

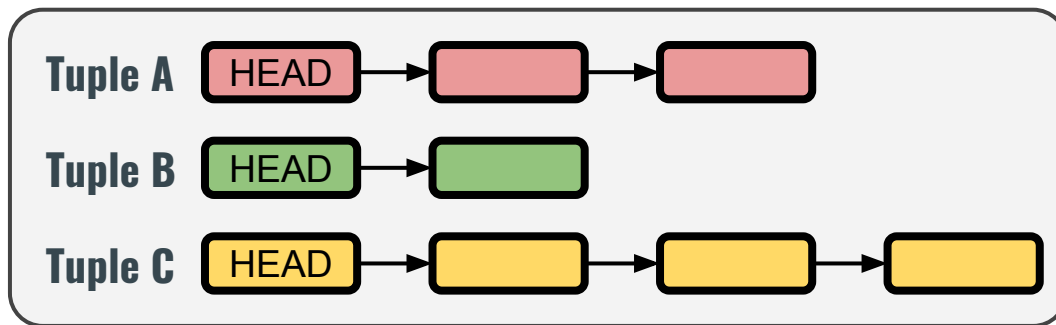
# An Experimental Evaluation of In-Memory Multi-Version Concurrency Control

Yingjun Wu, Joy Arulraj, Jiexi Lin, Ran Xian, Andrew Pavlo



# MULTI-VERSION CONCURRENT CONTROL

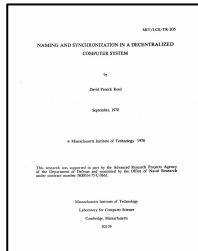
- ❑ Avoid read-write conflicts
- ❑ Support time-travel queries
- ❑ Enable snapshot isolation



**VERSION CHAINS**

# A BRIEF HISTORY OF MVCC

## 1979: FIRST MENTION



DAVID  
**REED**

Naming and synchronization in a decentralized computer system  
*Ph.D. Thesis, 1979*



# A BRIEF HISTORY OF MVCC

- **1979: FIRST MENTION**
- **1981: FIRST IMPLEMENTATION**



**InterBase/Firebird**

# A BRIEF HISTORY OF MVCC

- **1979: FIRST MENTION**
- **1981: FIRST IMPLEMENTATION**
- **1984:** Oracle
- **1985:** Postgres



# A BRIEF HISTORY OF MVCC

- **1979: FIRST MENTION**
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- **1984:** Oracle
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- **2001:** MySQL-InnoDB



# A BRIEF HISTORY OF MVCC

- **1979: FIRST MENTION**
- **1981: FIRST IMPLEMENTATION**
- **1984:** Oracle
- **1985:** Postgres
- **2001:** MySQL-InnoDB
- **2010-2017:** Hyrise, Hekaton, MemSQL, SAP HANA, NuoDB, HyPer...

# A BRIEF HISTORY OF MVCC

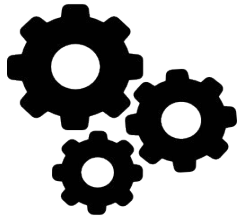
ORACLE®



Search for the best MVCC scheme  
for multi-core main-memory DBMSs



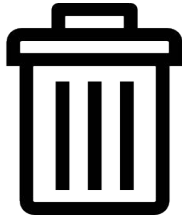
# DESIGN DECISIONS OF MVCC



**Concurrency Control Protocol**



**Version Storage**









**Garbage Collection**



**Index Management**

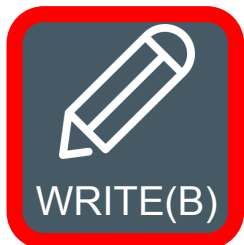
# CONCURRENCY CONTROL PROTOCOL

SCHEME	DBMS
Timestamp Ordering (MVTO)	N/A
Optimistic Concurrency Control (MVOCC)	  Microsoft SQL Server  MEMSQL
Two-phase Locking (MV2PL)	ORACLE  SAP HANA 
Serialization Certifier	 PostgreSQL

# CONCURRENCY CONTROL PROTOCOL

## ❏ Approach #1: Timestamp Ordering (MVT0)

- ❏ The DBMS assigns transactions timestamps that determine serial order.



	TXN-ID	READ-TS	BEGIN-TS	END-TS
A <sub>x</sub>	0	Tid	10	20
B <sub>x</sub>	Tid	17	15	30
B <sub>x+1</sub>	Tid	0	-	-

# CONCURRENCY CONTROL PROTOCOL

## ❏ Approach #3: Optimistic Concurrency Control (MVOC)

- ❏ Transactions optimistically access physical versions before validating the read-set consistency.

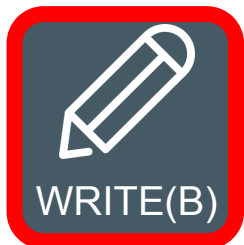


	TXN-ID	BEGIN-TS	END-TS
A <sub>x</sub>	0	10	20
B <sub>x</sub>	Tid	15	30
B <sub>x+1</sub>	Tid	-	-

# CONCURRENCY CONTROL PROTOCOL

## ❑ Approach #3: Two-Phase Locking (MV2PL)

- ❑ Transactions acquire appropriate lock on physical version before they can read/write a logical tuple.



	TXN-ID	READ-CNT	BEGIN-TS	END-TS
A <sub>x</sub>	0	2	10	20
B <sub>x</sub>	Tid	0	15	30
B <sub>x+1</sub>	Tid	0	-	-

# CONCURRENCY CONTROL PROTOCOL










## ❏ Approach #4: Serialization Certifier

- ❏ The DBMS maintains a serialization graph for detecting and removing “dangerous structures” formed by concurrent transactions.



