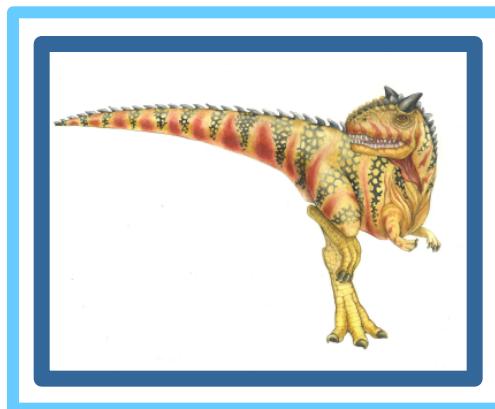
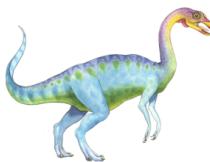


Chapter 2: Operating-System Structures





Operating System Structure

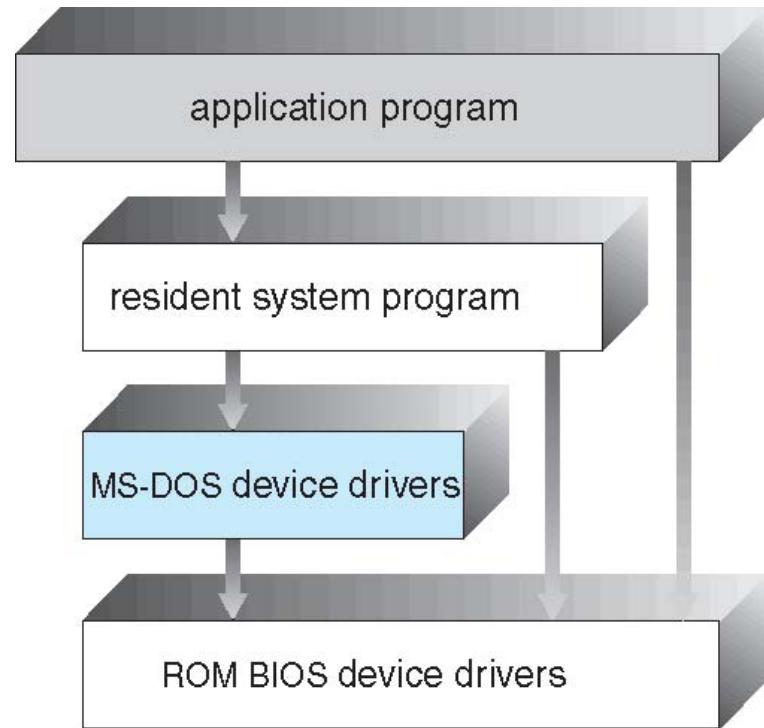
- General-purpose OS is very large program
- Types of structures are,
 - Simple – MS-DOS
 - Monolithic – UNIX
 - Layered – an abstraction
 - Microkernel – Mach
 - Modules





Simple Structure -- MS-DOS

- MS-DOS – written to provide the most functionality in the least space
 - Not divided into modules
 - Although MS-DOS has some structure, its interfaces and levels of functionality are not well separated





