- Let x[n], $0 \le n \le N-1$, denote a length-N time-domain sequence
- Let X[k], $0 \le k \le N-1$, denote the coefficients of the N-point orthogonal transform of x[n]

 A general form of the orthogonal transform pair is of the form

$$\mathcal{X}[k] = \sum_{n=0}^{N-1} x[n] \psi^*[k,n], \quad 0 \le k \le N-1 \quad \text{Analysis equation}$$

$$x[n] = \sum_{k=0}^{N-1} \mathcal{X}[k] \psi[k,n], \quad 0 \le n \le N-1 \quad \text{Synthesis equation}$$

• $\psi[k,n]$, called the basis sequences, are also length-N sequences

 In the class of transforms to be considered in this course, the basis sequences satisfy the condition

$$\frac{1}{N} \sum_{n=0}^{N-1} \psi[k,n] \psi^*[\ell,n] = \begin{cases} 1, & \ell = k \\ 0, & \ell \neq k \end{cases}$$

 Basis sequences satisfying the above condition are said to be orthogonal to each other

• To verify the inverse transform expression

$$x[n] = \sum_{k=0}^{N-1} X[k] \psi[k,n], \quad 0 \le n \le N-1$$

we substitute it into

$$\chi[k] = \sum_{n=0}^{N-1} x[n] \psi^*[k,n], \quad 0 \le k \le N-1$$

The substitution yields

$$\sum_{n=0}^{N-1} x[n] \psi^*[\ell, n] = \sum_{n=0}^{N-1} \left(\frac{1}{N} \sum_{k=0}^{N-1} \chi[k] \psi[k, n] \right) \psi^*[\ell, n]$$

$$= \sum_{k=0}^{N-1} \chi[k] \left(\frac{1}{N} \sum_{n=0}^{N-1} \psi[k, n] \psi^*[\ell, n] \right) = \chi[\ell]$$

Energy Preservation Property-

$$\sum_{n=0}^{N-1} |x[n]|^2 = \frac{1}{N} \sum_{k=0}^{N-1} |\mathcal{X}[k]|^2$$

- An important consequence of the orthogonality of the basis sequences
- More commonly known as the Parseval's theorem

- <u>Definition</u> The simplest relation between a length-N sequence x[n], defined for $0 \le n \le N-1$, and its DTFT $X(e^{j\omega})$ is obtained by uniformly sampling $X(e^{j\omega})$ on the ω -axis between $0 \le \omega < 2\pi$ at $\omega_k = 2\pi k/N$, $0 \le k \le N-1$
- From the definition of the DTFT we thus have

$$X[k] = X(e^{j\omega})\Big|_{\omega = 2\pi k/N} = \sum_{n=0}^{N-1} x[n]e^{-j2\pi kn/N},$$

$$0 \le k \le N-1$$

- Note: X[k] is also a length-N sequence in the frequency domain
- The sequence X[k] is called the **discrete** Fourier transform (DFT) of the sequence x[n]
- Using the notation $W_N = e^{-j2\pi/N}$ the DFT is usually expressed as:

$$X[k] = \sum_{n=0}^{N-1} x[n]W_N^{kn}, \ 0 \le k \le N-1$$

• The inverse discrete Fourier transform (IDFT) is given by

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn}, \quad 0 \le n \le N-1$$

• To verify the above expression we multiply both sides of the above equation by $W_N^{\ell n}$ and sum the result from n = 0 to n = N-1

resulting in

$$\sum_{n=0}^{N-1} x[n] W_N^{\ell n} = \sum_{n=0}^{N-1} \left(\frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn} \right) W_N^{\ell n}$$

$$= \frac{1}{N} \sum_{n=0}^{N-1} \sum_{k=0}^{N-1} X[k] W_N^{-(k-\ell)n}$$

$$= \frac{1}{N} \sum_{k=0}^{N-1} X[k] \left(\sum_{n=0}^{N-1} W_N^{-(k-\ell)n} \right)$$

