NATIONAL CRIMINAL DATABASE SYSTEM

A Project Report

Submitted by

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Roll Nos.: B006, B015, B018, B019

Place: Mumbai

Date: April 1, 2020

CERTIFICATE				
This is to certify that the project entitled "Note the bonafide work carried out by Sarvesh Siddhant Burse of B.Tech. Computer Engine during the IV semester of the academic year requirements for the Course Database Managements."	Agrawal, Naman Bhansali, Raj Bora, eering, MPSTME (NMIMS), Mumbai, 2019-2020 in partial fulfillment of the			
	Prof. Ishani Saha Internal Mentor			
Examiner 1	Examiner 2			

ABSTRACT

The present world is technology driven as it is employed by almost all domains in performing their operations efficiently. In the case of police department, the need for good record-keeping and information sharing practice is very significance in modern times. Not only do good records provide crucial internal information, police now need to communicate within department or department-to-department across the country, in order to protect the citizens. The proposed system applies to all Police stations across the country and looks into the subject of Centralized Crime Records Management. It is well understood that crime prevention, detection and conviction of criminals depend on a highly responsive backbone of information management. The efficiency of the police and the effectiveness with which it tackles crime depend on what quality of information it can derive from its existing records and how fast it can have access to it. It will be implemented by digitalising the present records, which accesses information across all records in the state thus helping speedy and successful completion to cases. This project has distributed architecture, with centralized storage of the database.

The application for the storage of the data uses the constructs of SQL server and all the user interfaces have been designed using Python GUI. SQL Queries are used to provide data relating criminals' bio, crimes, and other data like case details, policemen investigating, etc. Different types of users will have different level of view and access to the database. System Administrator will use it to add police users to the system. Various types of police will have different level of access to the System. Civilian will have access to the Data Analytic part and complaint registration option.

The project will have deep impact on the police working and aid in their investigation and record keeping.

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INTRODUCTION
For the academic project of the course Database Management System, we chose domain of Police Department and the problem pertaining to records management. Our project aims to improve the record access and updation system by the police and at the same time provide analyzed data from the crime and criminal records. The proposed system will help ensure quick action by the force unlike the current situation where they have to wait for confirmation of data from other police departments. Now they will have to simply login into the system and check info.
PROBLEM STATEMENT
There are inconsistencies in records management by the Police Department, which along with mis - communication causes lapses in their working. The current offline record management causes delay in the updating of criminal records and paperwork. This project tries to reduce a obstacle faced by the department in their performance. We will find a solution to this using Database Management System.

FUNCTIONAL REQUIREMENTS OF THE SYSTEM

Following functional requirements are expected from the system:

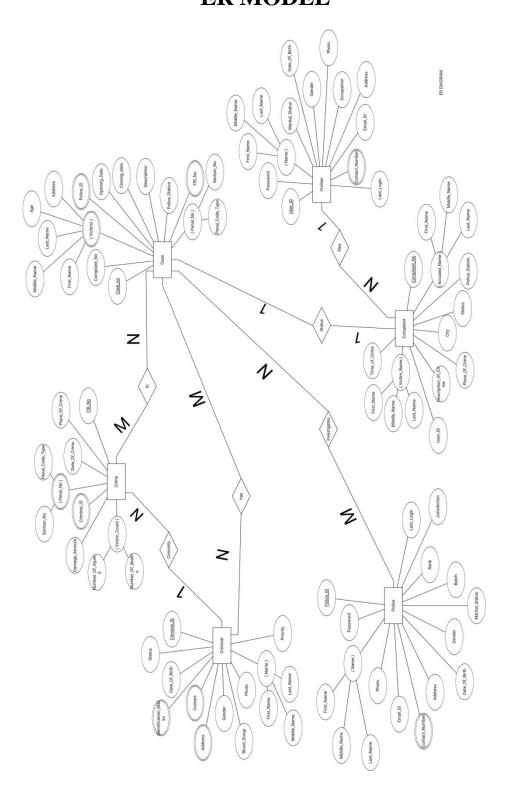
- 1. To create records of criminals, crime and cases by the authorized personnel.
- 2. To Alter the records as and when need arises by the police personnel.
- 3. To provide limited access to civilians to share data with them.

USERS OF THE SYSTEM

- 1. System Administrator: Person managing the database and Police users in the Police Department.
- 2. High Ranked Officer (A.C.P.): These users have access and privilege to alter the records.
- 3. Low Ranked Officer (Constable): These users only have access to the records.
- 4. Civilian: These users have option to file complaints as well as view analyzed crime data.

