of the point p and the segment p1->p2.
// 1 if point is on the segment; 0 if not on the line; -1
     if on the line but not on the segment
double s=triarea(p, p1, p2);
if(s>EPS) return(0);
double sg=(p.x-p1.x)*(p.x-p2.x);
if(sg>EPS) return(-1);
sg=(p.v-p1.v)*(p.v-p2.v):
if(sg>EPS) return(-1);
return(1):
int lineAndCircle(PT& oo, double r, PT& p1, PT& p2, PT& r1,
    PT& r2){
// returns -1 if there is no intersection
// returns 1 if there is only one intersection
```

```
closestpt(p1,p2,oo,m);
 PT v = p2-p1;
 v.normalize();
 double r0=dist(oo. m):
 if(r0>r+EPS) return -1;
 if(fabs(r0-r)<EPS){</pre>
 r1=r2=m;
 return 1;
 double dd = sqrt(r*r-r0*r0);
 r1 = m-v*dd: r2 = m+v*dd:
return 0:
}
int CAndC(PT o1, double r1, PT o2, double r2, PT &q1, PT& q2
    ){
 // intersection of two circles
 // -1 if no intersection or infinite intersection
 // 1 if only one point
 double r=dist(o1,o2);
 if(r1<r2) { swap(o1,o2); swap(r1,r2); }</pre>
 if(r<EPS) return(-1):</pre>
 if(r>r1+r2+EPS) return(-1);
 if(r<r1-r2-EPS) return(-1);</pre>
 PT v = o2-o1; v.normalize();
 q1 = o1+v*r1;
 if(fabs(r-r1-r2) < EPS || fabs(r+r2-r1) < EPS)</pre>
 { q2=q1; return(1); }
 double a=angle(r2, r, r1);
 q2=q1;
 rotate(o1, q1, a, q1);
 rotate(o1, q2, -a, q2);
return 0:
int pAndPoly(vector<PT> pv, PT p){
// the relation of the point and the simple polygon
 // 1 if p is in pv; 0 outside; -1 on the polygon
 int i, j;
 int n=pv.size();
 pv.push_back(pv[0]);
 for(i=0;i<n;i++) if(pAndSeg(pv[i], pv[i+1], p)==1) return</pre>
 for(i=0;i<n;i++) pv[i] = pv[i]-p;</pre>
 p.x=p.y=0.0;
 double a, y;
 while(1){
 a=(double)rand()/10000.00:
```

```
j=0;
for(i=0;i<n;i++){
  rotate(p, pv[i], a, pv[i]);
  if(fabs(pv[i].x)<EPS) j=1;
}
if(j==0){
  pv[n]=pv[0];
  j=0;
  for(i=0;i<n;i++) if(pv[i].x*pv[i+1].x < -EPS){
   y=pv[i+1].y-pv[i+1].x*(pv[i].y-pv[i+1].y)/(pv[i].x-pv[i+1].x);
   if(y>0) j++;
  }
  return(j%2);
}
return 1;
}
```

13 Giant Step Baby Step

```
// Giant Step - Baby Step for discrete log
// find x with a^x = b mod MOD
// Find one soln can be changed to find all
// O(root(MOD)*log(MOD)) can be reduced with unordered map
    or array
11 solve(11 a,11 b,11 MOD){
   int n=(int)sqrt(MOD+.0)+1;
   11 an=1,cur;
   forn(i,n) an=(an*a)%MOD;
   vector<pair<11,int> > vals;
   forsn(i,1,n+1){
    vals.pb(mp(cur,i));
       cur=(cur*an)%MOD;
   cur=b:
   sort(all(vals));
   forn(i,n+1){
    auto in=lower_bound(all(vals),mp(cur,-1))-vals.begin();
    if(in!=vals.size() && vals[in].fi==cur){
     ll ans=n*(ll)vals[in].se-i:
     if(ans<MOD) return ans;</pre>
    }
       cur=(cur*a)%MOD:
   return -1:
```

14 Hashtable

```
struct hashtable{
v64 hash1, hash2, inv1, inv2;
11 MOD1=MOD.MOD2=MOD+2:
ll pr1=31,pr2=37;
void create(string &p){
 int len=p.size():
 hash1.resize(len); hash2.resize(len);
 inv1.resize(len):inv2.resize(len):
 ll p1=1,p2=1;
 int i=0;
 while(p[i]){
  hash1[i]= (i==0)? 0:hash1[i-1]:
  hash2[i]= (i==0)? 0:hash2[i-1]:
  hash1[i] = (hash1[i] +p[i] *p1)%MOD1;
  hash2[i] = (hash2[i] + p[i] * p2) %MOD2;
  p1=p1*pr1%MOD1;
  p2=p2*pr2%MOD2;
  i++;
 11 iv1=inv(pr1,MOD1),iv2=inv(pr2,MOD2);
 inv1[0]=1.inv2[0]=1:
 forsn(i.1.len){
  inv1[i]=inv1[i-1]*iv1%MOD1;
  inv2[i]=inv2[i-1]*iv2%MOD2:
p64 gethash(int l,int r){
 11 ans1=hash1[r-1];
 if(1!=0) ans1+=MOD1-hash1[1-1]:
 ll ans2=hash2[r-1]:
 if(1!=0) ans2+=MOD2-hash2[1-1];
 ans1=ans1*inv1[1]%MOD1:
 ans2=ans2*inv2[1]%MOD2;
 return mp(ans1,ans2);
};
```

15 Heavy Light Decomposition

```
struct SegTree{
  v32 T,lazy;
  int N,MX;
  void clear(int n,int mx){
  N=n,MX=mx;
  T.assign(4*N,0);
```

```
lazv.assign(4*N.0):
 void build(int a[],int v,int tl,int tr){
 if(tl==tr){
  T[v]=a[t1];
  }else{
  int tm=(tl+tr)>>1,lf=v<<1,rt=lf^1;;</pre>
  build(a,lf,tl,tm);
  build(a,rt,tm+1,tr);
  T[v]=min(T[lf],T[rt]);
 void push(int v){
  int lf=v<<1,rt=lf^1;</pre>
 T[lf]=(T[lf]+lazy[v]);
 lazy[lf]=(lazy[lf]+lazy[v]);
 T[rt]=(T[rt]+lazy[v]);
 lazy[rt]=(lazy[rt]+lazy[v]);
 lazv[v]=0:
 }
 void update(int v,int tl,int tr,int l,int r,int val){
  if(l>r or tl>r or tr<l) return;</pre>
  if(1<=t1 && tr<=r){
  T[v]=T[v]+val:
  lazy[v]=(lazy[v]+val);
  }else{
  if(tl==tr) return:
  push(v):
   int tm=(tl+tr)>>1,lf=v<<1,rt=lf^1;;</pre>
   update(lf.tl.tm.l.r.val):
   update(rt,tm+1,tr,l,r,val);
  T[v]=max(T[lf],T[rt]);
 int query(int v,int tl,int tr,int l,int r){
 if(1>r) return MX:
  if(1<=t1 && tr<=r) return T[v]:
  push(v):
  int tm=(tl+tr)>>1,lf=v<<1,rt=lf^1;</pre>
 return max(query(lf,tl,tm,l,min(r,tm)),query(rt,tm+1,tr,
       \max(1, tm+1), r)):
 int a(int 1.int r){
 return query(1,0,N-1,1,r);
 void u(int l.int r.int val){
 update(1,0,N-1,1,r,val);
} st:
struct hld{
```

```
int n.t:
v32 sz,in,out,root,par,depth;
vv32 g;
SegTree tree:
void dfs_sz(int v=0,int p=0){
sz[v]=1:
for(auto &u: g[v]){
 if(u==p) continue;
 dfs_sz(u,v);
 sz[v] += sz[u];
 if(sz[u]>sz[g[v][0]]) swap(u, g[v][0]);
void dfs_hld(int v=0,int p=0){
   in[v]=t++:
   par[v]=p;
   depth[v]=depth[p]+1;
   for(auto u: g[v]){
    if(u==p) continue;
       root[u] = (u == g[v][0] ? root[v]:u);
       dfs_hld(u,v);
   out[v]=t;
void pre(vv32 &v){
g=v;n=v.size();t=0;
sz.assign(n,0);in.assign(n,0);out.assign(n,0);
root.assign(n.0):par.assign(n.0):depth.assign(n.0):
depth[0]=-1;
dfs sz():dfs hld():
tree.clear(n,-MOD);
template <class BinaryOperation>
void processPath(int u,int v,BinaryOperation op){
for(;root[u]!=root[v];v=par[root[v]]){
 if(depth[root[u]] > depth[root[v]]) swap(u,v);
 op(in[root[v]],in[v]);
if(depth[u]>depth[v]) swap(u,v);
op(in[u],in[v]);
void modifyPath(int u,int v,const int &value){
   processPath(u.v.[this.&value](int l.int r){tree.u(l.r.
        value):}): // [1.r]
 void modifvSubtree(int u.const int &value){
  tree.u(in[u],out[u]-1,value);
 int queryPath(int u,int v){
   int res=-MOD:
```

16 Hopcraft Karp