Monolithic Structure -- UNIX

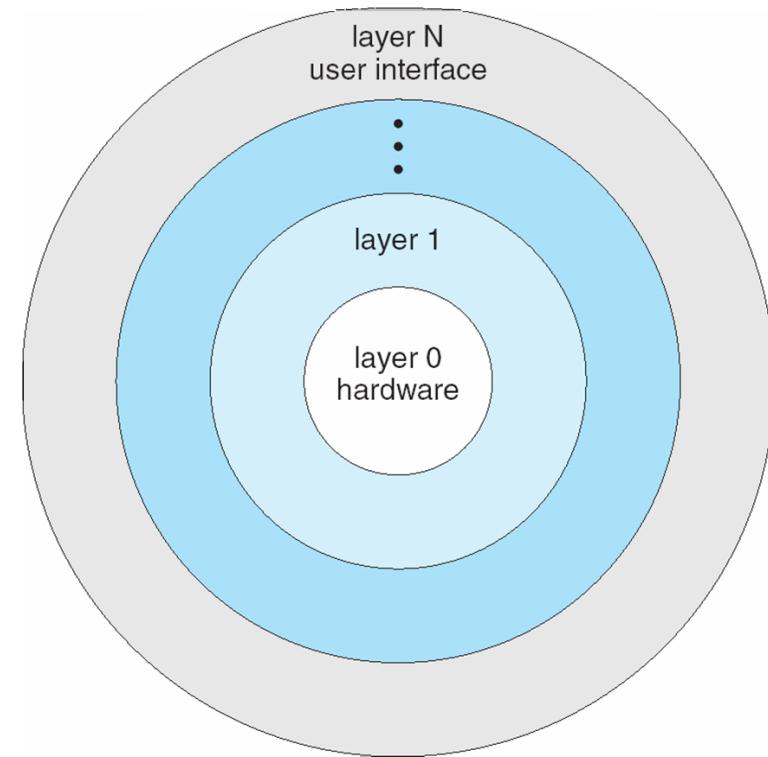
- UNIX – limited by hardware functionality, the original UNIX operating system had limited structuring.
- The UNIX OS consists of two separable parts
 - Systems programs
 - The kernel
 - Consists of everything below the system-call interface and above the physical hardware
 - Provides the file system, CPU scheduling, memory management, and other operating-system functions; a large number of functions for one level





Layered Approach

- The operating system is divided into a number of layers (levels), each built on top of lower layers. The bottom layer (layer 0), is the hardware; the highest (layer N) is the user interface.
- With modularity, layers are selected such that each uses functions (operations) and services of only lower-level layers