	TXN-ID	BEGIN-TS	END-TS
A <sub>x</sub>	0	10	20
B <sub>x</sub>	Tid	15	30
B <sub>x+1</sub>	Tid	-	-

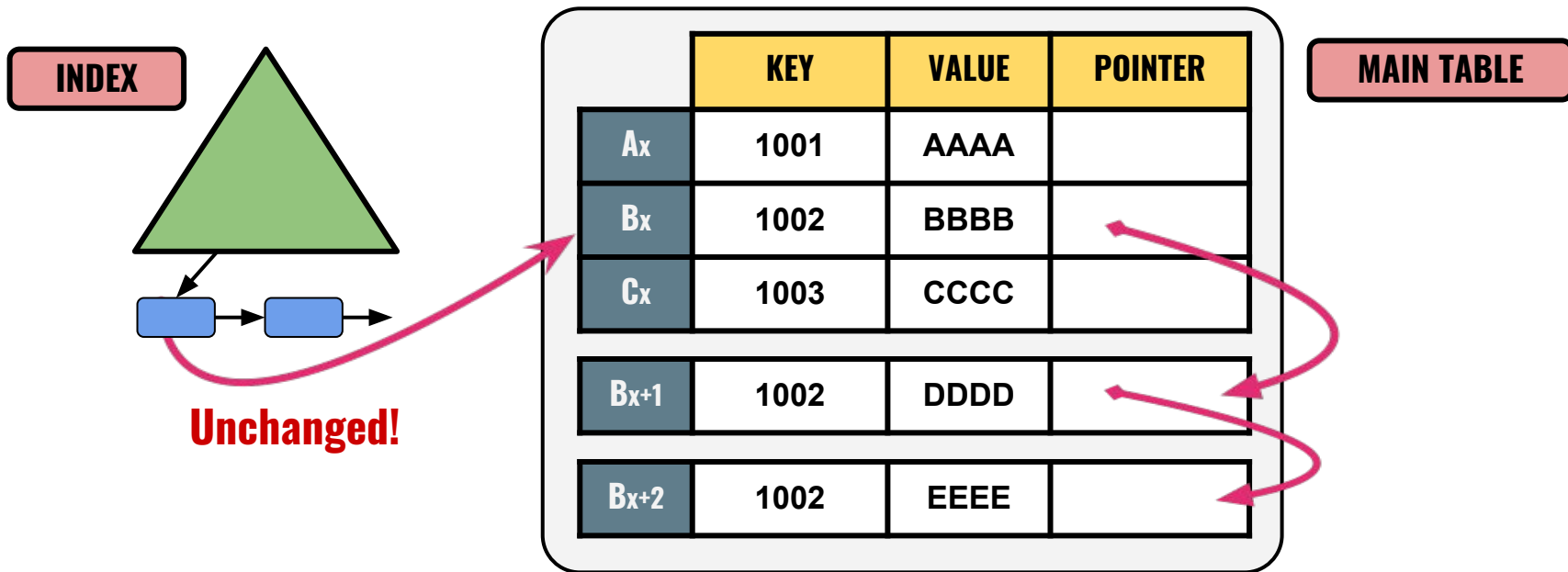
# VERSION STORAGE

SCHEME	DBMS
Append-Only	 PostgreSQL  HYRISE  Microsoft SQL Server  NUODB  MEMSQL
Time-Travel	 SAP HANA
Delta	 ORACLE  MySQL  HyPer

# VERSION STORAGE

## ❏ Approach #1: Append-Only Storage (Oldest-to-Newest)

- ❏ New versions are appended to the same table space.

