Posets whose persistence modules are always interval decomposable and homological invariants

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Joint work with

Toshitaka Aoki (Kobe), Emerson G. Escolar (Kobe)

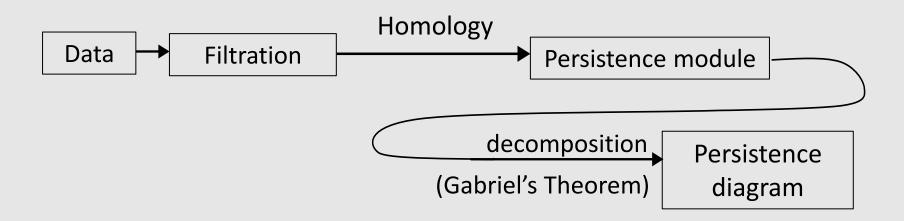
Preprint Summand-injectivity of interval approximations and monotonicity of interval global dimension. Toshitaka Aoki, Emerson G. Escolar, Shunsuke Tada. arXiv:2308.14979.

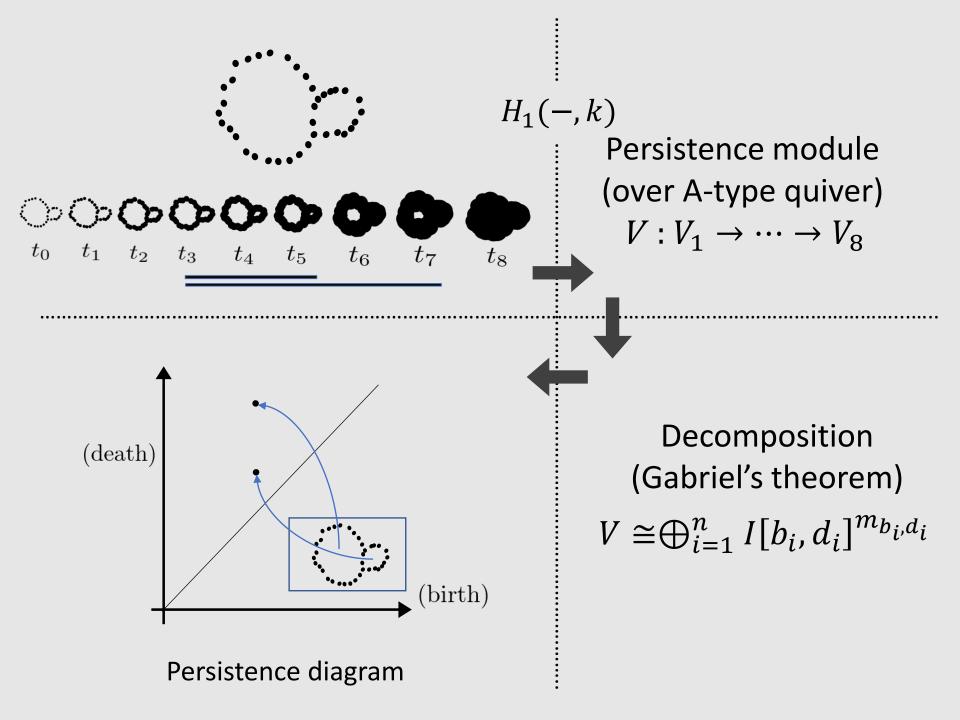
Contents

- Persistent homology?
- ■Three results

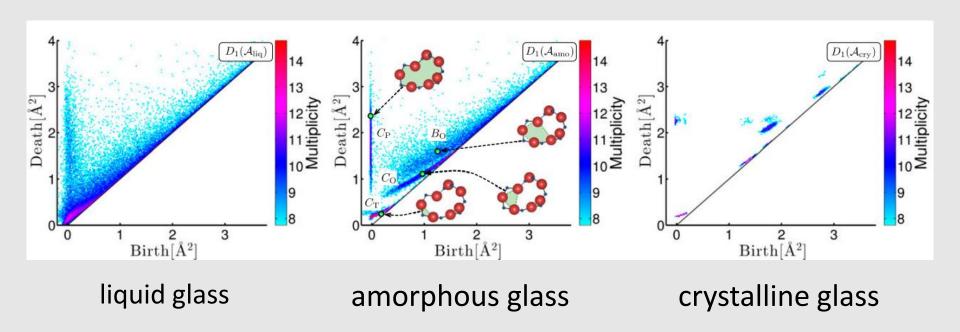
Persistent homology

Focus on the "persistence" of the shape (connected components, holes or voids) of data.





