

MAS 433: Cryptography

Lecture 8

Block Cipher (Part 4, Modes of Operation)

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Lecture Outline

- Classical ciphers
- Symmetric key encryption
 - One-time pad & information theory
 - Block cipher
 - DES, Double DES, Triple DES
 - AES
 - **Modes of Operation**
 - Attacks
 - Stream cipher
- Hash function and Message Authentication Code
- Public key encryption
- Digital signature
- Key establishment and management
- Introduction to other cryptographic topics

Recommended Reading

- CTP Section 3.7
- HAC Section 7.2.2
- Wikipedia:
 - Modes of operation
http://en.wikipedia.org/wiki/Block_cipher_modes_of_operation
 - Ciphertext Stealing
http://en.wikipedia.org/wiki/Ciphertext_stealing

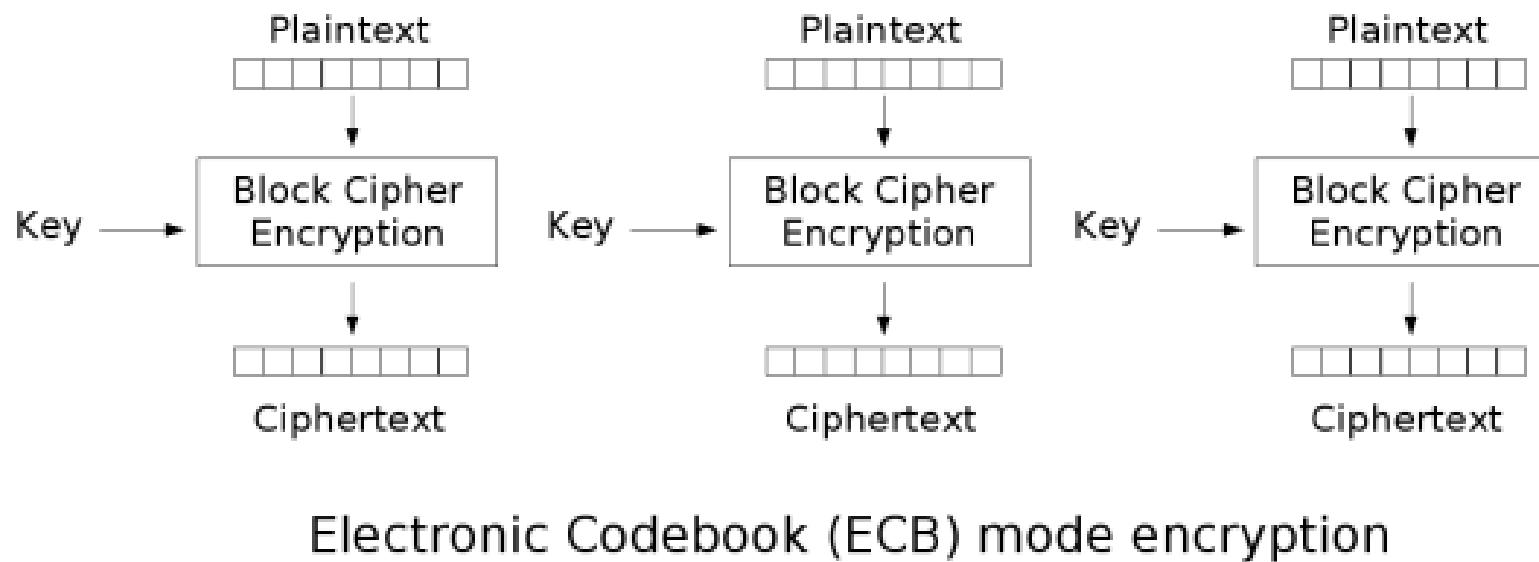
Block cipher

- Fixed block size
 - DES: 64 bits
 - AES: 128 bits
- How to encrypt many messages using the same key in a secure way?

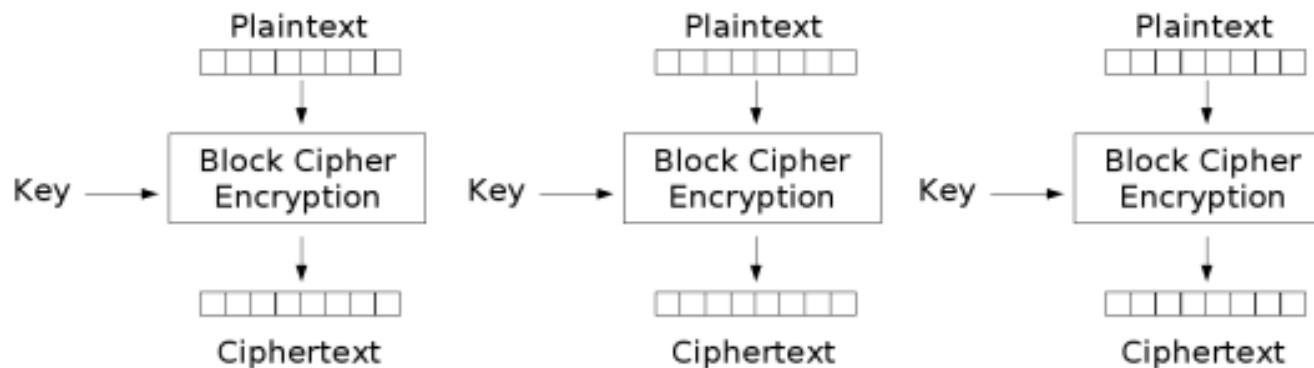
Block Cipher Modes of Operation

- NIST Special Publication 800-38A (2001)
<http://csrc.nist.gov/publications/nistpubs/800-38a/sp800-38a.pdf>
- Five encryption modes are recommended
 - Electronic Codebook (ECB)
 - Cipher Block Chaining (CBC)
 - Cipher Feedback (CFB)
 - Output Feedback (OFB)
 - Counter (CTR)

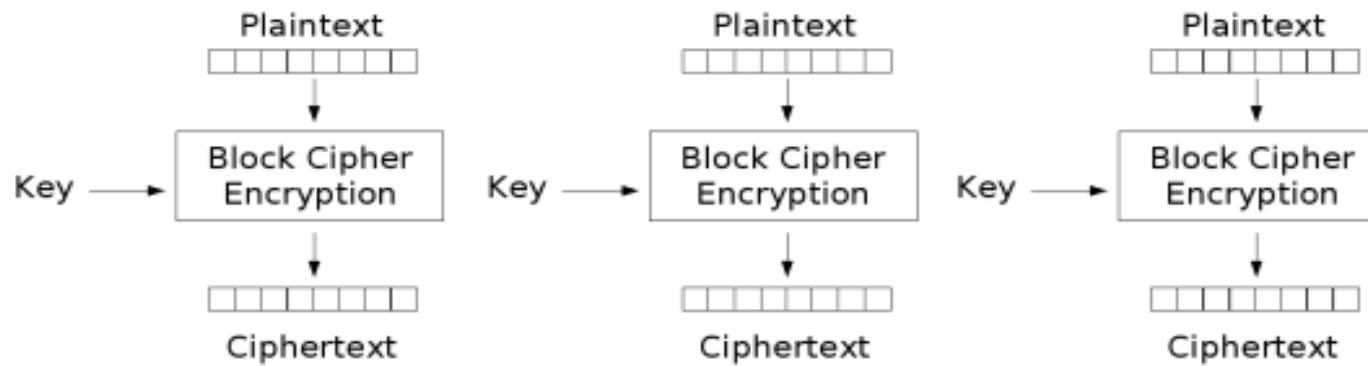
Electronic Codebook (ECB)



Electronic Codebook (ECB) (contd.)



