### **DESIGN & ANALYSIS OF ALGORITHM**

PCC-CS501

## DESIGN & ANALYSIS OF ALGORITHM SCHEDULE ----TOPIC WISE

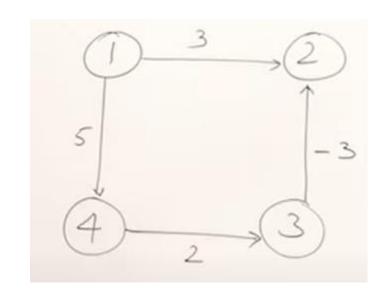
	Topic	Sub Topic
1	INTRODUCTION	DESIGN OF ALGORITHM, ANALYSIS OF ALGORITHM,
		ALGORITHM PROPERTIES
2	FRAMEWORK FOR ALGORITHM	HOW TO COUNT EXECUTION TIME OF ALGORITHM, INPUT INSTANCES
	ANALYSIS	
3	ASYMPTOTIC NOTATION	BEST CASE, AVERAGE CASE, WORST CASE
4	SOLVING RECURRENCE RELATION	SUBSTITUTION METHOD, MASTER THEOREM
5	ALGORITHM DESIGN TECHNIQUES	DIVIDE & CONQUER, GREEDY, DYNAMIC PROGRAMMING,
		BACKTRACKING,
6	DISJOINT SET MANIPULATION	UNION FIND
7	NETWORK FLOW PROBLEM	FORD FULKERSON ALGORITHM
8	NP COMPLETENESS	NP,NP HARDALGORITHM
9	APPROXIMATION ALGORITHM	COMPLEXITY ANALYSIS OF NP COMPETE PROBLEM

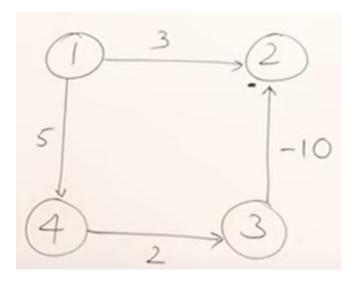
SINGLE SOURCE SHORTEST PATH ALGORITHM

### DRAWBACK

### DIJKSTRA ALGORITHM

 May or May not Support Negative weight edge.

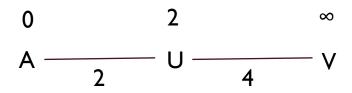


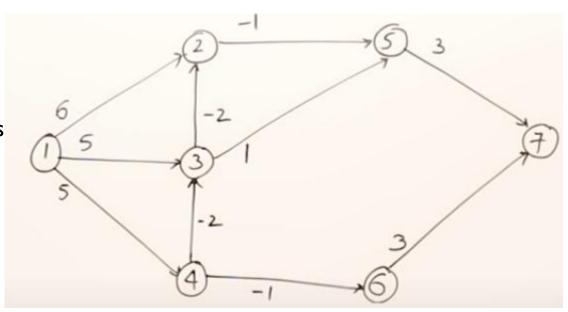


#### DYNAMIC PROGRAMMING

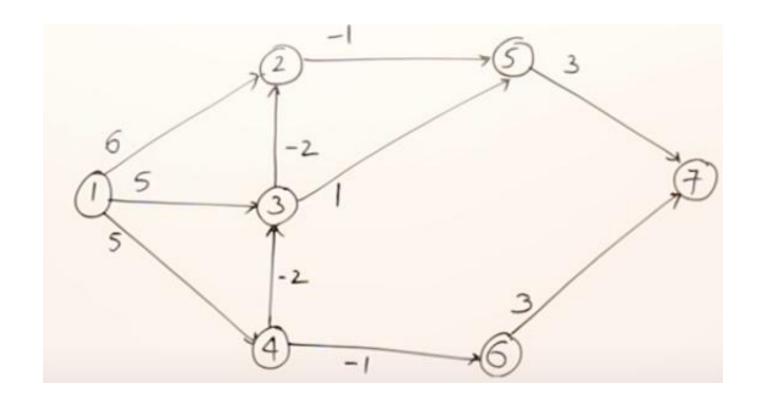
- Applicable on directed / undirected graph
- Optimization Problem
- Rule of Relaxation

It is similar to Dijkstra's algorithm but it can work with graphs in which edges can have negative weights.