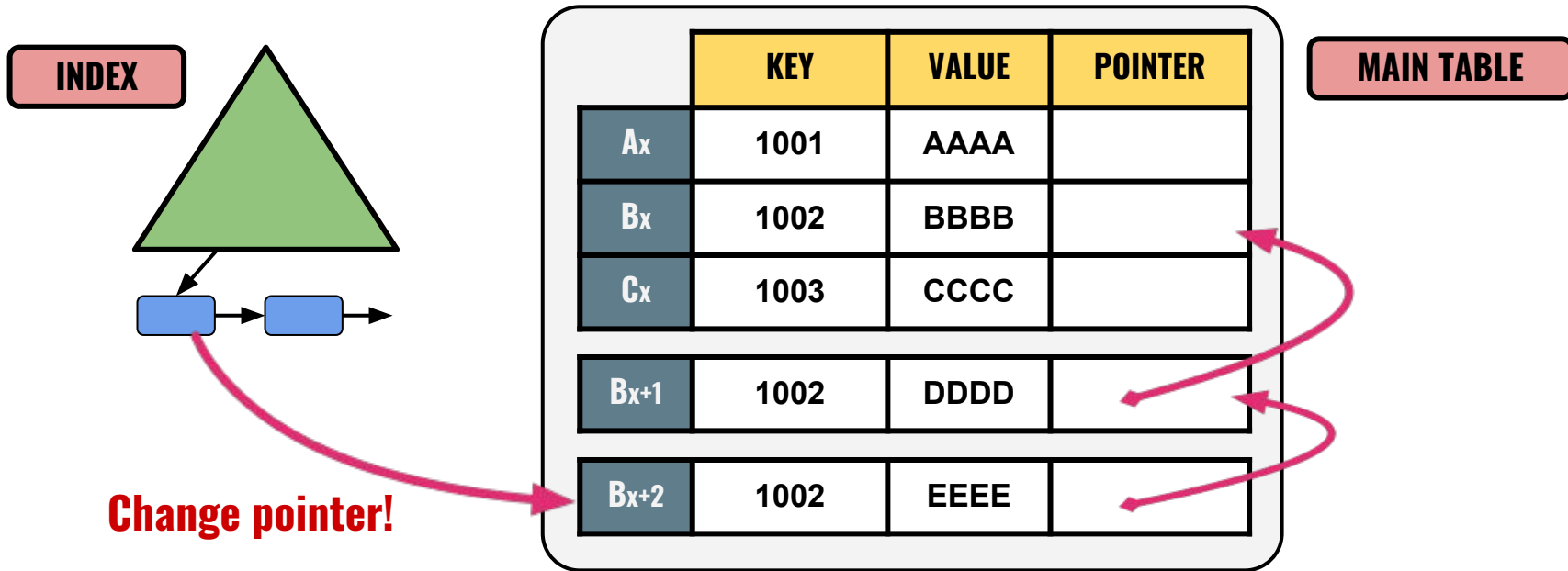




# VERSION STORAGE

## ❏ Approach #1: Append-Only Storage (Newest-to-Oldest)

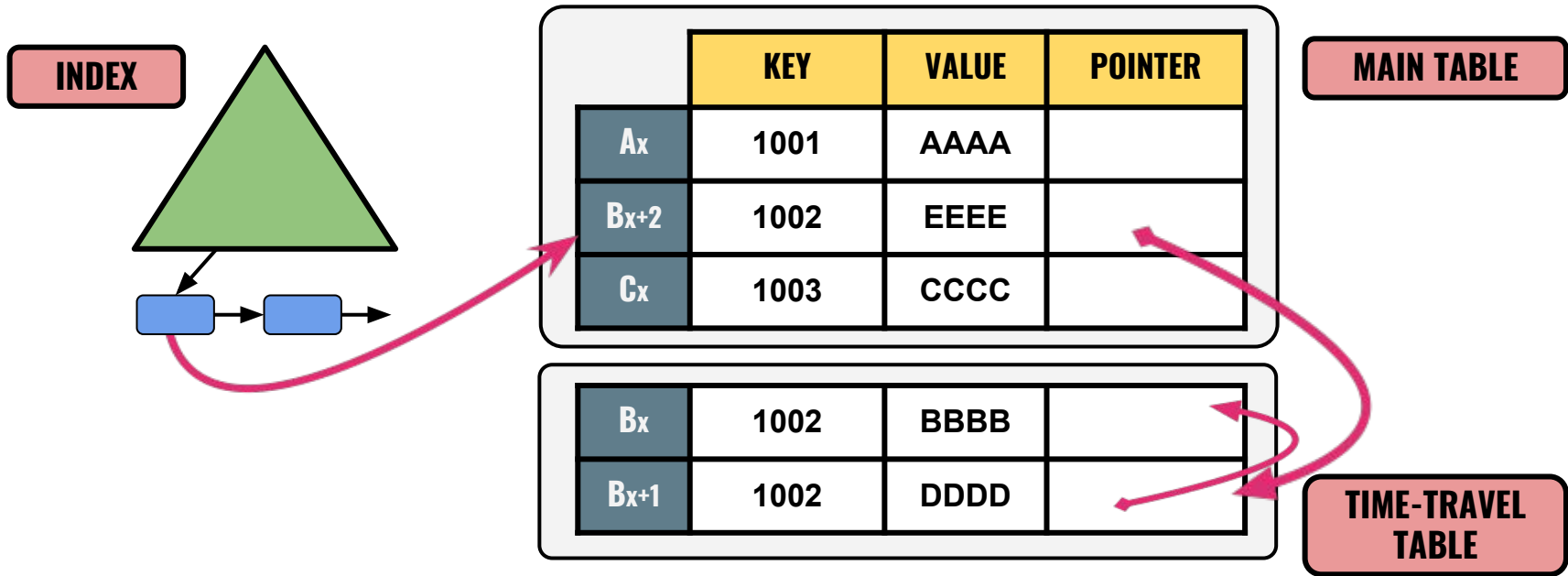
- ❏ New versions are appended to the same table space.



# VERSION STORAGE

## ❏ Approach #2: Time-Travel Storage

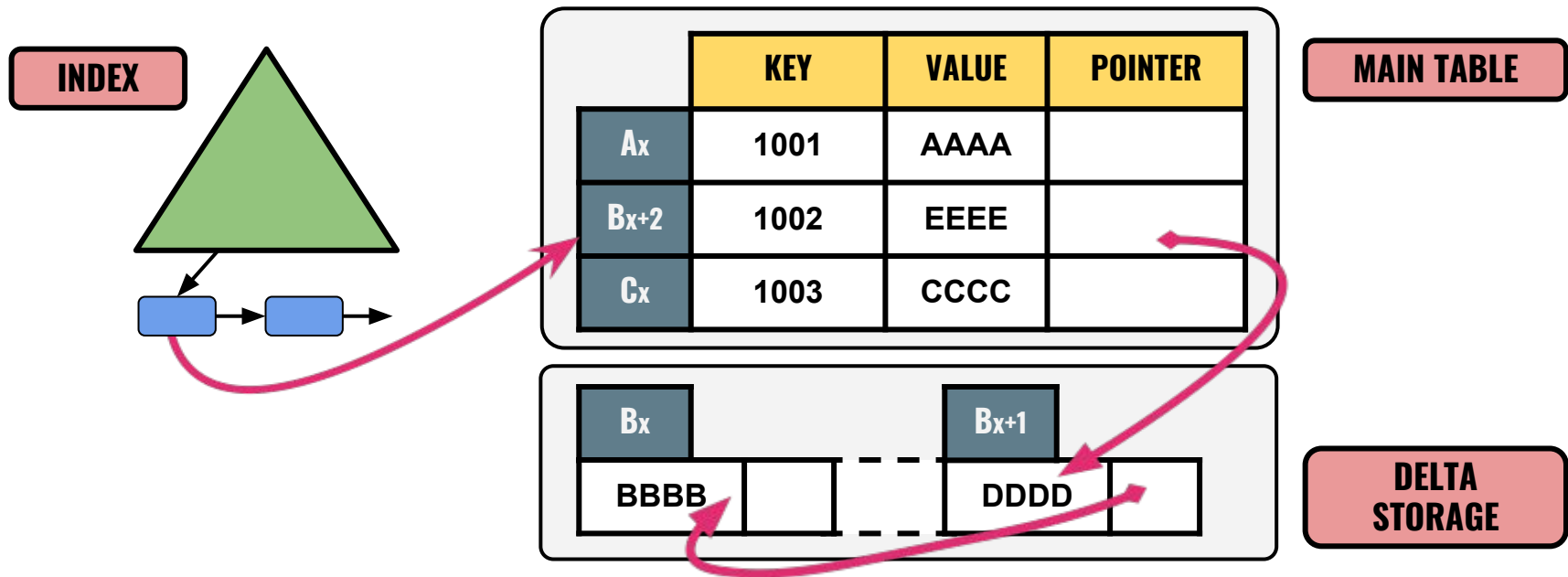
- ❏ Old versions are copied to separate table space.



