

Another broad class of policies for MDPs are **index policies**

- Suppose state and action spaces decompose as $S = S_1 \times S_2 \times \dots \times S_n$ and $A = \{a_1, a_2, \dots, a_k\}$. Then an index policy comprises a set of fns $\phi_1, \phi_2, \dots, \phi_m$ with $\phi_i : S_i \rightarrow \mathbb{R}$ (**indices**)
s.t. $A(s) = \operatorname{argmax}_{k \in [m]} \{\phi_k(s_i)\}$ for $s = (s_1, s_2, \dots, s_m)$
- In other words, for each 'part' of the state, we compute a fn, and then 'act on the part' with the highest fn value

Eg - (single machine scheduling with discounting)

- There are m jobs, with each job i having known processing time t_i and reward r_i upon completion
Want to schedule them on a single machine to maximize discounted sum of rewards (discount factor β)
- If $S_m \subseteq [m]$ = set of remaining jobs, then
$$V(S_m) = \max_{j \in S_m} [r_j \beta^{t_j} + \beta^{t_j} V(S_m \setminus j)]$$
- If $m = 2$, $V(\{1, 2\}) = \max[r_1 \beta^{t_1} + r_2 \beta^{t_1+t_2}, r_2 \beta^{t_2} + r_1 \beta^{t_1+t_2}]$
 \Rightarrow we first serve $\operatorname{argmax}_{i \in \{1, 2\}} \{r_i \beta^{t_i} / 1 - \beta^{t_i}\}$
- This problem has an easy soln via an **interchange argument**
Suppose order is $i_1 \rightarrow i_2 \rightarrow \dots \rightarrow j \rightarrow k \rightarrow \dots \rightarrow i_m$
 \Rightarrow Reward is of form $R_1 + \beta^{T+t_j} r_j + \beta^{T+t_j+t_m} r_m + R_2$
by previous argument, comparing $i_1 \rightarrow i_2 \rightarrow \dots \rightarrow j \rightarrow k \rightarrow \dots \rightarrow i_m$ and $i_1 \rightarrow i_2 \rightarrow \dots \rightarrow k \rightarrow j \rightarrow \dots \rightarrow i_m$ is same as comparing $r_{i_1} \beta^{t_{i_1}} / 1 - \beta^{t_{i_1}}$ and $r_{i_m} \beta^{t_{i_m}} / 1 - \beta^{t_{i_m}}$
 \Rightarrow OPT policy = serve jobs in decreasing order of $r_j \beta^{t_j} / 1 - \beta^{t_j}$

Bayesian Multi-Armed Bandits

- Now we consider a vast generalization of the above
 - There are m 'arms', where each arm i is a **Markov Chain** $X_i[t]$ on state S_i .
 - In each round, we can **play a single arm**, i.e., $A = [m]$
 - If $A[t] = i$, then $R(t) = R_i(X_i[t])$ and $X_i[t] \rightarrow X_i[t+1]$
 - (**Non-restlessness**) The state of arm i changes only when it is 'played' (i.e., $X_i[t+1] = X_i[t]$ if $i \notin A[t] \neq i$)
 - (**Discounted infinite-horizon objective**) $\max \sum_{t=0}^{\infty} \beta^t (\sum_{i=1}^m \mathbb{1}_{\{A[t]=i\}} R_i(X_i[t]))$
- Eg (MAB with Beta-Bernoulli priors).-
 - m actions, where action i gives reward $Beta(P_i)$
 - P_i unknown, but assume $P_i \sim Beta(N_i, S_i)$ (prior)
 - If we play action k , then posterior = $\begin{cases} Beta(N_k + S_k) & \text{if } R_i = 1 \\ Beta(N_k, S_k) & \text{if } R_i = 0 \end{cases}$
 - Aim : $\max \sum_{t=0}^{\infty} \beta^t (\sum_{i=1}^m \mathbb{1}_{\{A[t]=i\}} R_i[t])$

The 1.5 arm problem - Suppose we just have 2 arms

1) MC $X[t]$ with reward $R_i(X_i[t])$

2) Constant reward arm with reward γ

No opt policy is a stopping problem: play arm 1 till some time T (**stopping time**), then play 2 forever (since no new info)

$$\Rightarrow R = \sup_{T \geq 0} \mathbb{E} \left[\sum_{t=0}^{T-1} R_i(X_i[t]) \beta^t + \beta^T \gamma / (1-\beta) \right]$$