Electronic Codebook (ECB) mode encryption

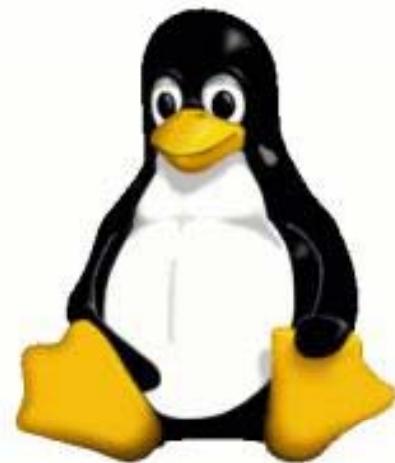


Electronic Codebook (ECB) mode encryption

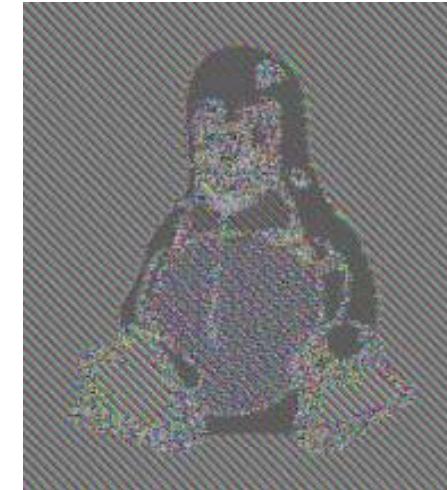
Electronic Codebook (ECB) (contd.)

- ECB mode, same key + same plaintext lock
=> the same ciphertext block
 - If this property is undesirable in an application, **ECB mode should not be used**
 - Example: Data with high redundancy
 - Uncompressed image file
 - Example: A secret key being used to encrypt too many data
 - there is a lot of redundancy in many documents, so the plaintext space in the ECB mode is limited
 - » Example: For AES, the input space size is 2^{128} . However, if each byte is used to represent an English letter, and if we assume that each letter carries about 1.5-bit information, then a 128-bit message block contains only about $16 \times 1.5 = 24$ -bit information.

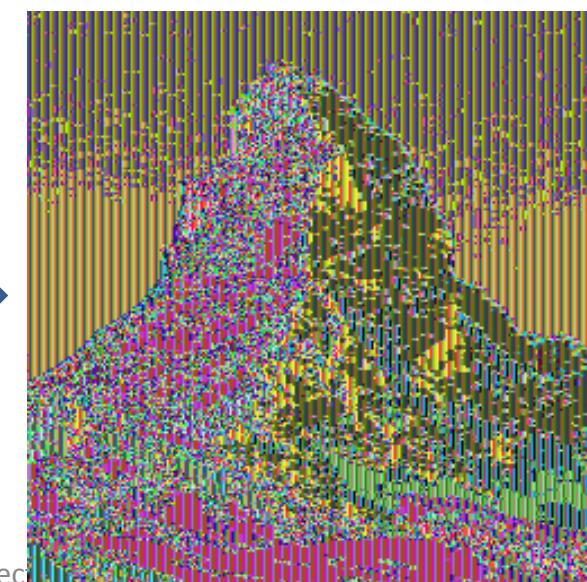
Electronic Codebook (ECB) (contd.)