Example in material science



Hiraoka, Y., Nakamura, T., Hirata, A., Escolar, E. G., Matsue, K., & Nishiura, Y. (2016). Hierarchical structures of amorphous solids characterized by persistent homology. *Proceedings of the National Academy of Sciences*, *113*(26), 7035-7040.

Zigzag persistence module / Interval decomposability

- Carlsson, Gunnar, and Vin De Silva. "Zigzag persistence." Foundations of computational mathematics 10 (2010): 367-405.
- Botnan, Magnus, and Michael Lesnick. "Algebraic stability of zigzag persistence modules." Algebraic & geometric topology 18.6 (2018): 3133-3204.
- McDonald, R Neuhausler, R Robinson, M Larsen, L Harrington, H Bruna, M "Zigzag persistence for coral reef resilience using a stochastic spatial model." Journal of the Royal Society, Interface volume 20 issue 205 20230280-(23 Aug 2023).

... and others.

Question

P: Posets of type A



Modules over *P* are interval decomposable (by Gabriel's theorem).

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$$\bullet \longleftrightarrow \bullet \longleftrightarrow \bullet \longleftrightarrow \bullet$$

$$V \cong \bigoplus_{i=1}^n I[b_i, d_i]^{m_{b_i, d_i}}$$

P: ?



Modules over *P* are interval decomposable.

Results

- (1)Classification of posets
- (2) Direct summand injectivity
- (3) Monotonicity

Persistence module (1/2)

• Let P be a finite partially ordered set (poset). (we see it as a category by $a \le b \Leftrightarrow \exists ! a \rightarrow b$)

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Example.

 $V_1 \rightarrow V_2 \rightarrow V_3$ is a module over $P: 1 \rightarrow 2 \rightarrow 3$, where each V_i is a finite dimensional k-vector space.

Intervals (2/2)

A full subposet *I* of *P* is called *interval* if *I* is
(1) connected (the Hasse diagram of *I* is connected),
(2) convex (x ≤ y ≤ z, and x, z ∈ *I* imply y ∈ *I*).

•

•

$$I \subset P: \qquad \begin{array}{c} \bullet & \bullet & \bullet \\ & \searrow \nearrow & \nearrow \\ & & \searrow \nearrow \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\$$

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- For an interval I of P, the *interval module* k_I is defined by $k_I(p) \coloneqq k$ for $p \in I$, otherwise $k_I(p) \coloneqq 0$, $k_I(a \to b) \coloneqq \mathrm{id}_k$ for a, $b \in I$, otherwise 0.

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- A module is *interval decomposable* if the module decomposes into interval modules.

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$$V \cong \bigoplus_{i=1}^n I[b_i, d_i]^{m_{b_i, d_i}}$$

P: ?



Modules over *P* are interval decomposable.

Theorem 1 [Aoki-Escolar-T]

Let *P* be a connected finite poset. The following are equivalent.

- (a) Every module over *P* is interval decomposable.
- (b) The Hasse diagram of *P* is one of the following form:

$$1 \leftrightarrow 2 \leftrightarrow \cdots \leftrightarrow n$$

$$0 \qquad 0 \qquad 1$$

$$1' \rightarrow 2' \rightarrow \cdots \rightarrow n$$

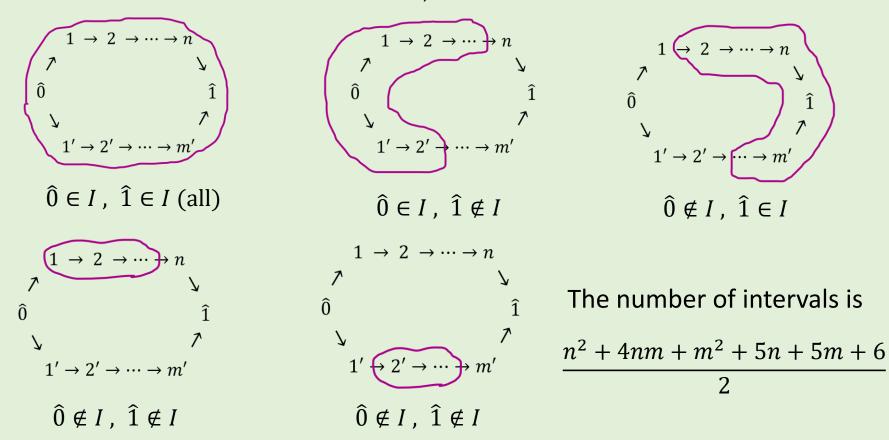
$$1' \rightarrow 2' \rightarrow \cdots \rightarrow m'$$

