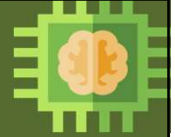


Elective Course

Course Code: CS4103

Autumn 2025-26



## Lecture #17

# Artificial Intelligence for Data Science

**Week-5: CONSTRAINT SATISFACTION PROBLEM (CSP) [Part-III]**

(Solving CSPs)

Course Instructor:**Dr. Monidipa Das**

Assistant Professor

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Indian Institute of Science Education and Research Kolkata, India 741246

## Solving CSPs: Improving Efficiency

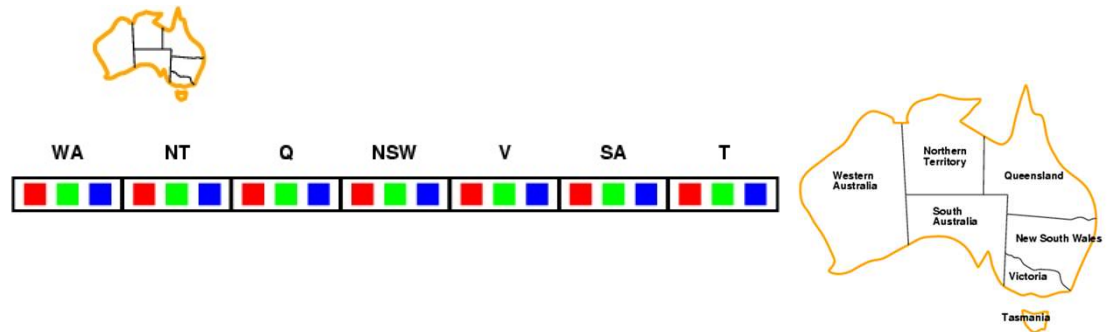


- Which variable should be assigned next?
- In what order should its values be tried?
- Can we **detect** inevitable failures early?
- Can we take advantage of the **problem structure**?

# Forward checking



- Idea:
  - Keep track of remaining legal values for unassigned variables
  - Terminate search when any variable has no legal values

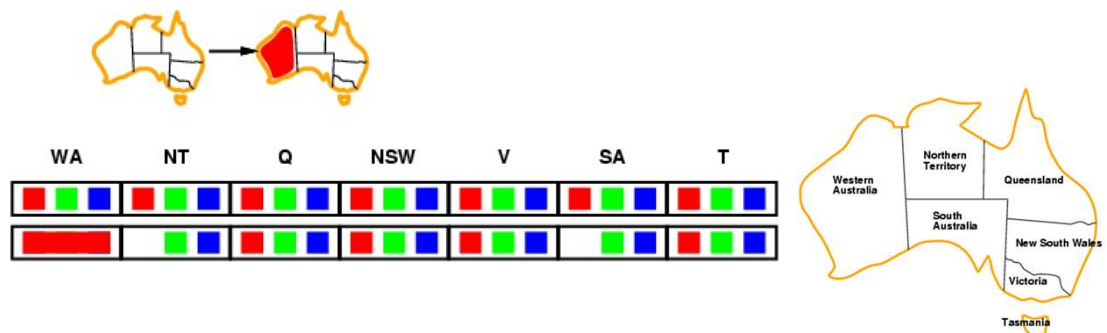


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# Forward checking



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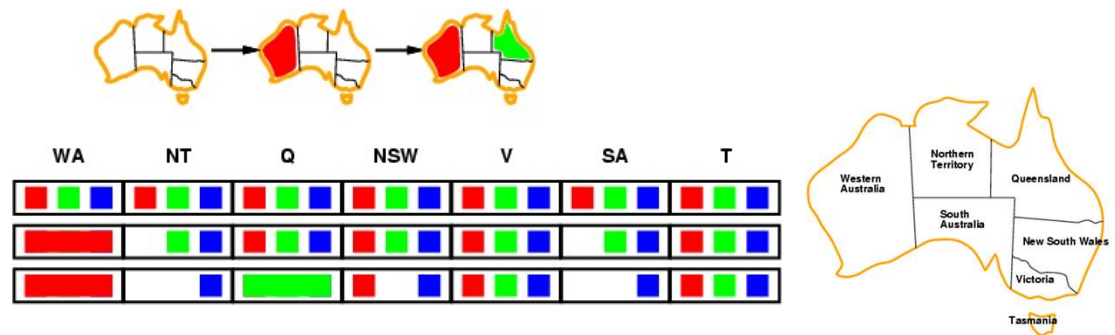