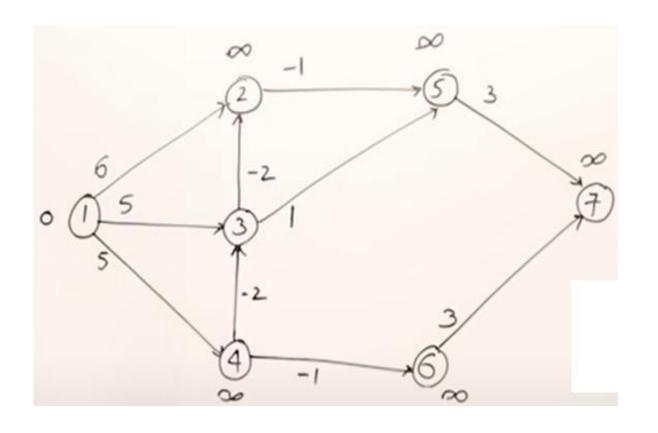
- Start vertex : I
- Visit |V|-I time all the edges.



## BELLMAN FORD ALGORITHM RELAX THE EDGES (1,2)(1,3)(1,4)(2,5)(3,2)(3,5)(4,3)(4,6)(5,6)(6,7)

Start vertex : I

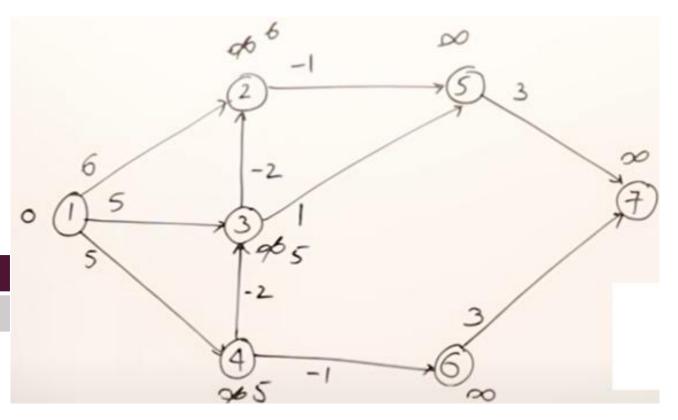
1	2	3	4	5	6	7
0	∞	∞	∞	∞	∞	∞



## BELLMAN FORD ALGORITHM RELAX THE EDGES (1,2)(1,3)(1,4)(2,5)(3,2)(3,5)(4,3)(4,6)(5,7)(6,7)

Start vertex : I

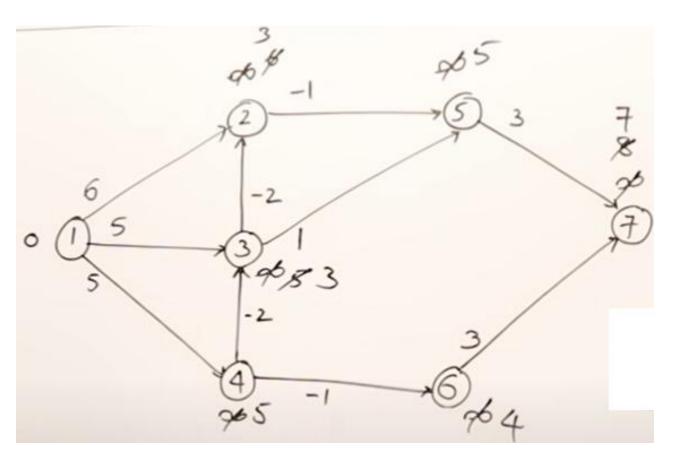
I	2	3	4	5	6	7
0	6	5	5	∞	∞	∞



## BELLMAN FORD ALGORITHM RELAX THE EDGES (1,2)(1,3)(1,4)(2,5)(3,2)(3,5)(4,3)(4,6)(5,7)(6,7)

Start vertex : I

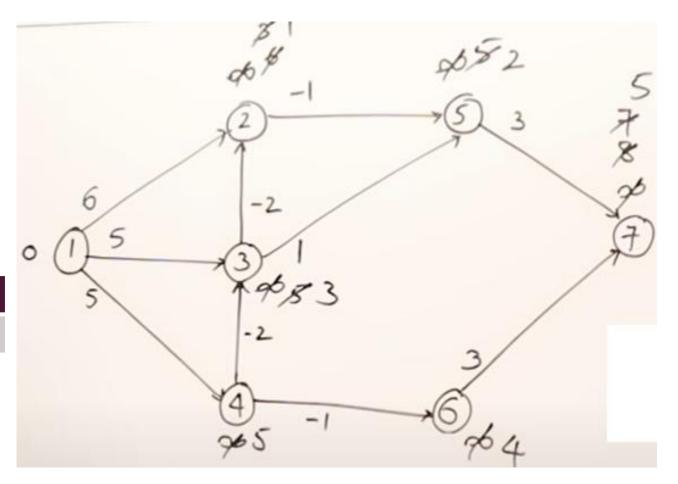
1	2	3	4	5	6	7
0	3	3	5	5	4	7



## BELLMAN FORD ALGORITHM RELAX THE EDGES (1,2)(1,3)(1,4)(2,5)(3,2)(3,5)(4,3)(4,6)(5,7)(6,7)

After 2<sup>nd</sup> visit.

