

Lipschitz constant on a relation means that relation is a function ↗ if F is contraction relation
 L Lipschitz constant on relation $F \subseteq V \times FV$ ↗ $\|u-v\|_2 \leq L \|x-y\|_2$ ↗ if F is nonexpansive

$$\Rightarrow (x=y \Rightarrow \|u-v\|_2 \leq L \cdot 0 = 0 \Rightarrow u=v)$$

↪ same input argument results in a single value, so F is a function

* Basic properties of Lipschitz function : As a result, nonexpansive and contraction mapping they are function too

(F Lipschitz constant $L \geq 1$, \tilde{F} Lipschitz constant $\tilde{L} \geq 1$) $\Rightarrow L \tilde{L}$ Lipschitz constant of $F \circ \tilde{F}$
 \Leftrightarrow if F and \tilde{F} are nonexpansive $\Leftrightarrow L \tilde{L} \leq 1$ (if $L \tilde{L} > 1$)

• A ↗ if $F \circ \tilde{F}$ has Lipschitz constant $|L \tilde{L}| \leq 1$
 \Leftrightarrow if F and \tilde{F} are nonexpansive $\Leftrightarrow L \tilde{L} \leq 1$

↪ also, $0F + (1-\theta)F$ has Lipschitz constant $|\theta|L + |1-\theta|\tilde{L} \leq \theta + (1-\theta) \leq 1$, so $0F + (1-\theta)F$ is nonexpansive

↪ nonexpansive nonexpansive \Leftrightarrow nonexpansive

↪ if $\theta \in (0,1)$ ↗ one of F , or \tilde{F} contraction ↗ $0 \leq \theta \leq 1$ [as we will get 1 only both $L=1, \tilde{L}=1$]

↪ \downarrow ↗

↪ $0F + (1-\theta)F$ is a contraction mapping then? ↗ strict convex combination of one expansive mapping and one contraction mapping is contraction

* Fixed point set of nonexpansive operator F ↗ $\{x \mid Fx = x\}$, $\{x \mid Fx \leq x\} = (I-F)^{-1}(0)$ ↗ So fixed point set of a nonexpansive operator is essentially 0 passed through the resolvent of that operator

↪ Note the scaling factor of the operator is -1

↪ Fixed point of contraction mapping F is singleton ↗ $F^{-1}(0) = \{x\}$

* Example: 1) Affine function $F(x) = Ax + b$ has Lipschitz constant $L = \|A\|_2 = \max_{\lambda}(\lambda)$ # For any affine function the maximum singular value of the associated matrix is the Lipschitz constant

$$\begin{aligned} \|F(x)-F(y)\|_2 &\leq \|Ax+bx-Ay-by\|_2 = \|A(x-y)\|_2 \\ &\Leftrightarrow \max_{\lambda} \frac{\|A(x-y)\|_2}{\|x-y\|_2} \leq L \\ &\Leftrightarrow \max_{\lambda} \frac{\|Ax\|_2}{\|x\|_2} \leq L \\ &\Leftrightarrow \max_{\lambda} \frac{\|Ax\|_2}{\|x\|_2} = \sqrt{\max(\lambda)} \end{aligned}$$

So, L can be taken as $\sqrt{\max(\lambda)}$, if this is < 1 then we can have contraction (nonexpansive) mapping.

2) Differentiable function:

A differentiable function is Lipschitz iff the Jacobian norm of the function is globally bounded by the associated Lipschitz constant.

$$F: \mathbb{R}^n \rightarrow \mathbb{R}^m \text{ Lipschitz with parameter } L \Leftrightarrow \forall x \quad \|DF(x)\|_2 \leq L$$

Proof: (\Leftarrow)

$\forall x, \|DF(x)\|_2 \leq L$

define $g(t) = (F(x)-F(y))^T F(x+(1-t)y)$ ↗ mean value theorem: $\exists c \in (x,y) \quad g'(c) = \frac{f(x)-f(y)}{x-y}$

this is a continuous function in t ↗ so by mean value theorem $\exists c \in (x,y) \quad g'(c) = \frac{g(x)-g(y)}{x-y} = 0$

minim $g'(t) = \frac{d}{dt} (F(x)-F(y))^T F(x+(1-t)y)$

$$= (F(x)-F(y))^T \frac{d}{dt} (F(x+(1-t)y)) \frac{d}{dt} ((1-t)y) \quad \text{Chain rule} \quad \frac{d}{dt} g(x) = D_g(x) \cdot D_x g(x)$$