3

SYSTEM DESIGN & CONSTRAINTS ER MODEL



RELATIONAL MODEL

CIVILIAN(<u>USERID</u>, PASSWORD, FNAME, MNAME, LNAME, DOB, GENDER, MARITALSTATUS, EMAILID, OCCUPATION, ADDRESS, LASTLOGIN, PHOTO, CONTACT NUMBER)

POLICE (<u>POLICEID</u>, PASSWORD, FNAME, MNAME, LNAME, LASTLOGIN, EMAILID, JURISDICTION, ADDRESS, GENDER, DOB, BATCH, RANK, MARITALSTATUS, PHOTO, CONTACT)

COMPLAINT (COMPLAINT NO, PLACEOFCRIME, TIMEOFCRIME, CRIMEDESCRIPTION, CITY, POLICESTATION, VFNAME, VMNAME, VLNAME, AFNAME, AMNAME, ALNAME, STATUS, USERID)

CRIMINAL(<u>CRIMINALID</u>, FNAME, MNAME, LNAME, DOB, BLOODGROUP, STATUS, PRIORITY, GENDER, PHOTO, CONTACT ADDRESS, IDENTIFICATIONMARKS)

CASE (<u>CASENO</u>, PENALCODETYPE, SECTIONNUMBER, POLICESTATION, DESCRIPTION, OPENDATE, CLOSEDATE, COMPLAINT_NO, POLICEID, VFNAME, VMNAME, VLNAME, VAGE, VADDRESS, FIRNO)

CRIME (<u>FIRNO</u>, DAMAGEAMOUNT, INJURED, DEATHS, DATEOFCRIME, PLACEOFCRIME, CRIMINALID, PENALCODETYPE, SECTIONNUMBER)

CONSTRAINTS

Constraints are the rules enforced on the data columns of a table. These are used to limit the type of data that can go into a table. This ensures the accuracy and reliability of the data in the database.

Constraints used on the Database:

1. Domain Constraint - Domain constraint defines the domain or set of values for an attribute. It specifies that the value taken by the attribute must be the atomic value from its domain.

For Example, Contact Number will take integer/number value and not any text.

2. Tuple Uniqueness Constraint - Tuple Uniqueness constraint specifies that all the tuples must be necessarily unique in any relation.

For Example, there cant be two records for same criminal_id in criminal table.

3. Key Constraint - Key constraint specifies that in any relation-

All the values of primary key must be unique. The value of primary key must not be NULL.

For Example, PoliceId, primary key in Police Entity, will be unique to each user.

4. Entity Integrity Constraint - Entity integrity constraint specifies that no attribute of primary key must contain a null value in any relation. This is because the presence of null value in the primary key violates the uniqueness property.

For Example, UserID cant be NULL in Civilian table or it will be tedious to fetch the data efficiently.

5. Referential Integrity Constraint - This constraint is enforced when a foreign key references the primary key of a relation. It specifies that all the

values taken by the foreign key must either be available in the relation of the primary key or be null.

For Example, we can file a case only if complaint is registered as complaint_no is foreign key in Case table.

SQL Constraints:

- 1. NOT NULL Constraint Ensures that a column cannot have NULL value. For example, POLICEID can't be NULL in Police entity.
- 2. UNIQUE Constraint Ensures that all values in a column are different. For example, email id should be unique in Civilian entities.
- 3. PRIMARY Key Uniquely identifies each row/record in a database table.
- 4. FOREIGN Key Uniquely identifies a row/record in any of the given database table.
- 5. CHECK Constraint The CHECK constraint ensures that all the values in a column satisfies certain conditions.
- 6. DEFAULT Constraint Provides a default value for a column when none is specified.

Mapping Cardinalities used in the ER Diagram:

- > Many to One: When entities in one entity set can take part only once in the relationship set and entities in other entity set can take part more than once in the relationship set, cardinality is many to one. For Example: Criminal Crime Relationship: One criminal does many crimes.
- > Many to Many: When entities in all entity sets can take part more than once in the relationship cardinality is many to many. For Example: Police Case Relationship: Many Police investigates many cases.
- > One to One: When each entity in each entity set can take part only once in the relationship, the cardinality is one to one.

# Participation Constraints are :	8
> Total Participation – Each entity is involved in the relationship. Total participation is represented by double lines.	
> Partial participation — Not all entities are involved in the relationship. Partial participation is represented by single lines.	