ECB mode
encryption

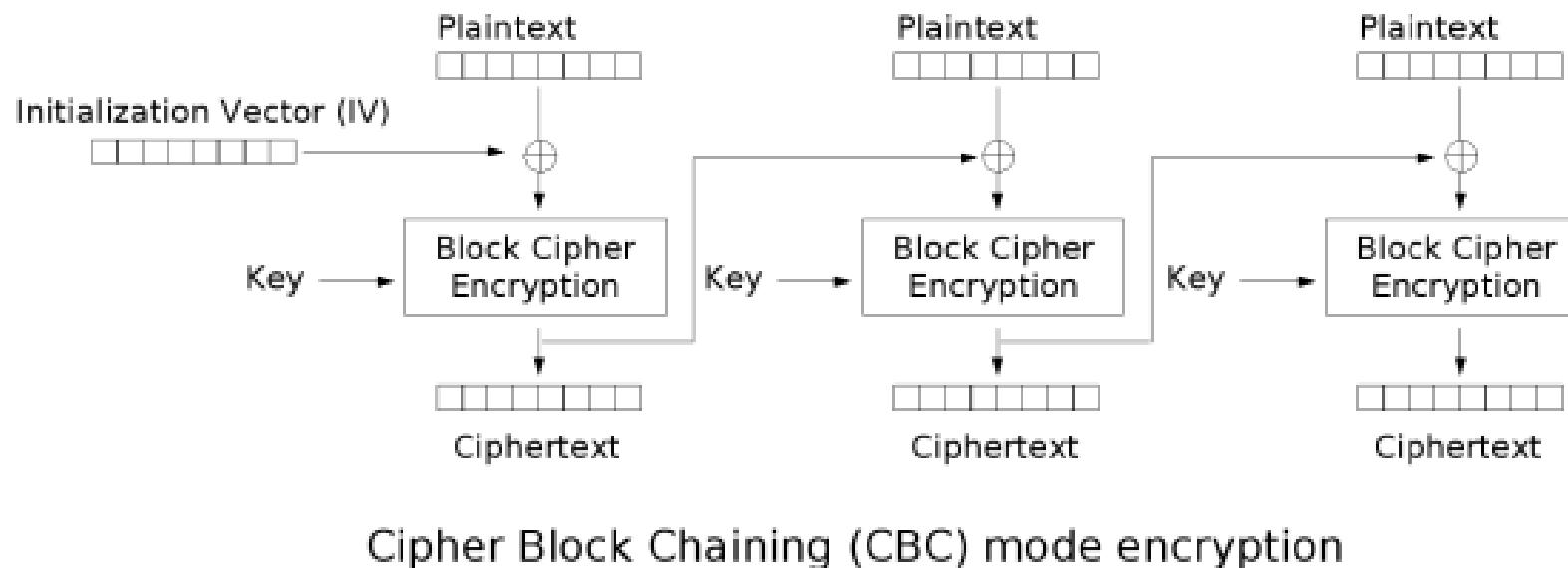


ECB mode
encryption

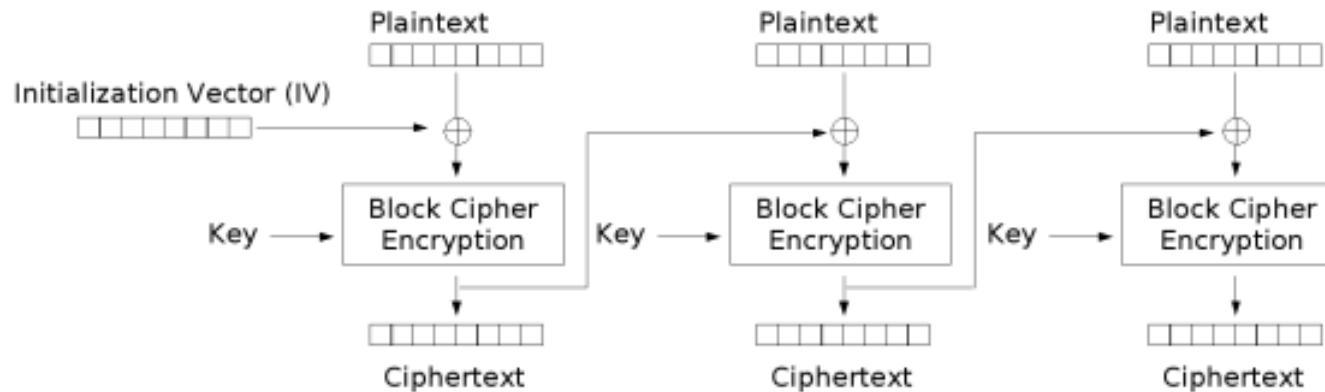


Cipher Block Chaining (CBC) Mode

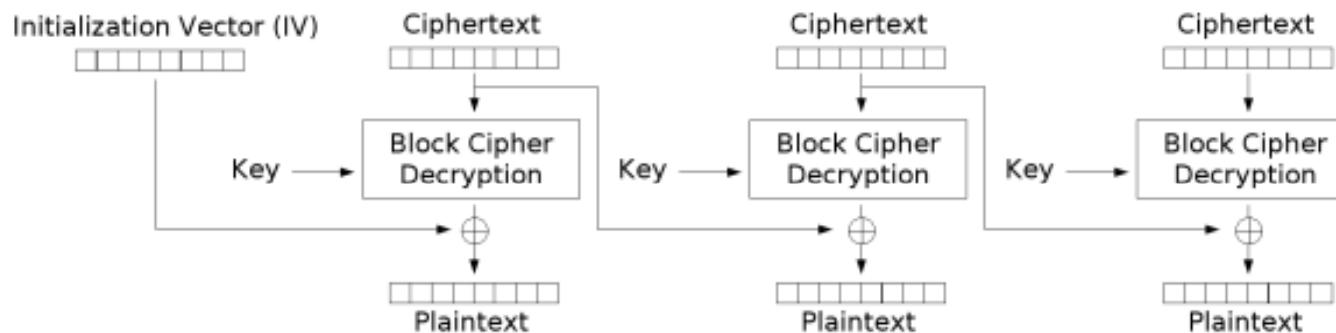
- Invented by IBM in 1976
- Initialization vector (also called “nonce”)
 - need not be secret (normally sent/stored together with ciphertext)
 - but must be unpredictable for CBC mode



Cipher Block Chaining (CBC) Mode (contd.)



Cipher Block Chaining (CBC) mode encryption $C_i = E_K(P_i \oplus C_{i-1}), C_0 = IV$



Cipher Block Chaining (CBC) mode decryption $P_i = D_K(C_i) \oplus C_{i-1}, C_0 = IV$

Cipher Block Chaining (CBC) Mode (contd.)

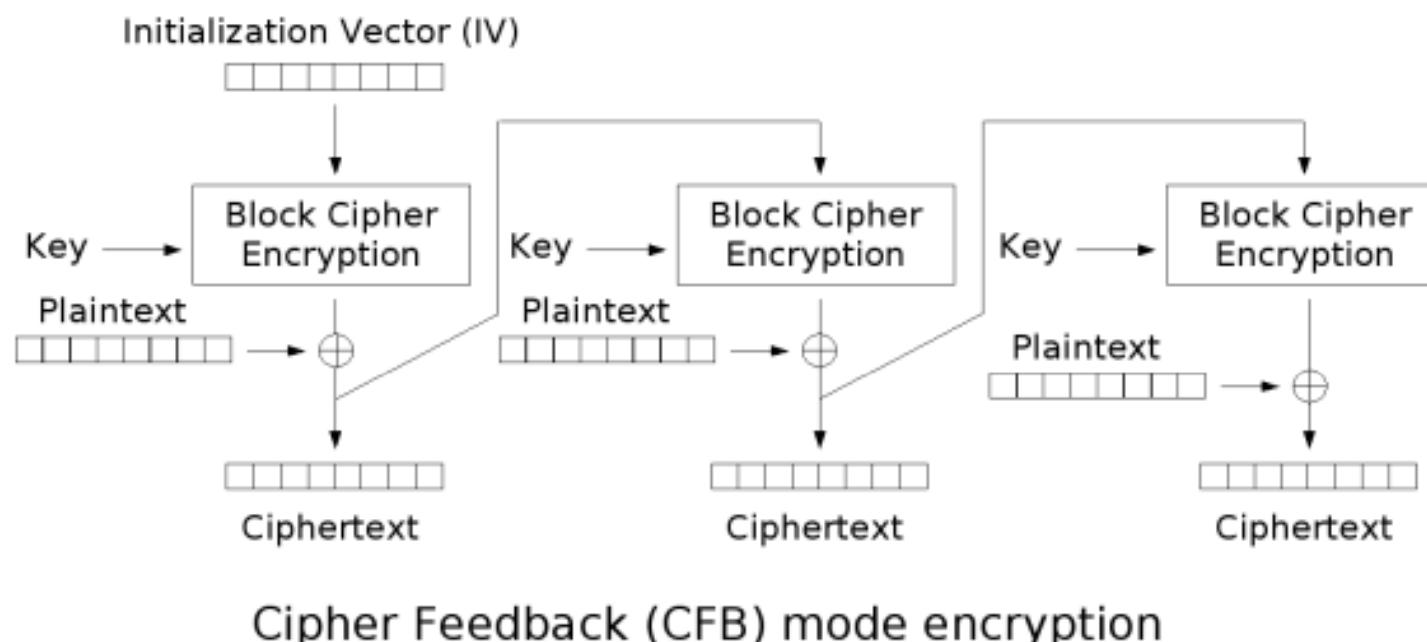
- CBC mode, the same key, the same plaintext block at different locations => different ciphertext blocks
 - The security of CBC is not that sensitive to the security of IV (If two IVs happen to be the same for the same key, encryption would not fail completely.)
 - The most reliable encryption mode!
 - The most commonly used encryption mode!