$$= (F(x)-F(y))^T D(F(x+(1-t)y)) (x-y)$$

$$\therefore g'(t) = (F(x)-F(y))^T D(F(x+(1-t)y)) (x-y)$$

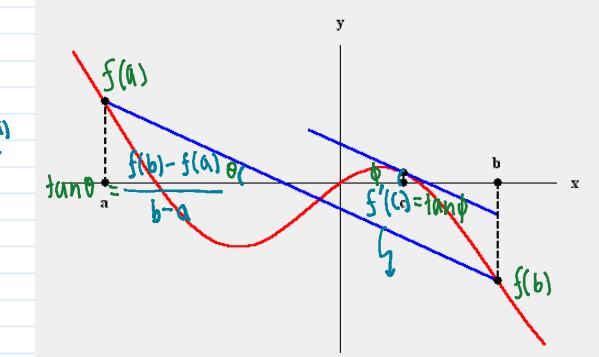
$\|g'(t)\|_2 = (F(x)-F(y))^T F(x+(1-t)y)$

$$= \|F(x)-F(y)\|^2$$

$\therefore g'(t) = (F(x)-F(y))^T D(F(x+(1-t)y)) (x-y) \leq \|x-y\|_2 \|D(F(x+(1-t)y))\|_2 \leq \|x-y\|_2 \|F(x)-F(y)\|_2$

By Cauchy-Schwarz inequality: $|ab| \leq \|a\|_2 \|b\|_2$

$\leq L \|F(x)-F(y)\|_2 \|x-y\|_2$



$$\|F(x)-F(y)\|_2 \leq L \|F(x)-F(y)\|_2 \|x-y\|_2$$

$$\Rightarrow \|F(x)-F(y)\|_2 \leq L \|x-y\|_2 \quad (\text{Q.E.D.})$$

(\Rightarrow)

F has a Lipschitz parameter L ,

Want to prove $\|DF(x)\|_2 \leq L$

$$\Leftrightarrow \max_{v \neq 0} \frac{\|DF(x)v\|_2}{\|v\|_2} \leq L$$

$$\Leftrightarrow \forall v \neq 0 \quad \frac{\|DF(x)v\|_2}{\|v\|_2} \leq L$$

$$\Leftrightarrow \forall v \quad \frac{\|DF(x)v\|_2}{\|v\|_2} \leq L \quad [\because v=0 \text{ is the equality } 0=0 \text{ case}]$$

Now: $\forall h \quad \|DF(x)v\|_2 = \lim_{h \rightarrow 0} \frac{1}{h} \|F(x+hv) - F(x)\|_2$

$\leq L \|x+hv - x\|_2 = L \|hv\|_2 = L \|h\|_2 \|v\|_2 \leq L \|v\|_2$ ↗ If $v \neq 0$ then $\|v\|_2 \neq 0$

* Projections: We will prove, $\Pi_C(z)$ is a non-expansive operator $\Leftrightarrow \frac{\|\Pi_C(z) - \Pi_C(y)\|}{\|z-y\|} \leq 1$

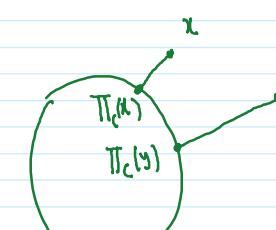
% If there are two lovers in real world and one projection set symbolizing the fantasy world, then the distance between the lovers will be smaller in the fantasy world than that of the real world.

$$\Pi_C(z) = \arg\min_{x \in C} \frac{1}{2} \|z-x\|^2 \quad \text{Always exists and unique}$$

↓

point in C closest to x

$$\text{Optimality criterion for } \Pi_C(z) \quad \left(\begin{array}{l} \nabla \Pi_C(z)^T (y-z) \geq 0 \\ \forall y \in C \end{array} \right) \quad \text{is nonexpansive nature}$$



Optimality criterion for $\forall y \in \mathcal{Y}$ is y^* is y^*
This is the first-order optimality condition.
In our case the problem is:

$$\begin{aligned} & \min_{z \in \mathcal{Z}} f_c(z) = \|z - x\|_2^2 \\ & \text{subject to } z \in \mathcal{C} \end{aligned}$$

$$\therefore \nabla f_c(z) = (z - x)$$

$$\nabla f_c(z^*) = (\Pi_c(x) - x)$$

$$\text{So, the optimality condition is } \forall z \in \mathcal{Z} \quad \nabla f_c(z^*)^\top (z - z^*) = (\Pi_c(x) - x)^\top (z - \Pi_c(x)) \geq 0$$

$$(\Pi_c(x) - x)^\top (z - \Pi_c(x)) \geq 0$$

similarly when we are taking projection of y :