```
// Max matching
//1 indexed Hopcroft-Karp Matching in O(E sgrtV)
struct Hopcroft_Karp{
static const int inf = 1e9;
int n:
vector<int> matchL, matchR, dist;
vector<vector<int> > g;
Hopcroft Karp(int n):n(n).matchL(n+1).matchR(n+1).dist(n+1)
     g(n+1)
void addEdge(int u, int v){
 g[u].pb(v);
bool bfs(){
 queue<int> q;
 for(int u=1:u<=n:u++){
 if(!matchL[u]){
   dist[u]=0;
   q.push(u);
  }else dist[u]=inf:
 dist[0]=inf:
 while(!q.empty()){
  int u=q.front();
  q.pop();
  for(auto v:g[u]){
   if(dist[matchR[v]] == inf){
    dist[matchR[v]] = dist[u] + 1;
    q.push(matchR[v]):
 }
 return (dist[0]!=inf):
bool dfs(int u){
 if(!u) return true:
 for(auto v:g[u]){
```

```
if(dist[matchR[v]] == dist[u]+1 &&dfs(matchR[v])){
   matchL[u]=v:
   matchR[v]=u:
   return true:
  }
 dist[u]=inf;
 return false;
int max_matching(){
 int matching=0;
 while(bfs()){
  for(int u=1;u<=n;u++){</pre>
   if(!matchL[u])
    if(dfs(u)) matching++;
  }
 }
 return matching;
}
};
```

17 Hungarian Algorithm

```
struct Hungarian{
 //Important: cost matrix a[1..n][1..m]>=0,n<=m (works with
       negative costs)
 // O(V^3) Use p to find matching of 1..m
 vv64 a:
 v64 u,v;
 v32 p,way;
 int n.m:
 Hungarian(int n, int m): n(n), m(m), u(n+1,0), v(m+1,0), p(m+1,0)
      +1,0), way (m+1,0), a (n+1,v64(m+1,0)){}
 void addEdge(int u,int v,ll val){
   a[u][v]=val;
 11 solveAssignmentProblem(){
   for(int i=1:i<=n:++i){</pre>
     p[0]=i;
     int j0=0;
     v64 minv(m+1,2e17+10);
     vector<bool> used(m+1,0);
     dof.
       used[i0] = true:
       int i0=p[j0];
       ll delta=2e17+10;
       int j1=0;
       for(int j=1;j<=m;++j){</pre>
```

```
if(!used[j]){
           11 cur=a[i0][j]-u[i0]-v[j];
           if(cur<minv[j]){</pre>
             minv[j]=cur;
             wav[i]=i0;
           if(minv[j]<delta){</pre>
             delta=minv[j];
             j1=j;
       for(int j=0;j<=m;++j){</pre>
         if(used[j]){
           u[p[i]]+=delta;
           v[i]-=delta;
         }else minv[j]-=delta;
       }
       j0=j1;
     }while(p[j0]!=0);
       int j1=way[j0];
       p[j0]=p[j1];
       j0=j1;
     }while(j0!=0);
   return -v[0];
};
```

18 Linear Sieve

```
int mu[LIM],is_com[LIM];
v32 pr;
void sieve(){
    mu[1]=1;
    forsn(i,2,LIM){
        if(!is_com[i]) pr.pb(i),mu[i]=-1;
        forstl(it,pr){
            if(it*i>=LIM) break;
            is_com[i*it]=1;
            if(i%it==0){
                mu[i*it]=0;
                break;
        }else{
                mu[i*it]=mu[i]*mu[it];
        }
}
```

```
}
```

19 Lowest Common Ancestor

```
vv32 v:
v32 tin,tout,dist;
vv32 up;
int 1:
void dfs(int i,int par,int lvl){
   tin[i]= ++t:
   dist[i]= lvl:
   up[i][0] = par;
   forsn(j,1,1+1) up[i][j] = up[up[i][j-1]][j-1];
   forstl(it,v[i]) if(it!=par) dfs(it,i,lvl+1);
   tout[i] = ++t;
bool is_ancetor(int u, int v){
   return tin[u] <=tin[v] && tout[u] >=tout[v]:
int lca(int u, int v){
   if (is_ancetor(u, v)) return u;
   if (is_ancetor(v, u)) return v;
   rforn(i,1) if(!is_ancetor(up[u][i], v)) u=up[u][i];
   return up[u][0];
int get_dis(int u,int v){
   int lcauv=lca(u.v):
   return dist[u]+dist[v]-2*dist[lcauv];
void preprocess(int root){
   tin.resize(n);
   tout.resize(n):
   dist.resize(n);
   t=0:
   l=ceil(log2((double)n));
   up.assign(n,v32(1+1));
   dfs(root,root,0);
```

20 Lucas Theorem

```
//Lucas Theorem: Find (n Choose m) mod p for prime p and
    large n,m. in O(log(m*n))
// nCm mod p by lucas theorem for large n,m >=0
```

21 Manacher

```
Manacher
// Given a string s of length N, finds all palindromes as
     its substrings.
// p[0][i] = half length of longest even palindrome around
// p[1][i] = longest odd at i (half rounded down i.e len 2*x
    +1).
//Time: O(N)
void manacher(const string& s){
int n=s.size();
v32 p[2] = \{v32(n+1), v32(n)\};
forn(z,2) for(int i=0,1=0,r=0;i<n;++i){
int t=r-i+!z;
if(i<r) p[z][i]=min(t,p[z][1+t]);</pre>
int L=i-p[z][i],R=i+p[z][i]-!z;
while (L>=1 && R+1<n && s[L1]==s[R+1]) p[z][i]++,L--,R++;
if(R>r) 1=I..r=R:}}
```

22 Matrix

```
int MOD1=MOD;
inline ll add(ll a,ll b){
  return (a+b)%MOD1;
}
inline ll mult(ll a,ll b){
  return a*b%MOD1;
}
struct matrix{
  int arr[105][105]={0};
```

```
int SZ:
void reset(int sz){
 S7=sz:
 //memset(arr.0.sizeof(arr)):
void makeiden(int sz){
 reset(sz):
 for(int i=0;i<SZ;i++){</pre>
  arr[i][i]=1:
matrix operator +(const matrix &o)const{
 matrix res;
 res.reset(SZ):
 for(int i=0:i<SZ:i++){</pre>
  for(int j=0;j<SZ;j++){</pre>
  res.arr[i][j]=add(arr[i][j],o.arr[i][j]);
  }
 }
 return res:
matrix operator *(const matrix &o)const{
 matrix res:
 res.reset(SZ):
 for(int i=0:i<SZ:i++){</pre>
  for(int j=0;j<SZ;j++){</pre>
   res.arr[i][j]=0;
   for(int k=0:k<SZ:k++){</pre>
    res.arr[i][j]=add(res.arr[i][j],mult(arr[i][k],o.arr[k][
  }
 return res;
matrix mpower(matrix a,int sz,ll b){
matrix res:
res.makeiden(sz);
while(b){
 if(b&1){
  res=res*a;
 a=a*a;
 b>>=1;
return res;
```

23 Merge Sort Tree

```
// Merge sort Tree
const int MAXN=1e5+5;
v32 T[4*MAXN]; // nlogn memory
void build(int a[],int v,int tl,int tr){
if(tl==tr){
T[v]=v32(1.a[t1]):
}else{
 int tm=(tl+tr)>>1:
 build(a, v << 1, t1, tm);
 build(a,(v<<1)^1,tm+1,tr);
 merge(all(T[v<<1]),all(T[(v<<1)^1]),back_inserter(T[v]));</pre>
 // built in combine in sorted order (2pointer)
// number of numbers <=x in [1,r]</pre>
int querv(int v.int tl.int tr.int l.int r.int x){
if(1>r) return 0;
if(1<=t1 && tr<=r){</pre>
 return upper_bound(all(T[v]),x)-T[v].begin();
int tm=(tl+tr)>>1:
return query(v<<1,tl,tm,l,min(r,tm),x)+query((v<<1)^1,tm+1,</pre>
     tr,max(1,tm+1),r,x);
// Number of distinct integers in [1,r]
void convert(int a[],int n){ // b store next occ index
m32 m; // Can be replaced by vv32 in small numbers
rforn(i,n-1){
 auto it=m.find(a[i]);
 if(it==m.end()) b[i]=MOD;
 else b[i]=it->se;
 m[a[i]]=i;
build(b.1.0.n-1):
inline int q(int 1,int r){ // no. of val in [1,r] with nxt
return (r-l+1)-query(1,0,n-1,1,r,r);
```

24 Miller Rabin