Cipher Feedback (CFB) Mode

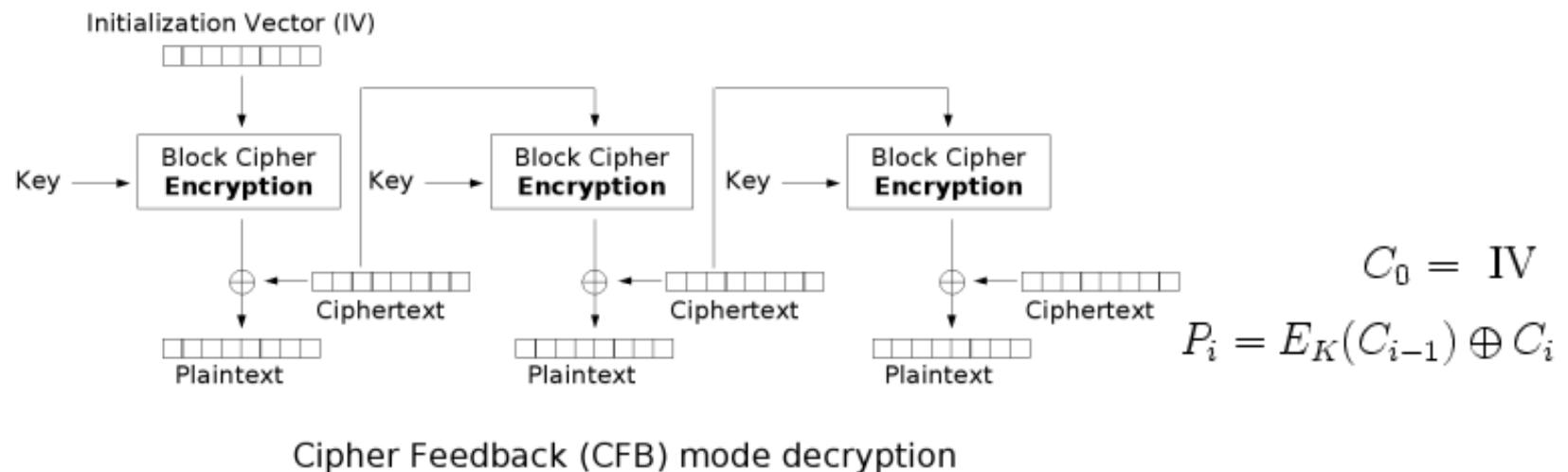
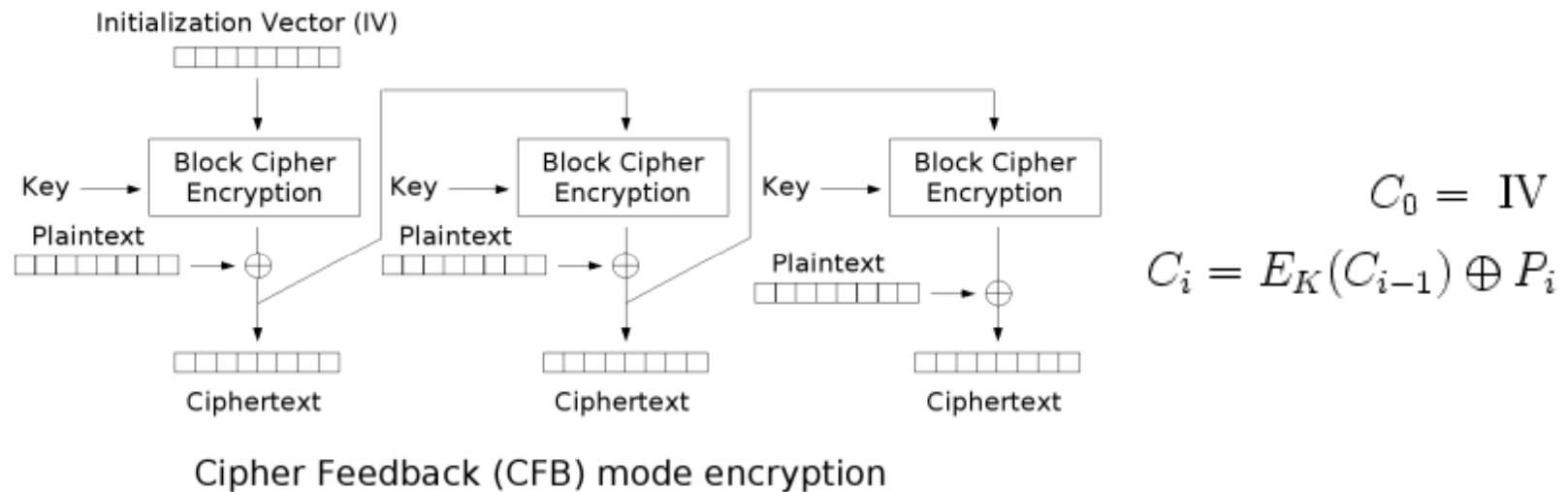
- A simplified version of CFB:

$$C_0 = \text{IV}$$

$$C_i = E_K(C_{i-1}) \oplus P_i$$

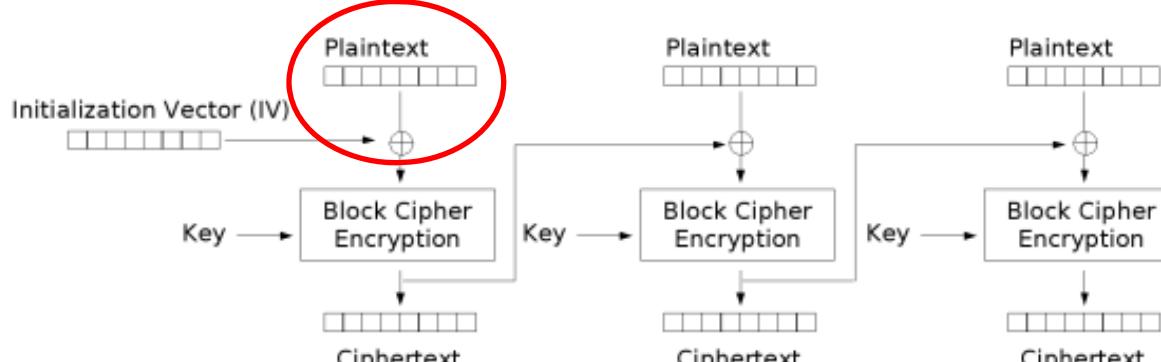


Cipher Feedback (CFB) Mode (contd.)