$$\forall z \in \mathcal{Z} \quad (\Pi_c(y) - y)^\top (z - \Pi_c(y)) \geq 0$$

$$z = \Pi_c(y)$$

$$(\Pi_c(x) - x)^\top (\Pi_c(y) - \Pi_c(x)) = \frac{1}{2} \|x - \Pi_c(x)\|_2^2 + \frac{1}{2} \|y - \Pi_c(y)\|_2^2 \geq 0$$

$$(\Pi_c(y) - y)^\top (\Pi_c(x) - \Pi_c(y)) \geq 0$$

$$(\Pi_c(y) - y)^\top (\Pi_c(x) - \Pi_c(y)) + (\Pi_c(x) - \Pi_c(x))^\top (\Pi_c(x) - \Pi_c(y)) \geq 0$$

$$= ((\Pi_c(y) - y) + x - \Pi_c(x))^\top (\Pi_c(x) - \Pi_c(y))$$

$$= ((x - y)^\top - (\Pi_c(x) - \Pi_c(y)))^\top (\Pi_c(x) - \Pi_c(y))$$

$$= (x - y)^\top (\Pi_c(x) - \Pi_c(y)) - \|\Pi_c(x) - \Pi_c(y)\|_2^2 \geq 0$$

$$\rightarrow \|\Pi_c(x) - \Pi_c(y)\|_2^2 \leq (\Pi_c(x) - \Pi_c(y))^\top (x - y) \leq \|\Pi_c(x) - \Pi_c(y)\|_2 \|x - y\|_2$$

By Cauchy-Schwarz
 $a^\top b \leq \|a\| \|b\|$

$$\therefore \|\Pi_c(x) - \Pi_c(y)\|_2 \leq \|x - y\|_2$$

So $\Pi_c(\cdot)$ is a nonexpansive operator

Similarly, overprojection operator $B_C = 2\Pi_C - I$ on $C \subseteq \mathbb{R}^n$ is nonexpansive.

Caution: nonexpansive and contraction mapping are function, but monotone or strongly monotone operator can be nontrivial relations.

A monotone operator always split out vector of same dimension as the input argument.

* Monotone operators: def: monotone operators definitions and related

$$\text{relation: } R \text{ is monotone} \Leftrightarrow \forall (x, y) \in R \quad (f(x) - f(y))^\top (x - y) \geq 0 \quad \text{if } f \text{ is } F \text{ then } R \text{ is } F \text{ relation}$$

(in the increasing sense) $\Leftrightarrow \forall (x, y) \in R \quad f(x) - f(y) \geq 0 \quad \text{if } f \text{ is } F \text{ then } R \text{ is } F \text{ relation}$

def: monotone operator

def: monotone relation

def: monotone mapping

def: monotone function

def: maximal monotone

def: strongly monotone

def: Lipschitz continuous

def: maximal Lipschitz continuous

$\forall x \in \text{int}(C) \quad N_C(x) = \{0\}$

$\forall x \in C \quad N_C(x) \neq \text{nontrivial}$ # The intuition behind this is that remember at the boundary of any function, the subdifferential might not exist. One strange function in this regard is the function of \mathbb{R}^n set.

Why $N_C(x)$ matters? Because this is the subdifferential mapping of the convex indicator function $\gamma_C(x) = \begin{cases} 0, & x \in C \\ \infty, & x \notin C \end{cases}$, i.e., $N_C = \gamma_C$.

Makes sense because remember, for any convex function the subdifferential always exists in the int(dom f). At the boundary the subdifferential may not exist, i.e., $\gamma_C(x) \in N_C(x)$ might not exist, as a result $N_C(x|_{\partial C}) = \gamma_C(x|_{\partial C}) = \text{nontrivial}$

[eq: subdifferential is monotone mapping]

proof: $\gamma_C(x) = \begin{cases} 0, & x \in C \\ \infty, & x \notin C \end{cases}$, want to determine, $N_C(x) = N_\gamma(x)$

By definition of subgradient, $\exists_{y \in \text{dom}(\gamma_C)} \forall_{y \in \text{dom}(\gamma_C)} \gamma_C(y) \geq \gamma_C(x) + g_x^T(y-x) \Leftrightarrow \min_{y \in \text{dom}(\gamma_C)} \gamma_C(y) = \gamma_C(x) + g_x^T(y-x) \Leftrightarrow \forall_{y \in \text{dom}(\gamma_C)} \gamma_C(y) \geq \gamma_C(x) + g_x^T(y-x)$

* Case 1: $x \in C$, then $\forall y \in C \quad 0 \geq \gamma_C(y) + g_x^T(y-x) \Leftrightarrow \forall y \in C \quad g_x^T(y-x) \leq 0$

* Case 2: $x \notin C \Rightarrow \exists y \in C \quad 0 > \gamma_C(y) + g_x^T(y-x) \Rightarrow \text{no finite } g_x \text{ can exist} \Rightarrow N_C(x) = \emptyset$

Not the most rigorous proof ...