```
using u64 = uint64_t;
using u128 = __uint128_t;
```

```
u64 binpower(u64 base, u64 e, u64 mod) {
   u64 result = 1:
   base %= mod:
   while (e) {
       if (e & 1)
           result = (u128)result * base % mod;
       base = (u128)base * base % mod:
       e >>= 1:
   return result:
}
bool check_composite(u64 n, u64 a, u64 d, int s) {
   u64 x = binpower(a, d, n);
   if (x == 1 | | x == n - 1)
       return false:
   for (int r = 1; r < s; r++) {</pre>
       x = (u128)x * x % n:
       if (x == n - 1)
           return false;
   return true;
};
bool MillerRabin(u64 n) { // returns true if n is prime,
     else returns false.
   if (n < 2)
       return false;
   int r = 0:
   u64 d = n - 1;
   while ((d & 1) == 0) {
       d >>= 1;
       r++:
   for (int a: {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
        37}) {
       if (n == a)
           return true:
       if (check_composite(n, a, d, r))
           return false:
   return true;
```

25 Min Cost Max Flow

```
// Mincost Maxflow : O(E^2)
// [Hell-Johnson MinCostMaxFlow using Dijstra with potential
     & Fibonnaci Heap]
// Negative cost cycles are not supported.
struct MCMF{
 struct Edge{
   int u.v.rind:
   FLOW cap,flow;
   COST cost:
 };
 int N;
 vector<COST> pot,dist;
 vector<vector<Edge> > v;
 vector<pair<int,int> > par;
 MCMF(int n): N(n),dist(n),v(n),par(n){}
 void AddEdge(int to,int from,int cap,int cost){
   if(to==from){
     assert(cost>=0);
     return:
   int i1=v[to].size(),i2=v[from].size();
   v[to].push_back({to,from,i2,cap,0,cost});
   v[from].push_back({from,to,i1,0,0,-cost});
 void setpi(int s){
   pot.assign(N,CINF);
   pot[s]=0;
   int ch=1,ite=N;
   COST cur, nw;
   while(ch-- && ite--){
    for(int i=0:i<N:++i){</pre>
      if(pot[i]!=CINF){
        cur=pot[i]:
        for(auto &e: v[i]){
          if(e.cap>0 && (nw=cur+e.cost)<pot[e.v]){</pre>
            pot[e.v]=nw; ch=1;
      }
     }
   assert(ite>=0); // Else negative cycle
 bool path(int s.int t){
   fill(dist.begin(),dist.end(),CINF);
   dist[s]=0:
   __gnu_pbds::priority_queue<pair<COST,int> > pq;
   vector<decltype(pq)::point_iterator> its(N);
```

```
pq.push({0,s});
 COST curr.val:
 int node.cnt:
 bool ok=0:
  while(!pq.empty()){
   tie(curr.node)=pg.top():
   pq.pop();
   curr=-curr;
   if(curr!=dist[node]) continue:
   curr+=pot[node];
   if(node==t) ok=1:
   cnt=0:
   for(auto &e: v[node]){
     if(e.cap>e.flow && (val=curr+e.cost-pot[e.v]) < dist[e.</pre>
       dist[e.v]=val;
       par[e.v] = make_pair(node, cnt);
       if(its[e.v] == pq.end()) its[e.v] = pq.push({-val,e.v})
           else pq.modify(its[e.v],{-val,e.v});
     }
     ++cnt:
 for(int i=0:i<N:++i){</pre>
   pot[i]=min(pot[i]+dist[i],FINF);
 return ok:
pair<FLOW.COST> SolveMCMF(int s.int t.FLOW need=FINF.bool
    neg=0){
 FLOW tot=0,cflow=0; COST tcost=0;
 if(s==t) return {tot.tcost}:
 if(!neg) pot.assign(N,0);
 else setpi(s);
  int cntr=0:
  while(path(s,t) && need>0){
   cflow=need:
   for(int node=t,u,ind;node!=s;node=u){
     u=par[node].first:
     ind=par[node].second:
     cflow=min(cflow,v[u][ind].cap-v[u][ind].flow);
    tot+=cflow: need-=cflow:
   for(int node=t,u,ind,rind;node!=s;node=u){
     u=par[node].first:
     ind=par[node].second;
     rind=v[u][ind].rind:
     v[u][ind].flow+=cflow:
     v[node][rind].flow-=cflow:
```

```
}
return {tot,tcost};
}
```

26 Nearest Pair of Points

```
vector<pt> t;
void rec(int 1, int r) {
   if (r - 1 \le 3) {
       for (int i = 1: i < r: ++i) {</pre>
           for (int j = i + 1; j < r; ++j) {
              upd_ans(a[i], a[j]);
       sort(a.begin() + 1, a.begin() + r, cmp v()):
   int m = (1 + r) >> 1;
   int midx = a[m].x:
   rec(1, m):
   rec(m, r);
   merge(a.begin() + 1, a.begin() + m, a.begin() + m, a.
        begin() + r, t.begin(), cmp_v());
   copy(t.begin(), t.begin() + r - 1, a.begin() + 1);
   int tsz = 0:
   for (int i = 1: i < r: ++i) {</pre>
       if (abs(a[i].x - midx) < mindist) {</pre>
           for (int j = tsz - 1; j >= 0 && a[i].y - t[j].y <</pre>
                 mindist: --i)
              upd_ans(a[i], t[j]);
           t[tsz++] = a[i]:
// In main, call as:
t.resize(n):
sort(a.begin(), a.end(), cmp_x());
mindist = 1E20;
rec(0, n);
```

27 Number Theoretic Transform

```
const int mod=998244353:
// 998244353=1+7*17*2^23 : g=3
// 1004535809=1+479*2^21 : g=3
// 469762049=1+7*2^26 : g=3
// 7340033=1+7*2^20 : g=3
// For below change mult as overflow:
// 10000093151233=1+3^3*5519*2^26 : g=5
// 1000000523862017=1+10853*1373*2^26 : g=3
// 100000000949747713=1+2<sup>2</sup>9*3*73*8505229 : g=2
// For rest find primitive root using Shoup's generator
// root_pw: power of 2 >= maxn, Mod-1=k*root_pw => w =
template<long long Mod.long long root pw.long long primitive
struct NTT{
inline long long powm(long long x,long long pw){
 if(abs(pw)>Mod-1) pw%=(Mod-1);
 if(pw<0) pw+=Mod-1;</pre>
 ll res=1:
 while(pw){
  if(pw&1LL) res=(res*x)%Mod:
  pw>>=1:
  x=(x*x)\Mod;
 return res:
 inline 11 inv(11 x){
    return powm(x,Mod-2);
}
11 root,root_1;
 NTT(){
 root=powm(primitive,(Mod-1)/root_pw);
 root_1=inv(root);
 void ntt(vector<long long> &a,bool invert){
 int n=a.size();
 for(long long i=1, j=0; i<n; i++){</pre>
  long long bit=n>>1;
  for(;j&bit;bit>>=1) j^=bit;
  j^=bit;
  if(i<j) swap(a[i],a[j]);</pre>
  for(long long len=2;len<=n;len<<=1){</pre>
  long long wlen= invert ? root_1:root;
  for(long long i=len;i<root_pw;i<<=1) wlen=wlen*wlen%Mod;</pre>
  for(long long i=0;i<n;i+=len){</pre>
```

```
long long w=1:
  for(long long j=0; j<len/2; j++){</pre>
   long long u=a[i+j], v=a[i+j+len/2]*w%Mod;
   a[i+i] = u+v < Mod ? u+v:u+v-Mod:
   a[i+j+len/2] = u-v>=0 ? u-v:u-v+Mod;
   w=w*wlen%Mod:
  }
 }
if(invert){
 ll n 1=inv(n):
 for(long long &x: a) x=x*n 1%Mod:
vector<long long> multiply(vector<long long> const& a,
    vector<ll> const& b){
vector<long long> fa(a.begin(),a.end()),fb(b.begin(),b.end
      ()):
int n=1:
 while(n<a.size()+b.size()) n<<=1:</pre>
point(fa,1,n);
point(fb.1.n):
 for(int i=0;i<n;++i) fa[i]=fa[i]*fb[i]%Mod;</pre>
coef(fa):
return fa;
void point(vector<long long> &A,bool not_pow=1,int atleast
    =-1){
if(not_pow){
 if(atleast==-1){
  atleast=1:
  while(atleast<A.size()) atleast<<=1;</pre>
 A.resize(atleast,0);
ntt(A.0):
void coef(vector<long long> &A,bool reduce=1){
ntt(A,1);
if(reduce) while(A.size() and A.back()==0) A.pop back();
void point_power(vector<long long> &A,long long k){
for(long long &x: A) x=powm(x,k);
void coef_power(vector<long long> &A,int k){
while(A.size() and A.back()==0) A.pop_back();
int n=1;
 while(n<k*A.size()) n<<=1:</pre>
point(A.1.n):
 point_power(A,k);
```

