

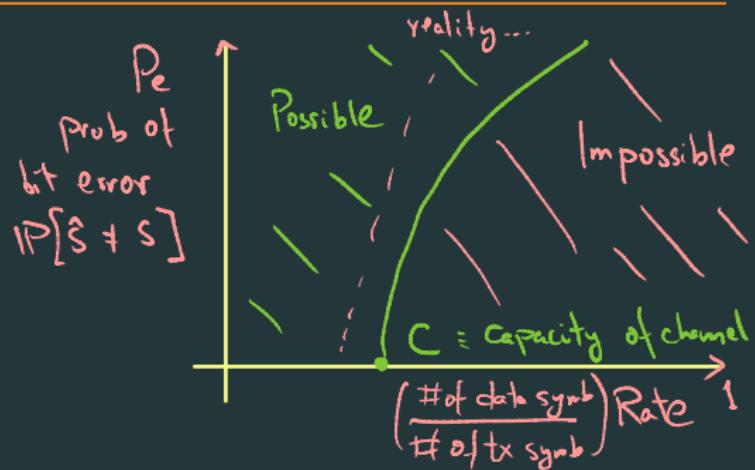


ORIE 4742 - Info Theory and Bayesian ML

Chapter 5: Channel Coding

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entropy: basic properties

rv X taking values $\mathcal{X} = \{a_1, a_2, \dots, a_k\}$, with pmf $\mathbb{P}[X = a_i] = p_i$

Shannon's entropy function

- outcome $X = a_i$ has *information content*: $h(a_i) = \log_2 \left(\frac{1}{p_i} \right)$
- random variable X has *entropy*: $H(X) = \mathbb{E}[h(X)] = \sum_{i=1}^k p_i \log_2 \left(\frac{1}{p_i} \right)$

- only depends on distribution of X (i.e., $H(X) = H(p_1, p_2, \dots, p_k)$)
- $H(X) \geq 0$ for all X
- if $X \sim \text{uniform}$ on \mathcal{X} , then $H(X) = \log_2 |\mathcal{X}|$; else, $H(X) \leq \log_2 |\mathcal{X}|$
- if $X \perp\!\!\!\perp Y$, then $H(X, Y) = H(X) + H(Y)$
where joint entropy $H(X, Y) \triangleq \sum_{(x,y)} \underbrace{p(x,y)}_{p(x)p(y)} \underbrace{\log_2 1/p(x,y)}_{p(x)p(y)}$

conditional entropy

conditional entropy

$$\text{for any rvs } X, Y: H(X|Y) = \sum_{y \in \mathcal{Y}} p(y) H(X|Y=y)$$
$$= \sum_{y \in \mathcal{Y}} p(y) \sum_{x \in \mathcal{X}} p(x|y) \log_2(1/p(x|y))$$

Consider $X \sim \left(\underbrace{a_1, a_2}_{Y=1}, \underbrace{a_3, a_4}_{Y=0} \right)$ w.p. (P_1, P_2, P_3, P_4) , $Y = \mathbb{1}_{\{X \in \{a_1, a_2\}\}}$

$\downarrow Y=1 \Rightarrow X \in \{a_1, a_2\}$ w.p. $\left(\frac{P_1}{P_1+P_2}, \frac{P_2}{P_1+P_2} \right)$

$\downarrow Y=0 \Rightarrow X \in \{a_3, a_4\}$ w.p. $\left(\frac{P_3}{P_3+P_4}, \frac{P_4}{P_3+P_4} \right)$

$$H(X) = H(Y) + \left[\begin{array}{l} (P_1+P_2) H_2 \left(\frac{P_1}{P_1+P_2} \right) \\ + (P_3+P_4) H_2 \left(\frac{P_3}{P_3+P_4} \right) \end{array} \right] H(X|Y)$$