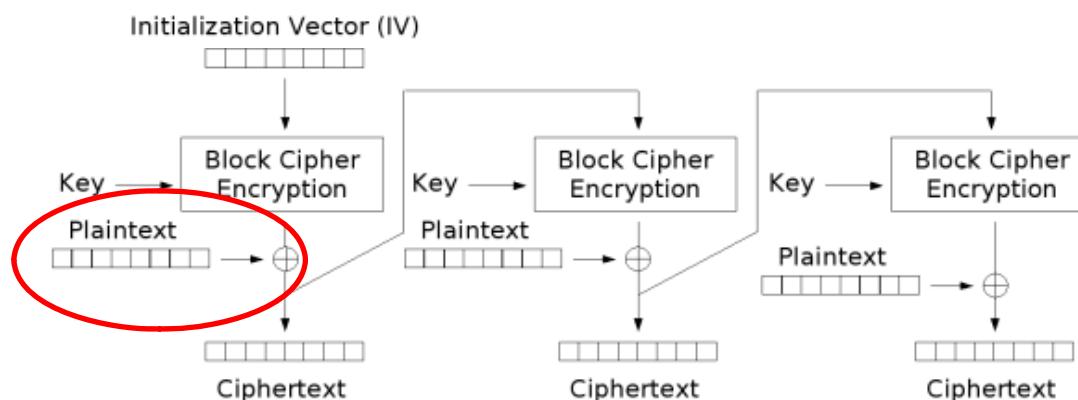


Cipher Feedback (CFB) Mode (contd.)

- Compare CBC & CFB



Cipher Block Chaining (CBC) mode encryption



Cipher Feedback (CFB) mode encryption

Similarity:
Each ciphertext block is used in the encryption of next block

Difference:
The plaintext

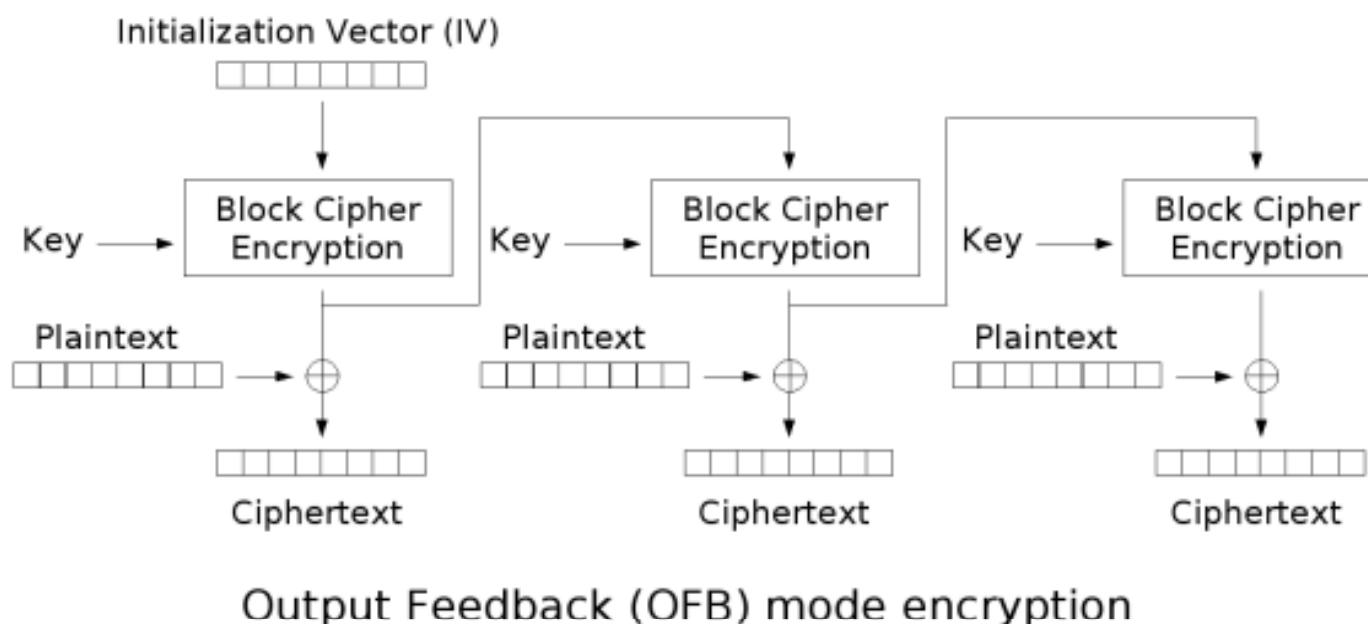
Output Feedback (OFB) Mode

- A simplified version of OFB:

