ORIE 4520 - Stochastics at Scale

Instructor: Siddhartha Banerjee

Semester: Fall 2015

August 26, 2015

Essential Course Information

Instructor

Prof. Siddhartha Banerjee Office: 229 Rhodes Hall

E-mail: sbanerjee@cornell.edu

Website: people.orie.cornell.edu/sbanerjee/

Office hours: MW 2:30pm-3:30pm (immediately after class)

Teaching Assistant
 Anna Srapionyan

E-mail: as3348@cornell.edu

Essential Course Information (contd.)

Lectures and Recitations

Course Number: ORIE 4520 Class time: MWF 1:25-2:15pm Class location: Phillips 403

Recitation time/location: To be decided

(Recitation time on schedule: Tuesay, 2:55-4:10pm)

Course Communication:

Website: http://people.orie.cornell.edu/sbanerjee/ orie4520f15.html

I will use BlackBoard for all announcements (search for ORIE

4520)

 Basic probability (at the level of ORIE 3500): Random variables, conditional probability and expectation, common probability distributions and their properties (binomial, geometric, exponential, Poisson); simulations.

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- Algorithms and graph theory: asymptotic (Big O) notation, basic algorithms (sorting, searching), LP
- Mathematical maturity

What is 'scaling'??

A warmup example: Balls in Bins



Courtesy: www.fixturescloseup.com

Suppose you throw m balls into n bins uniformly at random (u.a.r.)

• Assume n is very very large. Think of number of balls m(n) as a function of n.

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 Answer.
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 (The 'Birthday Paradox')
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 Answer. ⊕ (n log n)

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Answer: $\Theta(\sqrt{n})$

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Takeaway: In large stochastic systems, simple questions have 'interesting' answers

Balls in Bins: One final twist

We throw m balls into n bins uniformly at random (u.a.r.)

• If we choose m=n, how many balls are there in the most-loaded bin?

Answer: Maximum load is $\Theta\left(\frac{\log n}{\log\log n}\right)$

The power of two choices

Suppose instead we do the following:

For each ball, choose 2 bins u.a.r., and drop ball in less-loaded bin.

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Takeaway: In large stochastic systems, small changes can lead to dramatic outcomes

A (tentative) list of topics

- First unit: Intro to randomized algorithms and scaling
 - Tools: Tail inequalities (the Chernoff bound), randomized rounding, random walks
 - Examples: Sorting, median finding, graph algorithms (min and max cut, centrality), routing problems

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- Third unit: Threshold phenomena in large stochastic systems
 - Tools: Birth-death chains, branching processes, fluid approximations
 - Examples: Power of two choices, random graphs, epidemics

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Back to Administrivia

Course Material

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- References for the third unit:
 - Networks, Crowds and Markets (Sections V, VI) by D. Easley and J. Kleinberg
 - Epidemics and Rumours in Complex Networks by M. Draief and L. Massoulié.

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Coursework and Grading

Homework:

8 homeworks – weekly until the prelim, and biweekly after that. Homeworks due on Friday 12pm.

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One prelim: 90 min in-class exam, held during recitation hours (tentatively, during the week of 19th to 23rd October)

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• Grading:

Homeworks (45%) – max $\{6 \times 5\% + 2 \times 10\%, 45\}$ Prelim (25%), Project (25%+5%).