CS 225

**Data Structures** 

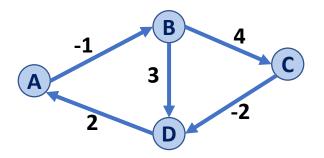
April 30 — Floyd-Warshall's Algorithm
Wade Fagen-Ulmschneider

Floyd-Warshall's Algorithm is an alterative to Dijkstra in the presence of negative-weight edges (not negative weight cycles).

```
FloydWarshall(G):
     Let d be a adj. matrix initialized to +inf
     foreach (Vertex v : G):
       d[v][v] = 0
     foreach (Edge (u, v) : G):
       d[u][v] = cost(u, v)
10
11
12
     foreach (Vertex u : G):
13
       foreach (Vertex v : G):
14
         foreach (Vertex w : G):
15
           if d[u, v] > d[u, w] + d[w, v]:
16
             d[u, v] = d[u, w] + d[w, v]
```

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FloydWarshall(G):
     Let d be a adj. matrix initialized to +inf
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         foreach (Vertex w : G):
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           if d[u, v] > d[u, w] + d[w, v]:
16
             d[u, v] = d[u, w] + d[w, v]
```

	Α	В	С	D
A				
В				
С				
D				



```
12 foreach (Vertex u : G):
13 foreach (Vertex v : G):
14 foreach (Vertex k : G):
15 if d[u, v] > d[u, k] + d[k, v]:
16 d[u, v] = d[u, w] + d[w, v]
```

	Α	В	С	D
Α	0	-1	∞	∞
В	∞	0	4	3
С	∞	∞	0	-2
D	2	∞	∞ B	0
	<b>C</b>			

#### Let us consider k=A:





$$C \longrightarrow B + \infty$$
 vs.  $C \longrightarrow A \longrightarrow B + \infty$ 

$$C \longrightarrow D$$
 -2 VS.  $C \longrightarrow A \longrightarrow D$  + $\propto$ 

$$D \longrightarrow B + \infty$$
 vs.  $D \longrightarrow B \longrightarrow B$ 

$$D \longrightarrow C + \infty$$
 VS.  $D \longrightarrow B \longrightarrow C$ 

	Α	В	С	D
A	0	-1	∞	∞
В	∞	0	4	3
С	∞	∞	0	-2
D	2	∞	∞ (3)	0
-1 B 4 C				

	Α	В	С	D
A	0	-1	∞	∞
В	∞	0	4	3
С	∞	∞	0	-2
D	2	1	∞ •	0
		-1	3 B	<b>C</b>

#### Let us consider k=B:



vs.  $A \rightarrow B \rightarrow C$ 



 $VS. \qquad A \longrightarrow B \longrightarrow D$ 



vs.  $(C) \rightarrow (B) \rightarrow (A)$ 



vs.  $C \rightarrow B \rightarrow D$ 



vs.  $D \rightarrow B \rightarrow A$ 



vs.  $(D) \rightarrow (B) \rightarrow (C)$ 

	Α	В	С	D
A	0	-1	∞	∞
В	∞	0	4	3
С	∞	∞	0	-2
D	2	1	∞ •	0
		-1	B	4 C

Running Time?

```
FloydWarshall(G):
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       d[u][v] = cost(u, v)
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     foreach (Vertex u : G):
13
       foreach (Vertex v : G):
14
         foreach (Vertex w : G):
           if d[u, v] > d[u, w] + d[w, v]:
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16
             d[u, v] = d[u, w] + d[w, v]
```

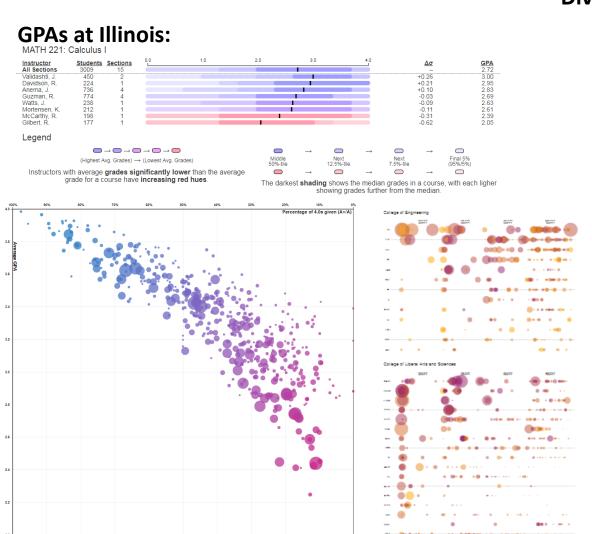
### Final Exam Review Session

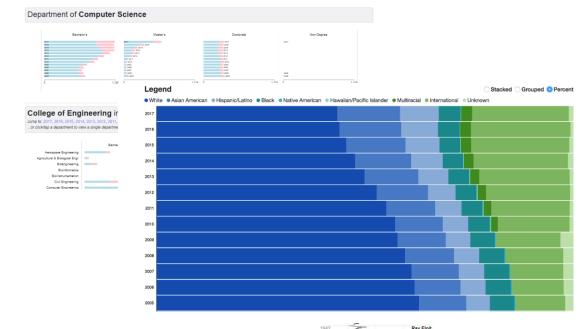
- Implementations
  - Edge List
  - Adjacency Matrix
  - Adjacency List
- Traversals
  - Breadth First
  - Depth First
- Minimum Spanning Tree
  - Kruskal's Algorithm
  - Prim's Algorithm
- Shortest Path
  - Dijkstra's Algorithm
  - Floyd-Warshall's Algorithm

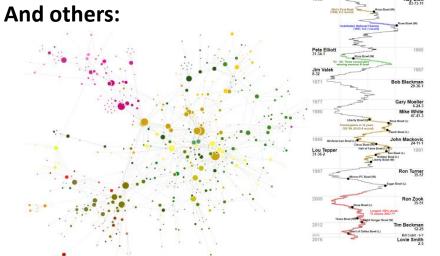
...and this is just the beginning. The journey continues to CS 374!

# My Passion: Data Discovery

#### **Diversity at Illinois:**







### Final Exam Review Session

 CS 225 will be doing an in-lecture final exam review session during Wednesday's lecture.

### **ICES Forms**

My promise: I will read the back of every ICES form. Please take the time to give feedback on the course.

Thanks for an amazing semester!