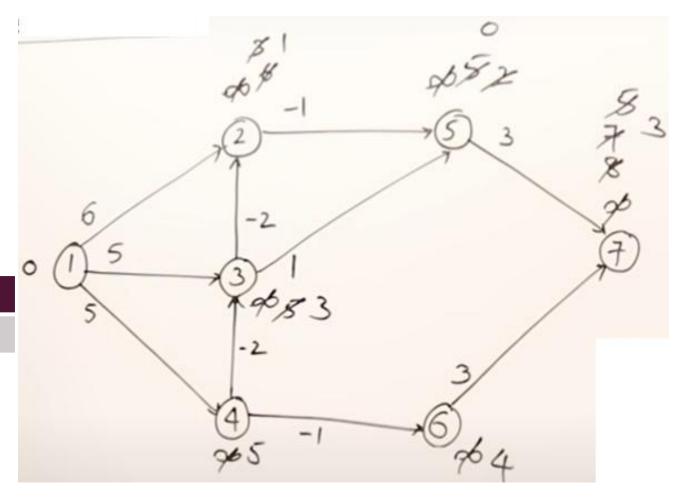
1	2	3	4	5	6	7
0	1	3	5	2	4	5



## BELLMAN FORD ALGORITHM RELAX THE EDGES (1,2)(1,3)(1,4)(2,5)(3,2)(3,5)(4,3)(4,6)(5,6)(6,7)

After 3<sup>rd</sup> visit.

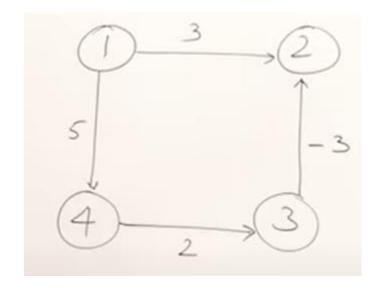
1	2	3	4	5	6	7
0	1	3	5	0	4	3

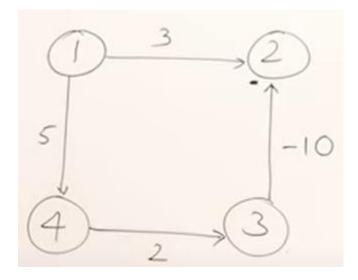


#### **ANALYSIS**

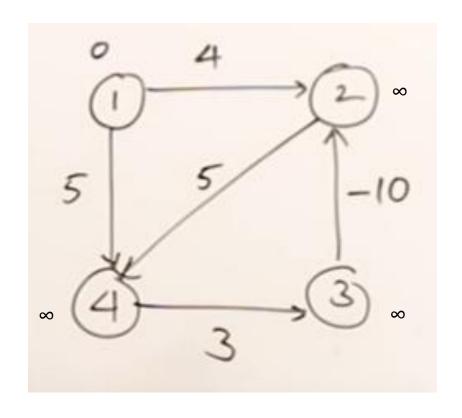
 $(|V|-I) \to O(|V||E|) \to O(n^2)$ 

■ For complete Graph  $E= n(n-1)/2 \dots O(|V||E|) \rightarrow O(n^3)$ 





# BELLMAN FORD ALGORITHM (1,2)(1,4)(2,4)(3,2)(4,3)



|V|-I

After 3<sup>rd</sup> time visit

1	2	3	4
0	-2	6	3

**NEGATIVE WEIGHT CYCLE** 

#### **NEXT CLASS**

■ Travelling Salesman Problem.