

## Disjoint Sets

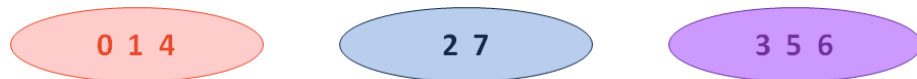
Let **R** be an equivalence relation. We represent R as several disjoint sets. Two key ideas from Monday:

- Each element exists in exactly one set.
- Every set is an equitant representation.
  - Mathematically:  $4 \in [0]_R \rightarrow 8 \in [0]_R$
  - Programmatically: `find(4) == find(8)`

## Building Disjoint Sets:

- Maintain a collection  $S = \{s_0, s_1, \dots, s_k\}$
- Each set has a representative member.
- ADT:

```
void makeSet(const T &t);
void union(const T &k1, const T &k2);
T &find(const T &k);
```



[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]

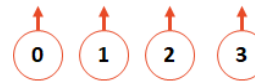
**Operation:** `find(k)`

**Operation:** `union(k1, k2)`

## Implementation #2:

- Continue to use an array where the index is the key
- The value of the array is:
  - 1, if we have found the representative element
  - The index of the parent**, if we haven't found the rep. element

## Impl #2 (continued):



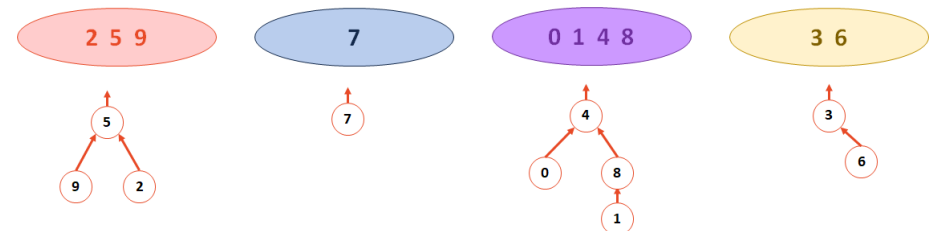
[0]	[1]	[2]	[3]

[0]	[1]	[2]	[3]

[0]	[1]	[2]	[3]

[0]	[1]	[2]	[3]

## Example:



4	8	5	6	-1	-1	-1	-1	4	5
[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]

...where is the error in this table?

## Implementation – DisjointSets::find

DisjointSets.cpp (partial)	
1	<code>int DisjointSets::find(int i) {</code>
2	<code>    if ( s[i] &lt; 0 ) { return i; }</code>
3	<code>    else { return _find( s[i] ); }</code>
4	<code>}</code>

What is the running time of `find`?

Structure which is similar to a Linked List  
 $O(h) == O(n)$

What is the ideal UpTree?

One root node with every other nodes as  
 its children  
 $O(1)$

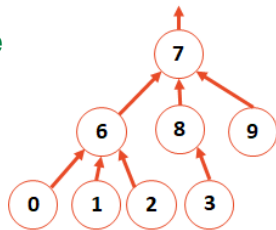
## Implementation – DisjointSets::union

DisjointSets.cpp (partial)	
1	void DisjointSets::union(int r1, int r2) {
2	
3	
4	}

How do we want to union the two UpTrees?

### Building a Smart Union Function

if use 7 as new root, most of the element do not change height



if using 4 as the new root, then the max height does not change



The implementation of this visual model is the following:

6	6	6	8	-1	10	7	-1	7	7	4	5
[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]	[10]	[11]

What are possible strategies to employ when building a “smart union”?

Union by Height ?  
Union by Size ?

## Smart Union Strategy #1: Union by height

**Idea:** Keep the height of the tree as small as possible!

Metadata at Root:

After union( 4, 7 ):

6	6	6	8		10	7		7	7	4	5
[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]	[10]	[11]

## Smart Union Strategy #2: Union by size

**Idea:** Minimize the number of nodes that increase in height.  
(Observe that the tree we union have all their nodes gain in height.)

Metadata at Root:

After union( 4, 7 ):

6	6	6	8		10	7		7	7	4	5
[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]	[10]	[11]

## Smart Union Implementation:

DisjointSets.cpp (partial)	
1	void DisjointSets::unionBySize(int root1, int root2) {
2	int newSize = arr_[root1] + arr_[root2];
3	
4	if ( arr_[root1] < arr_[root2] ) {
5	arr_[root2] = root1; arr_[root1] = newSize;
6	} else {
7	arr_[root1] = root2; arr_[root2] = newSize;
8	}
9	}

## CS 225 – Things To Be Doing:

1. Theory Exam 3 is on-going
2. MP6 released; Extra Credit deadline on Monday, April 9<sup>th</sup>
3. lab\_heaps released today
4. Daily POTDs are ongoing!