# VERSION STORAGE

## ❏ Approach #3: Delta Storage

- ❏ The original values of the modified attributes are copied into a separate delta space.



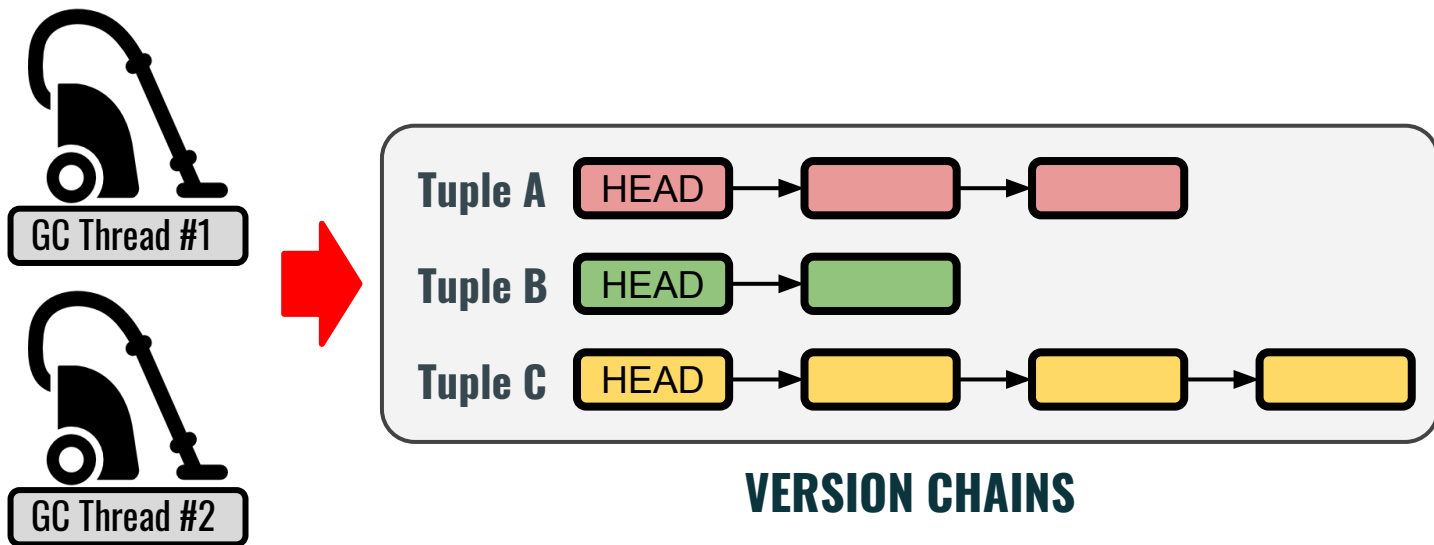
# GARBAGE COLLECTION

SCHEME	DBMS
Tuple-Level	 <p>ORACLE<sup>®</sup> MySQL<sup>®</sup> PostgreSQL<sup>®</sup> SAP HANA<sup>®</sup> MEMSQL<sup>®</sup> HYRISE<sup>®</sup> Microsoft<sup>®</sup> SQL Server<sup>®</sup> HEKATON<sup>®</sup> NUODB<sup>®</sup></p>
Transaction-Level	 <p>HyPer</p>

