1 Complexity Zoo

$1.1 \quad TIME[f(n)]$

Informally: problems that can be solved in f(n) time.

Definition 1.1. Given some function $f: \mathbb{N} \to \mathbb{N}$, TIME[f(n)] are the set of problems solvable within O(f(n)) atomic steps on a deterministic Turing machine. Where n is the size of the input.

$1.2 \quad NTIME[f(n)]$

Informally: problems that can be solved nondeterministically in f(n) time.

Definition 1.2. Given some function $f: \mathbb{N} \to \mathbb{N}$, NTIME[f(n)] are the set of problems solvable within O(f(n)) atomic steps on a nondeterministic Turing machine.

$1.3 \quad SPACE[f(n)]$

Informally: problems that can be solved in f(n) space.

Definition 1.3. Given some function $f: \mathbb{N} \to \mathbb{N}$, SPACE[f(n)] are the set of problems solvable using a tape of length O(f(n)) on a deterministic Turing machine. Where n is the size of the input.

1.4 NSPACE[f(n)]

Informally: problems that can be solved non-deterministically in f(n) space.

Definition 1.4. Given some function $f: \mathbb{N} \to \mathbb{N}$, NSPACE[f(n)] are the set of problems solvable using a tape of length O(f(n)) on a non-deterministic Turing machine. Where n is the size of the input.

1.5 P

Informally: all problems that can be solved in polynomial time.

Definition 1.5.

$$\mathbf{P} = \bigcup_{k \ge 0} \mathrm{TIME}[n^k]$$

Descriptive Complexity definitions:

Definition 1.6.

$$\mathbf{P} = FO(LFP)$$

(First Order logic extended with the Least Fixed Point operator, with successor. A high level, handwavy description of the LFP operator is the added ability to recursively define FO formulas.)

Definition 1.7.

$$\mathbf{P} = SO(Horn)$$

(Second Order logic restricted with Horn. SO logic allows you to quantify over subsets/relations/functions on the domain, and Horn means all 'clauses' are really implications with literal in the conclusion and all literals positive.)

Circuit Complexity definition:

Definition 1.8.

P = Set of problems that can be solved by a polynomial-time uniform family of boolean circuits

Notable Problems in \mathbf{P} :

- 2-SAT
- 2-Colourability
- Reachability

1.6 NP

Informally: all problems that can be solved in nondeterministic polynomial time.

Definition 1.9.

$$\mathbf{NP} = \bigcup_{k \ge 0} \mathrm{NTIME}[n^k]$$

In terms of a verifier:

Informally: The set of decision problems where a solution can be verified in polynomial time.

Descriptive Complexity Definition:

Definition 1.10.

$$\mathbf{NP} = \mathrm{SO}\exists$$

(Existential Second Order)

Notable Problems in \mathbf{NP} :

- SAT
- 3-Colourability
- TSP
- Subset sum

1.7 FPT

Informally, the set of problems that can be solved in polynomial time for some fixed parameter.

Definition 1.11. The set of problems that can be parameterised by k and can be solved in $f(k)n^c$, where f(x) is only dependent on k, and c is an independent constant.

P is contained within **FPT**. If a problem is in **FPT**, then for any fixed k that problem is in **P**. **FPT** is also known as $\mathbf{W}[\mathbf{0}]$

Notable Problems in \mathbf{FPT} :

• Vertex Cover



- 1.8 W[1]
- 1.9 W[2]
- 1.10 W[i]
- 1.11 FPTAS
- 1.12 PTAS
- 1.13 L
- 1.14 NL
- 1.15 PSPACE
- 1.16 coNP
- 1.17 Σ_2^p
- 1.18 Σ_i^p
- 1.19 Π_2^p
- 1.20 Π_i^p
- 1.21 PH
- **1.22** P^{SAT}
- 1.23 NP^{SAT}
- 1.24 P/poly
- 1.25 P-Uniform
- 1.26 EXP
- 1.27 NC
- 1.28 NC_0
- 1.29 NC_1
- 1.30 NC_2
- 1.31 NC_i
- **1.32** AC_i
- 1.33 AC_0
- **1.34** AC_1
- 1.35 BPP
- 1.36 RP
- 1.37 co-RP
- 1.38 ZPP
- 1.39 APX
- 1.40 PO

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