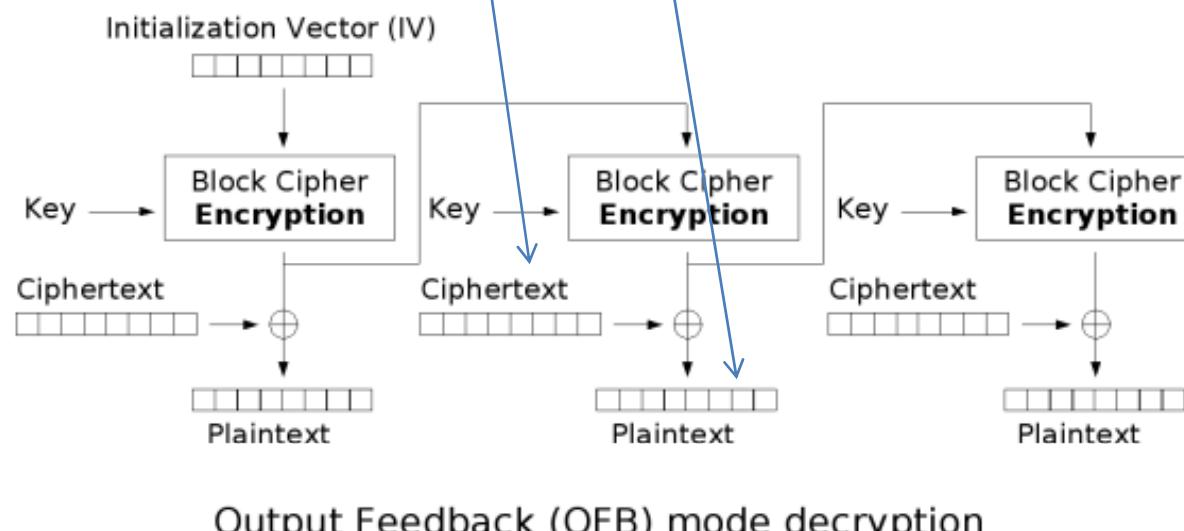
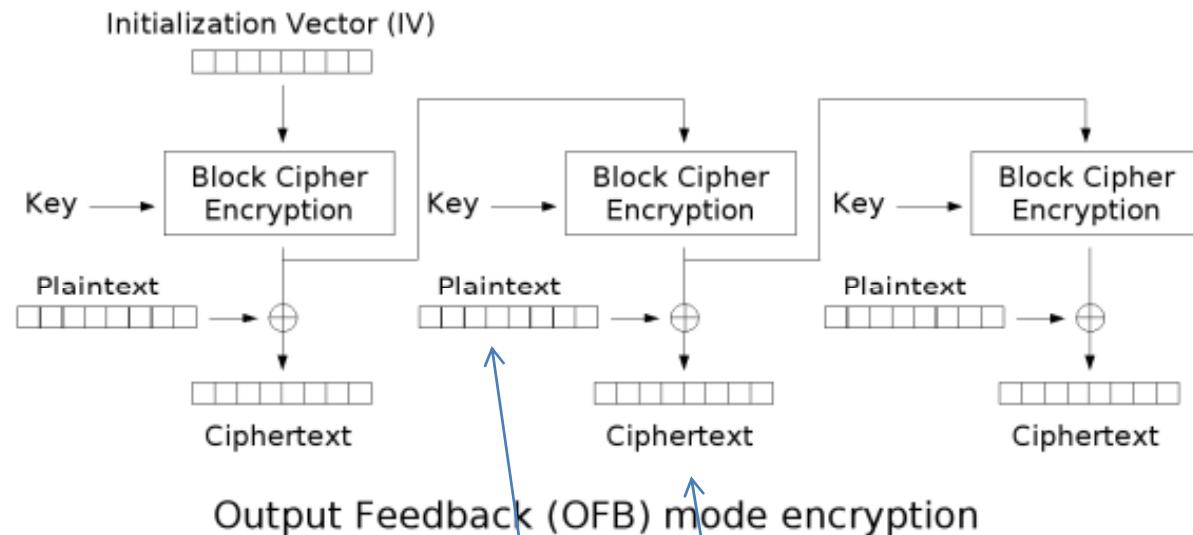
$$O_0 = \text{IV}$$

$$O_i = E_K(O_{i-1})$$

$$C_i = P_i \oplus O_i$$

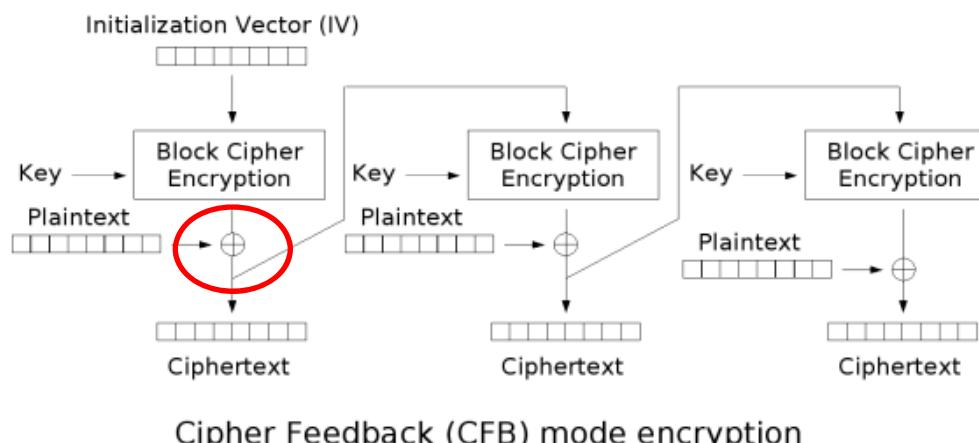
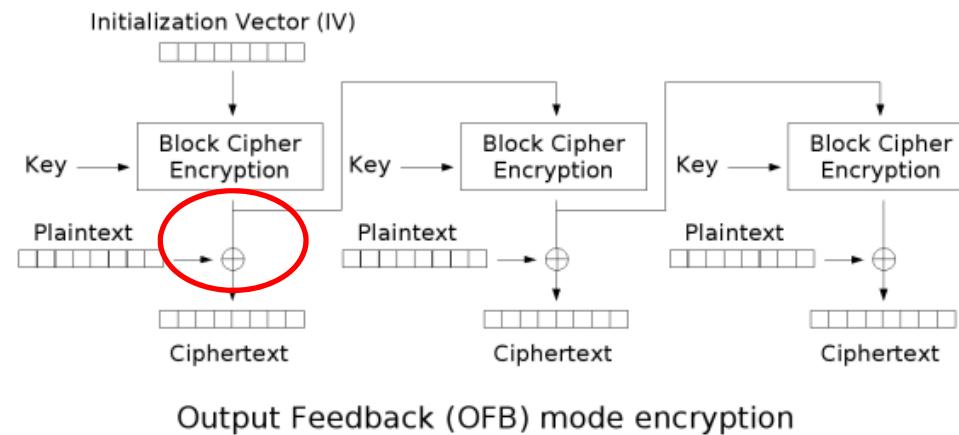


Output Feedback (OFB) Mode (cond.)



Output Feedback (OFB) Mode (cond.)