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$$C_{n,m} : \underline{bipath\ poset}\ of\ size\ (n,m)\ (\underline{commutative\ cycle})$$

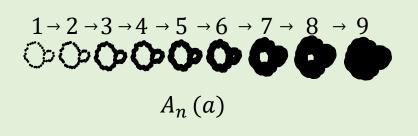
Theorem 1 [Aoki-Escolar-T]

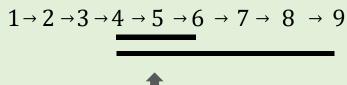
In particular, the intervals in $C_{n,m}$ are following forms.

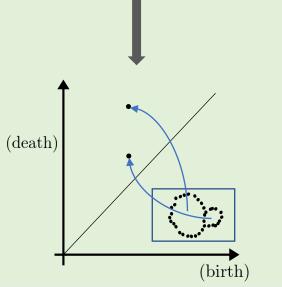


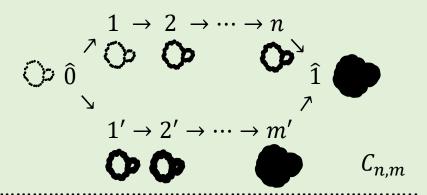
^{*} We introduce the sketch of proof after in Theorem 3.

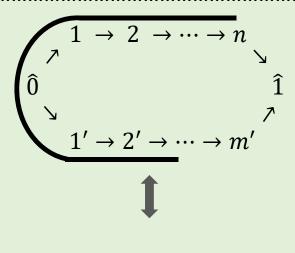
Comparison between type A and $C_{n,m}$





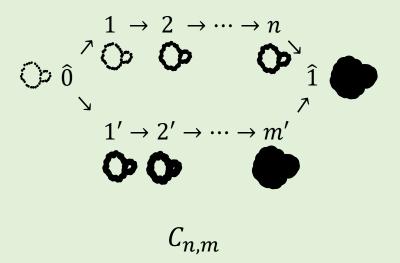


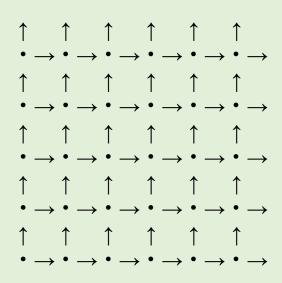






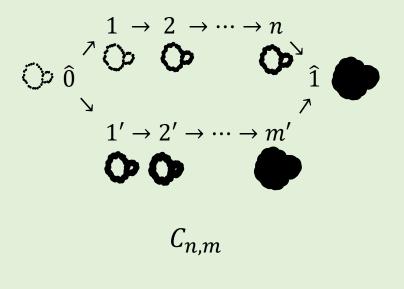
Idea for application of $C_{n,m}$

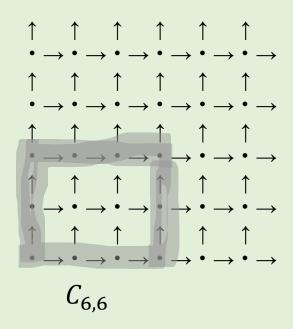




- We can get a part of information of multi-parameter persistence modules by restricting multi-filtration to $\mathcal{C}_{n,m}$ (like fibered barcode?).
- By using two functions, we can see the robustness of shape of data in terms of the two functions .

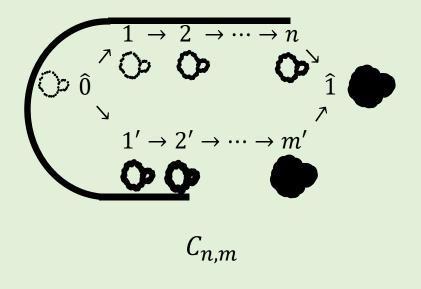
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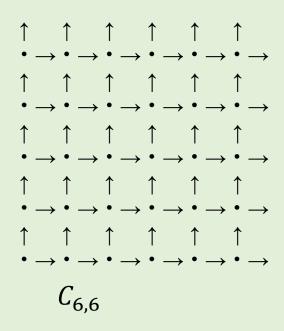




