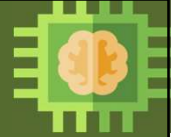


Elective Course

Course Code: CS4103

Autumn 2025-26



## Lecture #18

# Artificial Intelligence for Data Science

## Week-5: CONSTRAINT SATISFACTION PROBLEM (CSP) [Part-IV]

(Solving CSPs)

Course Instructor:

Dr. Monidipa Das

Assistant Professor

Department of Computational and Data Sciences

Indian Institute of Science Education and Research Kolkata, India 741246

# Backtracking Search for Solving CSP



```

function BACKTRACKING-SEARCH(csp) returns a solution, or failure
  return BACKTRACK({ }, csp)

function BACKTRACK(assignment, csp) returns a solution, or failure
  if assignment is complete then return assignment
  var ← SELECT-UNASSIGNED-VARIABLE(csp)
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do
    if value is consistent with assignment then
      add {var = value} to assignment
      inferences ← INFERENCE(csp, var, value)
      if inferences ≠ failure then
        add inferences to assignment
        result ← BACKTRACK(assignment, csp)
        if result ≠ failure then
          return result
      remove {var = value} and inferences from assignment
  return failure
  
```

# Comparison of CSP Algorithms



Comparison on the consistency check counted over five runs

Problem	Backtracking	BT+MRV	Forward Checking	FC+MRV
USA	(> 1,000K)	(> 1,000K)	2K	60
<i>n</i> -Queens	(> 40,000K)	13,500K	(> 40,000K)	817K
Zebra	3,859K	1K	35K	0.5K

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## Street Puzzle



1 2 3 4 5

$N_i = \{\text{English, Spaniard, Japanese, Italian, Norwegian}\}$

$C_i = \{\text{Red, Green, White, Yellow, Blue}\}$

$D_i = \{\text{Tea, Coffee, Milk, Fruit-juice, Water}\}$

$J_i = \{\text{Painter, Sculptor, Diplomat, Violinist, Doctor}\}$

$A_i = \{\text{Dog, Snails, Fox, Horse, Zebra}\}$

The Englishman lives in the Red house

The Spaniard has a Dog

The Japanese is a Painter

The Italian drinks Tea

The Norwegian lives in the first house on the left

The owner of the Green house drinks Coffee

The Green house is on the right of the White house

The Sculptor breeds Snails

The Diplomat lives in the Yellow house

The owner of the middle house drinks Milk

The Norwegian lives next door to the Blue house

The Violinist drinks Fruit juice

The Fox is in the house next to the Doctor's

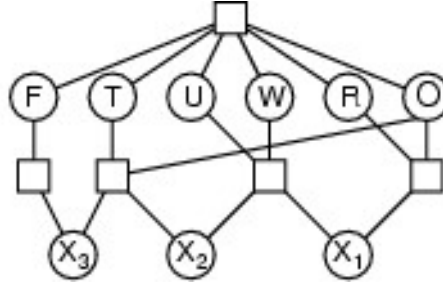
The Horse is next to the Diplomat's

Who owns the Zebra?  
Who drinks Water?

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# Cryptarithmic Puzzle



$$\begin{array}{r} \text{TWO} \\ + \text{TWO} \\ \hline \text{FOUR} \end{array}$$


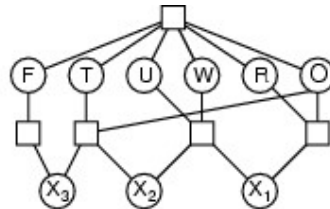
- **Variables:**  $F T U W R O X_1 X_2 X_3$
- **Domains:**  $\{0,1,2,3,4,5,6,7,8,9\}$
- **Constraints:**  $\text{Alldiff}(F, T, U, W, R, O)$

$$\begin{aligned} O + O &= R + 10 \cdot X_1 \\ X_1 + W + W &= U + 10 \cdot X_2 \\ X_2 + T + T &= O + 10 \cdot X_3 \\ X_3 &= F, T \neq 0, F \neq 0 \end{aligned}$$

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## Solution



$$\begin{array}{r} \text{TWO} \\ + \text{TWO} \\ \hline \text{FOUR} \end{array}$$


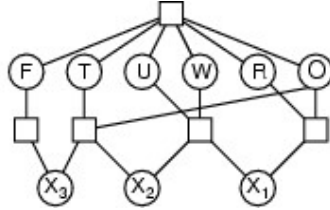
$$\begin{aligned} O + O &= R + 10 \cdot X_1 \\ X_1 + W + W &= U + 10 \cdot X_2 \\ X_2 + T + T &= O + 10 \cdot X_3 \\ X_3 &= F, T \neq 0, F \neq 0 \\ \text{Alldiff}(F, T, U, W, R, O) \end{aligned}$$

