

Optimal Stopping

- Setting - State space $S \cup \{ \text{Stopped} \}$
 - Action space $A = \{0, 1\}$
 - If $A_t = 0$, then $S_{t+i} = \text{'Stopped'}$ for all $i \geq 1$
 - Costs $C_t(x, 0) = c(x)$ (continuation cost)
 $\xrightarrow{x \neq \text{Stopped}} C_t(x, 1) = s(x)$ (stopping cost)
 $c_t(\text{Stopped}, \cdot) = 0$

• If we play this over a horizon T , then HJB eqn

$$V_t(x) = \min \left[c(x) + \mathbb{E} \left[V_{t+1}(X_{t+1}) \mid A_{t+1}^{x_{t+1}}, \begin{array}{l} a=1, \text{ie, continue} \\ a=0 \end{array} \right], s(x) \right] \quad x \in S$$

$$V_t(\text{Stopped}) = 0, \quad V_T(x) = 0 \quad \forall x$$

- 1-Step Look-Ahead Rule (1-sLA) - Stop whenever stopping now is better than stopping in next step, ie, $\forall t$
 $A_t = 1$ under the event
 $T_t \triangleq \{ s(x_t) \leq c(x_t) + \mathbb{E}[s(x_{t+1}) \mid X_t = x, A_t = 1] \}$

Thm (Chow & Robbins '61) - Suppose T_t is monotone,
i.e. $T_0 \subseteq T_1 \subseteq T_2 \dots$ a.s. Then 1-sla is optimal

Pf - Induction on $t \in \{T-1, T-2, \dots, 0\}$. For $t=T-1$, its clear we should stop if $s(x_t) \leq c(x_t) \Leftrightarrow x_t \in T$. Now if its true for $t \geq k$, then at time $t=k-1$, if $x_t \notin T$, then its better to continue till $t=k$; if $x_t \in T$, then $x_{t+1} \in T$ and by induction hypothesis, its better to stop.

Eg (Selling an asset) - X_1, X_2, \dots, X_T are iid offers for an asset with $\text{Var}(X_t) < \infty$; Per day holding cost c_i also, can recall the best past offer anytime.

$$- s(x_t) = -\max_{t' \leq t} \{x_{t'}\} \quad c(x_t) + E[s(x_{t+1})] = c - E[\max\{x_{t+1}, M_t\} | X_1, \dots, X_t]$$

$$\Rightarrow s(x_t) - c(x_t) - E[s(x_{t+1})] = E[(x_{t+1} - M_t)^+ | M_t] - c$$

decreasing as fn of M_t

$\Rightarrow 1-s_{la}$ is optimal!

Note - Not true without recall (ie, for the 'prophet inequality')

Corollary - If the set of 'allowed stopping points' is determined by X_t , then also we can define monotone stopping rules

Eg - Last-Success and Bruss' Odds Policy ($q_n = 1-p_n$)

- Sequence of T independent events $X_1, X_2, \dots, X_T, X_k \sim \text{Ber}(p_k)$
- Objective : Maximize probability of stopping at the 'last success' i.e., $R(X_t=1, A_t=0) = \text{IP}\{X_{t+k}=0 \forall k \geq 1\}$

- Let $\theta_i = p_i/q_i$ ('odds' for i^{th} event)
and $O_i = \sum_{t>i} \theta_t$

Thm (Bruss 2000) - The optimal stopping policy is to stop after the last time t with $I_t=1$ and $O_t \geq 1$

$$\begin{aligned} \text{Pf} - \text{Let } T_t &= \{X_t=1, \text{ and stopping at next success is worse}\} \\ &= \{X_t=1 \text{ and } q_{t+1}q_{t+2} \dots q_T > p_t + q_{t+2} - q_t + \dots + q_{t-1} - p_t\} \\ &= \{X_t=1, 1 > \theta_{t+1} + \theta_{t+2} + \dots + \theta_T\} \end{aligned}$$

Clearly T_t is monotone \Rightarrow opt policy is stop after last t st $\sum_{t \geq t} \theta_t \geq 1$

Corollary - For secretary problem, stop after last t st. $H_t \geq 1$

Optimal Stopping in Infinite Horizon Settings

- Above we assumed T was finite; can 1-sla work if $T \rightarrow \infty$

Eg - $X_t = X_{t-1} + Z_t$, with $Z_t = \pm 1 \text{ w.p } 1/2$ (symmetric RW)

$$c(X_t) = 0 \quad \forall t, \quad s(X_t) = e^{-X_t}$$

$$-s(X_t) = e^{-X_t}, \quad c(X_t) + \mathbb{E}[s(X_{t+1})] = e^{-X_t}(e + e^{-1})/2 = 1.5e^{-X_t} > s(X_t)$$

moreover $T_t = \mathbb{Z}$ which is monotone

- However X_t is **recurrent** \Rightarrow eventually hits every $x \in \mathbb{Z}$
 \Rightarrow opt policy is to wait forever (with cost 0)...