From the identity

$$\sum_{n=0}^{N-1} W_N^{-(k-\ell)n} = \begin{cases} N, \text{ for } k-\ell=rN, & r \text{ an integer} \\ 0, & \text{otherwise} \end{cases}$$

it follows then that the only non-zero term in

$$\sum_{n=0}^{N-1} W_N^{-(k-\ell)n}$$

is obtained when $k = \ell$ as $0 \le k, \ell \le N-1$

Hence

$$\sum_{n=0}^{N-1} x[n] W_N^{\ell n} = \frac{1}{N} \sum_{k=0}^{N-1} X[k] \left(\sum_{n=0}^{N-1} W_N^{-(k-\ell)n} \right)$$
$$= \frac{1}{N} \cdot X[\ell] \cdot N = X[\ell]$$

• Example - Consider the length-N sequence

$$x[n] = \begin{cases} 1, & n = 0 \\ 0, & 1 \le n \le N - 1 \end{cases}$$

• Its *N*-point DFT is given by

$$X[k] = \sum_{n=0}^{N-1} x[n]W_N^{kn} = x[0]W_N^0 = 1$$

$$0 \le k \le N-1$$

• Example - Consider the length-N sequence

$$y[n] = \begin{cases} 1, & n = m \\ 0, & 0 \le n \le m - 1, m + 1 \le n \le N - 1 \end{cases}$$

• Its *N*-point DFT is given by

$$Y[k] = \sum_{n=0}^{N-1} y[n]W_N^{kn} = y[m]W_N^{km} = W_N^{km}$$

$$0 < k < N-1$$

- Example Consider the length-N sequence defined for $0 \le n \le N-1$ $g[n] = \cos(2\pi rn/N), 0 \le r \le N-1$
- Using a trigonometric identity we can write

$$g[n] = \frac{1}{2} \left(e^{j2\pi rn/N} + e^{-j2\pi rn/N} \right)$$
$$= \frac{1}{2} \left(W_N^{-rn} + W_N^{rn} \right)$$

• The N-point DFT of g[n] is thus given by

$$G[k] = \sum_{n=0}^{N-1} g[n] W_N^{kn}$$

$$= \frac{1}{2} \left(\sum_{n=0}^{N-1} W_N^{-(r-k)n} + \sum_{n=0}^{N-1} W_N^{(r+k)n} \right),$$

$$0 \le k \le N-1$$

Making use of the identity

$$\sum_{n=0}^{N-1} W_N^{-(k-\ell)n} = \begin{cases} N, \text{ for } k - \ell = rN, r \text{ an integer} \\ 0, \text{ otherwise} \end{cases}$$

we get

$$G[k] = \begin{cases} N/2, & \text{for } k = r \\ N/2, & \text{for } k = N - r \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le k \le N-1$$

The DFT samples defined by

$$X[k] = \sum_{n=0}^{N-1} x[n] W_N^{kn}, \quad 0 \le k \le N-1$$

can be expressed in matrix form as

$$\mathbf{X} = \mathbf{D}_N \mathbf{x}$$

where

$$\mathbf{X} = \begin{bmatrix} X[0] & X[1] & \cdots & X[N-1] \end{bmatrix}^T$$

$$\mathbf{x} = \begin{bmatrix} x[0] & x[1] & \cdots & x[N-1] \end{bmatrix}^T$$

and \mathbf{D}_N is the $N \times N$ **DFT** matrix given by

$$\mathbf{D}_{N} = \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & W_{N}^{1} & W_{N}^{2} & \cdots & W_{N}^{(N-1)} \\ 1 & W_{N}^{2} & W_{N}^{4} & \cdots & W_{N}^{2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & W_{N}^{(N-1)} & W_{N}^{2(N-1)} & \cdots & W_{N}^{(N-1)^{2}} \end{bmatrix}$$

Likewise, the IDFT relation given by

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn}, \ 0 \le n \le N-1$$

can be expressed in matrix form as

$$\mathbf{x} = \mathbf{D}_N^{-1} \mathbf{X}$$

where \mathbf{D}_{N}^{-1} is the $N \times N$ **IDFT** matrix

where

$$\mathbf{D}_{N}^{-1} = \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & W_{N}^{-1} & W_{N}^{-2} & \cdots & W_{N}^{-(N-1)} \\ 1 & W_{N}^{-2} & W_{N}^{-4} & \cdots & W_{N}^{-2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & W_{N}^{-(N-1)} & W_{N}^{-2(N-1)} & \cdots & W_{N}^{-(N-1)^{2}} \end{bmatrix}$$