```
coef(A);
}
vector<long long> power(vector<long long> a,ll p){
  while(a.size() and a.back()==0) a.pop_back();
  vector<long long> res;
  res.pb(1);
  while(p){
   if(p&1) res=multiply(res,a);
    a=multiply(a,a);
   p/=2;
  }
  return res;
}
NTT<mod,1<<20,3> ntt;
```

28 Ordered Set

29 Polynomial

```
namespace algebra{
  const int inf = 1e9;
  const int magic = 500; // threshold for sizes to run the
      naive algo
namespace fft {
  const int maxn = 1 << 18;
  typedef double ftype;
  typedef complex<ftype> point;
```

```
point w[maxn]:
const ftype pi = acos(-1);
bool initiated = 0;
void init() {
if(!initiated) {
 for(int i = 1: i < maxn: i *= 2) {</pre>
  for(int j = 0; j < i; j++) {
   w[i + j] = polar(ftype(1), pi * j / i);
 initiated = 1;
template<tvpename T>
void fft(T *in, point *out, int n, int k = 1) {
if(n == 1) {
 *out = *in:
} else {
 n /= 2:
 fft(in, out, n, 2 * k);
 fft(in + k, out + n, n, 2 * k):
 for(int i = 0: i < n: i++) {</pre>
  auto t = out[i + n] * w[i + n];
  out[i + n] = out[i] - t:
  out[i] += t;
template<typename T>
void mul_slow(vector<T> &a, const vector<T> &b) {
 vector<T> res(a.size() + b.size() - 1);
 for(size t i = 0: i < a.size(): i++) {</pre>
 for(size_t j = 0; j < b.size(); j++) {</pre>
  res[i + i] += a[i] * b[i]:
 }}
 a = res;
template<typename T>
void mul(vector<T> &a, const vector<T> &b) {
if(min(a.size(), b.size()) < magic) {</pre>
 mul slow(a, b):
 return:}
 init():
 static const int shift = 15, mask = (1 << shift) - 1;</pre>
 size t n = a.size() + b.size() - 1:
 while( builtin popcount(n) != 1) n++;
 a.resize(n);
 static point A[maxn], B[maxn];
 static point C[maxn], D[maxn];
 for(size_t i = 0; i < n; i++) {</pre>
 A[i] = point(a[i] & mask, a[i] >> shift):
 if(i < b.size()) B[i] = point(b[i] & mask, b[i] >> shift)
 else B[i] = 0;
```

```
fft(A, C, n): fft(B, D, n):
  for(size t i = 0: i < n: i++) {</pre>
  point c0 = C[i] + conj(C[(n - i) \% n]);
  point c1 = C[i] - conj(C[(n - i) \% n]);
  point d0 = D[i] + conj(D[(n - i) \% n]);
  point d1 = D[i] - conj(D[(n - i) \% n]);
  A[i] = c0 * d0 - point(0, 1) * c1 * d1;
  B[i] = c0 * d1 + d0 * c1:
  fft(A, C, n); fft(B, D, n);
  reverse(C + 1, C + n):
  reverse(D + 1, D + n):
  int t = 4 * n;
  for(size t i = 0: i < n: i++) {</pre>
  int64_t A0 = llround(real(C[i]) / t);
  T A1 = llround(imag(D[i]) / t);
  T A2 = llround(imag(C[i]) / t);
  a[i] = A0 + (A1 << shift) + (A2 << 2 * shift);
 return: }}
template<typename T>
T bpow(T x, size_t n) {
return n ? n % 2 ? x * bpow(x, n - 1) : bpow(x * x, n / 2)
      : T(1):
template<typename T>
T bpow(T x, size_t n, T m) {
return n ? n % 2 ? x * bpow(x, n - 1, m) % m : bpow(x * x
     % m, n / 2, m) : T(1);
template<int m>
struct modular{
int64 t r:
modular() : r(0) {}
 modular(int64 t rr) : r(rr) \{if(abs(r) >= m) r \% = m : if(r)\}
     < 0) r += m:
 modular inv() const {return bpow(*this, m - 2);}
modular pow(int64_t k) const {return bpow(*this,k);}
modular operator * (const modular &t) const {return (r * t
      .r) % m:}
modular operator / (const modular &t) const {return *this
      * t.inv():}
modular operator += (const modular &t) {r += t.r: if(r >=
     m) r -= m; return *this;}
modular operator -= (const modular &t) {r -= t.r; if(r <</pre>
     0) r += m: return *this:}
modular operator + (const modular &t) const {return
     modular(*this) += t:}
modular operator - (const modular &t) const {return
     modular(*this) -= t:}
```

HTB

```
modular operator *= (const modular &t) {return *this = *
     this * t:}
modular operator /= (const modular &t) {return *this = *
     this / t:}
bool operator == (const modular &t) const {return r == t.r
bool operator != (const modular &t) const {return r != t.r
operator int64_t() const {return r;}
}:
template<int T>
istream& operator >> (istream &in. modular<T> &x) {
return in >> x.r;
template<typename T>
struct poly {
vector<T> a:
void normalize() { // get rid of leading zeroes
 while(!a.empty() && a.back() == T(0)) a.pop_back();}
polv(){}
poly(T a0) : a{a0}{normalize();}
poly(vector<T> t) : a(t){normalize();}
poly operator += (const poly &t) {
 a.resize(max(a.size(), t.a.size()));
 for(size_t i = 0; i < t.a.size(); i++) {</pre>
  a[i] += t.a[i]:
 normalize():
 return *this;
poly operator -= (const poly &t) {
 a.resize(max(a.size(), t.a.size()));
 for(size t i = 0: i < t.a.size(): i++) {</pre>
  a[i] -= t.a[i];
 normalize():
 return *this:
poly operator + (const poly &t) const {return poly(*this)
polv operator - (const polv &t) const {return polv(*this)
poly mod_xk(size_t k) const { // get same polynomial mod x
 k = min(k, a.size());
 return vector<T>(begin(a), begin(a) + k);
poly mul_xk(size_t k) const { // multiply by x^k
 polv res(*this):
 res.a.insert(begin(res.a), k, 0);
```

```
return res:
poly div_xk(size_t k) const { // divide by x^k, dropping
    coefficients
 k = min(k, a.size());
 return vector<T>(begin(a) + k, end(a));
poly substr(size_t l, size_t r) const { // return mod_xk(r
    ).div xk(1)
1 = min(1, a.size());
r = min(r, a.size()):
 return vector<T>(begin(a) + 1, begin(a) + r);
poly inv(size_t n) const { // get inverse series mod x^n
poly ans = a[0].inv():
 size_t a = 1;
 while(a < n)  {
 poly C = (ans * mod_xk(2 * a)).substr(a, 2 * a);
 ans -= (ans * C).mod xk(a).mul xk(a):
return ans.mod xk(n):
poly operator *= (const poly &t) {fft::mul(a, t.a);
    normalize(); return *this;}
poly operator * (const poly &t) const {return poly(*this)
polv reverse(size t n. bool rev = 0) const { // reverses
    and leaves only n terms
 polv res(*this):
 if(rev) { // If rev = 1 then tail goes to head
 res.a.resize(max(n, res.a.size()));
 std::reverse(res.a.begin(), res.a.end());
 return res.mod xk(n):
pair<poly, poly> divmod_slow(const poly &b) const { //
    when divisor or quotient is small
 vector<T> A(a);
 vector<T> res:
 while(A.size() >= b.a.size()) {
 res.push_back(A.back() / b.a.back());
 if(res.back() != T(0)) {
  for(size_t i = 0; i < b.a.size(); i++) {</pre>
   A[A.size() - i - 1] = res.back() * b.a[b.a.size() - i
        - 1]:
 A.pop_back();
```

```
std::reverse(begin(res), end(res));
 return {res. A}:
pair<poly, poly> divmod(const poly &b) const { // returns
     quotiend and remainder of a mod b
 if(a.size() < b.a.size()) {</pre>
 return {poly{0}, *this};
 int d = a.size() - b.a.size();
 if(min(d, b.a.size()) < magic) {</pre>
 return divmod slow(b):
 poly D = (reverse(d + 1) * b.reverse(d + 1).inv(d + 1)).
     mod xk(d + 1).reverse(d + 1, 1):
 return {D. *this - D * b}:
poly operator / (const poly &t) const {return divmod(t).
     first:}
poly operator % (const poly &t) const {return divmod(t).
poly operator /= (const poly &t) {return *this = divmod(t)
     .first:}
poly operator %= (const poly &t) {return *this = divmod(t)
poly operator *= (const T &x) {
for(auto &it: a) {
 it. *= x:
 normalize();
 return *this:
poly operator /= (const T &x) {
for(auto &it: a) {
 it /= x;
 normalize():
 return *this:
poly operator * (const T &x) const {return poly(*this) *=
poly operator / (const T &x) const {return poly(*this) /=
T operator [](int idx) const {
return idx >= (int)a.size() || idx < 0 ? T(0) : a[idx];</pre>
T& coef(size t idx) { // mutable reference at coefficient
return a[idx];
bool operator == (const poly &t) const {return a == t.a;}
bool operator != (const poly &t) const {return a != t.a;}
```

HTB 17

