0

(demand, distrib) (Distri) (Distri) (Drift)

- Assumptions

i) Di are independent

ii) Di realized before Di+1 (sequential)

- Controls - Protection levels (nested)

· Xj = pnotection bool for face classes j+1, j+2,..., in

- Let $S_j = Capacity$ available when fare class j arrives $(S_1 = C)$

then

$$\frac{R_{j}(s_{j},x_{j},D_{i})}{V_{j}(s_{j})} = \max_{x_{j} \in \{0,1,...,s_{j}\}} \frac{F_{j}\left[P_{j} \min \left\{D_{j}, s_{j}-x_{j}\right\} + V_{j+1}(s_{j+1})\right]}{X_{j} \in \{0,1,...,s_{j}\}}$$

where Siti = max {Si-Di, Xi}

- Question - Why protection levels?

- Alternative - Let's solve an 'easier' problem

· Assumption 1 - We know Dj before allocating class jeats

· Assumption 2 - Can choose exact number of seats for j

 $- \bigvee_{j} \left(S_{j} \right) = \left[\mathbb{E}_{j} \left[\max \left(P_{j} \cdot \min \left(x, D_{j} \right) + \bigvee_{j+1} \left(S_{j+1} \right) \right] \right]$ $\max \left(S_{j} - x, S_{j} - D_{j} \right)$ $= S_{j} - \min \left(x, D_{j} \right)$

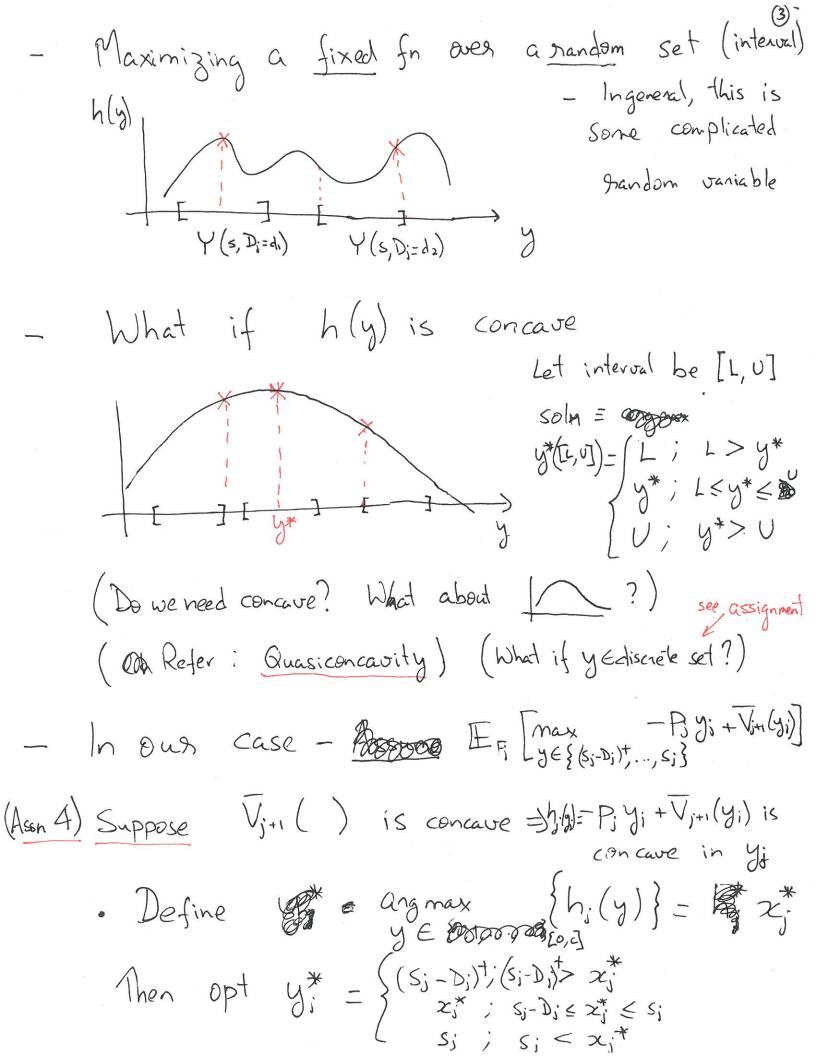
- Easier state variable - Let $y_i = S_i - min(x, D_i)$ final desired state

 $V_{j}\left(S_{j}\right) = \mathbb{E}_{F_{j}}\left[\max_{y_{j} \in \{(S_{j}-D_{j})^{+}, \dots, S_{j}\}} P_{j}\left(S_{j}-Y_{j}\right) + V_{j+1}\left(Y_{j}\right)\right]$ $\frac{1}{y_{j}}\left(S_{j}-Y_{j}\right) + V_{j+1}\left(Y_{j}\right)$

$$= \sum_{j=1}^{n} \left(S_{j} \right) = P_{j} S_{j} + \left[\sum_{j=1}^{n} \left(\sum_{j=1}^{n} \left(-P_{j} y_{j} + \sum_{j=1}^{n} \left(y_{j} \right) \right) \right] + \left(y_{j} \right) \right]$$

What can we say about the solution?

· Assumption 3 - Consider continuous D;



given Al- Oracle access' to Di A2 - Control = Exact # of seats of fave P; A3 - Continuous D; A4 - Concave V; (.) then optimal capacity allocation => Protection levels ?} - I where $x = \frac{1}{x \in [0, c]} = \frac{1}{x \in [0, c]}$ * Let's gremove the assumptions. A4, then A1 and A2 * Concavity of Vi () - Induction on appropriate n, n-1, ..., 1

 $- \overline{V_n(x_n|D_n)} \stackrel{\Delta}{=} \max_{(x_n-D_n)^+ \in \mathcal{Y} \leq x_n} \left[-P_n(x_n-D_n)^+ + P_n x_n \right]$ $(x_n-D_n)^+ \leq \mathcal{Y} \leq x_n \quad (\text{or } P_n \cdot \min\{x_n,D_n\})$

- concave in x Y Dn Vn(xn/Dn) - Pn Dn $- \overline{V}_n(x_n) = \mathbb{E}_{F_n} \left[\overline{V}_n(x_n | D_n) \right]$ = Concave in Xn

(linear combination of concavefus)

- Assume
$$V_{j+1}(x_{j+1})$$
 is concave in x_{j+1} .

Let $F_j(y) = \max_{x \in \mathbb{R}} \left[F_j(y) \right] + F_j(x_j)$
 $V_j(x_j|D_j) = \max_{x \in \mathbb{R}} \left[F_j(y) \right] + F_j(x_j)$

Sufficient to show this is concave

- max ($g(x)$)

Let $x = ay \max_{x \in \mathbb{R}} g(x)$
 $(x-a) \leq y \leq x$
 $x = ay \max_{x \in \mathbb{R}} g(x)$

Then $y = ay = ay = ay$
 $x = ay$

Notes - this works for discrete domands (A3) Protection levels aneoptimal even if we control exact what about A1? Claim - It does not matter...