

AVL TREES

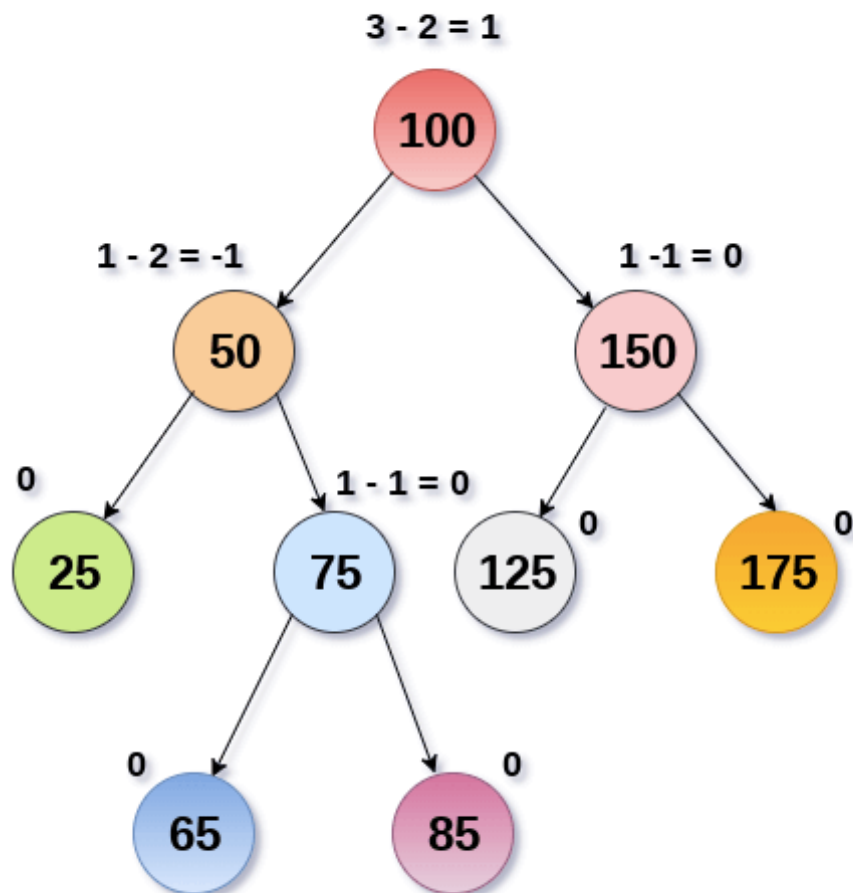
Balance Factor (k) = height (left(k)) - height (right(k))

If balance factor of any node is 1, it means that the left sub-tree is one level higher than the right sub-tree.

If balance factor of any node is 0, it means that the left sub-tree and right sub-tree contain equal height.

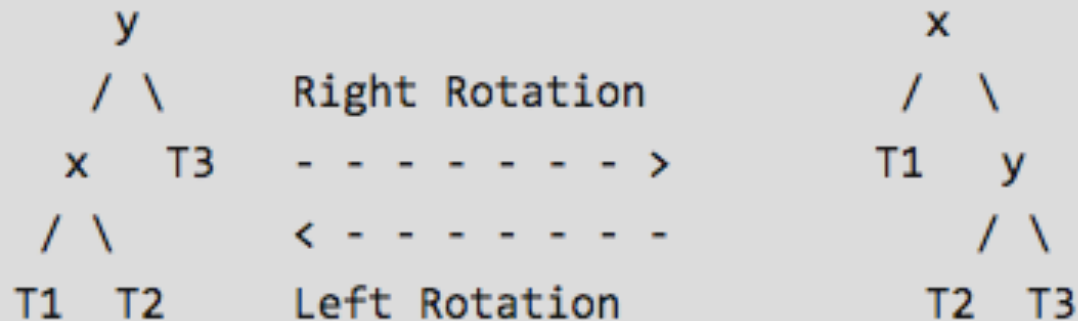
If balance factor of any node is -1, it means that the left sub-tree is one level lower than the right sub-tree.

An AVL tree is given in the following figure. We can see that, balance factor associated with each node is in between -1 and +1. therefore, it is an example of AVL tree.