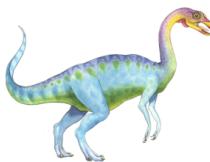




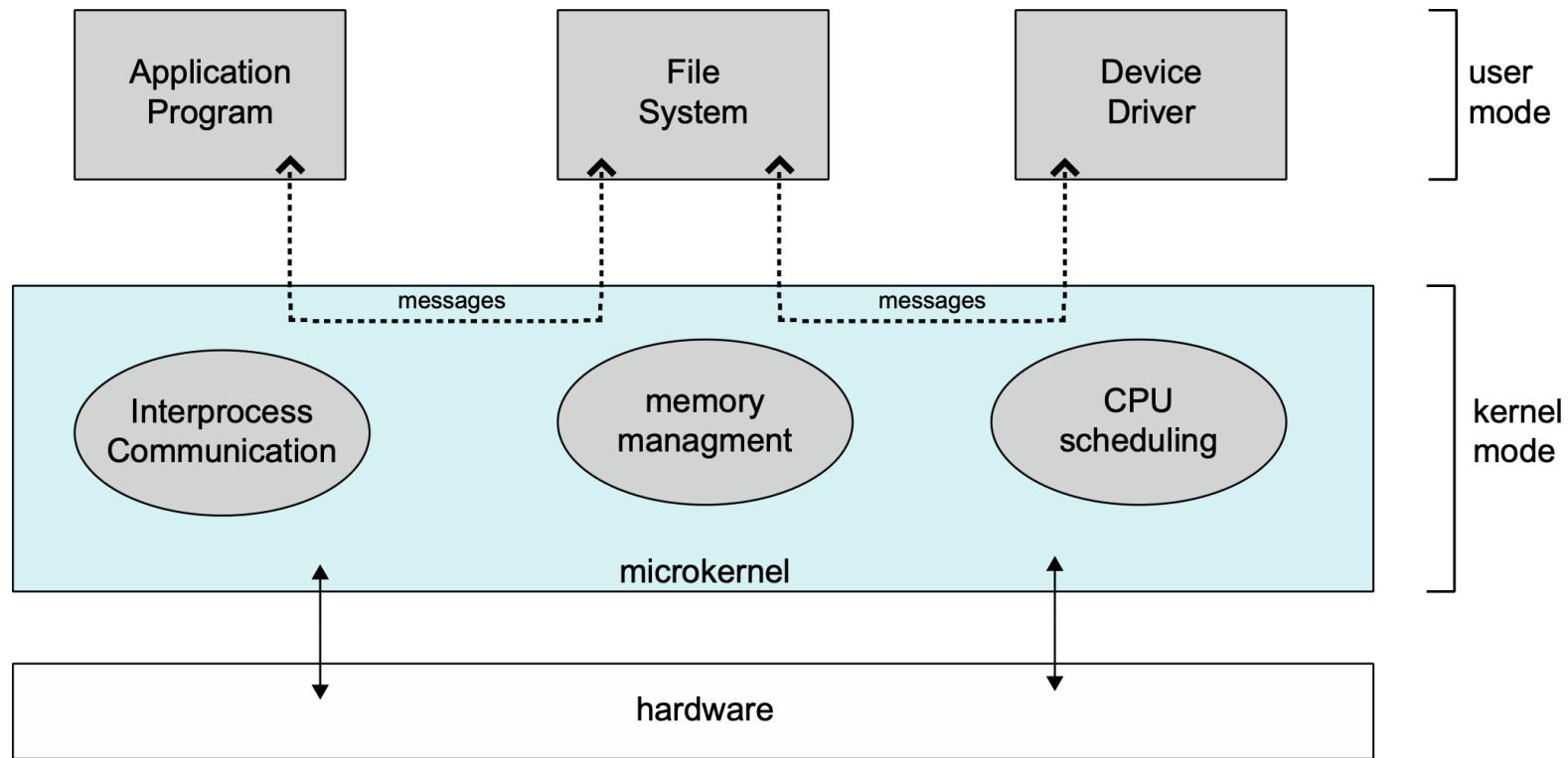
Microkernel System Structure

- Moves as much from the kernel into user space
- **Mach** example of **microkernel**
 - Mac OS X kernel (**Darwin**) partly based on Mach
- Communication takes place between user modules using **message passing**
- Benefits:
 - Easier to extend a microkernel
 - Easier to port the operating system to new architectures
 - More reliable (less code is running in kernel mode)
 - More secure
- Detriments:
 - Performance overhead of user space to kernel space communication