# GARBAGE COLLECTION

## ❏ Approach #1: Tuple-level

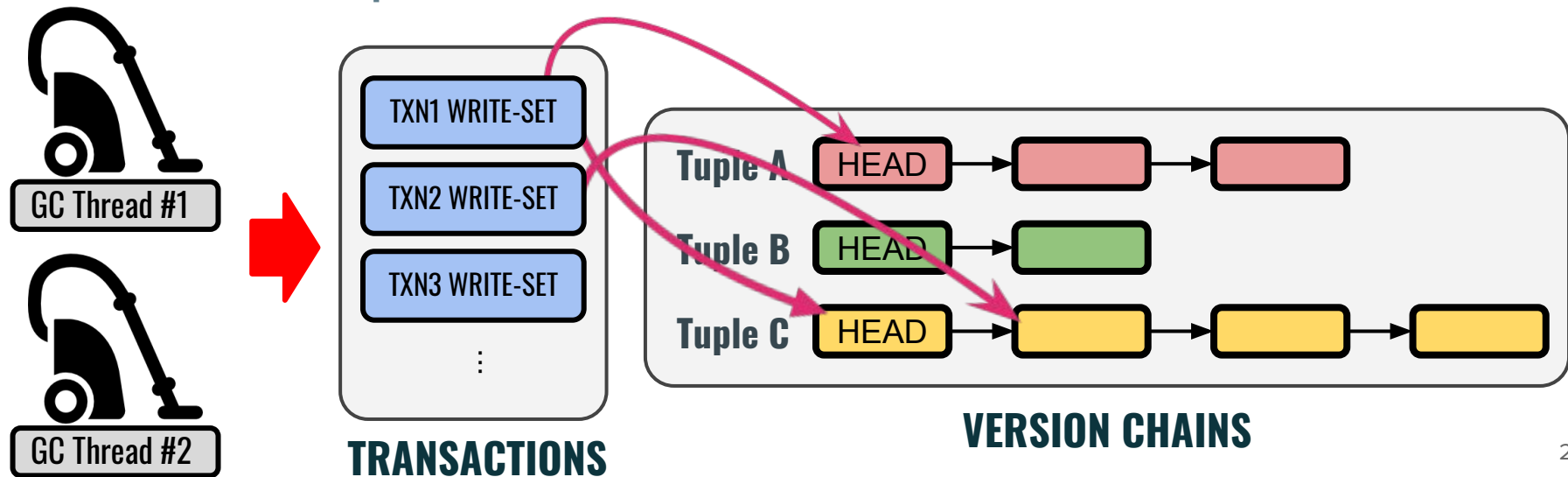
- ❏ Find old versions by examining tuples directly.




# GARBAGE COLLECTION

## ❏ Approach #2: Transaction-level

- ❏ Transactions keep track of their old versions so the DBMS does not have to scan tuples to determine visibility.



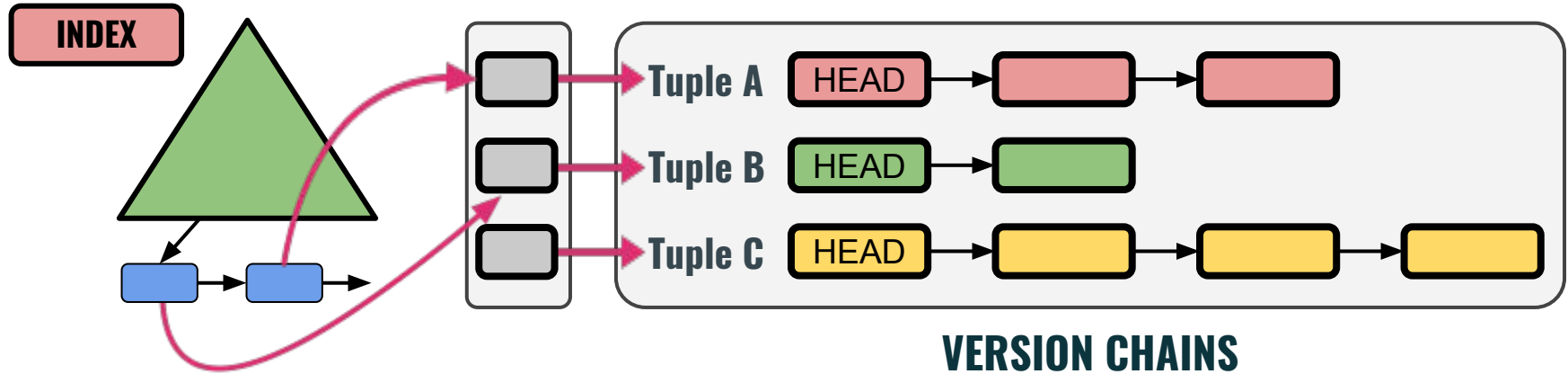
# INDEX MANAGEMENT

SCHEME	DBMS
Logical Pointers	    
Physical Pointers	   

# INDEX MANAGEMENT

## ❏ Approach #1: Logical Pointers

- ❏ Use a fixed identifier per tuple that does not change.

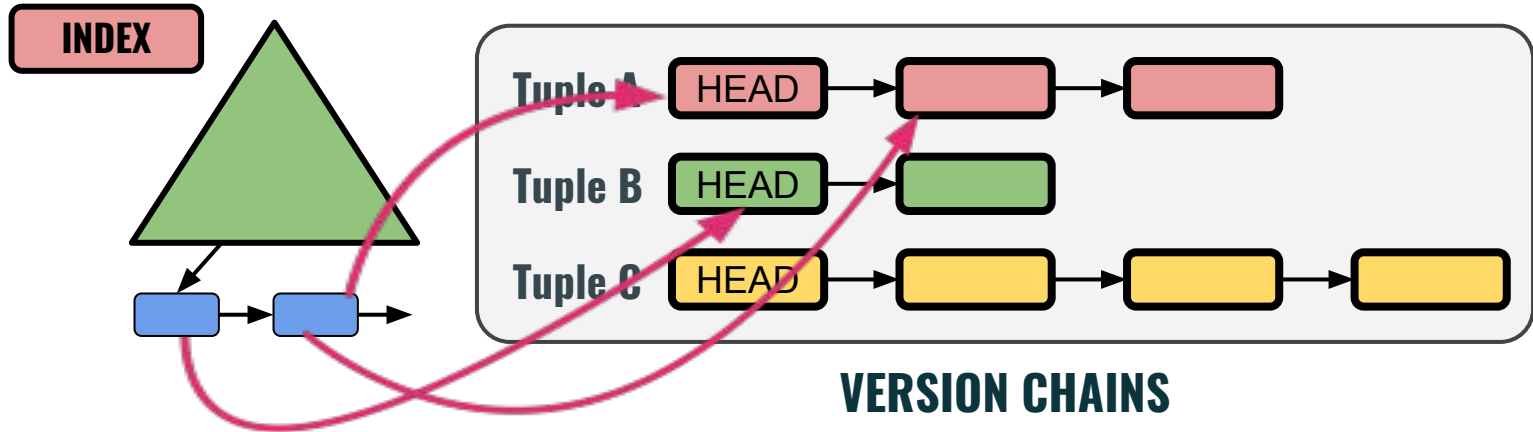




# INDEX MANAGEMENT

## ❏ Approach #2: Physical Pointers

- ❏ Use the physical address to the version chain head.