* saddle subdifferential:

$f: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$ s.t. $\{(\cdot, \cdot)\}$ is convex in \mathbb{R}^n , concave in \mathbb{R}^m

$F(x, y) = \{\text{saddle subdifferential relation}\}$

$$= \left[\begin{array}{c} \partial_x f(x, y) \\ \partial_y f(x, y) \end{array} \right] \quad \forall_{x, y} \quad \partial_x f(x, y) \neq \emptyset, \partial_y f(x, y) \neq \emptyset$$

* Why this is called saddle subdifferential?

Because $(x, y) \in F \Leftrightarrow (x, y) \in \text{dom} f(x, y) \Leftrightarrow \begin{cases} \exists_{x \in \text{dom} f(x, y)} \exists_{y \in \text{dom} f(x, y)} f(x, y) \leq f(x, \tilde{y}) \\ \exists_{y \in \text{dom} f(x, y)} \exists_{x \in \text{dom} f(x, y)} f(x, y) \leq f(\tilde{x}, y) \end{cases}$

(convex in x) \quad (concave in y)

(convex in y) \quad (concave in x)

$\therefore \text{Set of } F \text{ is the saddle points of } f.$

Saddle subdifferential relation for CCP function $f(x, y)$ is maximal

* KKT operator: # Manuscript version

$$\left(\begin{array}{c} \nabla f(x) \\ \nabla_h f(x, \lambda, \nu) \\ \nabla_v f(x, \lambda, \nu) \end{array} \right) \quad \begin{array}{l} \text{affine} \\ \text{if } \lambda \neq 0 \\ -\infty \end{array}$$

$$L(x, \lambda, \nu) = \begin{cases} f(x) + \sum_{i=1}^m \lambda_i h_i(x) + \nu_i v_i(x), & \text{if } \lambda \neq 0 \\ -\infty, & \text{otherwise} \end{cases}$$

$T(x, \lambda, \nu) = \{\text{KKT operator}\} = \left[\begin{array}{c} \partial_x L(x, \lambda, \nu) \\ -f(x) + \nabla_h f(x, \lambda, \nu) \\ \lambda \end{array} \right]$

so $\partial_x L(x, \lambda, \nu) = 0$ || vanishing gradient of Lagrangian
 $\nabla_h f(x, \lambda, \nu) = 0$ || primal inequality feasibility
 $\lambda = 0$ || dual variable feasibility
 $\nabla_v f(x, \lambda, \nu) = 0$ || ensures that dual feasibility condition holds

$T(x, \lambda, \nu)$ is special case of saddle subdifferential $\hat{f} = \text{monotone}$

$\exists (x^*, \lambda^*, \nu^*) \in T(x^*, \lambda^*, \nu^*) \Leftrightarrow (x^*, \lambda^*, \nu^*)$ solves the optimality problem

* KKT operator (slide version)

$\nabla f(x)$

$Ax = b$

$L(x, y) = f(x) + y^T(Ax - b)$

KKT operator: $F(x, y) = \left[\begin{array}{c} \partial_x L(x, y) \\ -\partial_y L(x, y) \end{array} \right]$ # gives vanishing gradient of Lagrangian (one of the KKT conditions)
gives LHS of primal feasibility (another KKT condition)

$$F(0, 0) = \left[\begin{array}{c} \partial_x L(0, 0) \\ -\partial_y L(0, 0) \end{array} \right]$$

so if $F(x^*, y^*) = 0 \Leftrightarrow (x^*, y^*)$: optimal primal-dual pair

$$\Leftrightarrow (x^*, y^*) \in F$$

$$\Leftrightarrow (0, (A^T y^*)) \in F^{-1}$$

$$\Leftrightarrow (x^*, y^*) \in F^{-1}(0)$$

So optimal primal-dual pair will belong to the output set of the inverse KKT operator with 0 fed into it!