Microkernel System Structure





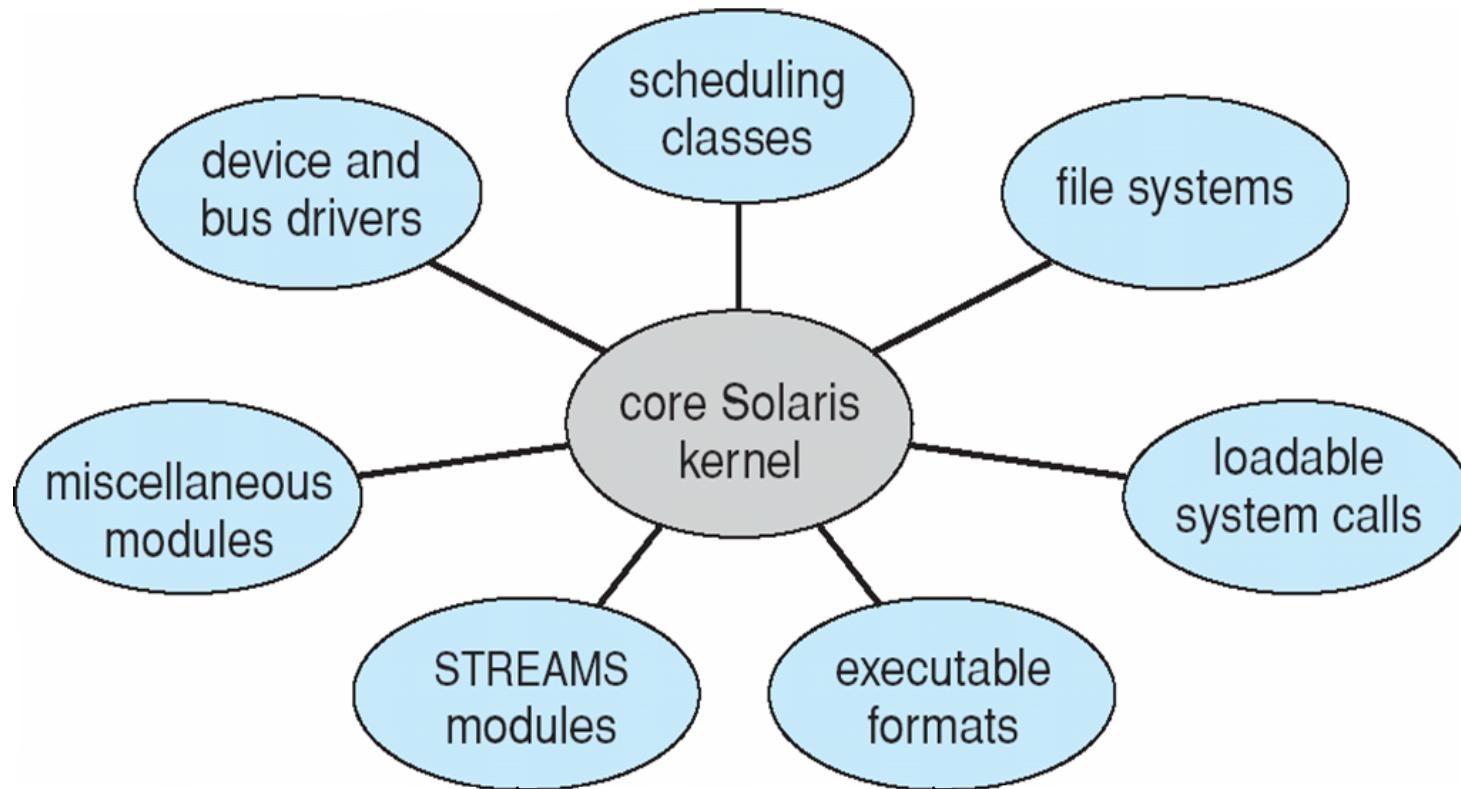
Modules

- Many modern operating systems implement **loadable kernel modules** — **LKMs**
 - Uses object-oriented approach
 - The core services are designed with the kernel
 - Additional services are added as LKMs
 - Overall, similar to layers but with more flexible
 - Linux, Solaris, etc





Solaris Modular Approach





Hybrid Systems

- Most modern operating systems are actually not one pure model
 - Hybrid combines multiple approaches to address performance, security, usability needs
 - Linux and Solaris kernels in kernel address space, so monolithic, plus modular for dynamic loading of functionality
 - Windows mostly monolithic, plus microkernel
 - Eg: Apple Mac OS, iOS, Android





Operating-System Debugging

- **Debugging** is finding and fixing errors, or **bugs**
- OS generate **log files** containing error information
- Failure of an application can generate **core dump** file capturing memory of the process
- Operating system failure can generate **crash dump** file containing kernel memory



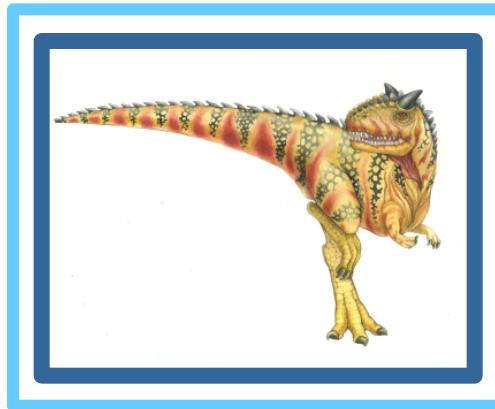


System Boot

- When power initialized on system, execution starts at a fixed memory location
 - Firmware ROM used to hold initial boot code
- Operating system must be made available to hardware so hardware can start it
 - Small piece of code – **bootstrap loader**, stored in **ROM** locates the kernel, loads it into memory, and starts it
 - Sometimes two-step process where **boot block** at fixed location loaded by ROM code, which loads bootstrap loader from disk