• Note:

$$\mathbf{D}_N^{-1} = \frac{1}{N} \mathbf{D}_N^*$$

DFT Computation Using MATLAB

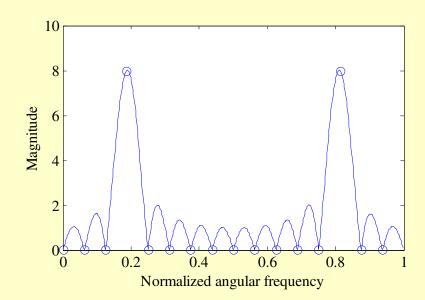
- The functions to compute the DFT and the IDFT are fft and ifft
- These functions make use of FFT algorithms which are computationally highly efficient compared to the direct computation
- Programs 5_1.m and 5_2.m illustrate the use of these functions

DFT Computation Using MATLAB

• Example - Program 5_3.m can be used to compute the DFT and the DTFT of the sequence

$$x[n] = \cos(6\pi n/16), 0 \le n \le 15$$

as shown below



• indicates DFT samples

• The *N*-point DFT X[k] of a length-*N* sequence x[n] is simply the frequency samples of its DTFT $X(e^{j\omega})$ evaluated at *N* uniformly spaced frequency points

$$\omega = \omega_k = 2\pi k/N$$
, $0 \le k \le N-1$

• Given the *N*-point DFT X[k] of a length-*N* sequence x[n], its DTFT $X(e^{j\omega})$ can be uniquely determined from X[k]

Thus

$$X(e^{j\omega}) = \sum_{n=0}^{N-1} x[n]e^{-j\omega n}$$

$$= \sum_{n=0}^{N-1} \left[\frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn} \right] e^{-j\omega n}$$

$$= \frac{1}{N} \sum_{k=0}^{N-1} X[k] \sum_{n=0}^{N-1} e^{-j(\omega - 2\pi k/N)n}$$

 To develop a compact expression for the sum S, let

$$r = e^{-j(\omega - 2\pi k/N)}$$

- Then $S = \sum_{n=0}^{N-1} r^n$
- From the above

$$rS = \sum_{n=1}^{N} r^n = 1 + \sum_{n=1}^{N-1} r^n + r^N - 1$$
$$= \sum_{n=0}^{N-1} r^n + r^N - 1 = S + r^N - 1$$

• Or, equivalently,

$$S - rS = (1 - r)S = 1 - r^{N}$$

Hence

$$S = \frac{1 - r^{N}}{1 - r} = \frac{1 - e^{-j(\omega N - 2\pi k)}}{1 - e^{-j[\omega - (2\pi k/N)]}}$$

$$= \frac{\sin\left(\frac{\omega N - 2\pi k}{2}\right)}{\sin\left(\frac{\omega N - 2\pi k}{2N}\right)} \cdot e^{-j[(\omega - 2\pi k/N)][(N-1)/2]}$$

Therefore

$$= \frac{1}{N} \sum_{k=0}^{N-1} X[k] \frac{\sin\left(\frac{\omega N - 2\pi k}{2}\right)}{\sin\left(\frac{\omega N - 2\pi k}{2N}\right)} \cdot e^{-j[(\omega - 2\pi k/N)][(N-1)/2]}$$

- Consider a sequence x[n] with a DTFT $X(e^{j\omega})$
- We sample $X(e^{j\omega})$ at N equally spaced points $\omega_k = 2\pi k/N, \ 0 \le k \le N-1$ developing the N frequency samples $\{X(e^{j\omega_k})\}$
- These *N* frequency samples can be considered as an *N*-point DFT *Y*[*k*] whose *N*-point IDFT is a length-*N* sequence *y*[*n*]