$H(X|Y)$

the chain rule

the chain rule (information content)

for any rvs X, Y and realizations x, y :

$$\underbrace{h(x, y)}_{\log_2 \left(\frac{1}{P(x,y)} \right)} = h(x) + h(y|x) = \underbrace{h(y)}_{\log_2 \left(\frac{1}{P(y)} \right)} + \underbrace{h(x|y)}_{\log_2 \left(\frac{1}{P(x|y)} \right)}$$

$$\begin{aligned} \log_2 \left(P(x,y) \right) &= \log_2 \left(P(x) P(y|x) \right) \\ &= \log_2 \left(P(x) \right) + \log_2 \left(P(y|x) \right) \end{aligned}$$

the chain rule

the chain rule (entropy)

for any rvs X, Y :

$$H(X, Y) = \underbrace{H(X) + H(Y|X)}_{\sum_{(x,y)} P(x,y) \log_2 \frac{1}{P(x,y)}} = H(Y) + \underbrace{H(X|Y)}_{\sum_{y} P(y) \log_2 \frac{1}{P(y)}} = \sum_{(x,y)} P(x,y) \log_2 \frac{1}{P(x,y)}$$
$$\frac{1}{P(y)P(x|y)}$$

$$\underline{\text{Note}} : H(X) - H(X|Y) = H(Y) - H(Y|X)$$

mutual information

(information gain ...)

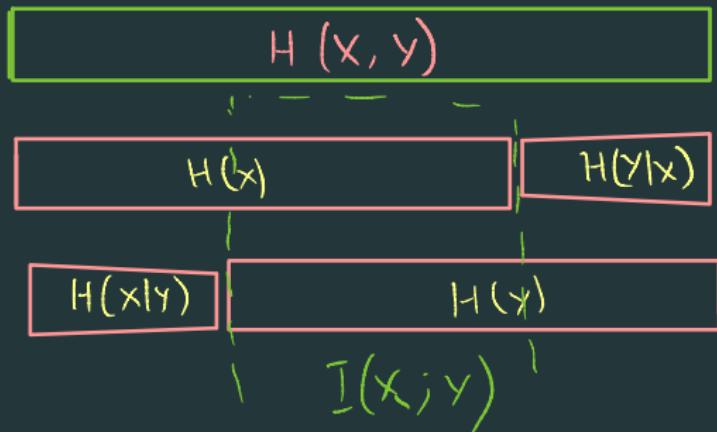
mutual information

for any rvs X, Y :

$$I(X; Y) = H(X) - H(X|Y) = H(Y) - H(Y|X)$$

moreover, given any other conditioning rv Z

$$I(X; Y|Z) = H(X|Z) - H(X|Y, Z) = H(Y|Z) - H(Y|X, Z)$$



blank

Alt defn - $I(x; y) = D_{KL}\left(P(x,y) \parallel P(x)P(y)\right)$

≥ 0 (by Gibbs' Ineq)

want to encode use to encode

LHS - $RHS \equiv \sum_{(x,y)} p(x,y) \log_2 \left(\frac{p(x,y)}{p(x)p(y)} \right) = \sum_{x,y} p(y)p(x|y) \log_2 \left(\frac{p(x,y)p(y)}{p(x)p(y)} \right)$

$= \underbrace{\sum p(y) \log_2 \left(\frac{1}{p(y)} \right)} - \underbrace{\sum p(x,y) \log_2 \left(\frac{1}{p(y|x)} \right)}$

$I(x;y) = H(y) - H(y|x)$

example

| $P(x, y)$ | | x | | | | $P(y)$ |
|-----------|---|----------------|----------------|----------------|----------------|---------------|
| | | 1 | 2 | 3 | 4 | |
| y | 1 | $\frac{1}{8}$ | $\frac{1}{16}$ | $\frac{1}{32}$ | $\frac{1}{32}$ | $\frac{1}{4}$ |
| | 2 | $\frac{1}{16}$ | $\frac{1}{8}$ | $\frac{1}{32}$ | $\frac{1}{32}$ | $\frac{1}{4}$ |
| | 3 | $\frac{1}{16}$ | $\frac{1}{16}$ | $\frac{1}{16}$ | $\frac{1}{16}$ | $\frac{1}{4}$ |
| | 4 | $\frac{1}{4}$ | 0 | 0 | 0 | $\frac{1}{4}$ |
| $P(x)$ | | $\frac{1}{2}$ | $\frac{1}{4}$ | $\frac{1}{8}$ | $\frac{1}{8}$ | |