End of Chapter 2



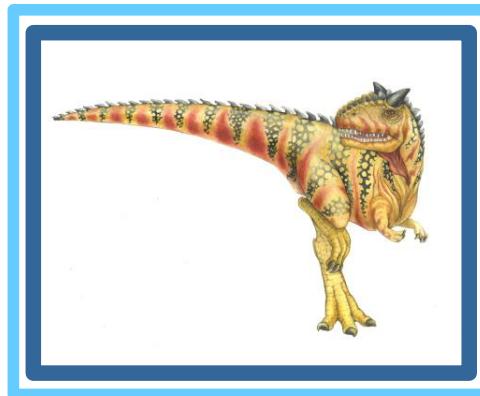


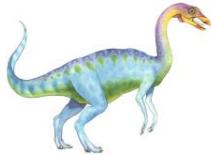
Quiz

- 1. The OS design requirements belong to two groups —> and**
- 2. What is the main principle to follow while implementing an OS?**
- 3. Different structures ?**
- 4. What are the disadvantages of layered approach?**
- 5. What is the communication type used in microkernel system structure?**
- 6. The process of finding and fixing of errors is called as?**
- 7. The small piece of code that help to locate the kernel is known as?**



Chapter 3: Processes

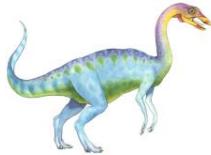




Chapter 3: Processes

- Process Concept
- Process Scheduling
- Operations on Processes
- Interprocess Communication
- Examples of IPC Systems
- Communication in Client-Server Systems

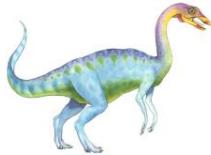




Objectives

- To introduce the notion of a process -- a program in execution, which forms the basis of all computation
- To describe the various features of processes, including scheduling, creation and termination, and communication
- To explore interprocess communication using shared memory and message passing
- To describe communication in client-server systems

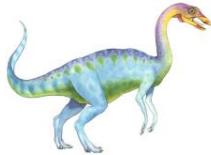




Process

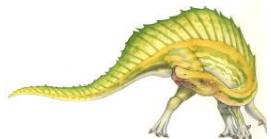
- Program is **passive** entity stored on disk (**executable file**), process is **active**
 - Program becomes process when executable file loaded into memory
- Execution of program started via GUI mouse clicks, command line entry of its name, etc
- One program can be several processes
 - Consider multiple users executing the same program