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# Forward checking



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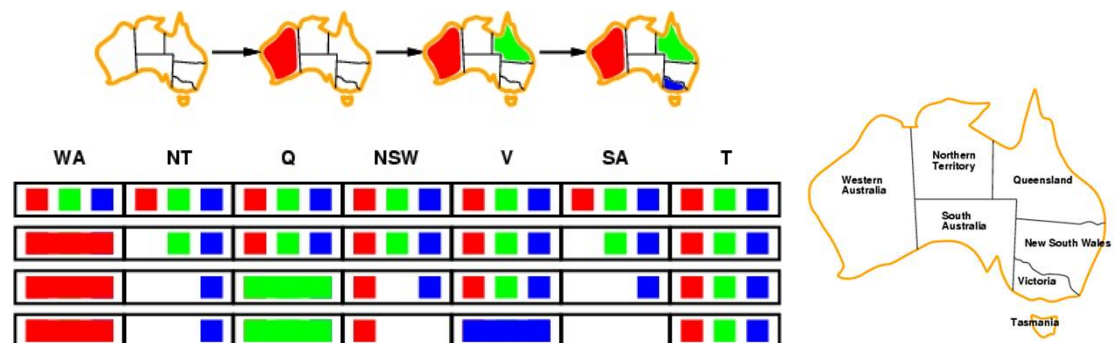


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# Forward checking

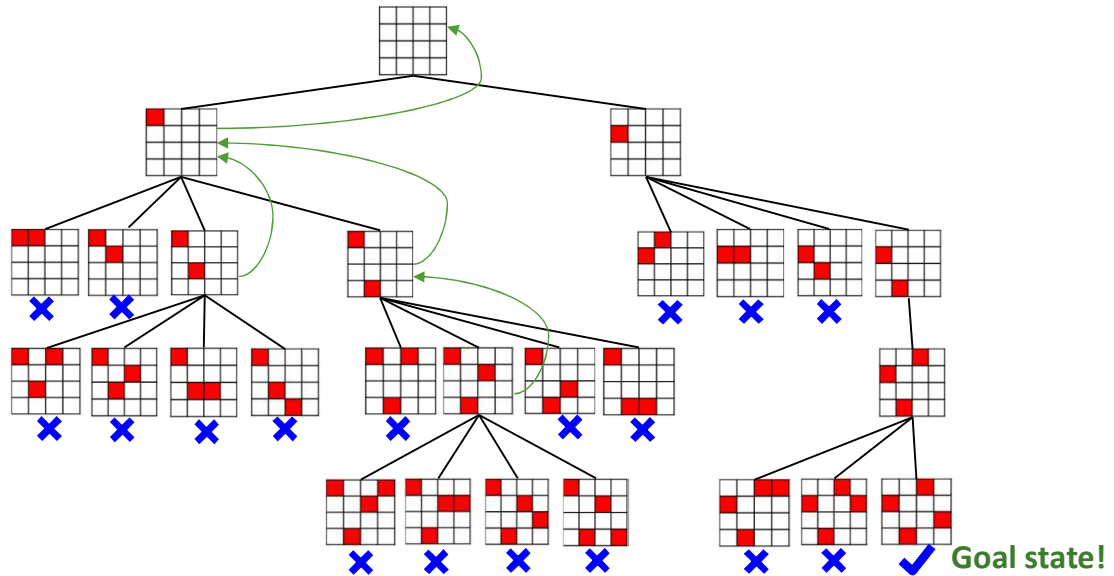


- Idea:
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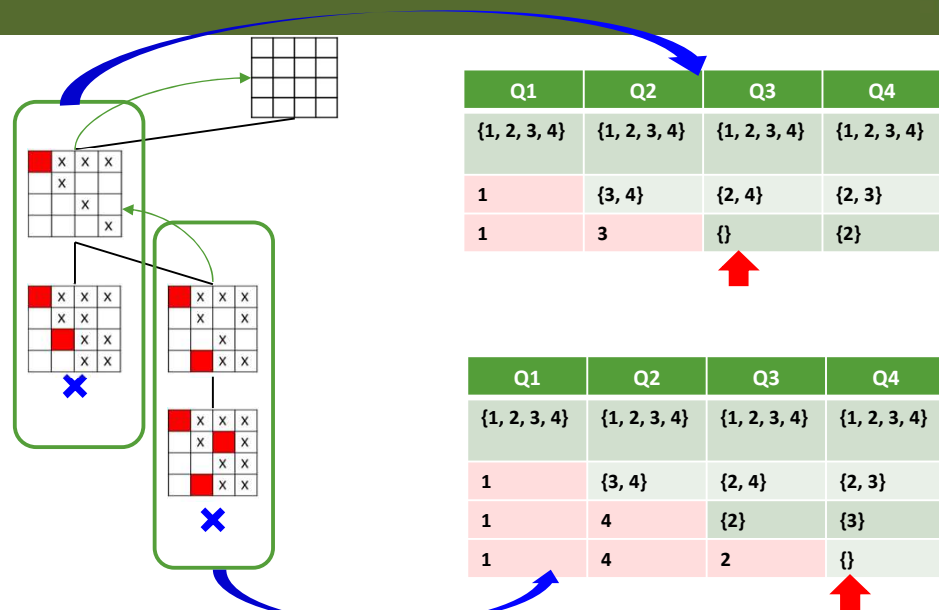
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## Backtracking Search Example: 4-Queens Problem