| $P(x)$ | | 1 | 2 | 3 | 4 | $P(y x)$ | |
|--------|---|---------------|---------------|---------------|---------------|----------|---------------|
| y | 1 | $\frac{1}{2}$ | $\frac{1}{4}$ | $\frac{1}{8}$ | $\frac{1}{8}$ | 1 | $\frac{1}{4}$ |
| | 2 | $\frac{1}{4}$ | $\frac{1}{2}$ | $\frac{1}{8}$ | $\frac{1}{8}$ | 2 | $\frac{1}{8}$ |
| | 3 | $\frac{1}{8}$ | $\frac{1}{4}$ | $\frac{1}{4}$ | $\frac{1}{4}$ | 3 | $\frac{1}{8}$ |
| | 4 | — | 0 | 0 | 0 | 4 | $\frac{1}{2}$ |

$$h(x) = 1 \cdot 2 + 3 \cdot 3, H(x) = \frac{7}{4}$$

$$h(y) = 2 \cdot 2 + 2 \cdot 2, H(y) = 2$$

$$H(x,y) = 2 \cdot \left(\frac{1}{4}\right) + 3\left(\frac{2}{8}\right) + 4\left(\frac{6}{16}\right) + 5\left(\frac{4}{32}\right) = \frac{27}{8}$$

$$H(x) + H(y) = \frac{30}{8} > H(x,y)$$

$$\left. \begin{aligned} H(x|y) &= \frac{1}{4} \cdot \left(\frac{7}{4} + \frac{7}{4} + 2 + 0 \right) = \frac{11}{8} \\ H(y|x) &= \frac{1}{2} \cdot \left(\frac{7}{4} \right) + \frac{1}{4} \left(\frac{3}{2} \right) + \frac{1}{8} \left(\frac{3}{2} \right) + \frac{1}{8} \left(\frac{3}{2} \right) = \frac{13}{8} \\ H(x) + H(y|x) &= \frac{7}{4} + \frac{13}{8} = \frac{27}{8} \\ H(y) + H(x|y) &= 2 + \frac{11}{8} = \frac{27}{8} \\ H(x) - H(x|y) &= \frac{7}{4} - \frac{11}{8} = \frac{3}{8} \\ H(y) - H(y|x) &= 2 - \frac{13}{8} = \frac{3}{8} \end{aligned} \right] = H(x,y) = I(x;y)$$

mutual information and KL divergence

mutual information

for any rvs X, Y : $I(X; Y) = H(X) - H(X|Y) = H(Y) - H(Y|X)$

$$= D_{KL} \left(P(x,y) \parallel P(x)P(y) \right)$$

= loss in encoding length when
we approximate $P(x,y)$ by
 $P(x)P(y)$

blank

Some properties

- Suppose $X \perp\!\!\!\perp Y$

$$H(X|Y) = H(X), H(Y|X) = H(Y) \leftarrow (\text{check...})$$

$$\Rightarrow I(X;Y) = 0$$

- Suppose $Y = f(X)$ (deterministic, f)

$$H(Y|X) = 0 \quad (H(X|Y) = ? \quad 0 \text{ if } f \text{ invertible})$$

$$\Rightarrow I(X;Y) = H(Y) \quad (= H(X) \text{ if } f \text{ invertible})$$

Data Processing Inequality

Markov Chain



X and Y are conditionally indep given Z

Data Proc^g ineq

$$I(X; Z) \leq I(X; Y)$$

= iff Z is a sufficient stats

- $X=Y, Z=0 \Rightarrow I(X; Z)=0, I(X; Y)=H(X)$
- $Y \perp\!\!\!\perp X, Z=Y \Rightarrow I(X; Y)=0, I(X; Z)=0$

visualizing mutual information