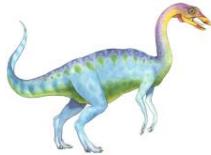




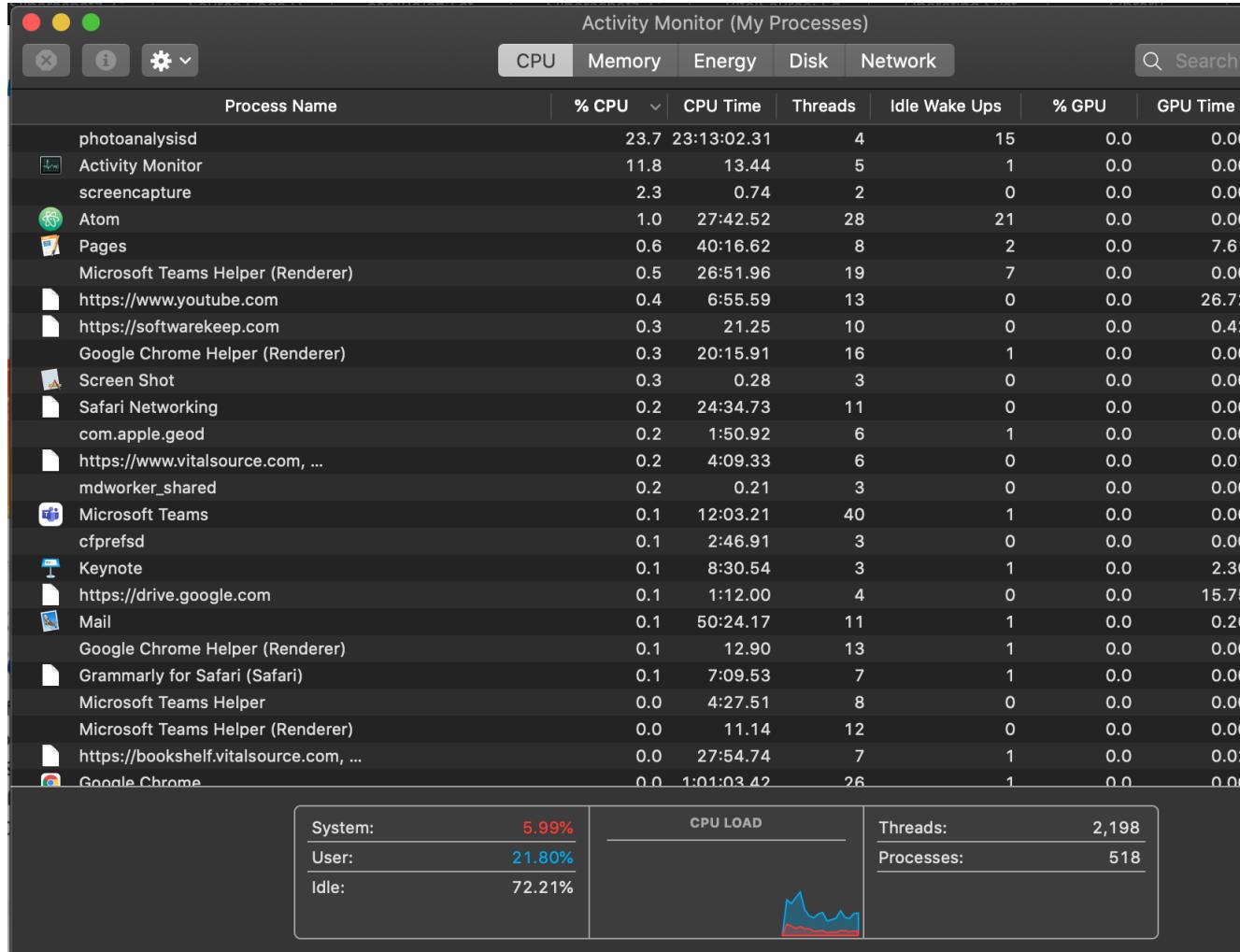
Process Vs Threads

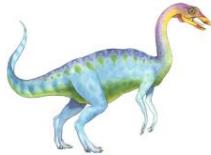
- An operating system executes a variety of programs
 - **Process** – a program in execution
 - **Thread** — Unit of execution within a process, A single process can have any number of threads.
-
- Task manager in Windows
 - Activity monitor in Mac





Process Vs Threads





Process State

- As a process executes, it changes **state**
 - **new**: The process is being created
 - **running**: Instructions are being executed
 - **waiting**: The process is waiting for some event to occur
 - **ready**: The process is waiting to be assigned to a processor
 - **terminated**: The process has finished execution



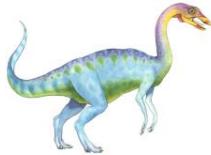
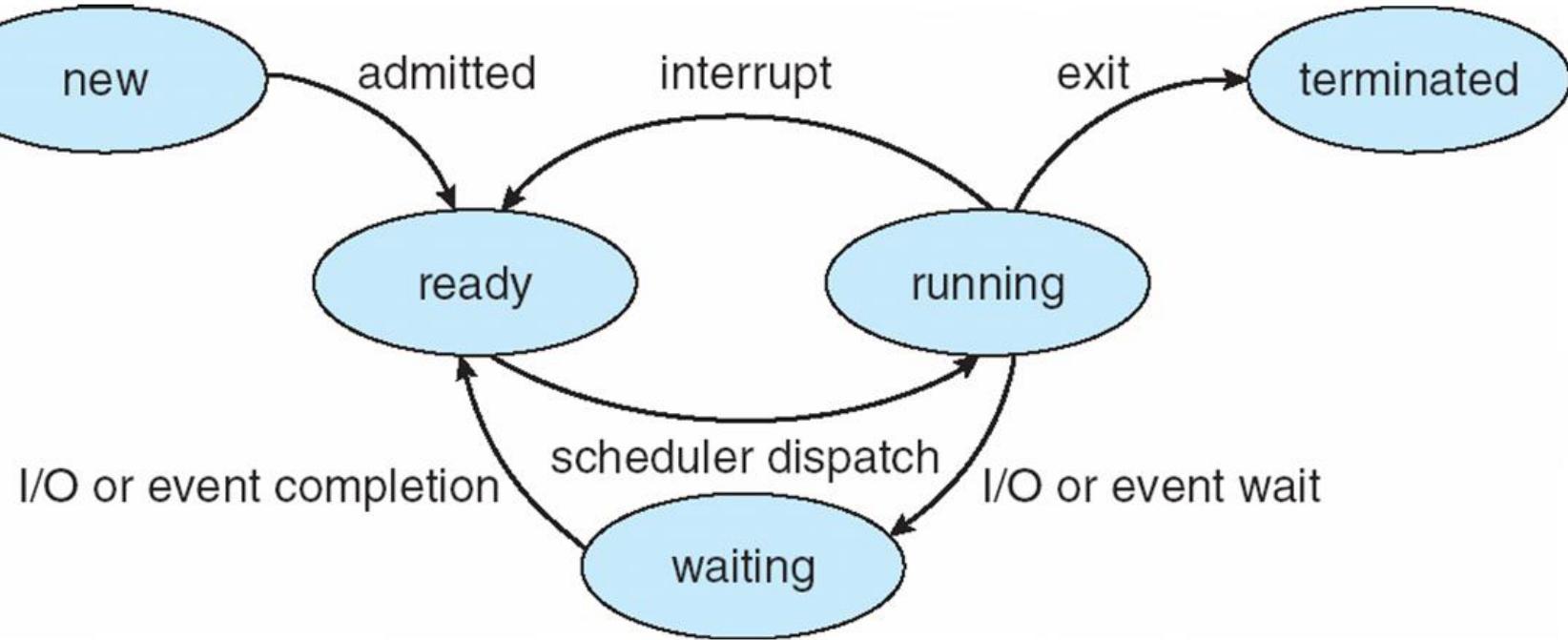