Rotation

T1, T2 and T3 are subtrees of the tree rooted with y (on the left side) or x (on the right side)



Keys in both of the above trees follow the following order

$\text{keys}(T1) < \text{key}(x) < \text{keys}(T2) < \text{key}(y) < \text{keys}(T3)$

So BST property is not violated anywhere.

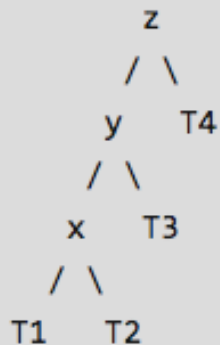
Balance your Binary Search Tree

Left Left Case

y is left child of z and x is left child of y

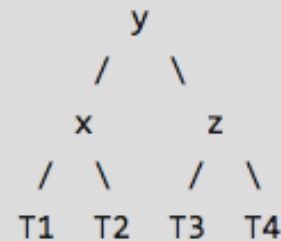
a) Left Left Case

T1, T2, T3 and T4 are subtrees.



Right Rotate (z)

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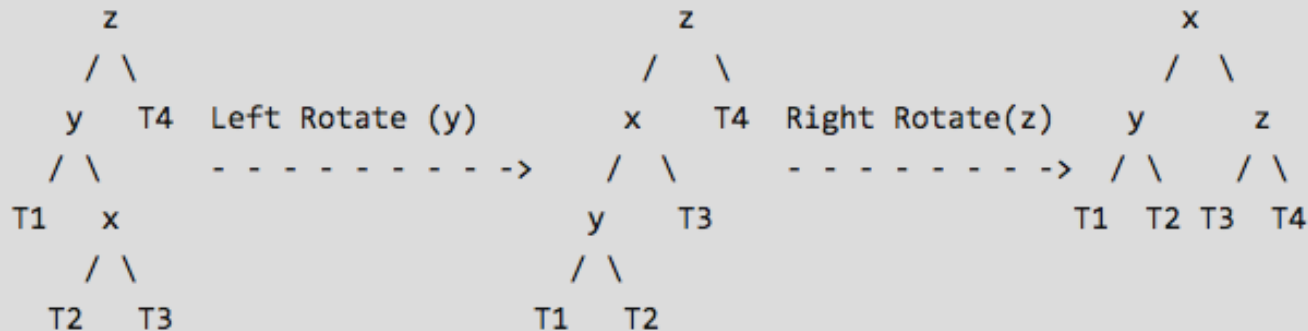