The **Gittins Index** of arm 1 in state $X_i[0]=2$ is the smallest constant reward γ' s.t. you are indifferent between playing arm 1 in state $X_i[0]$ and arm 2

• Formally for any $x \in S$:

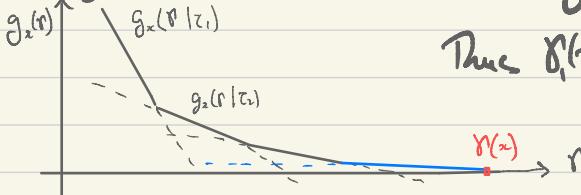
$$\gamma_i(x) = \sup \left\{ r : \frac{r}{1-\beta} \leq \sup_{T \geq 0} \mathbb{E} \left[\sum_{t=0}^{T-1} \beta^t R_i(x_{i[t]}) + \frac{\beta^T r}{1-\beta} \mid X_{i[0]} = x \right] \right\}$$

This can also be viewed as the **minimum per-horizon charge** for pulling arm i s.t. you are indifferent between pulling once or not when $X_{i[0]} = x$

$$\gamma_i(x) = \sup \left\{ r : 0 \leq \sup_{T \geq 0} \mathbb{E} \left[\sum_{t=0}^{T-1} \beta^t [R_i(x_{i[t]}) - r] \mid X_{i[0]} = x \right] \right\}$$

Note that $g_x(r)$ is decreasing and convex in r

- To see this, note that for a fixed sample path x_1, x_2, \dots and fixed T , $\sum_{t=0}^{T-1} \beta^t [\mathbb{E}[R_i(x_t)] - r]$ is linear decreasing
- Taking expectation preserves linearity, and sup over T makes it convex



Thus $\gamma_i(x)$ has a unique soln

Also for the optimal T , we have $\mathbb{E} \left[\sum_{t=0}^{T-1} \beta^t R_i(x_{i[t]}) - \gamma_i(x) \sum_{t=0}^{T-1} \beta^t \mid X_{i[0]} = x \right] = 0$
 Thus

$$\gamma_i(x) = \sup_{T > 0, \text{ stopping time}} \frac{\mathbb{E} \left[\sum_{t=0}^{T-1} \beta^t R_i(x_{i[t]}) \mid X_{i[0]} = x \right]}{\mathbb{E} \left[\sum_{t=0}^{T-1} \beta^t \mid X_{i[0]} = x \right]}$$

Gittins Index for arm 1
 in state x

Eg - Suppose $X[0] = X, X[1] = X, X[2] = \dots$ $= \begin{cases} M & \text{w.p } p \\ 0 & \text{ow} \end{cases}$ 'Collapsing' arm
 $R_i(x) = x$ (play one to learn)

$$\text{Then } \gamma(\text{'unknown'}) = \sup \left[\gamma \mid \gamma / 1 - \beta \leq pM / 1 - \beta + (1-p)\beta\gamma / 1 - \beta \right]$$

$$= \frac{pM}{1 - \beta(1-p)}$$

Eg - Single job with reward π_i , processing time t_i :

$\xrightarrow{\substack{\text{amount of} \\ \text{job processed}}}$ $R_i(0) = \sup_{T > 0} \frac{\pi_i \beta^{t_i} \prod_{j \neq i} \beta^{t_j}}{\sum_{t=1}^{T-1} \beta^t} = \left(\frac{\pi_i \beta^{t_i}}{1 - \beta^{t_i}} \right) (1 - \beta)$

the index we get via interchange

Thm (Gittins '79) - For finite arms $[m]$, and bounded rewards $R_i(x) \in [-c, c] \forall i \in [m], x \in S$:
A policy is optimal if and only if it always selects arm i at time t with highest Gittins index $\gamma_i(X:[t])$.

There are actually more general conditions for when an index policy is optimal. When is it not, though?

- Independence of irrelevant alternatives (IIA): A policy Π satisfies IIA if for any set of arms $[n]$ and $i \in [n]$, if $\Pi([n]) = i$, then $\Pi([n]) = i$ for any $[n] \supseteq [m]$.

An index policy is optimal if and only if opt policy is IIA.