- Compare OFB & CFB

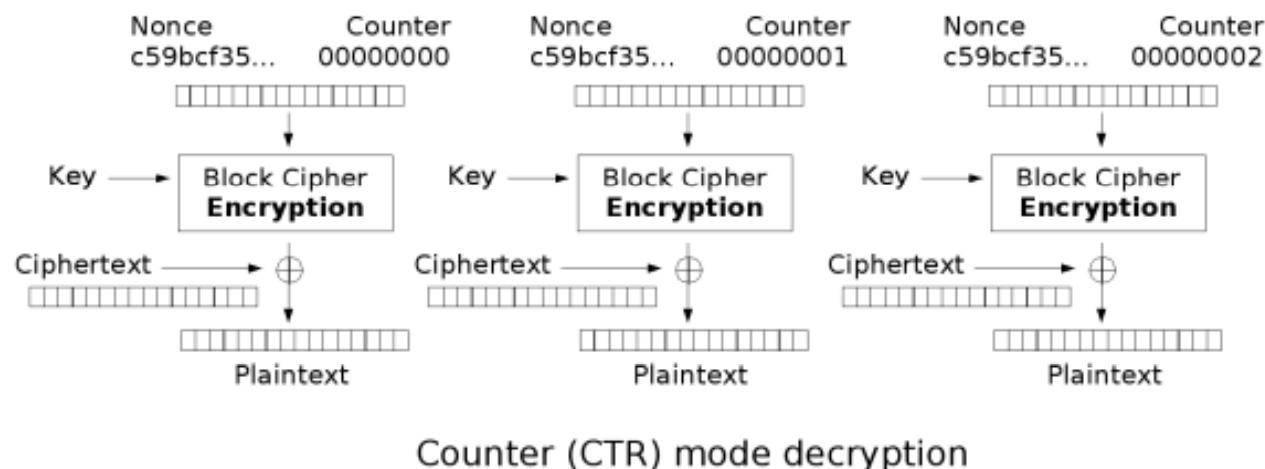
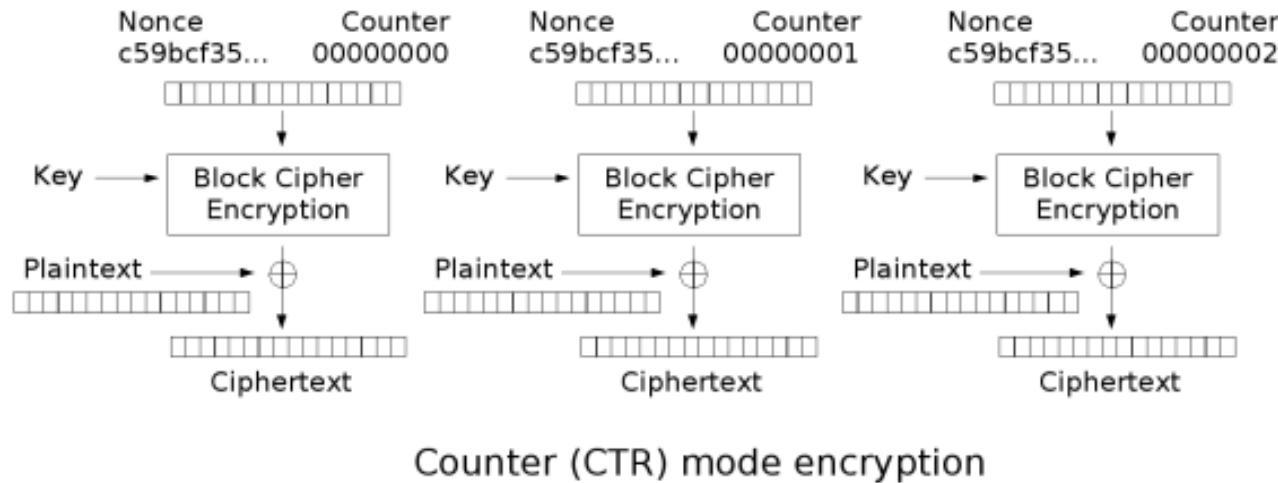


Counter (CTR) Mode

- For AES-CTR
 - The input of block cipher consists of 64-bit IV and 64-bit counter
 - IV: **different for each message**,
remains the same for each message
 - Counter: start from 0,
increased by 1 after each block
perform the same way for each message

IV (64 bits)	Counter (64-bits)
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Counter (CTR) Mode (contd.)

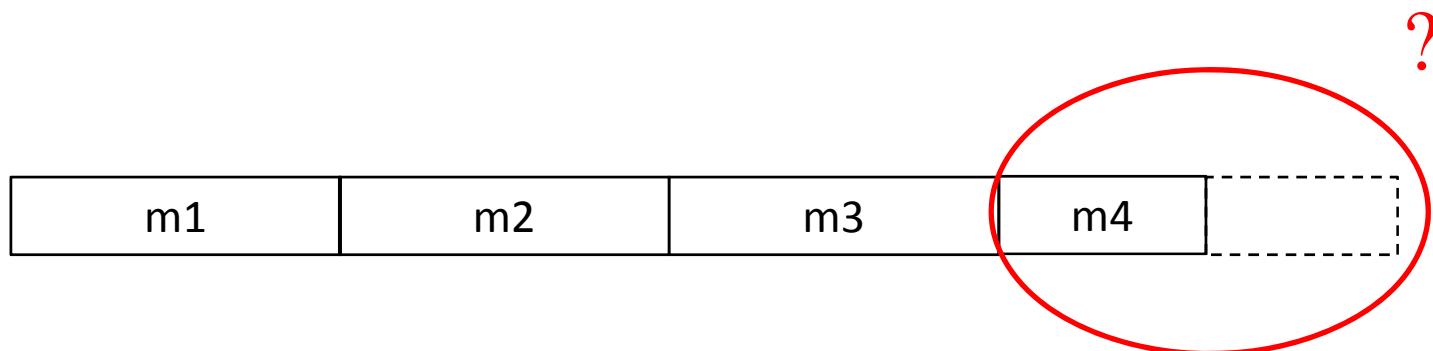


Counter (CTR) Mode (contd.)

- For each message, the counter should not repeat
 - i.e., the length of each message for AES-CTR should not be more than 2^{64} blocks

How to encrypt a partial block?

- If the message length is not the multiple of the block size of the block cipher
 - the last block is called partial block