X1	X2	X3	F	T	U	W	R	O
{0,1}	{0,1}	{0,1}	{1,2,3,4,5,6,7,8,9}	{1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}
{0,1}	{0,1}	{1}	{1}	{1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0, 2, 4, 6, 8}	{0,1,2,3,4,5,6,7,8,9}
{0,1}	{0,1}	1	1	{2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0, 2, 4, 6, 8}	{0,2,3,4,5,6,7,8,9}
{0,1}	0	1	1	{2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0, 2, 4, 6, 8}	{0, 2, 4, 6, 8}
0	0	1	1	{5, 6, 7}	{0, 4, 6, 8}	{0, 2, 3, 4}	{0, 4, 8}	{0, 2, 4}
0	0	1	1	{6, 7}	{4, 6, 8}	{2, 3, 4}	{4, 8}	{2, 4}
0	0	1	1	7	{6}	{2, 3}	{8}	{4}
0	0	1	1	7	{6}	{3}	{8}	{4}
0	0	1	1	7	6	3	8	4

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# Solution



$$\begin{array}{r} \text{TWO} \\ + \text{TWO} \\ \hline \text{FOUR} \end{array}$$


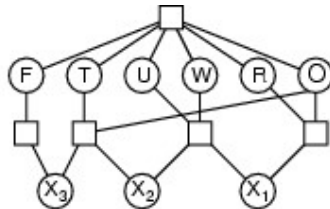
$$\begin{aligned} O + O &= R + 10 \cdot X_1 \\ X_1 + W + W &= U + 10 \cdot X_2 \\ X_2 + T + T &= O + 10 \cdot X_3 \\ X_3 &= F, T \neq 0, F \neq 0 \\ \text{Alldiff } (F, T, U, W, R, O) \end{aligned}$$

X1	X2	X3	F	T	U	W	R	O
{0,1}	{0,1}	{0,1}	{1,2,3,4,5,6,7,8,9}	{1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}
{0,1}	{0,1}	{1}	{1}	{1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0, 2, 4, 6, 8}	{0,1,2,3,4,5,6,7,8,9}
{0,1}	{0,1}	1	1	{2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0, 2, 4, 6, 8}	{0,2,3,4,5,6,7,8,9}
{0,1}	0	1	1	{2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0, 2, 4, 6, 8}	{0, 2, 4, 6, 8}
0	0	1	1	{5, 6, 7}	{0, 4, 6, 8}	{0, 2, 3, 4}	{0, 4, 8}	{0, 2, 4}
0	0	1	1	{6, 7}	{4, 6, 8}	{2, 3, 4}	{4, 8}	{2, 4}
0	0	1	1	6	{8}	{3}	{4}	{2}
					Failure!			

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# Another Solution



$$\begin{array}{r} \text{TWO} \\ + \text{TWO} \\ \hline \text{FOUR} \end{array}$$


$$\begin{aligned} O + O &= R + 10 \cdot X_1 \\ X_1 + W + W &= U + 10 \cdot X_2 \\ X_2 + T + T &= O + 10 \cdot X_3 \\ X_3 &= F, T \neq 0, F \neq 0 \\ \text{Alldiff } (F, T, U, W, R, O) \end{aligned}$$

X1	X2	X3	F	T	U	W	R	O
{0,1}	{0,1}	{0,1}	{1,2,3,4,5,6,7,8,9}	{1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}
{0,1}	{0,1}	{1}	{1}	{1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0,1,2,3,4,5,6,7,8,9}	{0, 2, 4, 6, 8}	{0,1,2,3,4,5,6,7,8,9}
{0,1}	{0,1}	1	1	{2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0, 2, 4, 6, 8}	{0,2,3,4,5,6,7,8,9}
{0,1}	1	1	1	{6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0,2,3,4,5,6,7,8,9}	{0, 4, 6, 8}	{3, 5, 7, 9}
0	1	1	1	{6}	{0, 4, 8}	{5, 7, 9}	{6}	{3}
1	1	1	1	{7,8}	{3,5,7}	{6,7,8}	{0, 4}	{5, 7}
1	1	1	1	{7,8}	{3,7}	{6,8}	0	{5}
1	1	1	1	{7}	{3}	{6}	0	5
1	1	1	1	7	3	6	0	5

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# Local Search Problem: Definition



- **Definition:** A **local search problem** consists of a:
  - **CSP:** a set of variables, domains for these variables, and constraints on their joint values. A node in the search space will be a **complete assignment** to all of the variables.
  - **Neighbor relation:** an edge in the search space will exist when the neighbor relation holds between a pair of nodes.
  - **Scoring function:**  $h(n)$ , judges cost of a node (want to minimize)
    - E.g. the number of constraints violated in node  $n$ .
    - E.g. the cost of a state in an optimization context.