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Idea for application of $C_{n,m}$



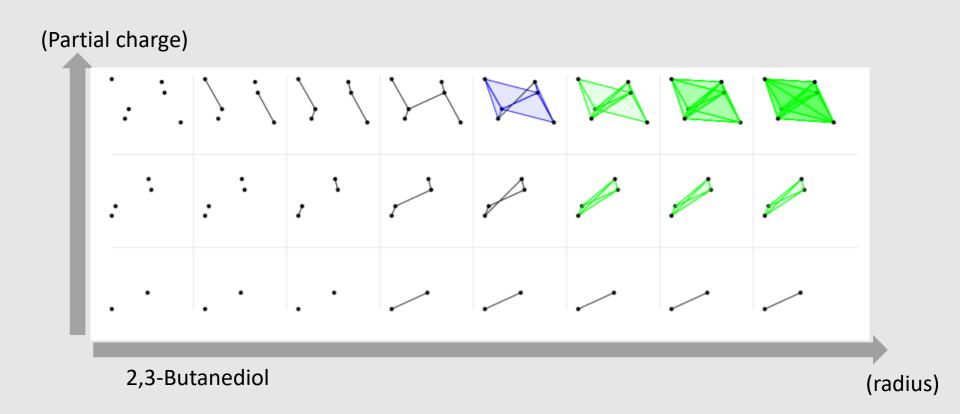


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Results

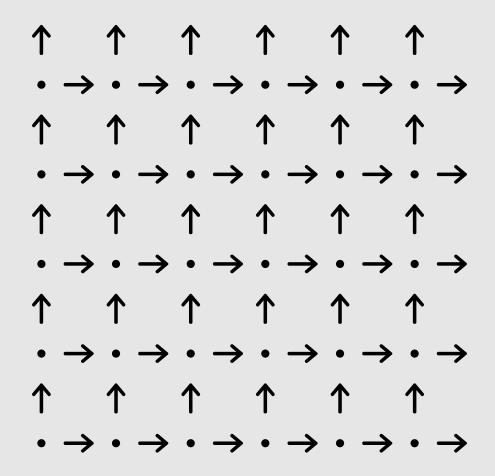
- (1)Classification of posets
- (2) Direct summand injectivity
- (3) Monotonicity

Multi-parameter persistent homology



Keller B, Lesnick M, Willke TL. Persistent Homology for Virtual Screening. ChemRxiv. Cambridge: Cambridge Open Engage; 2018

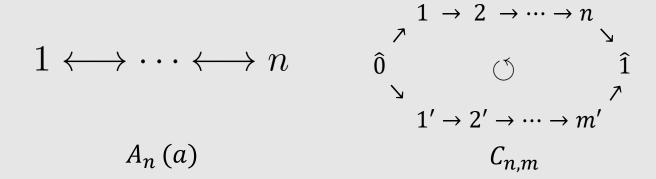
Multi-parameter persistent homology



It is difficult to classify all the indecomposable module (wild representation type)

Multi-parameter persistent homology

M.Buchet, Emerson G. Escolar "Every 1D Persistence Module is a Restriction of Some Indecomposable 2D Persistence Module" *Journal of Applied and Computational Topology*



(Modules are always interval decomposable)



We want to understand complex modules (obtained from data) over posets, using "good" modules (interval modules).

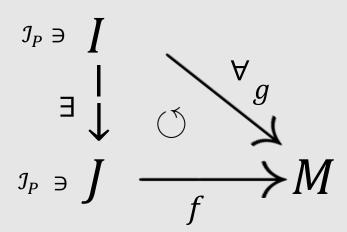
Interval approximation

- Benjamin Blanchette, Thomas Brüstle, and Eric J Hanson. "Homological approximations in persistence theory." *Canadian Journal of Mathematics*, pages 1–38, 2021.
- · Hideto Asashiba, Emerson G Escolar, Ken Nakashima, and Michio Yoshiwaki. "Approximation by interval-decomposables and interval resolutions of persistence modules." *Journal of Pure and Applied Algebra*, 227(10):107397, 2023.

Interval approximation (1/1)

- \mathcal{I}_P is the set of interval decomposable modules over P.
- *M* is a module over *P*.