# EVALUATION

## ❏ Benchmarks

- ❏ YCSB
- ❏ TPC-C

## ❏ Configuration

- ❏ 4X Intel Xeon E7-4820 (40 cores)
- ❏ 128 GB DRAM



# EVALUATION

## Concurrency Control Protocol



**MVTO**



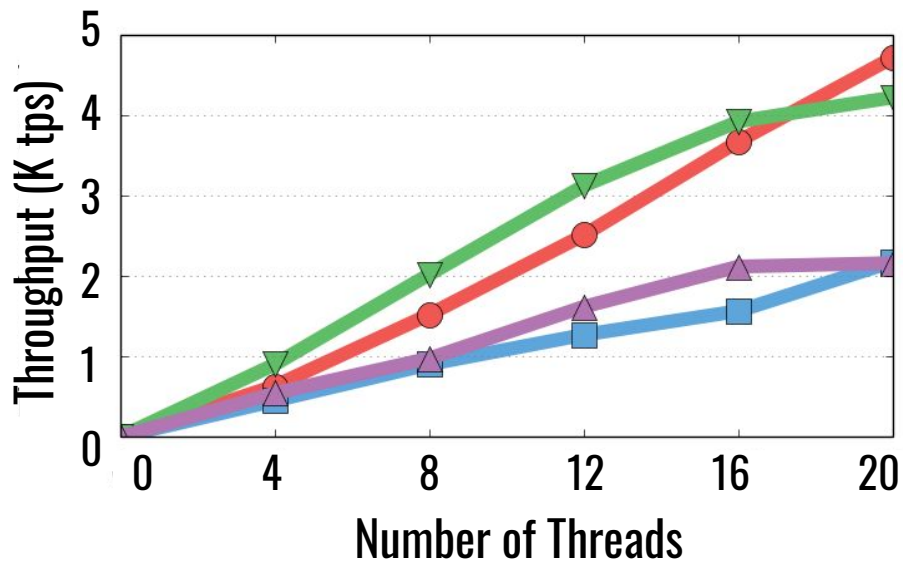
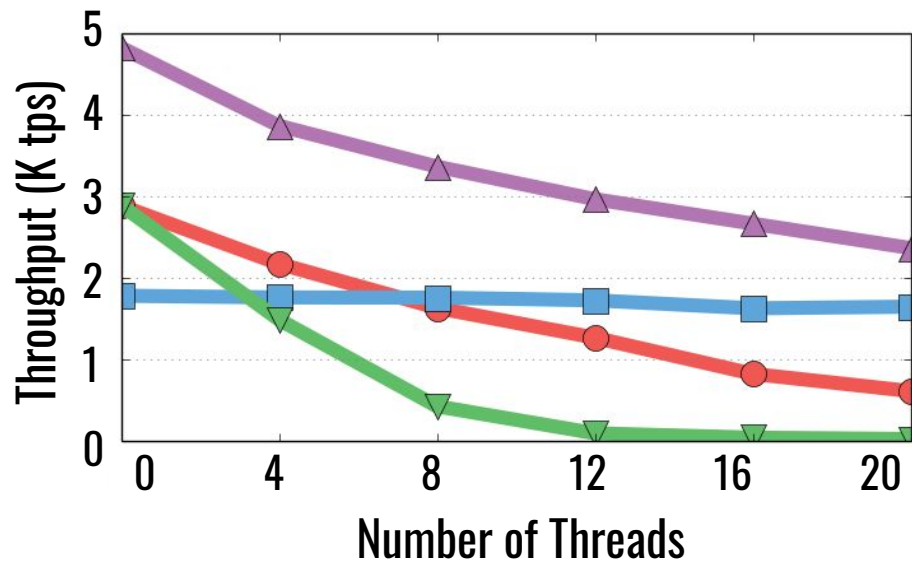
**MVOCC**



