Introduction: Fibonacci Numbers II

Daniel Kane

Department of Computer Science and Engineering University of California, San Diego

Algorithmic Toolbox Data Structures and Algorithms

Learning Objectives

- Produce a simple algorithm to compute Fibonacci numbers.
- Show that this algorithm is very slow.

Definition

$$F_n = \begin{cases} 0, & n = 0, \\ 1, & n = 1, \\ F_{n-1} + F_{n-2}, & n > 1. \end{cases}$$

Definition

$$F_n = \begin{cases} 0, & n = 0, \\ 1, & n = 1, \\ F_{n-1} + F_{n-2}, & n > 1. \end{cases}$$

Grows rapidly.

Computing Fibonacci numbers

Compute F_n

Input: An integer $n \geq 0$.

Output: F_n .

Algorithm

```
if n \le 1:
return n
```

Algorithm

```
FibRecurs(n)
```

```
\begin{array}{l} \text{if } n \leq 1 \colon \\ \text{return } n \\ \\ \text{else:} \\ \text{return FibRecurs}(n-1) + \text{FibRecurs}(n-2) \end{array}
```

Let T(n) denote the number of lines of code executed by FibRecurs(n).

If $n \leq 1$

```
 \begin{array}{l} \text{if } n \leq 1: \\ \text{return } n \\ \\ \text{else:} \\ \text{return FibRecurs}(n-1) + \text{FibRecurs}(n-2) \\ \end{array}
```

If $n \leq 1$

FibRecurs(n)

T(n) = 2.

```
if n \le 1:
return n
else:
return FibRecurs(n-1) + FibRecurs(n-2)
```

If $n \ge 2$

```
\label{eq:normalized} \begin{split} &\text{if } n \leq 1\colon \\ &\text{return } n \\ &\text{else:} \\ &\text{return } \text{FibRecurs}(n-1) + \text{FibRecurs}(n-2) \end{split}
```

If $n \ge 2$

```
 \begin{array}{l} \text{if } n \leq 1 \colon \\ \\ \text{return } n \\ \\ \text{else:} \\ \\ \text{return FibRecurs}(n-1) + \\ \text{FibRecurs}(n-2) \\ \end{array}
```

$$T(n) = 3$$

If $n \ge 2$

```
 \begin{array}{l} \text{if } n \leq 1 \colon \\ \\ \text{return } n \\ \\ \text{else:} \\ \\ \text{return FibRecurs}(n-1) + \\ \text{FibRecurs}(n-2) \\ \end{array}
```

$$T(n) = 3 + T(n-1) + T(n-2).$$

$$T(n) = \begin{cases} 2 & \text{if } n \leq 1 \\ T(n-1) + T(n-2) + 3 & \text{else.} \end{cases}$$

$$T(n) = \begin{cases} 2 & \text{if } n \leq 1 \\ T(n-1) + T(n-2) + 3 & \text{else.} \end{cases}$$

Therefore $T(n) \geq F_n$

$$T(n) = \begin{cases} 2 & \text{if } n \leq 1 \\ T(n-1) + T(n-2) + 3 & \text{else.} \end{cases}$$

Therefore
$$T(n) \geq F_n$$

$$T(100) \approx 1.77 \cdot 10^{21} \qquad (1.77 \text{ sextillion})$$

$$T(n) = \begin{cases} 2 & \text{if } n \leq 1 \\ T(n-1) + T(n-2) + 3 & \text{else.} \end{cases}$$

Therefore
$$T(n) \geq F_n$$

$$T(100) \approx 1.77 \cdot 10^{21}$$
 (1.77 sextillion)

Takes 56,000 years at 1GHz.