• Now
$$X(e^{j\omega}) = \sum_{\ell=-\infty}^{\infty} x[\ell] e^{-j\omega\ell}$$

• Thus
$$Y[k] = X(e^{j\omega_k}) = X(e^{j2\pi k/N})$$

$$= \sum_{\ell=-\infty}^{\infty} x[\ell] e^{-j2\pi k\ell/N} = \sum_{\ell=-\infty}^{\infty} x[\ell] W_N^{k\ell}$$

• An IDFT of *Y*[*k*] yields

$$y[n] = \frac{1}{N} \sum_{k=0}^{N-1} Y[k] W_N^{-kn}$$

• i.e.
$$y[n] = \frac{1}{N} \sum_{k=0}^{N-1} \sum_{\ell=-\infty}^{\infty} x[\ell] W_N^{k\ell} W_N^{-kn}$$

$$= \sum_{\ell=-\infty}^{\infty} x[\ell] \left[\frac{1}{N} \sum_{k=0}^{N-1} W_N^{-k(n-\ell)} \right]$$

Making use of the identity

$$\frac{1}{N} \sum_{n=0}^{N-1} W_N^{-k(n-r)} = \begin{cases} 1, & \text{for } r = n + mN \\ 0, & \text{otherwise} \end{cases}$$

we arrive at the desired relation

$$y[n] = \sum_{m=-\infty}^{\infty} x[n+mN], \quad 0 \le n \le N-1$$

• Thus y[n] is obtained from x[n] by adding an infinite number of shifted replicas of x[n], with each replica shifted by an integer multiple of N sampling instants, and observing the sum only for the interval $0 \le n \le N-1$

To apply

$$y[n] = \sum_{m=-\infty}^{\infty} x[n+mN], \quad 0 \le n \le N-1$$

to finite-length sequences, we assume that the samples outside the specified range are zeros

• Thus if x[n] is a length-M sequence with $M \le N$, then y[n] = x[n] for $0 \le n \le N-1$

- If M > N, there is a time-domain aliasing of samples of x[n] in generating y[n], and x[n] cannot be recovered from y[n]
- Example Let $\{x[n]\} = \{0 \ 1 \ 2 \ 3 \ 4 \ 5\}$
- By sampling its DTFT $X(e^{j\omega})$ at $\omega_k = 2\pi k/4$, $0 \le k \le 3$ and then applying a 4-point IDFT to these samples, we arrive at the sequence y[n] given by

$$y[n] = x[n] + x[n+4] + x[n-4], 0 \le n \le 3$$

• i.e.

$$\{y[n]\} = \{4 \quad 6 \quad 2 \quad 3\}$$



 $\{x[n]\}\$ cannot be recovered from $\{y[n]\}$

Numerical Computation of the DTFT Using the DFT

- A practical approach to the numerical computation of the DTFT of a finite-length sequence
- Let $X(e^{j\omega})$ be the DTFT of a length-N sequence x[n]
- We wish to evaluate $X(e^{j\omega})$ at a dense grid of frequencies $\omega_k = 2\pi k/M$, $0 \le k \le M-1$, where M >> N:

Numerical Computation of the DTFT Using the DFT

$$X(e^{j\omega_k}) = \sum_{n=0}^{N-1} x[n]e^{-j\omega_k n} = \sum_{n=0}^{N-1} x[n]e^{-j2\pi kn/M}$$

• Define a new sequence

$$x_e[n] = \begin{cases} x[n], & 0 \le n \le N - 1 \\ 0, & N \le n \le M - 1 \end{cases}$$

Then

$$X(e^{j\omega_k}) = \sum_{n=0}^{M-1} x_e[n]e^{-j2\pi kn/M}$$

Numerical Computation of the DTFT Using the DFT

- Thus $X(e^{j\omega_k})$ is essentially an M-point DFT $X_e[k]$ of the length-M sequence $x_e[n]$
- The DFT $X_e[k]$ can be computed very efficiently using the FFT algorithm if M is an integer power of 2
- The function freqz employs this approach to evaluate the frequency response at a prescribed set of frequencies of a DTFT expressed as a rational function in $e^{-j\omega}$