```
polv deriv() { // calculate derivative
 vector<T> res:
for(int i = 1; i <= a.size(); i++) {</pre>
 res.push back(T(i) * a[i]):
}
return res:
polv integr() { // calculate integral with C = 0
vector<T> res = {0}:
for(int i = 0; i <= a.size(); i++) {</pre>
 res.push back(a[i] / T(i + 1));
return res;
size_t leading_xk() const { // Let p(x) = x^k * t(x),
     return k
int res = 0:
 while(a[res] == T(0)) {
 res++:
}
 return res;
polv log(size_t n) { // calculate log p(x) mod x^n
assert(a[0] == T(1)):
return (deriv().mod_xk(n) * inv(n)).integr().mod_xk(n);
poly exp(size_t n) { // calculate exp p(x) mod x^n
if(a.emptv()) return T(1);
 assert(a[0] == T(0));
polv ans = T(1):
 size_t a = 1;
 while(a < n) {
 poly C = ans.log(2 * a).div_xk(a) - substr(a, 2 * a);
 ans -= (ans * C).mod_xk(a).mul_xk(a);
 a *= 2:
 return ans.mod_xk(n):
poly pow_slow(size_t k, size_t n) { // if k is small
return k ? k % 2 ? (*this * pow_slow(k - 1, n)).mod_xk(n)
      : (*this * *this).mod xk(n).pow slow(k / 2, n) : T
poly pow(size_t k, size_t n) { // calculate p^k(n) mod x^n
if(a.empty()) return *this;
if(k < magic) return pow_slow(k, n);</pre>
int i = leading_xk();
T i = a[i]:
poly t = div_xk(i) / j;
```

```
return bpow(j, k) * (t.log(n) * T(k)).exp(n).mul_xk(i * k | 30 Primitive Root.
      ).mod xk(n):
 }
};
template<typename T>
poly<T> operator * (const T& a. const poly<T>& b) {
 return b * a:
using namespace algebra;
const int mod = 998244353:
const int lim = 2e6+5:
typedef modular<mod> base;
typedef poly<base> polyn;
base fact[lim],invfact[lim];
void pre(){
fact[0]=1:
   invfact[lim-1]=367642781;
   for(int i=1;i<lim;++i) fact[i]=fact[i-1]*base(i);</pre>
   for(int i=lim-1:i>0:--i) invfact[i-1]=invfact[i]*base(i):
void example(){
   int t,n,k;
   cin>>t:
   while(t--){
    cin>>n>>k:
    if(k==1){
     base nk=1,N=n,Ans;
     for(int i=n-2;i>=0;--i){
      base ans=nk*invfact[i]*invfact[n-i]*invfact[n-2-i]:
      Ans+=ans:
      if(i) nk*=N;
     Ans=Ans*fact[n-2]*fact[n]/nk;
     cout<<Ans<<endl:
    }else{
     vector<base> a(n+1):
     for(int i=0;i<=n;++i){</pre>
      a[i]=base(i+1).pow(k)*invfact[i];
     polvn A(a), B(A, pow(n,n)):
     base Ans=(B[n-2]/base(n).pow(n-2))*fact[n-2];
     cout<<Ans<<endl:
   }
```

```
// Primitive root Exist for n=1,2,4,(odd prime power),2*(odd
      prime power)
// O(Ans.log(p).logp + sqrt(phi)) \le O((log p)^8 + root(p))
// Change phi when not prime
// Include powm (inverse)
ll phi cal(ll n){
ll result=n:
for(11 i=2:i*i<=n:++i){</pre>
 if(n%i==0){
  while(n%i==0) n/=i;
  result-=result/i:
if(n>1) result-=result/n:
return result;
11 generator(ll p){
v64 fact:
ll phi=p-1; // Call phi_cal if not prime
11 n=phi:
for(ll i=2:i*i<=n:++i){</pre>
 if(n%i==0){
  fact.push_back(i);
  while(n%i==0) n/=i:
}
if(n>1) fact.push_back(n);
for(11 res=2;res<=p;++res){</pre>
 bool ok=true:
 for(size t i=0:i<fact.size() && ok:++i)</pre>
  ok&=(powm(res,phi/fact[i],p)!=1);
 if(ok) return res;
return -1;
```

31 Push Relabel

```
//Push-Relabel Algorithm for Flows - Gap Heuristic,
    Complexity: O(V^3)
//To obtain the actual flow values, look at all edges with
    capacity > 0
//Zero capacity edges are residual edges
struct edge{
int from, to, cap, flow, index;
```

```
edge(int from, int to, int cap, int flow, int index):
 from(from), to(to), cap(cap), flow(flow), index(index) {}
}:
struct PushRelabel{
int n;
 vector<vector<edge> > g:
 vector<long long> excess;
 vector<int> height,active,count;
 queue<int> Q:
 PushRelabel(int n): n(n),g(n),excess(n),height(n),active(n)
      .count(2*n) {}
 void addEdge(int from, int to, int cap){
 g[from].push_back(edge(from,to,cap,0,g[to].size()));
 if(from==to) g[from].back().index++;
 g[to].push_back(edge(to,from,0,0, g[from].size()-1));
 void enqueue(int v){
 if(!active[v] && excess[v]>0){
  active[v]=true:
  Q.push(v);
 }
 void push(edge &e){
 int amt=(int)min(excess[e.from],(long long)e.cap - e.flow)
 if(height[e.from] <=height[e.to] || amt==0) return;</pre>
 e.flow += amt:
 g[e.to][e.index].flow -= amt:
 excess[e.to] += amt;
 excess[e.from] -= amt:
 enqueue(e.to);
 void relabel(int v){
 count[height[v]]--;
 int d=2*n:
 for(auto &it:g[v]){
  if(it.cap-it.flow>0) d=min(d, height[it.to]+1);
 height[v]=d;
 count[height[v]]++;
 enqueue(v):
 void gap(int k){
 for(int v=0;v<n;v++){</pre>
  if(height[v]<k) continue;</pre>
  count[height[v]]--:
  height[v]=max(height[v], n+1);
  count[height[v]]++;
  enqueue(v);
 }
```

```
void discharge(int v){
 for(int i=0; excess[v]>0 && i<g[v].size(); i++) push(g[v][</pre>
 if(excess[v]>0){
  if(count[height[v]]==1) gap(height[v]);
  else relabel(v);
 }
long long max_flow(int source, int dest){
 count[0] = n-1:
 count[n] = 1:
 height[source] = n;
 active[source] = active[dest] = 1;
 for(auto &it:g[source]){
  excess[source]+=it.cap;
  push(it):
 while(!Q.emptv()){
  int v=Q.front():
  Q.pop();
  active[v]=false:
  discharge(v);
 long long max_flow=0;
 for(auto &e:g[source]) max_flow+=e.flow;
 return max flow:
};
```

32 Rabin Miller

```
ull mulm(ull a,ull b,ull MOD){
  ull res=0;
  a%=MOD,b%=MOD;
  while(b){
    if(b&1LL) res=(res+a)%MOD;
    b>>=1;
    a=(a+a)%MOD;
}
  return res;
}
ull powm(ull x,ull pw,ull MOD){ //return x^pw % MOD
    x%=MOD;
  ull res=1;
  while(pw){
    if(pw&1LL) res=mulm(res,x,MOD);
    pw>>=1;
```

```
x=mulm(x.x.MOD):
 return res;
inline ull inv(ull x,ull MOD){
 return powm(x,MOD-2,MOD):
bool prime(ull p){
 if(p==2) return 1;
 if(p==1 || !(p&1LL)) return 0;
 ull s=p-1:
 while(!(s&1LL)) s/=2:
 forn(i,15){
   ull a=rand()%(p-1)+1,tmp=s;
   ull mod=powm(a,tmp,p);
   while(tmp!=p-1 && mod!=1 && mod!=p-1){
    mod=mulm(mod,mod,p);
     tmp<<=1;
   if(mod!=p-1 && !(tmp&1LL)) return 0;
 return 1:
```

33 Simplex

```
// Two-phase simplex algorithm for solving linear programs
    of the form
11
11
      maximize c^T x
11
      subject to Ax <= b
11
                  x >= 0
11
// INPUT: A -- an m x n matrix
        b -- an m-dimensional vector
11
        c -- an n-dimensional vector
        x -- a vector where the optimal solution will be
    stored
// OUTPUT: value of the optimal solution (infinity if
    unbounded
11
          above, nan if infeasible)
// To use this code, create an LPSolver object with A. b.
// arguments. Then, call Solve(x).
#include <iostream>
```