- However, we can avoid these pathologies with additional assumption
- Lemma - If $S = \sup_x s(x) < \infty$, $C = \inf_x c(x) > 0 \Rightarrow F_T(x) \rightarrow F(x)$ as $T \rightarrow \infty$
- Moreover a 1-sla is optimal if and only if T_t is monotone
(see Weber section 6.3)

Eg - (**Sequential Prob Ratio Test**) - Suppose $X_1, X_2, \dots \sim F$, where our prior is that $F = F_0 \text{ w.p } p_0, F = F_1 \text{ w.p } p_1 = 1 - p_0$.

- To get each additional sample, we pay cost δ
- If we stop and declare H_0 is true, we pay cost $s_0 L_t$ if wrong

- $L_t = p_1 \frac{\mathbb{P}[F=F_1]}{\mathbb{P}[F=F_0]}, \quad l_n = \frac{\mathbb{P}[F=F_1 | X_1, \dots, X_n]}{\mathbb{P}[F=F_0 | X_1, \dots, X_n]} = \frac{f_1(x_n)}{f_0(x_n)} l_{n-1}$ (likelihood ratios)
- We now want to consider a 1-sla policy

- Stopping at time $t \equiv$

$$S(L_t) = \min \left(\frac{c_0 L_t}{1 + L_t}, \frac{c_1}{1 + L_t} \right) \Rightarrow (1 + L_t) S(L_t) = \min(c_0 L_t, c_1)$$

- Continuing at t and stopping in next step

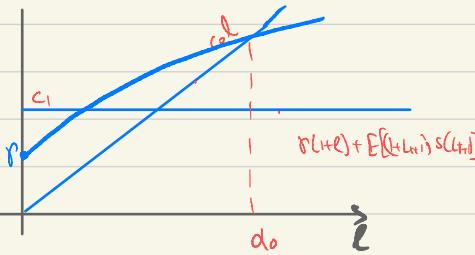
$$c(L_t) + \mathbb{E}[S(L_{t+1})|L_t] = r + \frac{1}{1+L_t} \mathbb{E}_{x \sim f_1} [S(L_{t+1})] + \frac{1}{1+L_t} \mathbb{E}_{x \sim f_0} [S(L_{t+1})]$$

$$\Rightarrow (1+L_t)(c(L_t) + \mathbb{E}[S(L_{t+1})|L_t]) = r(1+L_t) + \int L_t f_1(y) S(L_t f_1(y)/f_0(y)) dy \\ + \int f_0(y) S(L_t f_1(y)/f_0(y)) f_0(y) dy \\ = r(1+L_t) + \int (1+L_t f_1(y)/f_0(y)) S(L_{t+1}) f_0(y) dy \\ = r(1+L_t) + \mathbb{E}_{x \sim f_0} [(1+L_{t+1}) S(L_{t+1})]$$

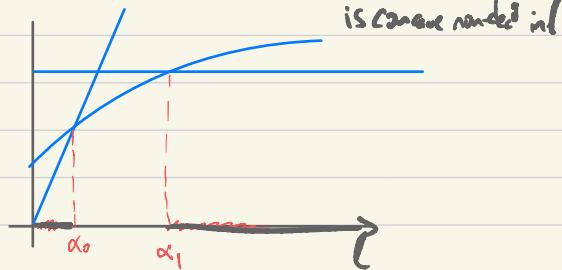
Thus the 1-sla rule corresponds to the following:

- 1) If $c_0 L_t \leq r(1+L_t) + \mathbb{E}_{f_0} [(1+L_{t+1}) S(L_{t+1})] \Rightarrow$ Stop and declare $F=F_0$
- 2) If $c_1 \leq r(1+L_t) + \mathbb{E}_{f_0} [(1+L_{t+1}) S(L_{t+1})] \Rightarrow$ Stop and declare $F=F_1$
- 3) Else continue

- Note $(1+l)S(l) = \min(c_0 l, c_1) \equiv$ concave, non-dec² in l
 $\Rightarrow \mathbb{E}_{F_0} [(1+L_{t+1}) S(L_{t+1})|L_t=l] = \mathbb{E}_{x \sim f_0} [\min(c_0 f_1(y)/f_0(y), c_1)]$



2 intervals



3 intervals

It can be seen from the plots that the 1-sla either has a single threshold l_0 , or 2 thresholds $l_0 < l_1$. This is clearly monotone \Rightarrow the 1-sla is optimal.