**MV2PL**



**SI+SSN**



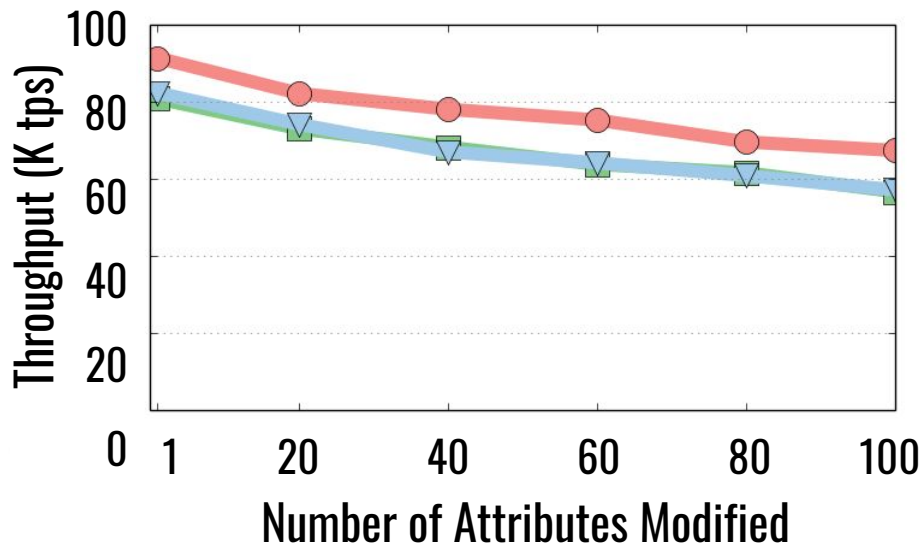
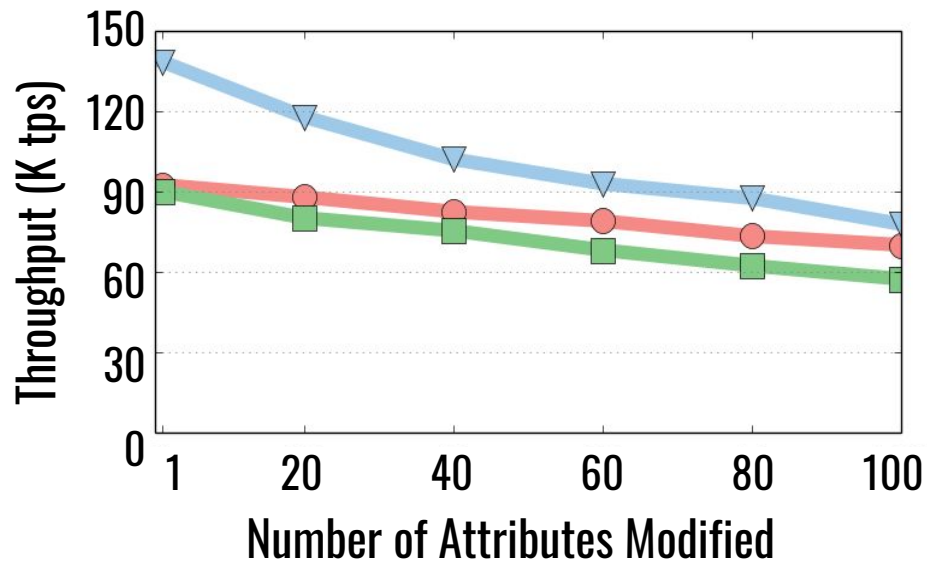
# EVALUATION

## Version Storage

●● Append-Only

■■ Time-Travel

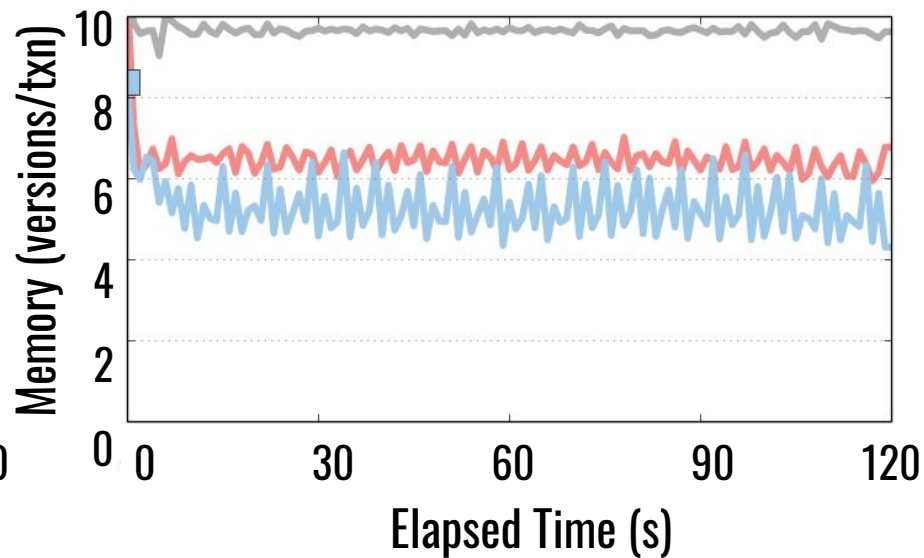
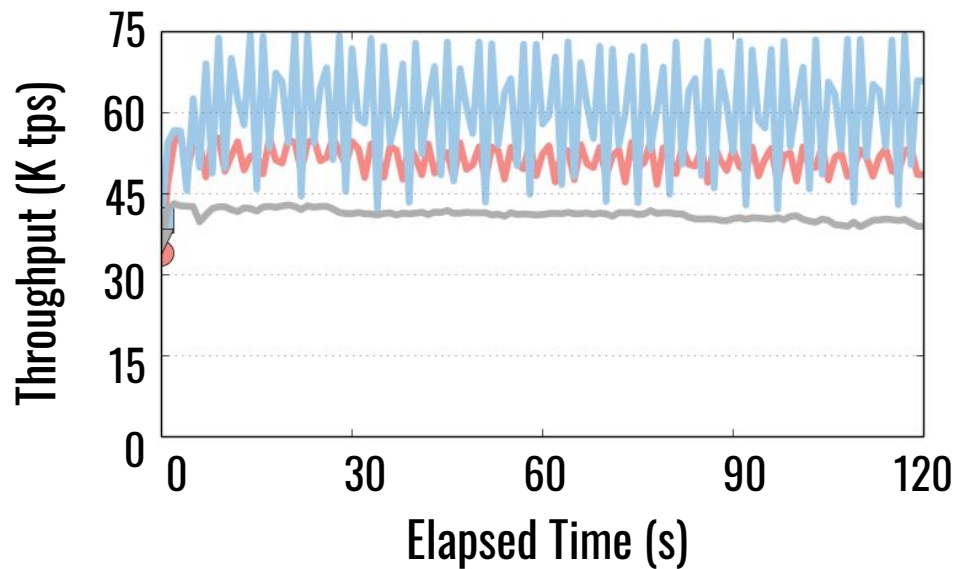
▼▼ Delta