```
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std:
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9:
struct LPSolver {
 int m. n:
 VI B, N;
 VVD D;
 LPSolver(const VVD &A, const VD &b, const VD &c) :
   m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2), VD(n + 1)
   for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D
        [i][j] = A[i][j];
   for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1;
         D[i][n + 1] = b[i];
   for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j];
   N[n] = -1; D[m + 1][n] = 1;
 void Pivot(int r. int s) {
   for (int i = 0; i < m + 2; i++) if (i != r)
    for (int j = 0; j < n + 2; j++) if (j != s)
      D[i][j] = D[r][j] * D[i][s] / D[r][s];
   for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] /= D[
        rl[sl:
   for (int i = 0: i < m + 2: i++) if (i != r) D[i][s] /= -D
        [r][s]:
   D[r][s] = 1.0 / D[r][s];
   swap(B[r], N[s]);
 bool Simplex(int phase) {
   int x = phase == 1 ? m + 1 : m;
   while (true) {
    int s = -1;
    for (int j = 0; j <= n; j++) {</pre>
      if (phase == 2 && N[j] == -1) continue;
      if (s == -1 || D[x][j] < D[x][s] || D[x][j] == D[x][s]
           ] && N[j] < N[s]) s = j;
     }
```

```
if (D[x][s] > -EPS) return true:
     int r = -1:
     for (int i = 0; i < m; i++) {</pre>
       if (D[i][s] < EPS) continue:</pre>
       if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] /
            D[r][s] ||
         (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s])
              && B[i] < B[r]) r = i;
     if (r == -1) return false;
     Pivot(r. s):
 }
 DOUBLE Solve(VD &x) {
   int r = 0:
   for (int i = 1: i < m: i++) if (D[i][n + 1] < D[r][n +
        11) r = i:
   if (D[r][n + 1] < -EPS) {
     Pivot(r. n):
     if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -</pre>
          numeric_limits<DOUBLE>::infinity();
     for (int i = 0; i < m; i++) if (B[i] == -1) {
       int s = -1:
       for (int j = 0; j <= n; j++)
         if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i</pre>
              [s] \&\& N[j] < N[s]) s = j;
       Pivot(i, s):
   if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity
   x = VD(n):
   for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][
   return D[m][n + 1]:
};
int main() {
  const int m = 4;
  const int n = 3:
 DOUBLE _A[m][n] = {
   \{ 6, -1, 0 \},
   \{-1, -5, 0\},\
   \{1, 5, 1\},\
   \{-1, -5, -1\}
 DOUBLE _b[m] = \{ 10, -4, 5, -5 \};
```

```
DOUBLE _c[n] = { 1, -1, 0 };

VVD A(m);

VD b(_b, _b + m);

VD c(_c, _c + n);

for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);

LPSolver solver(A, b, c);

VD x;

DOUBLE value = solver.Solve(x);

cerr << "VALUE: " << value << endl; // VALUE: 1.29032
 cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
 for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
 cerr << endl;
 return 0;
}</pre>
```

34 Stock Span

```
void stockspan(v32 &d,int n){
// smallest index j such that j>i and d[i]>d[i]
nxt.assign(n.n):
// largest index j such that j<i and d[j]>=d[i]
pre.assign(n,-1);
int stk[n+10],ptr=0;
forn(i,n){
 while(ptr && d[i]>d[stk[ptr-1]]) nxt[stk[--ptr]]=i;
 pre[i] = (ptr? stk[ptr-1]:-1);
 stk[ptr++]=i:
v32 nxt.pr: // nxt index with val < cur and prev index with
    val < cur
void stockspan(v32 &d,int n){
int stk[n+10].ptr=0:
pr.assign(n,-1);
forn(i,n){
 while(ptr && d[i]<=d[stk[ptr-1]]) --ptr;</pre>
 pr[i] = (ptr? stk[ptr-1]:-1);
 stk[ptr++]=i;
ptr=0;
nxt.assign(n.n):
forn(i,n){
 while(ptr && d[i]<d[stk[ptr-1]]) nxt[stk[--ptr]]=i;</pre>
 stk[ptr++]=i:
```

35 Suffix Array

```
struct SuffixArray{
v32 a:
string s:
SuffixArray(const string& _s): s(_s+'\0'){ // e.g. s="aba
     \0" will have a=[3.2.0.1]
 int N=s.size():
 vector<pair<11,int> > b(N);
 a.resize(N):
 for(int i=0:i<N:++i){</pre>
  b[i].first=s[i]:
  b[i].second=i;
 int q=8;
 while((1<<q)<N) q++;</pre>
 for(int moc=0;;moc++){
  sort(all(b)):
  a[b[0].second]=0:
  for(int i=1:i<N:++i){</pre>
   a[b[i].second]=a[b[i-1].second]+(b[i-1].first!=b[i].first
        );
  }
  if((1<<moc)>=N) break;
  for(int i=0:i<N:++i){</pre>
   b[i].first=(l1)a[i]<<q;
   if(i+(1<<moc)<N) b[i].first+=a[i+(1<<moc)];</pre>
   b[i].second=i:
 for(int i=0:i<N:++i) a[i]=b[i].second:</pre>
v32 lcp(){ // longest common prefixes:res[i]=lcp(a[i],a[i
     -1]) e.g. s="aba\0" will have res=[0.0.1.0]
 int n=a.size(),h=0;
 v32 inv(n).res(n):
 for(int i=0;i<n;++i) inv[a[i]]=i;</pre>
 for(int i=0:i<n:++i){</pre>
  if(inv[i]>0){
   int p0=a[inv[i]-1];
   while(s[i+h]==s[p0+h]) h++;
   res[inv[i]]=h:
   if(h>0) h--;
 }
 return res;
```

36 Suffix Automaton

} };

```
struct state {
   int len. link:
   map<char, int> next;
const int MAXLEN = 100000;
state st[MAXLEN * 2]:
int sz, last;
void sa init() {
   st[0].len = 0:
   st[0].link = -1:
   sz++;
   last = 0;
void sa extend(char c) {
   int cur = sz++:
   st[cur].len = st[last].len + 1;
   int p = last:
   while (p != -1 && !st[p].next.count(c)) {
      st[p].next[c] = cur;
      p = st[p].link;
   if (p == -1) {
      st[cur].link = 0:
   } else {
      int a = st[p].next[c]:
      if (st[p].len + 1 == st[q].len) {
          st[cur].link = q;
      } else {
          int clone = sz++;
          st[clone].len = st[p].len + 1;
          st[clone].next = st[a].next:
          st[clone].link = st[q].link;
          while (p != -1 && st[p].next[c] == q) {
              st[p].next[c] = clone;
              p = st[p].link;
          st[q].link = st[cur].link = clone;
   }
   last = cur;
```

37 Suffix Tree

```
string s;
int n:
struct node {
   int 1, r, par, link;
   map<char,int> next;
   node (int 1=0, int r=0, int par=-1)
      : 1(1), r(r), par(par), link(-1) {}
   int len() { return r - 1: }
   int &get (char c) {
       if (!next.count(c)) next[c] = -1;
       return next[c]:
   }
};
node t[MAXN]:
int sz:
struct state {
   int v, pos;
   state (int v, int pos) : v(v), pos(pos) {}
state ptr (0, 0);
state go (state st, int 1, int r) {
   while (1 < r)
       if (st.pos == t[st.v].len()) {
          st = state (t[st.v].get(s[1]), 0);
          if (st.v == -1) return st:
      }
       else {
          if (s[ t[st.v].l + st.pos ] != s[l])
              return state (-1, -1);
          if (r-1 < t[st.v].len() - st.pos)</pre>
              return state (st.v, st.pos + r-1);
          1 += t[st.v].len() - st.pos;
          st.pos = t[st.v].len();
      }
   return st;
int split (state st) {
   if (st.pos == t[st.v].len())
       return st.v;
```

```
if (st.pos == 0)
       return t[st.v].par;
   node v = t[st.v];
   int id = sz++:
   t[id] = node (v.1, v.1+st.pos, v.par);
   t[v.par].get(s[v.1]) = id:
   t[id].get( s[v.l+st.pos] ) = st.v;
   t[st.v].par = id;
   t[st.v].1 += st.pos;
   return id;
}
int get_link (int v) {
   if (t[v].link != -1) return t[v].link;
   if (t[v].par == -1) return 0;
   int to = get_link (t[v].par);
   return t[v].link = split (go (state(to,t[to].len()), t[v
        ].1 + (t[v].par==0), t[v].r));
void tree_extend (int pos) {
   for(::) {
       state nptr = go (ptr, pos, pos+1);
       if (nptr.v != -1) {
           ptr = nptr;
           return:
       int mid = split (ptr);
       int leaf = sz++:
       t[leaf] = node (pos, n, mid);
       t[mid].get( s[pos] ) = leaf;
       ptr.v = get_link (mid);
       ptr.pos = t[ptr.v].len();
       if (!mid) break:
}
void build tree() {
   for (int i=0; i<n; ++i)</pre>
       tree extend (i):
```

Template