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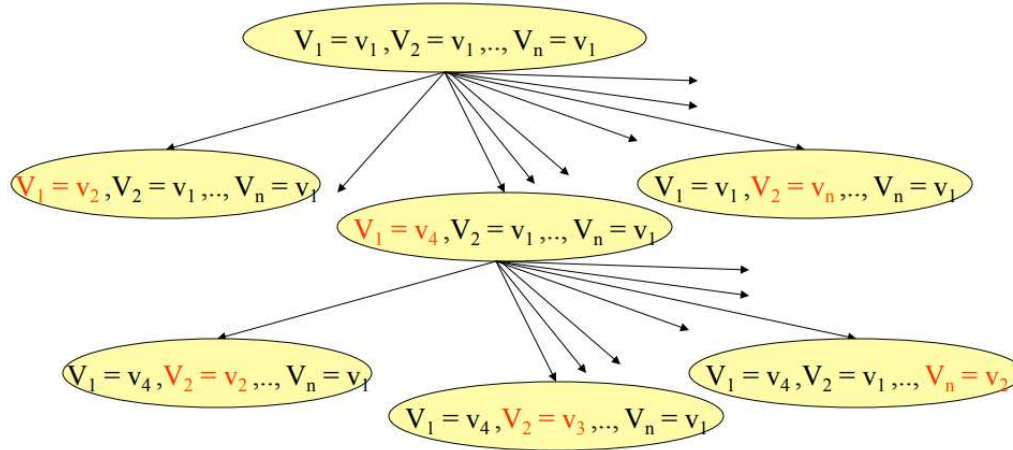
# Local Search for CSPs



- Hill-climbing, simulated annealing typically work with "complete" states, i.e., all variables assigned
- To apply to CSPs:
  - **allow states with unsatisfied constraints**
  - operators **reassign** variable values
- Variable selection: randomly select any conflicted variable
- Value selection by **min-conflicts** heuristic:
  - choose value that violates the fewest constraints
  - i.e., hill-climb with  $h(n)$  = total number of violated constraints

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# Search Space for Local Search



**Local search does NOT backtrack!**

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## Example: Sudoku as a local search problem



- **CSP:** usual Sudoku CSP
  - One variable per cell; domains  $\{1, \dots, 9\}$ ;
  - Constraints: each number occurs once per row, per column, and per 3x3 box
- **Neighbor relation:** value of a single cell differs
- **Scoring function:** number of constraint violations

1	8	1	4	8	3	4	3	5	2	8	1	4	8	3	4	3	5
7	9	3	6	2	8	1	4	7	7	9	3	6	2	8	1	4	7
4	6	5	7	1	2	8	5	6	4	6	5	7	1	2	8	5	6
3	3	7	3	1	4	1	9	3	3	3	7	3	1	4	1	9	3
8	5	7	8	2	2	9	7	8	8	5	7	8	2	2	9	7	8
5	4	4	3	7	8	7	6	2	5	4	4	3	7	8	7	6	2
4	8	7	1	2	8	5	3	6	4	8	7	1	2	8	5	3	6
1	1	7	5	9	3	4	2	8	1	1	7	5	9	3	4	2	8
7	5	8	4	8	6	7	3	5	7	5	8	4	8	6	7	3	5

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## Local Search for CSP: MIN-CONFLICTS Algorithm



**function** MIN-CONFLICTS(*csp*, *max\_steps*) **returns** a solution or failure

**inputs:** *csp*, a constraint satisfaction problem

*max\_steps*, the number of steps allowed before giving up

*current*  $\leftarrow$  an initial complete assignment for *csp*

**for** *i* = 1 to *max\_steps* **do**

**if** *current* is a solution for *csp* **then return** *current*

*var*  $\leftarrow$  a randomly chosen conflicted variable from *csp*.VARIABLES

*value*  $\leftarrow$  the value *v* for *var* that minimizes CONFLICTS(*var*, *v*, *current*, *csp*)

    set *var* = *value* in *current*

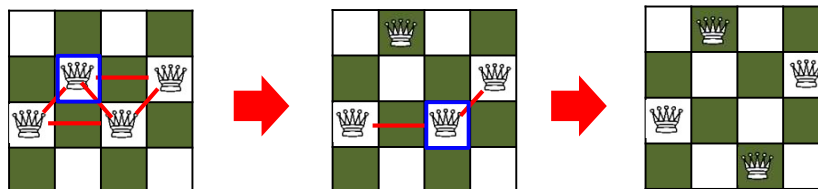
**return** failure

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## Example: 4-Queens



- **States:** 4 queens in 4 columns ( $4^4 = 256$  states)
- **Actions:** move queen in column
- **Goal test:** no attacks
- **Evaluation:**  $h(n)$  = number of attacks