An *interval approximation of M* is a morphism $f: J \to M$ with $J \in \mathcal{I}_P$ s. t. for any $g: I \to M$ with $I \in \mathcal{I}_P$ factor through f.



An *interval cover* of M is an interval approximation such that the number of direct summands of the domain is smaller than that of other interval approximations (uniquely determined).

Question

How can we calculate interval cover of any modules (easily)?

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Calculation of interval cover of a module M

- (1) We have an interval approximation $\bigoplus_{I \in \mathbb{I}(P)} k_I^{m_I} \to M$, where m_I is k-dimension of $\operatorname{Hom}_k(k_I, M)$ and $\mathbb{I}(P)$ is the set of all the intervals in the poset P.
- (2) We reduce the direct summands of the above interval approximation until we obtain the interval cover.

We give a helpful observation to the Question.

Theorem 2 [Aoki-Escolar-T]

Let P be a finite poset and M be a module over P. For any interval cover of M $f = (f_i) : \bigoplus_{i=1}^n k_{I_i} \to M,$

the following holds.

- (1) f is surjective.
- (2) Each $f_i: k_{I_i} \to M$ is injective.
- (3) For every $a \in P$, we have M(a) = 0 if and only if $(\bigoplus_{i=1}^{n} k_{I_i})(a) = 0$.

In particular, every k_{I_i} is given by an interval submodule of M.

Remark

Recently, [Asashiba, 2023, Proposition 4.8, arXiv:2307.06559] gave the essentially same result (see also [Blanchette-Brüstle-Hanson, Proposition 6.7, 2021, Canadian Journal of Mathematics, 1-38]).

Results

- (1)Classification of posets
- (2) Direct summand injectivity
- (3) Monotonicity

Resolution dimension (1/2)

An interval resolution of M is an exact sequence

$$0 \to J_m \xrightarrow{g_m} \cdots \to J_2 \xrightarrow{g_2} J_1 \xrightarrow{g_1} J \xrightarrow{f} M \to 0,$$

$$\cdots K_3^{\iota_3 \nearrow f_2 \nearrow K_2} K_2^{\iota_2 \nearrow f_1 \nearrow K_1} K_1^{\iota_1 \nearrow N}$$

then we say that the *interval resolution dimension of* M is m and write int-res-dim M=m.

Interval resolution global dimension (2/2)

interval resolution global dimension of P is
 int-res-gldim(P) := sup{int-res-dim(M) | M: modules over P}

* [Asashiba-Escolar-Nakashima-Yoshiwaki, 23] show that int-res-gldim(P) < ∞ for any poset finite poset P.

Interval resolution global dimension (2/2)

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Remark

int-res-gldim(P) is zero if and only if the Hasse diagram of P is either (i) or (ii), where $\sum_{i=1,\dots,n} 1_{i} + \sum_{j=1}^{n} 1_{j} + \sum_{i=1}^{n} 1_{j} + \sum_{i=1}^{n} 1_{i} + \sum_{j=1}^{n} 1_{j} + \sum_{i=1}^{n} 1_{i} + \sum_{j=1}^{n} 1_{j} + \sum_{j=1}^$

Theorem 3 [Aoki-Escolar-T]

Let P be a finite poset. For any full subposet Q of P, the following inequality holds.

 $int-res-gldim(Q) \le int-res-gldim(P)$.

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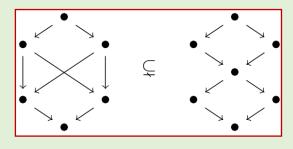
Poset

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Remark

The above monotonicity does not hold for (usual) global dimension in general [Igusa-Zacharia, 1990].



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Global dimension	3	>	2	
Interval global dimension	1	<	2	(over a field with two elements)

Theorem 1 [Aoki-Escolar-T]

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- (a) Every module over P is interval decomposable. $1 \rightarrow 2 \rightarrow \cdots \rightarrow n$
- (b) The Hasse diagram of P is $1 \longleftrightarrow \cdots \longleftrightarrow n$ or $\hat{0} \circlearrowleft \hat{1} C_{n,m}$. $A_n(a) \qquad \qquad 1' \to 2' \to \cdots \to m'$

Sketch of proof (a \Rightarrow b)

- The Hasse diagram of P does not have a vertex with degree 3(by Theorem 3).
- P is either A_n or \tilde{A}_m for some n and m.
- P must be A_n or $C_{n,m}$ for some n and m.