```
#pragma GCC optimize ("-02")
```

```
#pragma GCC optimize("Ofast")
// *#pragma GCC target("sse,sse2,sse3,sse4,popcnt,abm
     ,mmx,avx,tune=native")
// ~ #pragma GCC optimize("unroll-loops")
#include <bits/stdc++.h>
using namespace std:
#define fastio ios_base::sync_with_stdio(0);cin.tie(0);cout.
    tie(0)
#define pb push_back
#define mp make_pair
#define fi first
#define se second
#define all(x) x.begin(),x.end()
#define memreset(a) memset(a,0,sizeof(a))
#define testcase(t) int t:cin>>t:while(t--)
#define forstl(i,v) for(auto &i: v)
#define forn(i.e) for(int i=0:i<e:++i)</pre>
#define forsn(i,s,e) for(int i=s;i<e;++i)</pre>
#define rforn(i.s) for(int i=s:i>=0:--i)
#define rforsn(i.s.e) for(int i=s:i>=e:--i)
#define bitcount(a) __builtin_popcount(a) // set bits (add
    11)
#define ln '\n'
#define getcurrtime() cerr<<"Time = "<<((double)clock()/</pre>
    CLOCKS_PER_SEC) << end1
#define dbgarr(v,s,e) cerr<<#v<<" = "; forsn(i,s,e) cerr<<v[</pre>
    il<<". ": cerr<<endl
#define inputfile freopen("input.txt", "r", stdin)
#define outputfile freopen("output.txt", "w", stdout)
#define dbg(args...) { string _s = #args; replace(_s.begin()
    , _s.end(), ',', ''); \
stringstream _ss(_s); istream_iterator<string> _it(_ss); err
    ( it. args): }
void err(istream_iterator<string> it) { cerr<<endl; }</pre>
template<typename T, typename... Args>
void err(istream iterator<string> it. T a. Args... args) {
 cerr << *it << " = " << a << "\t"; err(++it, args...);</pre>
template<typename T1, typename T2>
ostream& operator <<(ostream& c,pair<T1,T2> &v){
c<<"("<<v.fi<<"."<<v.se<<")": return c:</pre>
template <template <class...> class TT, class ...T>
ostream& operator<<(ostream& out,TT<T...>& c){
   out<<"{ ";
   forstl(x,c) out<<x<" ";</pre>
   out<<"}"; return out;</pre>
typedef long long 11;
typedef unsigned long long ull;
```

```
typedef long double ld:
typedef pair<11,11> p64;
typedef pair<int,int> p32;
typedef pair<int,p32> p96;
typedef vector<11> v64;
typedef vector<int> v32:
typedef vector<v32> vv32;
typedef vector<v64> vv64;
typedef vector<p32> vp32;
typedef vector<p64> vp64;
typedef vector<vp32> vvp32;
typedef map<int.int> m32:
const int LIM=1e5+5,MOD=1e9+7;
const ld EPS = 1e-9:
mt19937 rng(chrono::steady_clock::now().time_since_epoch().
    count());
```

39 Topological Sort

```
int n: // number of vertices
vector<int> adj[LIM]; // adjacency list of graph
vector<bool> visited:
vector<int> ans:
void dfs(int v) {
   visited[v] = true;
   for (int u : adi[v]) {
       if (!visited[u])
           dfs(u):
    ans.push_back(v);
void topological_sort() {
   visited.assign(n, false);
    ans.clear():
   for (int i = 0; i < n; ++i) {</pre>
       if (!visited[i])
           dfs(i);
   reverse(ans.begin(), ans.end());
```

40 Treap

```
mt19937 rng(chrono::steady_clock::now().time_since_epoch().
     count()):
int getRand(int 1, int r){
    uniform_int_distribution<int> uid(1, r);
    return uid(rng);
}
struct Treap{
    struct data{
        int priority, val, cnt;
        data *1, *r;
        data(){
            val = 0, cnt = 0, l = NULL, r = NULL;
        data (int _val){
           val = _val, cnt = 1;
           1 = NULL, r = NULL;
           priority = getRand(1, 2e9);
       }
    typedef struct data* node;
    node head;
    Treap(): head(0) {}
    int cnt(node cur){
        return cur ? cur->cnt : 0:
    void updateCnt(node cur)
       if(cur)
            cur \rightarrow cnt = cnt(cur \rightarrow 1) + cnt(cur \rightarrow r) + 1;
    void push(node cur) //Lazy Propagation
    void combine(node &cur, node 1, node r)
       if(!1)
           cur = r;
            return;
       if(!r)
```

```
cur = 1:
       return;
   //Merge Operations like in segment tree
void reset(node &cur) //To reset other fields of cur
     except value and cnt
{
   if(!cur)
       return:
}
void operation(node &cur)
   if(!cur)
       return;
   reset(cur):
   combine(cur, cur->1, cur);
   combine(cur, cur, cur->r);
void splitPos(node cur, node &1, node &r, int k, int add
   if(!cur)
       return void(1 = r = 0):
   push(cur);
   int idx = add + cnt(cur->1):
   if(idx \le k)
       splitPos(cur->r, cur->r, r, k, idx + 1), l = cur;
       splitPos(cur\rightarrow 1, 1, cur\rightarrow 1, k, add), r = cur;
   updateCnt(cur):
   operation(cur):
void split(node cur, node &1, node &r, int k)
       return void(1 = r = 0);
   push(cur):
   int idx = cur -> val;
   if(idx \le k)
       split(cur->r, cur->r, r, k), l = cur;
       split(cur\rightarrow 1, 1, cur\rightarrow 1, k), r = cur;
   updateCnt(cur);
   operation(cur);
```

```
void merge(node &cur, node 1, node r)
   push(1);
   push(r):
   if(!1 || !r)
       cur = 1 ? 1 : r;
   else if(l->priority > r->priority)
       merge(1->r, 1->r, r), cur = 1;
       merge(r->1, 1, r->1), cur = r:
   updateCnt(cur);
   operation(cur);
void insert(int val)
   if(!head)
       head = new data(val);
       return:
   node 1, r, mid, mid2, rr;
   mid = new data(val);
   split(head, 1, r, val - 1);
   merge(1, 1, mid);
   split(r, mid2, rr, val);
   merge(head, 1, rr);
void erase(int val)
   node 1, r, mid;
   split(head, 1, r, val - 1);
   split(r, mid, r, val);
   merge(head, 1, r);
void inorder(node cur)
   if(!cur)
       return:
   push(cur);
   inorder(cur->1);
   cout<<cur->val<<" ":
   inorder(cur->r);
void inorder()
```

```
inorder(head);
   cout<<endl;
void clear(node cur)
   if(!cur)
       return:
   clear(cur->1), clear(cur->r);
   delete cur:
void clear()
   clear(head);
int find_by_order(int x) //1 indexed
   if(!x)
       return -1:
   x--;
   node 1, r, mid;
   splitPos(head, 1, r, x - 1);
   splitPos(r, mid, r, 0);
   int ans = -1:
   if(cnt(mid) == 1)
       ans = mid->val;
   merge(r, mid, r);
   merge(head, 1, r);
   return ans;
int order_of_key(int val) //1 indexed
   node 1, r, mid;
   split(head, 1, r, val - 1);
   split(r, mid, r, val);
   int ans = -1:
   if(cnt(mid) == 1)
       ans = 1 + cnt(1);
   merge(r, mid, r);
   merge(head, 1, r);
   return ans;
```

};

41 Tree Bridge

```
struct dsu{
   v32 par,rk;
   dsu(){}
   dsu(int n) {reset(n);}
   void reset(int n){
       rk.assign(n,0);
       par.resize(n);
       iota(all(par),0);
   int getpar(int i){
       return (par[i]==i)? i:(par[i]=getpar(par[i]));
   bool con(int i,int j){
       return getpar(i) == getpar(j);
   bool join(int i,int j){
       i=getpar(i), j=getpar(j);
       if(i==j) return 0;
       if(rk[i]>rk[j]) par[j]=i;
       else{
           par[i]=j;
           if(rk[i]==rk[j]) rk[j]++;
       return 1;
};
vector<vector<int> > getBridgeTree(vector<vector<int> > &v,
    int n){
int in[n]={0},low[n],ctm=0;
 vector<pair<int,int> > bridge;
 function<void(int,int)> dfs=[&](int u,int p){
 in[u]=low[u]=++ctm:
 for(auto &it: v[u]){
  if(it==p) continue;
  if(in[it]){
   if(low[it]<low[u]) low[u]=low[it];</pre>
  }else{
   dfs(it.u):
   if(low[it]<low[u]) low[u]=low[it];</pre>
   if(low[it]>in[u]) bridge.push_back({u,it});
   else d.join(u,it);
 }
for(int i=0;i<n;++i) if(!in[i]) dfs(i,i);</pre>
int par[n].id[n]:
for(int i=0;i<n;++i) par[i]=d.getpar(i),id[i]=-1;</pre>
```