Balance your Binary Search Tree

Left Right Case

y is left child of z and x is right child of y

b) Left Right Case

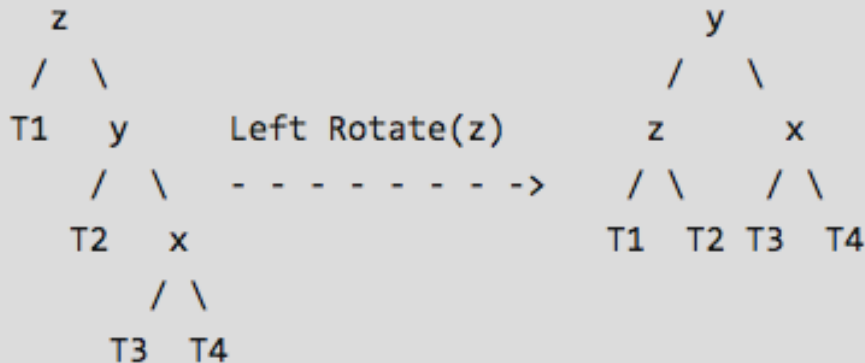


Balance your Binary Search Tree

Right Right Case

y is right child of z and x is right child of y

c) Right Right Case

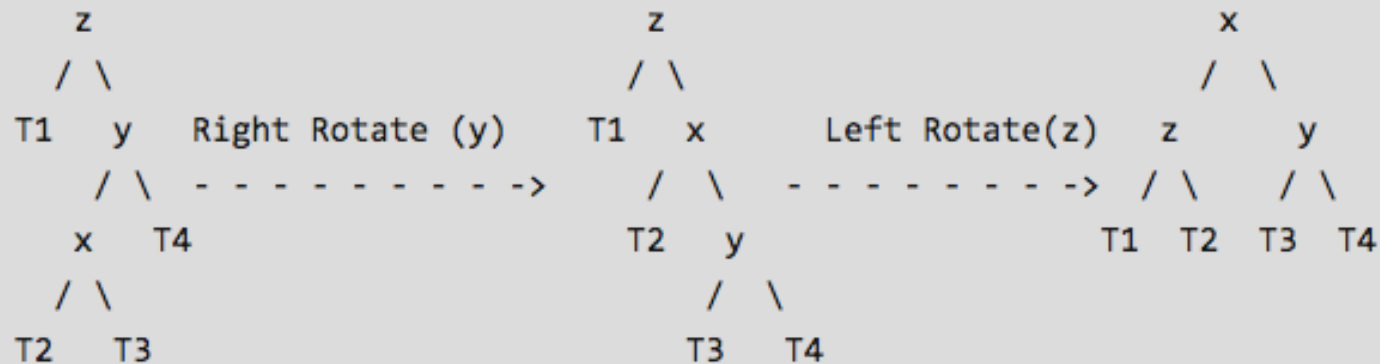


Balance your Binary Search Tree

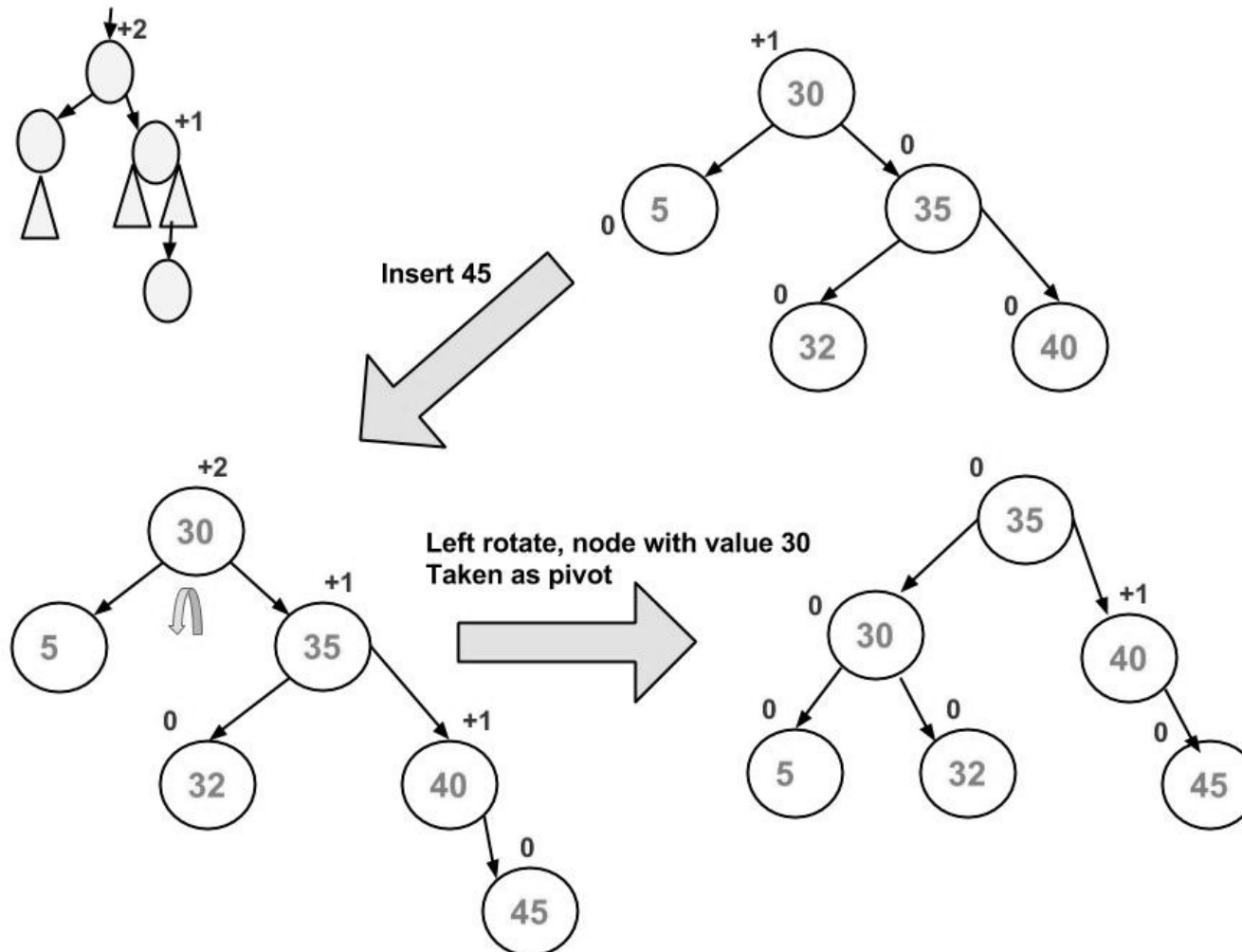
Right Left Case

y is right child of z and x is left child of y

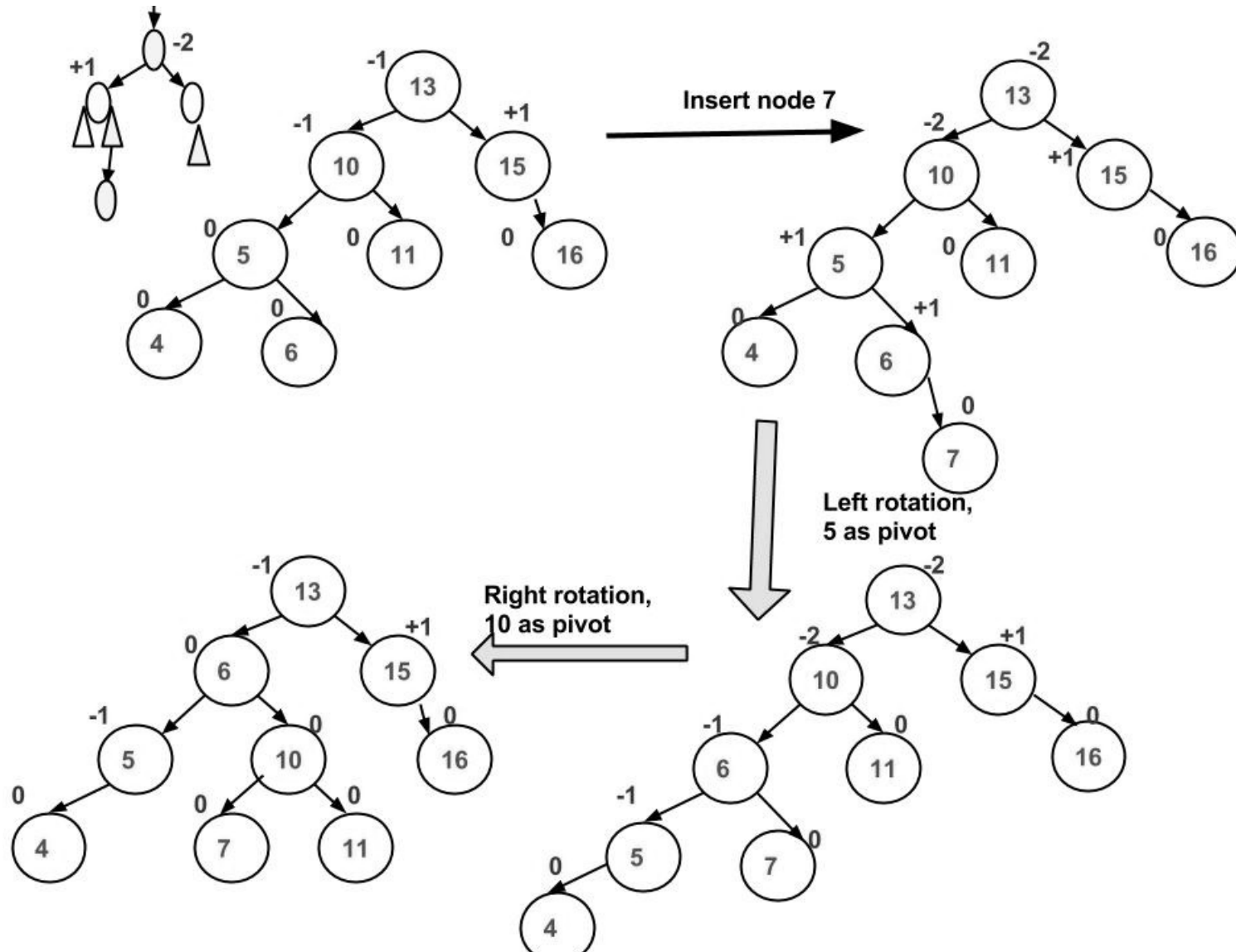
d) Right Left Case



Example



Example



Example