NORMALISATION TECHNIQUES

1NF

1. Case (CASENO, PENALCODETYPE, SECTION, POLICESTATION, DESCRIPTION,

OPENDATE, CLOSEDATE, COMPLAINT_NO, POLICEID, VFNAME, V MNAME, VLNAME, VAGE, VADDRESS, FIRNO)

Above relation is not in 1NF because it has-

FIRNO as multivalued attribute.

VICTIMS as multivalued and composite attribute.

POLICEID as multivalued attribute

So we decompose the relation as below-

- A. CASE1(CASENO, PENALCODETYPE, SECTION, POLICESTATION, DESCRIPTION, OPENDATE, CLOSEDATE, COMPLAINT_NO)
- B. CASE2(CASENO, POLICEID)
- C. CASE3(CASENO,VFNAME ,VMNAME, VLNAME,VAGE ,VADDRESS)
- D. CASE4(CASENO,FIRNO)

as the decomposition is lossless and dependency is preserved i.e.

attr(CASE1) attr(CASE2)= CASENO

attr(CASE1) attr(CASE3)= CASENO

attr(CASE1) attr(CASE4)= CASENO

CASENO is key of Case1()

==>attr(CASE1)Uattr(CASE2)Uattr(CASE3)Uattr(CASE4)=attr(CASE)

2.Crime(FIRNO, DAMAGEAMOUNT, INJURED, DEATHS, DATEOFCRIM E, PLACEOFCRIME, CRIMINALID, PENALCODETYPE, SECTIONNUMB ER)

Above relation is not in 1NF because it has-

CRIMINALID as multivalued attribute

PENALNO as multivalued and composite attribute

So we decompose the relation as below-

A. CRIME(FIRNO,DAMAGEAMOUNT,INJURED,DEATHS,DATEOFCR IME,PLACEOFCRIME)

B. CRIME2(FIRNO, CRIMINALID)

C. CRIME3(FIRNO, PENALCODETYPE, SECIONNUMBER)

as the decomposition is lossless and dependency is preserved i.e.

attr(CRIME) attr(CRIME2)= FIRNO

attr(CRIME) attr(CRIME3)=FIRNO

FIRNO is key of CRIME()

==>attr(CRIME)Uattr(CRIME2)Uattr(CRIME3)=attr(Crime)

3.Criminal(CRIMINALID ,FNAME ,MNAME ,LNAME,DOB ,BLOODGR OUP ,STATUS,PRIORITY ,GENDER ,PHOTO ,CONTACT,ADDRESS ,I DENTIFICATIONMARKS)

Above relation is not in 1NF because it has-

Address as multivalued attribute

Contact as multivalued attribute

Identification_Marks as multivalued attribute

Name as Composite attribute

So we decompose the relation as below-

A.CRIMINAL(CRIMINALID, FNAME, MNAME, LNAME, DOB, BLOOD GROUP, STATUS, PRIORITY, GENDER, PHOTO)

- B. CRIMINAL2(CRIMINALID, CONTACT)
- C. CRIMINAL3(CRIMINALID, ADDRESS)
- D. CRIMINAL4(CRIMINALID, IDENTIFICATION MARKS)

as the decomposition is lossless and dependency is preserved i.e.

attr(CRIMINAL) attr(CRIMINAL2)= CRIMINALID

attr(CRIMINAL) attr(CRIMINAL3)= CRIMINALID

attr(CRIMINAL) attr(CRIMINAL4)= CRIMINALID

CRIMINALID is key of CRIMINAL ()

==>attr(CRIMINAL)Uattr(CRIMINAL2)Uattr(CRIMINAL)Uattr(CRIMINAL4)=attr(Criminal)

4.Police(POLICEID ,PASSWORD ,FNAME ,MNAME ,LNAME,LASTLO GIN ,EMAILID,JURISDICTION, CONTACT ADDRESS ,GENDER, DOB ,BATCH ,RANK, MARITALSTATUS,PHOTO, CONTACT)

Above relation is not in 1NF because it has-

Contact Number as multivalued attribute

So we decompose the relation as below-

A.POLICE(POLICEID, PASSWORD, FNAME, MNAME, LNAME, LASTL OGIN, EMAILID, JURISDICTION, CONTACT ADDRESS, GENDER, DOB, BATCH, RANK, MARITAL STATUS, PHOTO)

B. POLICE1(POLICEID, CONTACT)

as the decomposition is lossless and dependency is preserved i.e.

attr(POLICE) attr(POLICE1)= POLICEID

POLICEID is key of POLICE ()

==> attr(POLICE)Uattr(POLICE1)=attr(Police)

