CS 225

Data Structures

March 5 — AVL Applications
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AVL Runtime Proof

On Friday, we proved an upper-bound on the height of an AVL tree is 2*lg(n) or O(lg(n)).

Summary of Balanced BST

AVL Trees

- Max height: 1.44 * lg(n)
- Rotations:

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AVL Trees

- Max height: 1.44 * lg(n)
- Rotations:

Zero rotations on find One rotation on insert O(h) == O(lg(n)) rotations on remove

Red-Black Trees

- Max height: 2 * lg(n)
- Constant number of rotations on insert, remove, and find

Why AVL?

Summary of Balanced BST

Pros:

- Running Time:

- Improvement Over:

- Great for specific applications:

Summary of Balanced BST

Cons:

- Running Time:

- In-memory Requirement:

C++ provides us a balanced BST as part of the standard library:

```
std::map<K, V> map;
```

```
V & std::map<K, V>::operator[]( const K & )
```

```
V & std::map<K, V>::operator[]( const K & )
std::map<K, V>::erase( const K & )
```

```
iterator std::map<K, V>::lower_bound( const K & );
iterator std::map<K, V>::upper_bound( const K & );
```

CS 225 -- Course Update

This weekend, the following grades were updated:

- mp1
- mp2*
- mp3*
- lab_inheritance
- lab_quacks
- lab_trees

Why do we care?

```
DFS dfs(...);
for ( ImageTraversal::Iterator it = dfs.begin(); it != dfs.end(); ++it ) {
    std::cout << (*it) << std::endl;
}</pre>
```

Why do we care?

```
DFS dfs(...);
for ( ImageTraversal::Iterator it = dfs.begin(); it != dfs.end(); ++it ) {
    std::cout << (*it) << std::endl;
}</pre>
```

```
1 DFS dfs(...);
2 for ( const Point & p : dfs ) {
3   std::cout << p << std::endl;
4 }</pre>
```

Why do we care?

```
1 DFS dfs(...);
2 for ( ImageTraversal::Iterator it = dfs.begin(); it != dfs.end(); ++it ) {
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```
1 DFS dfs(...);
2 for ( const Point & p : dfs ) {
3   std::cout << p << std::endl;
4 }</pre>
```

```
1 ImageTraversal & traversal = /* ... */;
2 for ( const Point & p : traversal ) {
3   std::cout << p << std::endl;
4 }</pre>
```

```
1 ImageTraversal *traversal = /* ... */;
2 for ( const Point & p : traversal ) {
3   std::cout << p << std::endl;
4 }</pre>
```

Every Data Structure So Far

	Unsorted Array	Sorted Array	Unsorted List	Sorted List	Binary Tree	BST	AVL
Find							
Insert							
Remove							
Traverse							

Q: Consider points in 1D: $p = \{p_1, p_2, ..., p_n\}$what points fall in [11, 42]?

Tree construction:

Balanced BSTs are useful structures for range-based and nearest-neighbor searches.

Q: Consider points in 1D: $\mathbf{p} = \{\mathbf{p}_1, \mathbf{p}_2, ..., \mathbf{p}_n\}$...what points fall in [11, 42]?

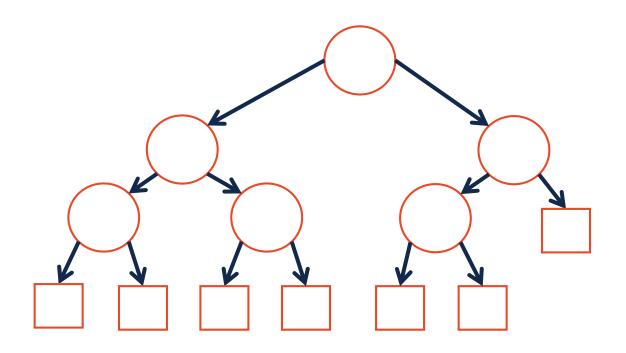


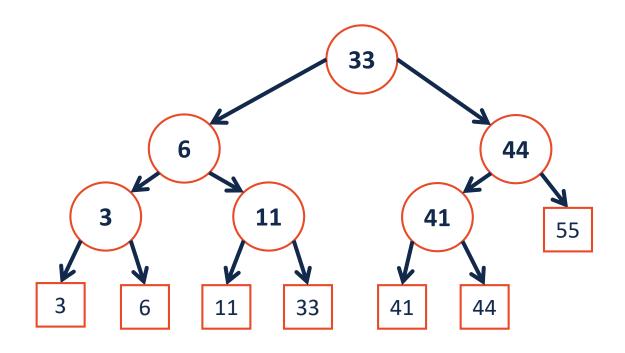
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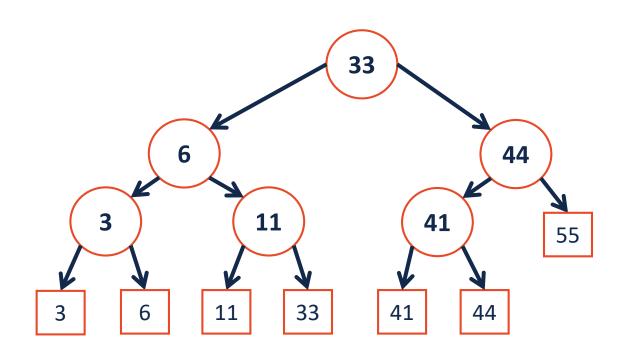
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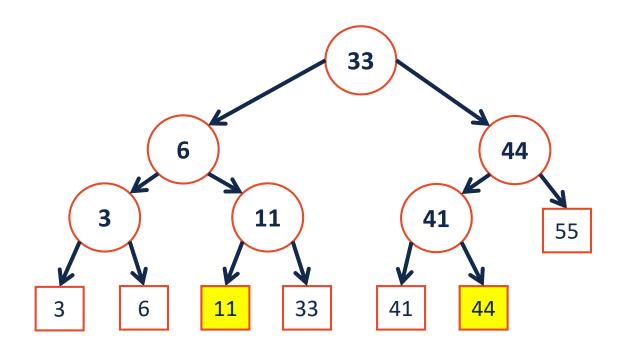
Tree construction:



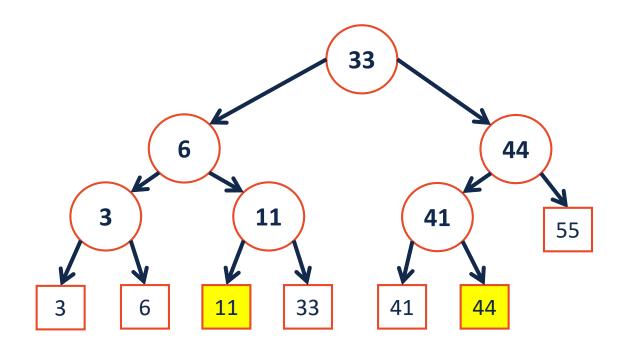


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Running Time



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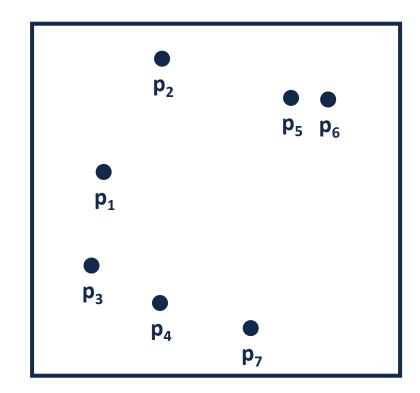




Consider points in 2D: $p = \{p_1, p_2, ..., p_n\}$.

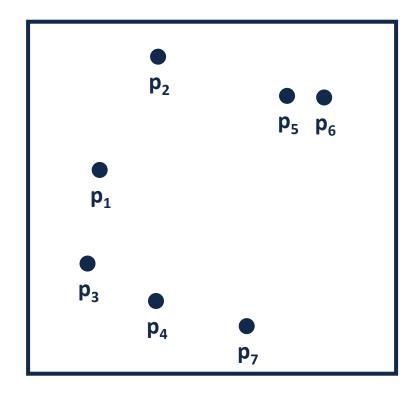
Q: What points are in the rectangle: $(x_1, y_1), (x_2, y_2)$]?

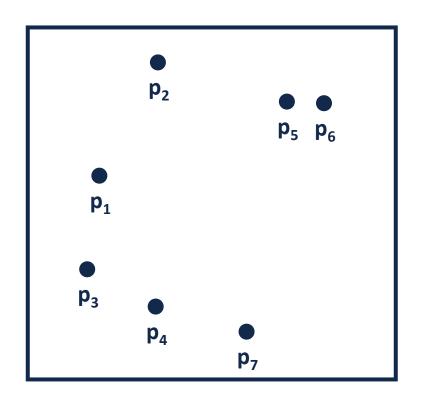
Q: What is the nearest point to (x_1, y_1) ?

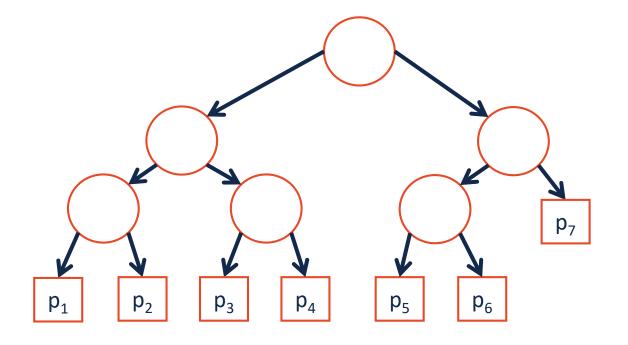


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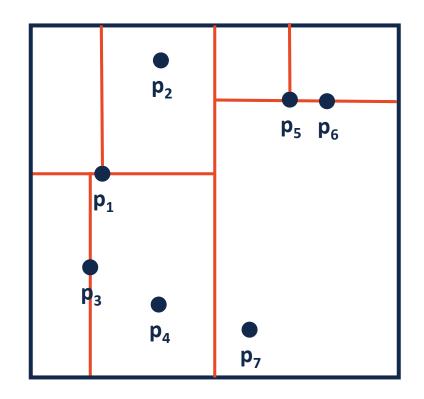
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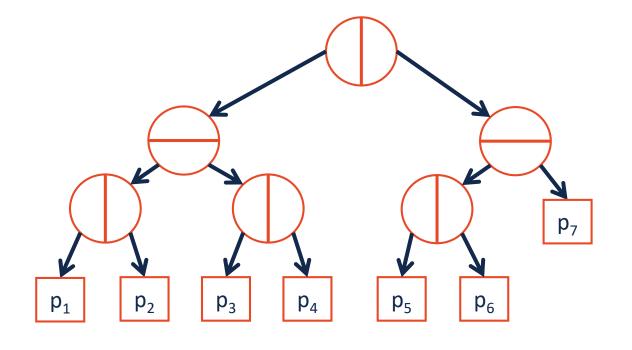






kD-Trees





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