Multiplier to residual mapping is very important as it has a connection with ADMM

* Multiplier to residual mapping. def: multiplier to residual mapping

$$\left(\begin{array}{c} \nabla f(x) \\ A \\ b \end{array} \right)$$

$L(x, y) = f(x) + y^T(Ax - b)$ more technically: $f(x) = b - A^T y \Leftrightarrow f(x) = b - A^T \text{argmin}_x L(x, y)$

def: $F(y) = b - A^T y$ if $x^* \in \text{argmin}_x L(x, y)$, because $f(x) = b - A^T y$ is strongly convex, so will have unique minimizer

This is a kind of multiplier to residual mapping because it takes the lagrange multiplier and outputs the residual that associated with the sub-optimality of x

$\nabla_x L(x, y) = 0 \Leftrightarrow \partial_x f(x) + A^T y = 0 \Leftrightarrow \nabla_x f(x) + A^T y = 0$

$\Leftrightarrow \exists_{\lambda \in \mathbb{R}^m} \nabla_x f(x) + A^T y = \lambda I \Leftrightarrow \lambda = -A^T y \Leftrightarrow x^* = -A^T y$

$\Leftrightarrow x^* = (A^T y)^{-1}(-A^T y) \Leftrightarrow x^* = (A^T y)^{-1}(-A^T y)$

$\Leftrightarrow x^* = (A^T y)^{-1}(-A^T y) \Leftrightarrow x^* = (A^T y)^{-1}(-A^T y)$ # but x^* is unique as it is the minimizer of a strongly convex function.

Alternative definition to "residual mapping operator":

$F(y) = b - A(\lambda^T y)^{-1}(-A^T y)$ Has a monotone operator, so multiplier to residual mapping is a monotone operator

proof: $\{\partial_x f(y)\}$ is monotone

$\Rightarrow \{\partial_x f(y)\}^{-1}$ is monotone

$\Rightarrow \forall_{y_1, y_2} \partial_x f^T(\partial_x f(y_1))^{-1}(\partial_x f(y_2)) \in \text{monotone}$

$\Rightarrow F: \text{monotone} \rightarrow A^T F(Ay) : \text{monotone}$

\downarrow $A^T F(Ay) = -A^T$

$\Rightarrow -A(\lambda^T y)^{-1}(-A^T y) \in \text{monotone}$

$\Rightarrow (b - A(\lambda^T y)^{-1}(-A^T y)) \in \text{monotone}$

addition of a constant vector to each
 element of a relation set will not change
 monotonicity of a relation; equivalent logic.
 $b \square$ is a monotone operator
 $-A(\lambda_{\text{mon}})^T(A^T b)$ is a monotone operator
 $\Rightarrow (b \square + (-A(\lambda_{\text{mon}})^T(A^T b)))$ is a monotone operator
 $\Rightarrow (\text{Lip} + (-A(\lambda_{\text{mon}})^T(A^T b)))$ is a monotone operator
 [Loc. sum of monotone is monotone.]
 Now we show:

$$\begin{aligned} F(y) &= b - A(\lambda_f)^T(-A^T y) \\ &= \partial_y [b^T y + f^*(-A^T y)] \end{aligned}$$

Proof:

$$\begin{aligned} \partial_y [b^T y + f^*(-A^T y)] &= b + \partial_y [f^*(-A^T y)] \\ &= b + \left[\begin{array}{c} \partial_x f^*(y) \\ \vdots \\ \partial_x f^*(y) \end{array} \right]_{x=A^T y} \quad / \text{ by chain rule: } \partial_x f(Ax) = A^T \partial_x f(x) |_{x=A^T y} \end{aligned}$$

Now

$$\begin{aligned} \partial_y f^*(y) &= \sup_x [f(x) + y^T x] \\ &= \sup_x [-f(x) + y^T x] = \underset{x}{\text{argmax}} \quad / \text{ recall, if } f = \sup_{x \in X} (f_x) \\ &\quad x^* = \underset{x}{\text{argmax}} (f(x)) \\ &= -f(x^*) + y^T x^* = \underset{x}{\text{argmax}} -f(x) + y^T x \\ &= \underset{x}{\text{argmin}} f(x) - y^T x \quad / \text{ find } \eta \in \partial f \Rightarrow \eta \in \partial f^* \\ &= (\partial f^{-1})(y) \end{aligned}$$

$\therefore \partial_y f^*(y) = (\partial f^{-1})(y)$

$$\therefore \partial_y (b^T y + f^*(-A^T y)) = b - \partial_y [(\partial f^{-1})(y)]|_{y=-A^T y}$$

$$= b - A(\partial f^{-1})(-A^T y) = b - A[\partial f^{-1}(y)]|_{y=-A^T y} = b - A(\partial f^{-1})(-A^T y) \quad \blacksquare$$