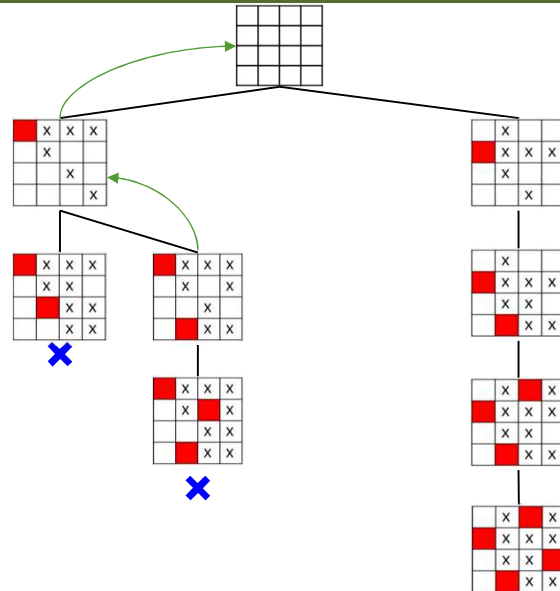
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## Forward checking: 4-Queens Problem



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# Forward checking: 4-Queens Problem



Q1	Q2	Q3	Q4
{1, 2, 3, 4}	{1, 2, 3, 4}	{1, 2, 3, 4}	{1, 2, 3, 4}
2	{4}	{1, 3}	{1, 3, 4}
2	4	{1}	{1, 3}
2	4	1	{3}
2	4	1	3

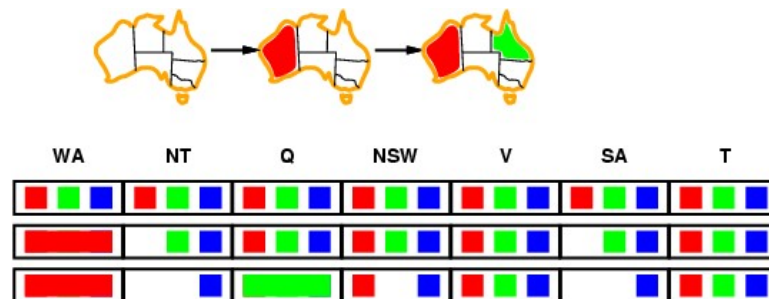
✓ Goal state!

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# Constraint propagation



- Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:



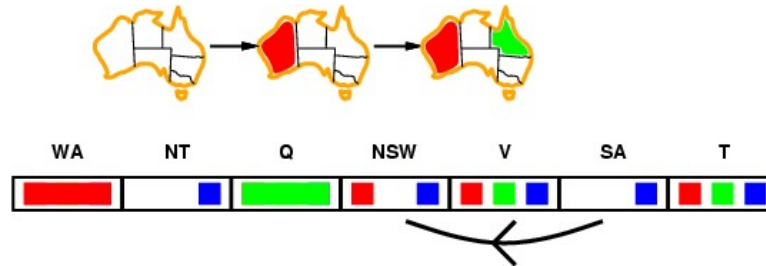
- NT and SA cannot both be blue!
- Constraint propagation repeatedly enforces constraints locally

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# Arc consistency



- Simplest form of propagation makes each arc **consistent**
- $X \rightarrow Y$  is consistent iff  
for **every** value  $x$  of  $X$  there is **some** allowed  $y$