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 \Re We prove the converse (interval decomposability of module over $C_{n,m}$) by using theory of special biserial algebra.

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$$1 \longleftrightarrow \cdots \longleftrightarrow n : A_n$$

m+1 $1 \longleftrightarrow \cdots \longleftrightarrow m : \tilde{A}_m$

 \aleph We prove the converse (interval decomposability of module over $C_{n,m}$) by using theory of special biserial algebra.

Summary

- (1) We classified finite posets whose modules are always interval decomposable.
- (2) We show that restriction of each direct summand of interval cover is injective.

(It makes calculation of interval cover easier.)

(3) We show the monotonicity of int-res-gldim. (This is used to show the first result.)

Discussion

- When $C_{n,m}$ is useful in TDA? Stability?
- Does int-res-gldim depend on the characteristic of fields?
- When do we have int-res-gldim(Q) = int-res-gldim(P) for $Q \subseteq P$?
- Can we calculate interval cover easily?
- Computation using GAP package QPA("Quiver and Path Algebras") and "pmgap" by E. G. Escolar to calculate modules over poset.

Thank you for your attention!



Our paper

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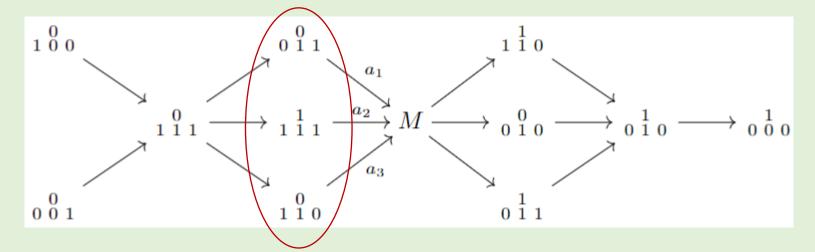
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Sketch of proof (Interval resolution global dimension $< \infty$)

- (0) Let G be a direct sum of all interval modules k_I over P. Note that all indecomposable projective (resp., injective) are interval modules.
- (1)[AENY, 23] shows that any submodule of interval module is an interval decomposable module.
- (2) Then, the endmorphism ring $\Gamma \coloneqq \operatorname{End}(G)$ is a left strongly quasihereditary algebra by [Ringel, 09, Theorem 5] (see also [Iyama, 03]). In particular, Γ has the finite global dimension.
- (3) We have int-res-gldim(P) = gldim (Γ) 2 < ∞ .([Erdmann-Holm-lyama-Schröer, 17])
 - See [AENY, 23] for the detail.

Example of interval cover and injectivity

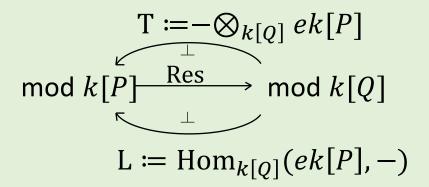


An interval cover of *M* is given by

For a full subposet Q of P, we have an isomorphism

$$k[Q] \cong ek[P]e$$

of *k*-algebras, where $e \coloneqq \Sigma_{x \in Q} e_x$. It induces adjoint functors



- Res preserves interval decomposability of modules.
- T and L do NOT preserve interval decomposability of modules in general.

We find a functor Θ that sends to interval modules over Q to interval modules over P by using T and L.

The functor ⊕

Using adjoint functors, we have

For a given module $M \in \text{mod } k[Q]$, let

$$\Theta(M) := \operatorname{Im}(\theta_M).$$

It gives rise to a functor Θ. It is called *intermediate extension* in [Kuhn, 94], and *prolongedment intermédiare* in [Beilison-Bernstein-Deligne, 82].

Proposition For a given interval I of Q, let k_I be the corresponding interval k[Q]-module. Then, we have

$$\Theta(k_I) \cong k_{\operatorname{conv}(I)},$$

where conv(I) is the smallest interval of P containing I.

The functor ⊕

We obtain a pair of functors

$$\operatorname{\mathsf{mod}} k[P] \ \ \operatorname{\mathsf{Res}} \ \ \operatorname{\mathsf{mod}} k[Q]$$

satisfying the following properties:

- (i) Res preserves interval decomposability of modules.
- (ii) Θ sends interval modules to interval modules by Proposition.
- (iii) $1_{\text{mod } k[Q]} \cong \text{Res} \circ \Theta$.

