

# 1 Complexity Zoo

## 1.1 TIME[f(n)]

Informally: problems that can be solved in  $f(n)$  time.

**Definition 1.1.** Given some function  $f : \mathbb{N} \rightarrow \mathbb{N}$ ,  $\text{TIME}[f(n)]$  are the set of problems solvable within  $O(f(n))$  atomic steps on a deterministic Turing machine. Where  $n$  is the size of the input.

## 1.2 NTIME[f(n)]

Informally: problems that can be solved nondeterministically in  $f(n)$  time.

**Definition 1.2.** Given some function  $f : \mathbb{N} \rightarrow \mathbb{N}$ ,  $\text{NTIME}[f(n)]$  are the set of problems solvable within  $O(f(n))$  atomic steps on a nondeterministic Turing machine.

## 1.3 SPACE[f(n)]

Informally: problems that can be solved in  $f(n)$  space.

**Definition 1.3.** Given some function  $f : \mathbb{N} \rightarrow \mathbb{N}$ ,  $\text{SPACE}[f(n)]$  are the set of problems solvable using a tape of length  $O(f(n))$  on a deterministic Turing machine. Where  $n$  is the size of the input.

## 1.4 NSPACE[f(n)]

Informally: problems that can be solved non-deterministically in  $f(n)$  space.

**Definition 1.4.** Given some function  $f : \mathbb{N} \rightarrow \mathbb{N}$ ,  $\text{NSPACE}[f(n)]$  are the set of problems solvable using a tape of length  $O(f(n))$  on a non-deterministic Turing machine. Where  $n$  is the size of the input.

## 1.5 P

Informally: all problems that can be solved in polynomial time.

**Definition 1.5.**

$$\mathbf{P} = \bigcup_{k \geq 0} \text{TIME}[n^k]$$

Descriptive Complexity definitions:

**Definition 1.6.**

$$\mathbf{P} = \text{FO}(\text{LFP})$$

(First Order logic extended with the Least Fixed Point operator, with successor. A high level, handwavy description of the LFP operator is the added ability to recursively define FO formulas.)

**Definition 1.7.**

$$\mathbf{P} = \text{SO}(\text{Horn})$$

(Second Order logic restricted with Horn. SO logic allows you to quantify over subsets/relations/functions on the domain, and Horn means all ‘clauses’ are really implications with literal in the conclusion and all literals positive.)

Circuit Complexity definition:

**Definition 1.8.**

$\mathbf{P}$  = Set of problems that can be solved by a polynomial-time uniform family of boolean circuits

Notable Problems in  $\mathbf{P}$ :

- 2-SAT
- 2-Colourability
- Reachability

## 1.6 NP

Informally: all problems that can be solved in nondeterministic polynomial time.

**Definition 1.9.**

$$\mathbf{NP} = \bigcup_{k \geq 0} \text{NTIME}[n^k]$$

In terms of a verifier:

Informally: The set of decision problems where a solution can be verified in polynomial time.

Descriptive Complexity Definition:

**Definition 1.10.**

$$\mathbf{NP} = \text{SO}\exists$$

(Existential Second Order)

Notable Problems in  $\mathbf{NP}$ :

- SAT
- 3-Colourability
- TSP
- Subset sum

## 1.7 FPT

Informally, the set of problems that can be solved in polynomial time for some fixed parameter.

**Definition 1.11.** The set of problems that can be parameterised by  $k$  and can be solved in  $f(k)n^c$ , where  $f(x)$  is only dependent on  $k$ , and  $c$  is an independent constant.

**P** is contained within **FPT**.

If a problem is in **FPT**, then for any fixed  $k$  that problem is in **P**.

**FPT** is also known as **W[0]**

Notable Problems in **FPT**:

- Vertex Cover

## 1.8 W[1]

**Definition 1.12.** The class of parametrized problems that admit a parametrized reduction to the following problem: Given a nondeterministic single-tape Turing machine, decide if it accepts within  $k$  steps.

N.B This is short acceptance

**Definition 1.13.** The class of parametrized problems that admit a parametrized reduction to the following problem: Given a Boolean circuit  $C$ , with a mixture of fanin-2 and unbounded-fanin gates. There is at most 1 unbounded-fanin gate along any path to the root, and the total depth (fanin-2 and unbounded-fanin) is constant. Does  $C$  have a satisfying assignment of Hamming weight  $k$ ?

N.B This is Weighted 3-SAT.

Notable Problems in **W[1]**:

- Short Acceptance
- Weighted 3-SAT
- Clique (of size  $k$ )
- Independent set (of size  $k$ )



1.9    **W[2]**  
 1.10   **W[i]**  
 1.11   **FPTAS**  
 1.12   **PTAS**  
 1.13   **L**  
 1.14   **NL**  
 1.15   **PSPACE**  
 1.16   **coNP**  
 1.17    $\Sigma_2^p$   
 1.18    $\Sigma_i^p$   
 1.19    $\Pi_2^p$   
 1.20    $\Pi_i^p$   
 1.21   **PH**  
 1.22    $P^{SAT}$   
 1.23    $NP^{SAT}$   
 1.24   **P/poly**  
 1.25   **P-Uniform**  
 1.26   **EXP**  
 1.27   **NC**  
 1.28    $NC_0$   
 1.29    $NC_1$   
 1.30    $NC_2$   
 1.31    $NC_i$   
 1.32    $AC_i$   
 1.33    $AC_0$   
 1.34    $AC_1$   
 1.35   **BPP**  
 1.36   **RP**  
 1.37   **co-RP**  
 1.38   **ZPP**  
 1.39   **APX**  
 1.40   **PO**  
 1.41   **PCP**