5. Civilian(USERID,PASSWORD,FNAME,MNAME,LNAME,DOB, GENDER, MARITALSTATUS, EMAILID,OCCUPATION, ADDRESS, LASTLOGIN, PHOTO, CONTACT)

Above relation is not in 1NF because it has-

Contact_Number as multivalued attribute

Name as composite attribute

So we decompose the relation as below-

A. CIVILIAN1(USERID,PASSWORD,FNAME,MNAME,LNAME,DOB, GENDER, MARITALSTATUS, EMAILID,OCCUPATION, ADDRESS, LASTLOGIN, PHOTO)

B. CIVILIAN2(USERID, CONTACT)

as the decomposition is lossless and dependency is preserved i.e.

attr(CIVILIAN) attr(CIVILIAN1)= USERID

USERID is key of CIVILIAN ()

And

attr(CIVILIAN)Uattr(CIVILIAN 1)=attr(Civilian)

2NF

All relations are in 2NF

3NF

All relations are in 3NF

BCNF

All relations are in BCNF

IMPLEMENTATION

HARDWARE & SOFTWARE REQUIREMENTS

- 1. Minimum 1500x800 Display size
- 2. Python 3.x
- 3. SQLite3

TOOLS & LIBRARIES

Front end developed with Python GUI Module and back end developed with SQLite Database.

Following Python modules were implemented in the project:

- > tkinter for GUI development
- > sqlite3 for database connectivity
- > matplotlib for data visualization
- > pillow for image handling in python
- > datetime for capturing login session details
- > uuid unique id generation for complaints
- > os for handling application pages

Following Database & related software were used for Database:

- > SQLite3 Database
- > DB Browser
- > SQLite Studio

APPLICATION - AT A GLANCE

Database:

Table CASE1



Table CASE2 and CASE4

	CASENO	POLICEID
	Filter	Filter
1	101	6
2	102	6
3	103	20
4	104	20
5	105	15
6	106	References POLIC Hold Ctrl+Shift ar
7	107	15
8	108	15
9	109	6
10	110	19
11	111	15
12	112	16

	CASENO	FIRNO
	Filter	Filter
1	101	503
2	102	506
3	103	509
4	104	512
5	105	502
6	106	505
7	107	508
8	108	511
9	109	501
10	110	504
11	111	510
12	112	507