- Given random initial state, can solve  $n$ -queens in almost constant time for arbitrary  $n$  with high probability (e.g.,  $n = 10,000,000$ )

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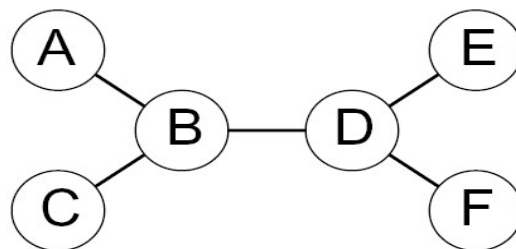
# Solving CSPs: Improving Efficiency



- Which variable should be assigned next?
- In what order should its values be tried?
- Can we **detect** inevitable failures early?
- Can we take advantage of the **problem structure**?

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## Tree-structured CSPs



**Theorem:** if the constraint graph has no loops, the CSP can be solved in  $O(nd^2)$  time

Compare to general CSPs, where worst-case time is  $O(d^n)$

This property also applies to logical and probabilistic reasoning:  
an important example of the relation between syntactic restrictions  
and the complexity of reasoning.

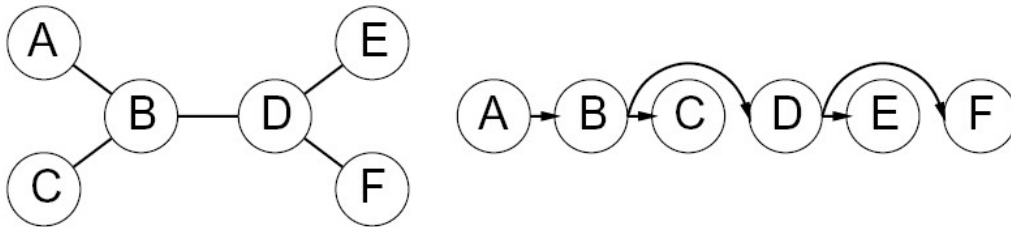
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# Algorithm for tree-structured CSPs



1. Choose a variable as root, order variables from root to leaves such that every node's parent precedes it in the ordering



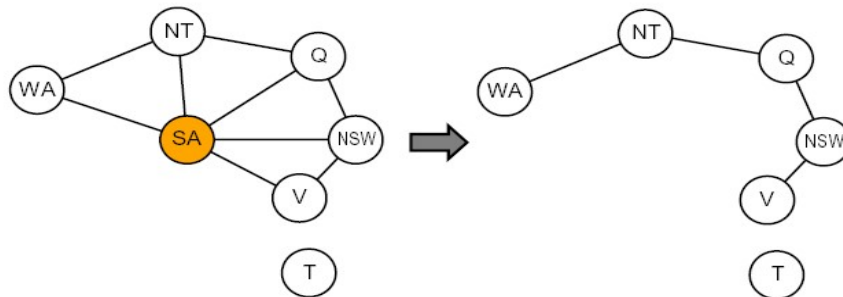
2. For  $j$  from  $n$  down to 2, apply  $\text{REMOVEINCONSISTENT}(\text{Parent}(X_j), X_j)$
3. For  $j$  from 1 to  $n$ , assign  $X_j$  consistently with  $\text{Parent}(X_j)$

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# Nearly tree-structured CSPs



**Conditioning:** instantiate a variable, prune its neighbors' domains

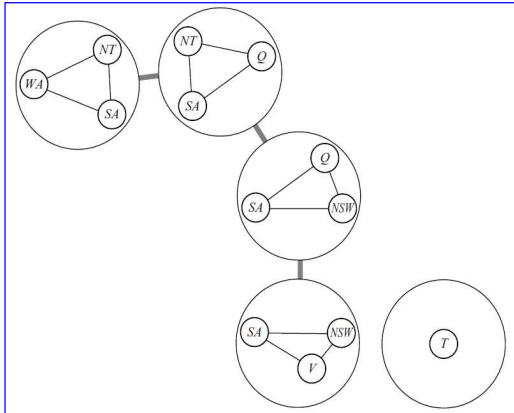


**Cutset conditioning:** instantiate (in all ways) a set of variables such that the remaining constraint graph is a tree

Cutset size  $c \Rightarrow$  runtime  $O(d^c \cdot (n - c)d^2)$ , very fast for small  $c$

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# Tree decomposition



- Every variable in original problem must appear in at least one subproblem
- If two variables are connected in the original problem, they must appear together (along with the constraint) in at least one subproblem
- If a variable occurs in two subproblems in the tree, it must appear in every subproblem on the path that connects the two

- Algorithm: solve for all solutions of each subproblem. Then, use the tree-structured algorithm, treating the subproblem solutions as variables for those subproblems.
- $O(nd^{w+1})$  where  $w$  is the *treewidth* (= one less than size of largest subproblem)

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## Questions?

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