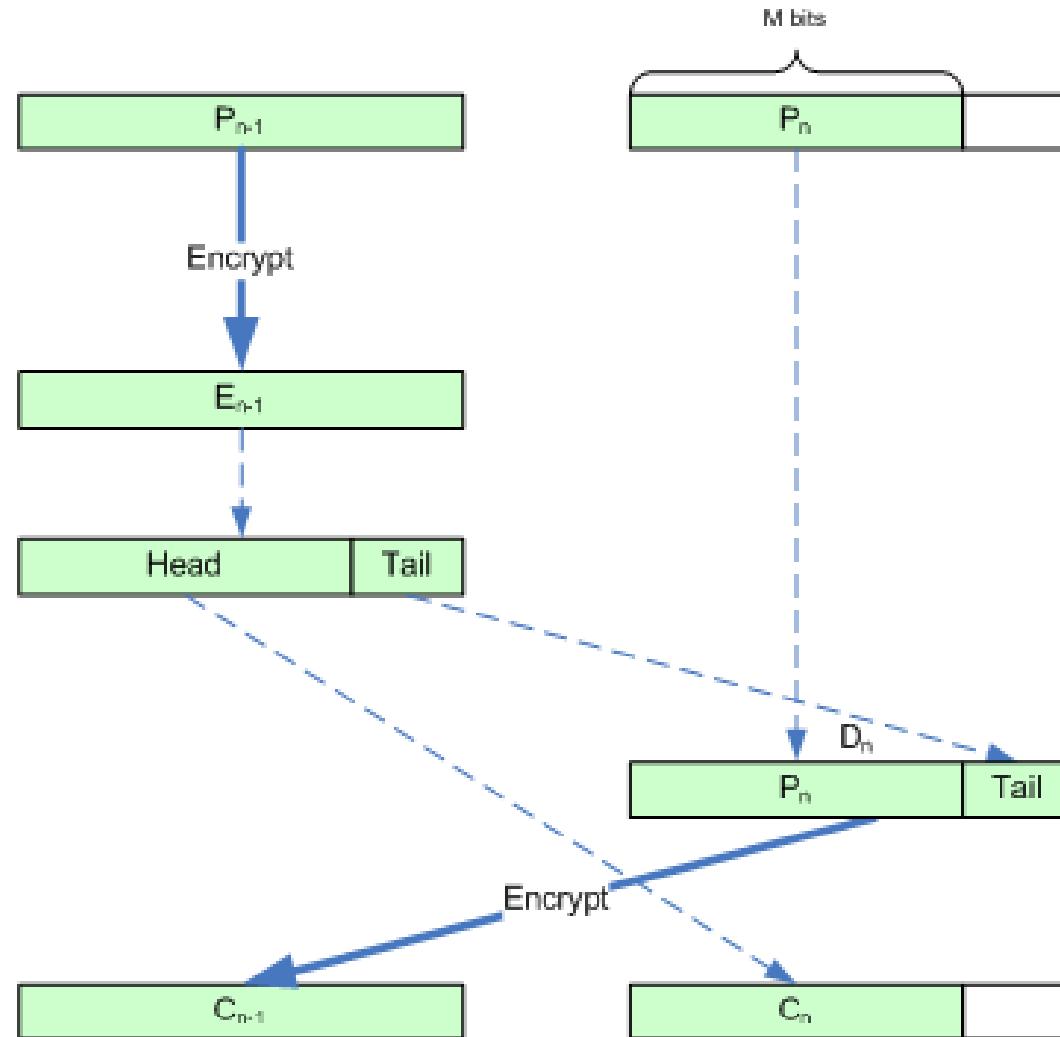
How to encrypt a partial block? (contd.)

- ECB & CBC
 - Straightforward encryption:
 - Pad the partial block to full block
 - ECB: padded with **random bits**
 - » Otherwise, the entropy of the partial block may be too small
 - CBC: padded with random bits or constant bits
 - Ciphertext length larger than plaintext length
 - The actual message length is sent together with the ciphertext
- CFB, OFB, CTR
 - NO partial block problem
 - Why?

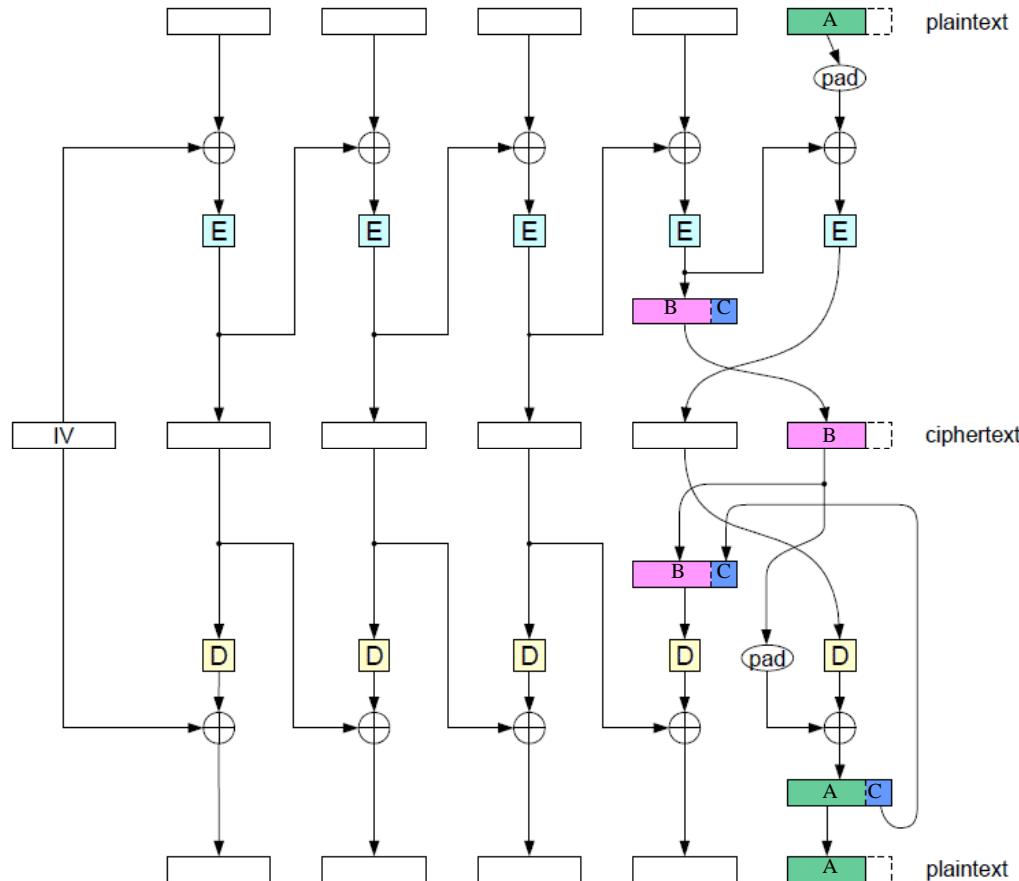
Ciphertext stealing

- Ciphertext stealing technique
 - Try to achieve:
 - Ciphertext length = Plaintext length
 - ECB ciphertext stealing
 - The plaintext should be more than one block
 - Otherwise, just use the padding method
 - CBC ciphertext stealing
 - Not necessary that plaintext is more than one block
 - If less than one full block, stealing from C_0 (IV).

ECB: Ciphertext Stealing



CBC: Ciphertext Stealing



data block
 block cipher (encryption/decryption)
 XOR
 zero-padding

<http://www.scythe.jp/memo/crypto-cts.html>

Summary

- Modes of operations
 - ECB: not strong
 - Parallel computation is possible
 - CBC: strong, the most commonly used
 - CFB
 - OFB: for the same key, all the IVs must be different
 - CTR: for the same key, all the IVs must be different
 - Parallel computation is possible
- Ciphertext stealing for encrypting the partial block
 - ECB
 - CBC
 - Not a problem for CFB, OFB & CTR