```
vector<vector<int> > g;
for(auto &it: bridge){
  it.first=par[it.first],it.second=par[it.second];
  if(id[it.first]==-1){
   id[it.first]=g.size();
    g.push_back(vector<int>(0));
  }
  if(id[it.second]==-1){
   id[it.second]=g.size();
    g.push_back(vector<int>(0));
  }
  g[id[it.first]].push_back(id[it.second]);
  g[id[it.second]].push_back(id[it.first]);
  }
  return g;
}
// dfs find all brideges and g contain bridge tree rest is
    diameter calculation
```

42 Z Algorithm

```
// Z Algorithm
// Z[i] is the length of the longest substring starting from
// which is also a prefix of S
// O(n)
void z_func(v32 &s,v32 &z){
int L=0.R=0:
int sz=s.size();
z.assign(sz,0);
forsn(i.1.sz){
 if(i>R){
  L=R=i:
  while(R<sz && s[R-L]==s[R]) R++;</pre>
  z[i]=R-L; R--;
 }else{
  int k=i-L;
  if(z[k]<R-i+1) z[i]=z[k];</pre>
  elsef
   while(R<sz && s[R-L]==s[R]) R++:</pre>
   z[i]=R-L; R--;
 }
}
```

43 Z Ideas

```
Grav codes:
Applications:
1. Gray code of n bits forms a Hamiltonian cycle on a
hypercube, where each bit corresponds to one dimension.
2. Gray code can be used to solve the Towers of Hanoi
problem. Let n denote number of disks. Start with Grav
code of length n which consists of all zeroes (G(0)) and
move between consecutive Grav codes (from G(i) to G(i+1)).
Let i-th bit of current Gray code represent n-th disk
(the least significant bit corresponds to the smallest
disk and the most significant bit to the biggest disk).
Since exactly one bit changes on each step, we can treat
changing i-th bit as moving i-th disk. Notice that there
is exactly one move option for each disk (except the
     smallest one)
on each step (except start and finish positions). There
are always two move options for the smallest disk but
there is a strategy which will always lead to answer: if n
is odd then sequence of the smallest disk moves looks
like
            ftrftr
                          ... where f is the initial rod, t
    is
the terminal rod and r is the remaining rod), and if
                  frtfrt
n is even:
int gray (int n) {
   return n ^ (n >> 1);
int rev_g (int g) {
 int n = 0:
 for (; g; g >>= 1)
   n ^= g;
 return n:
Enumerating all submasks of a bitmask:
for (int s=m: : s=(s-1)&m) {
... vou can use s ...
if (s==0) break:
Sparse Table:
int st[MAXN][K];
for (int i = 0: i < N: i++)</pre>
   st[i][0] = array[i];
```

```
for (int j = 1; j <= K; j++)
   for (int i = 0; i + (1 << j) <= N; i++)</pre>
       st[i][j] = min(st[i][j-1], st[i + (1 << (j - 1))][j -
To get minimum of a range:
int j = log[R - L + 1];
int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);</pre>
Divide and Conquer DP:
Some dynamic programming problems have a recurrence of this
dp(i,j) = minkj \{dp(i1,k)+C(k,j)\}
where C(k,j) is some cost function.
Say 1 in and 1 jm , and evaluating C takes O(1) time.
Straightforward evaluation of the above recurrence is O(nm2)
There are nm states, and m transitions for each state.
Let opt(i,j) be the value of k that minimizes the above
    expression.
If opt(i,j) opt (i,j+1) for all i,j, then we can apply
divide-and-conquer DP. This known as the monotonicity
The optimal "splitting point" for a fixed i increases as i
    increases.
int n:
long long C(int i, int j);
vector<long long> dp_before(n), dp_cur(n);
// compute dp_cur[1], ... dp_cur[r] (inclusive)
void compute(int 1, int r, int optl, int optr)
   if (1 > r)
       return:
   int mid = (1 + r) >> 1;
   pair<long long, int> best = {INF, -1}:
   for (int k = optl; k <= min(mid, optr); k++) {</pre>
       best = min(best, {dp_before[k] + C(k, mid), k});
   dp_cur[mid] = best.first;
   int opt = best.second:
```

```
compute(1, mid - 1, optl, opt);
   compute(mid + 1, r, opt, optr);
Lyndon factorization: We can get the minimum cyclic shift.
Factorize the string as s = w1w2w3...wn
string min_cyclic_string(string s) {
   s += s:
   int n = s.size():
   int i = 0, ans = 0:
   while (i < n / 2) {
       ans = i:
       int j = i + 1, k = i;
       while (j < n \&\& s[k] <= s[j]) {
          if (s[k] < s[i])
              k = i;
           else
              k++:
           j++;
       while (i <= k)
          i += j - k;
   return s.substr(ans, n / 2);
Rank of a matrix:
const double EPS = 1E-9;
int compute rank(vector<vector<double>> A) {
   int n = A.size();
   int m = A[0].size():
   int rank = 0:
   vector<bool> row_selected(n, false);
   for (int i = 0; i < m; ++i) {</pre>
       int i:
       for (i = 0; i < n; ++i) {
          if (!row_selected[j] && abs(A[j][i]) > EPS)
              break:
       if (i != n) {
           ++rank;
          row_selected[j] = true;
          for (int p = i + 1; p < m; ++p)
              A[i][p] /= A[i][i];
```

```
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```

```
for (int k = 0: k < n: ++k) {
               if (k != j && abs(A[k][i]) > EPS) {
                   for (int p = i + 1; p < m; ++p)
                       A[k][p] -= A[j][p] * A[k][i];
           }
       }
    return rank;
Determinant of a matrix:
const double EPS = 1E-9:
vector < vector<double> > a (n, vector<double> (n));
double det = 1:
for (int i=0: i<n: ++i) {</pre>
    int k = i:
   for (int j=i+1; j<n; ++j)</pre>
       if (abs (a[j][i]) > abs (a[k][i]))
           k = j;
    if (abs (a[k][i]) < EPS) {</pre>
        det = 0:
       break;
    swap (a[i], a[k]):
    if (i != k)
       det = -det:
    det *= a[i][i];
    for (int j=i+1; j<n; ++j)</pre>
       a[i][i] /= a[i][i]:
    for (int j=0; j<n; ++j)</pre>
       if (j != i && abs (a[j][i]) > EPS)
           for (int k=i+1: k<n: ++k)</pre>
               a[i][k] -= a[i][k] * a[i][i];
}
cout << det:
Generating all k-subsets:
vector<int> ans:
void gen(int n, int k, int idx, bool rev) {
    if (k > n | | k < 0)
       return:
    if (!n) {
       for (int i = 0; i < idx; ++i) {</pre>
```

```
if (ans[i])
              cout << i + 1:
       cout << "\n":
       return;
   ans[idx] = rev:
   gen(n-1, k-rev, idx+1, false);
   ans[idx] = !rev;
   gen(n-1, k-!rev, idx + 1, true):
void all combinations(int n, int k) {
   ans.resize(n):
   gen(n, k, 0, false);
Simpsons formula for integration:
const int N = 1000 * 1000; // number of steps (already
    multiplied by 2)
double simpson_integration(double a, double b){
   double h = (b - a) / N:
   double s = f(a) + f(b); // a = x_0 and b = x_2n
   for (int i = 1: i <= N - 1: ++i) { // Refer to final</pre>
        Simpson's formula
      double x = a + h * i;
      s += f(x) * ((i & 1) ? 4 : 2):
   s *= h / 3;
   return s:
Picks theorem:
Given a certain lattice polygon with non-zero area.
We denote its area by S, the number of points with integer
coordinates lying strictly inside the polygon by I and the
number of points lying on polygon sides by B.
Then, the Pick formula states: S=I + B/2 - 1
In particular, if the values of I and B for a polygon are
the area can be calculated in O(1) without even knowing the
    vertices.
Strongly Connected component and Condensation Graph:
```

```
vector < vector<int> > g, gr;
   vector<bool> used:
   vector<int> order, component;
   void dfs1 (int v) {
       used[v] = true:
       for (size_t i=0; i<g[v].size(); ++i)</pre>
          if (!used[ g[v][i] ])
              dfs1 (g[v][i]);
       order.push_back (v);
   void dfs2 (int v) {
       used[v] = true:
       component.push_back (v);
       for (size_t i=0; i<gr[v].size(); ++i)</pre>
           if (!used[ gr[v][i] ])
              dfs2 (gr[v][i]);
   }
   int main() {
       int n:
       ... reading n ...
       for (::) {
           int a. b:
           ... reading next edge (a,b) ...
           g[a].push_back (b);
           gr[b].push back (a):
       used.assign (n, false);
       for (int i=0; i<n; ++i)</pre>
          if (!used[i])
              dfs1 (i);
       used.assign (n, false);
       for (int i=0: i<n: ++i) {</pre>
          int v = order[n-1-i]:
           if (!used[v]) {
              dfs2 (v):
              ... printing next component ...
              component.clear():
      }
   }
FFT Matrices:
XOR FFT: 1 1 / 1 -1, AND FFT: 0 1/ 1 1, OR FFT: 1 1/ 1 0
```