Diagram of Process State





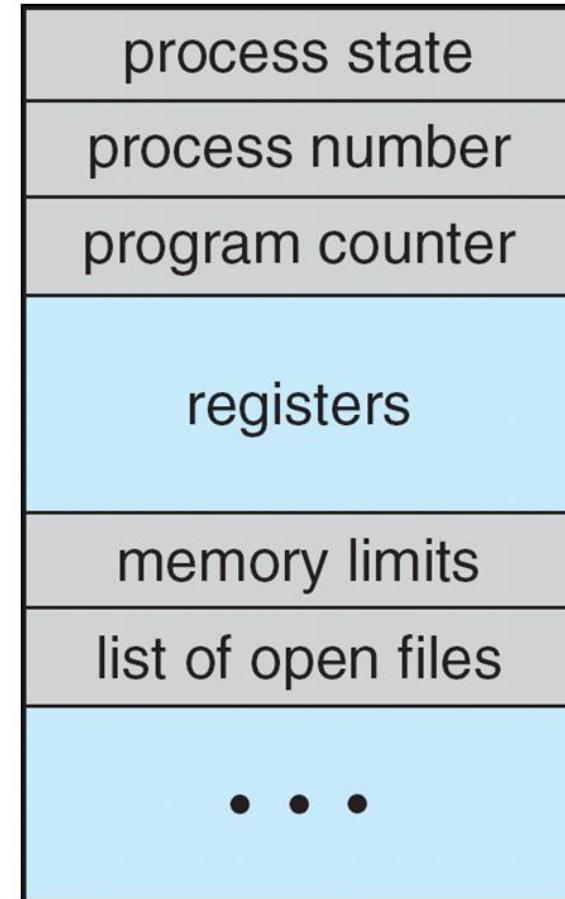
Process Control Block (PCB)

Each process can be represented in the OS by a PCB.

Information associated with each process

(also called **task control block**)

- Process state – running, waiting, etc
- Precess number — ID of the process
- Program counter – location of instruction to next execute, address of the next instruction in the code
- CPU registers – contents of all process-centric registers
- CPU scheduling information- priorities, scheduling queue pointers
- Memory-management information – memory allocated to the process
- Accounting information – CPU used, clock time elapsed since start, time limits
- I/O status information – I/O devices allocated to process, list of open files



Quiz

1. A single program can have multiple processes - T or F?
2. The unit of execution within a process is known as
3. In which state the process is waiting to be assigned to a processor?
4. The system does not do any useful task during
5. What is cascading termination?
6. Two methods for interprocess communication?