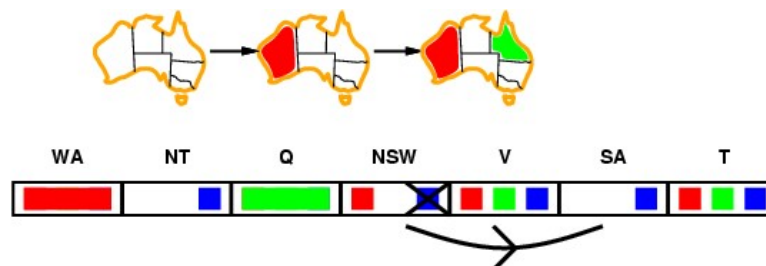


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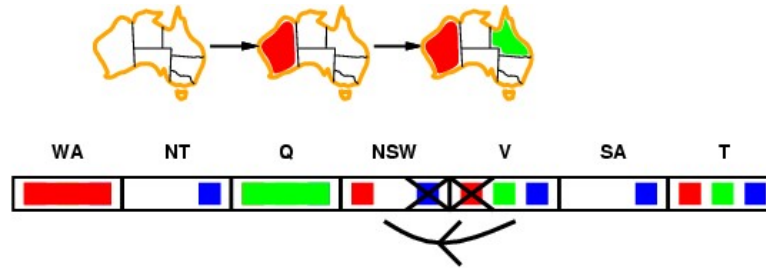


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# Arc consistency



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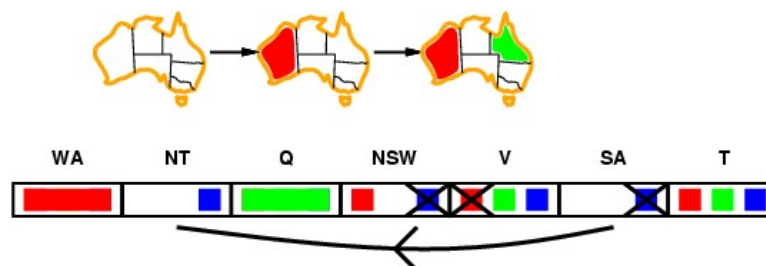
- If  $X$  loses a value, neighbors of  $X$  need to be rechecked

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# Arc consistency



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- $X \rightarrow Y$  is consistent iff  
for **every** value  $x$  of  $X$  there is **some** allowed  $y$



- If  $X$  loses a value, neighbors of  $X$  need to be rechecked
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment

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# Arc consistency Algorithm



```

function AC-3(csp) returns the CSP, possibly with reduced domains
  inputs: csp, a binary CSP with variables  $\{X_1, X_2, \dots, X_n\}$ 
  local variables: queue, a queue of arcs, initially all the arcs in csp

  while queue is not empty do
     $(X_i, X_j) \leftarrow \text{REMOVE-FIRST}(\textit{queue})$ 
    if RM-INCONSISTENT-VALUES( $X_i, X_j$ ) then
      for each  $X_k$  in NEIGHBORS[ $X_i$ ] do
        add  $(X_k, X_i)$  to queue

  function RM-INCONSISTENT-VALUES( $X_i, X_j$ ) returns true iff remove a value
    removed  $\leftarrow$  false
    for each  $x$  in DOMAIN[ $X_i$ ] do
      if no value  $y$  in DOMAIN[ $X_j$ ] allows  $(x, y)$  to satisfy constraint( $X_i, X_j$ )
        then delete  $x$  from DOMAIN[ $X_i$ ]; removed  $\leftarrow$  true
    return removed
  
```

Time complexity:  $O(n^2d^3)$

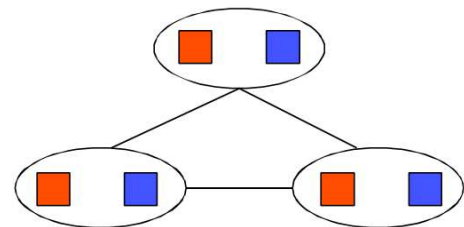
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# Limitations of Arc Consistency



- After running arc consistency:
  - Can have one solution left
  - Can have multiple solutions left
  - **Can have no solutions left (and not know it)**

What went wrong here?



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# $k$ -consistency



- A CSP is  $k$ -consistent if, for any set of  $k - 1$  variables, and for any consistent assignment to those variables, a consistent value can always be assigned to any  $k$ -th variable
  - 1-consistency is **node consistency**
  - 2-consistency is **arc consistency**
  - For binary constraint networks, 3-consistency is the same as **path consistency**
- Getting  $k$ -consistency requires time and space exponential in  $k$

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## Questions?

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