# EVALUATION

## Garbage Collection

**■ Tuple-Level**    **■ Transaction-Level**    **■ Disable**

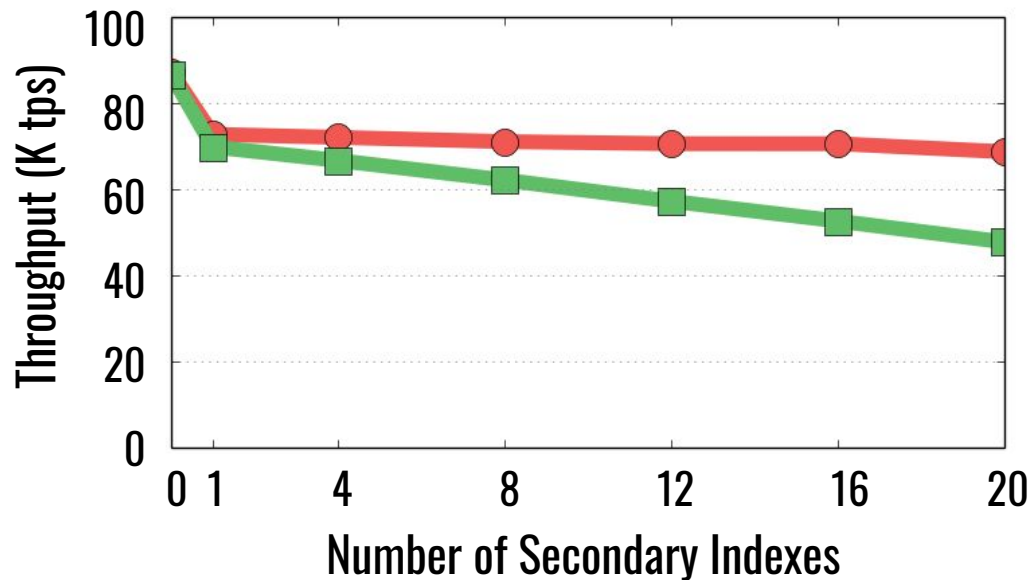


# EVALUATION

## Index Management

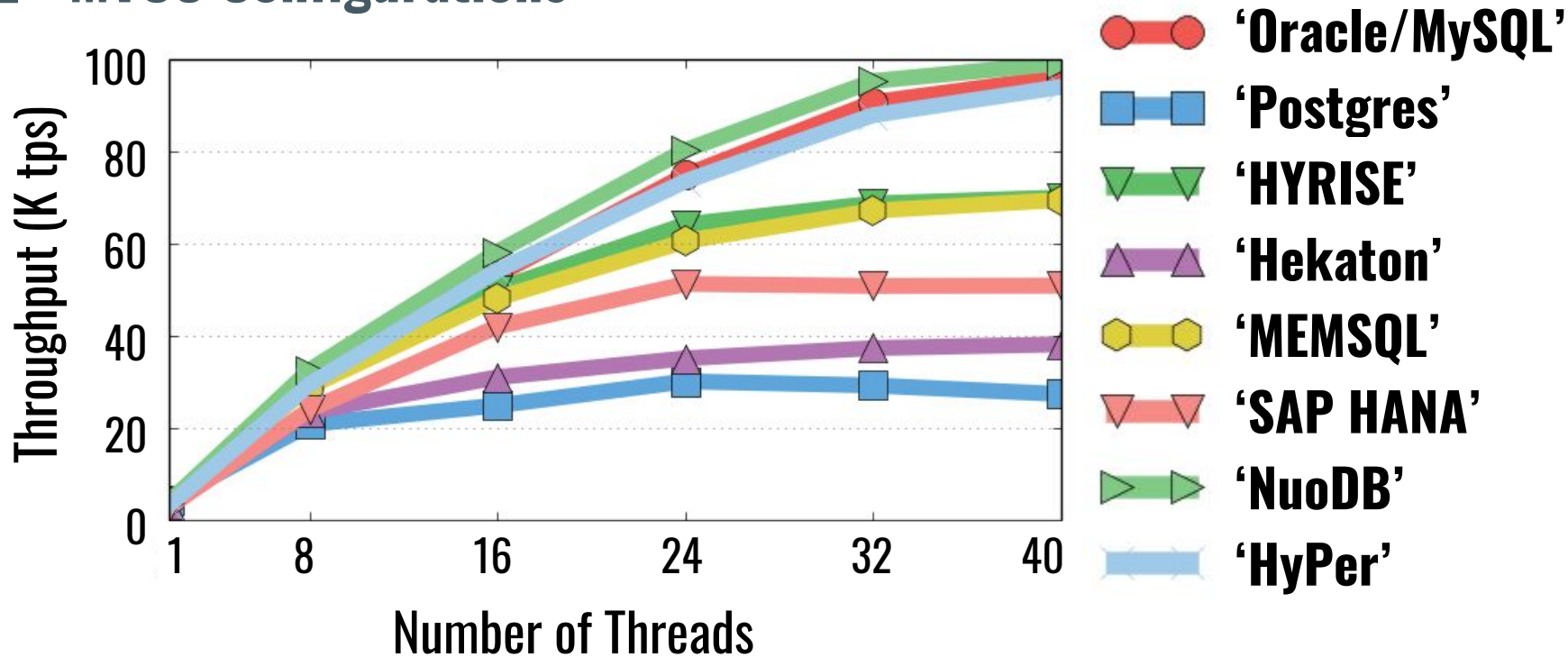
●● Logical Pointers

■ Physical Pointers



# EVALUATION

## ❑ MVCC Configurations



# CONCLUSION

- ❑ **Choosing the best MVCC scheme is challenging**
  - ❑ **Four design aspects**
  - ❑ **Multiple design decision combinations**
  - ❑ **Optimize for different objectives**



# END

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