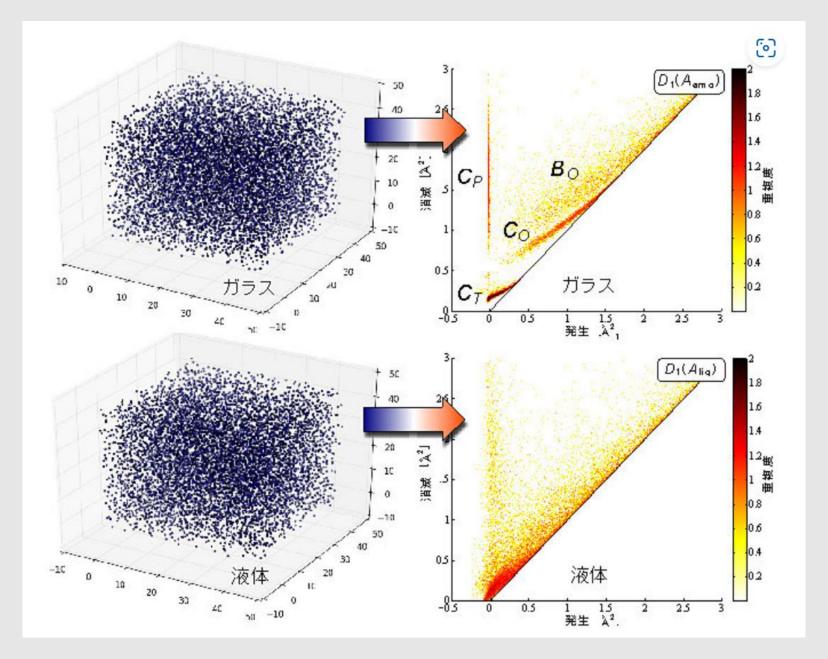
Proposition For any $M \in \text{mod } k[Q]$, we have the following inequality int-res-dim $(M) \leq \text{int-res-dim}(\Theta(M))$.

Since M is an arbitrary module, we obtain the desired inequality int-res-gldim $(k[Q]) \leq \text{int-res-gldim}(k[P])$.

Remark(Decomposition algorithm)

In [Aoki-Escolar-T], we prove the interval decomposability of modules over $C_{n,m}$ in a non-constructive way. (we used the knowledge of special biserial algebra)

Now, we have an algorithm for decomposition.



共同発表: ガラスの「形」を数学的に解明~トポロジーで読み解く無秩序の中の秩序~ (jst.go.jp)

Representations of posets (homological algebra)

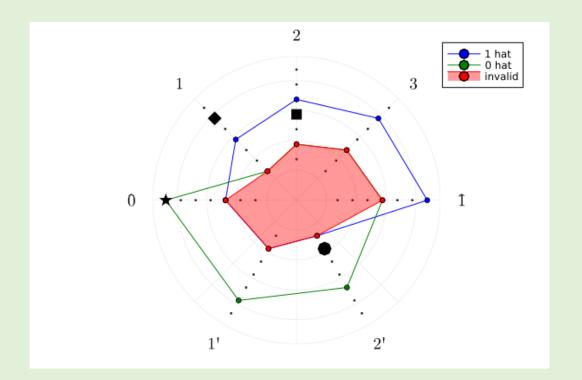
- · Magnus Bakke Botnan, Steffen Oppermann, and Steve Oudot. "Signed barcodes for multi-parameter persistence via rank decompositions and rank-exact resolutions." In International Symposium on Computational Geometry, 2021
- Benjamin Blanchette, Thomas Brüstle, and Eric J Hanson. "Homological approximations in persistence theory." *Canadian Journal of Mathematics*, pages 1–38, 2021.
- · Hideto Asashiba, Emerson G Escolar, Ken Nakashima, and Michio Yoshiwaki. "Approximation by interval-decomposables and interval resolutions of persistence modules." *Journal of Pure and Applied Algebra*, 227(10):107397, 2023.

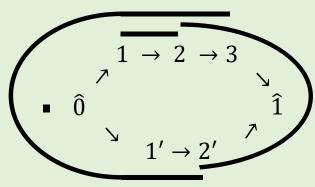
... and others.

We want to understand complex modules (obtained from data) over posets, using "good" modules (interval modules).

Interval approximation

Persistence diagram for $C_{n,m}$





- The longer persistence, the closer to the center.
- Place the points of $C_{n,m}$ on the circumference clockwise from 0hat.