Table CASE3

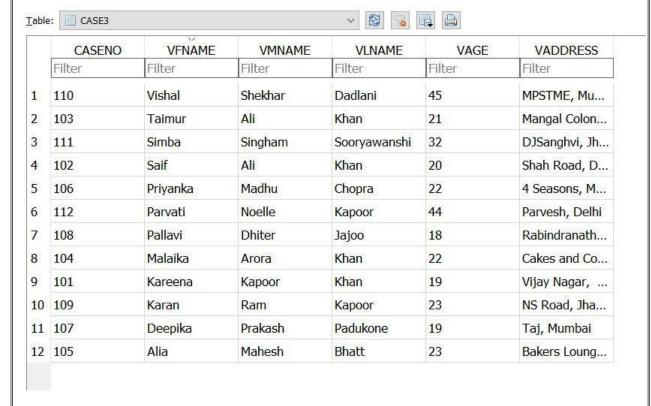


Table CRIME



Table CRIME2 & CRIME3

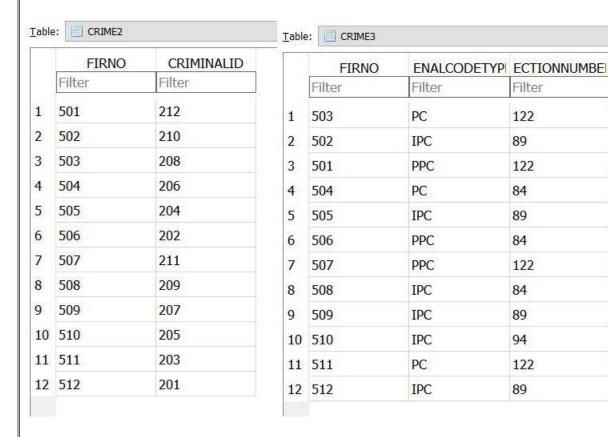


Table CRIMINAL

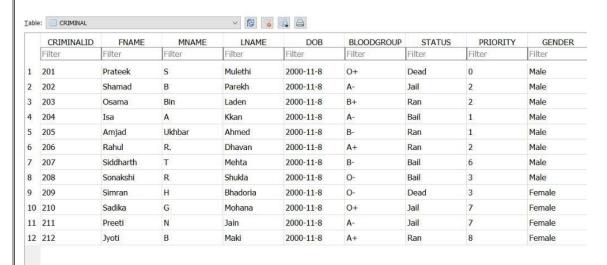


Table CRIMINAL1



Table CRIMINAL2 and CRIMINAL3



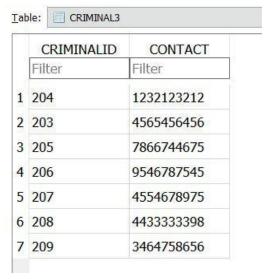


Table POLICE



Table POLICE1

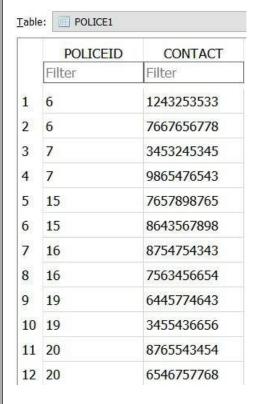


Table COMPLAINT



Table CIVILIAN1

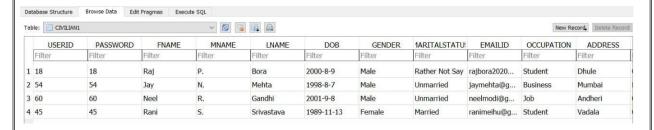
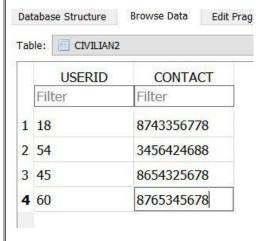


Table CIVILIAN2



GUI

Login Screen:

WELCOME TO POLICE PORTAL

Enter Below details to Login

Select User Type

USER ID

PASSWORD

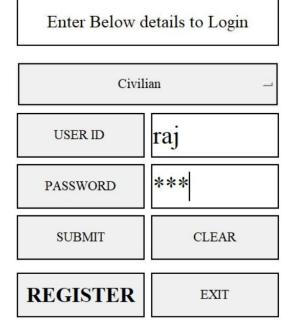
SUBMIT

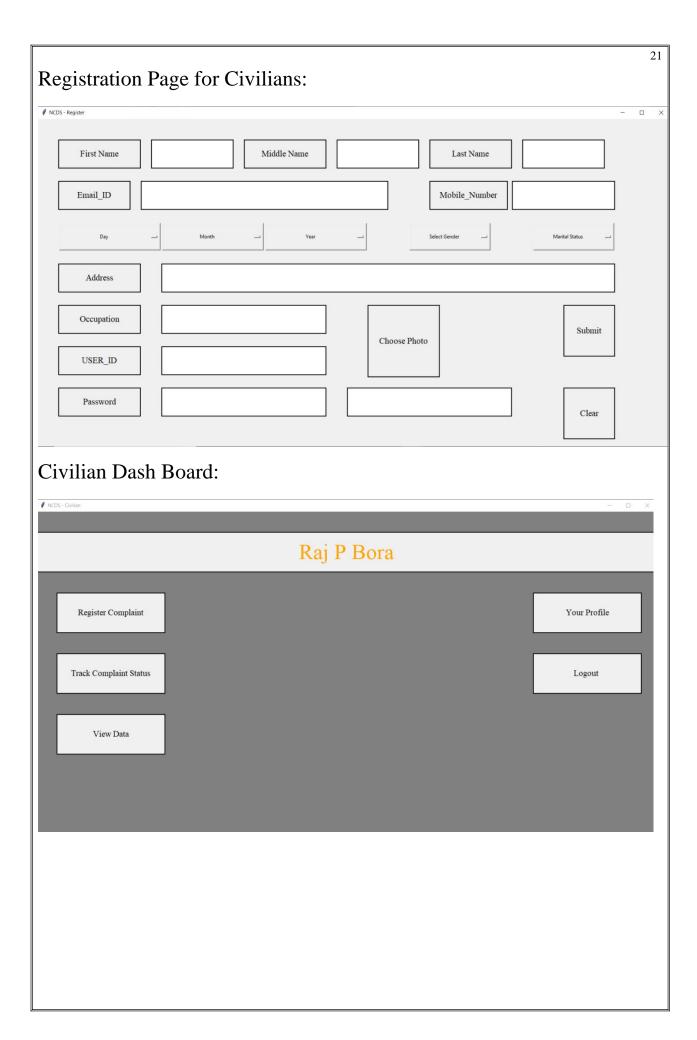
CLEAR

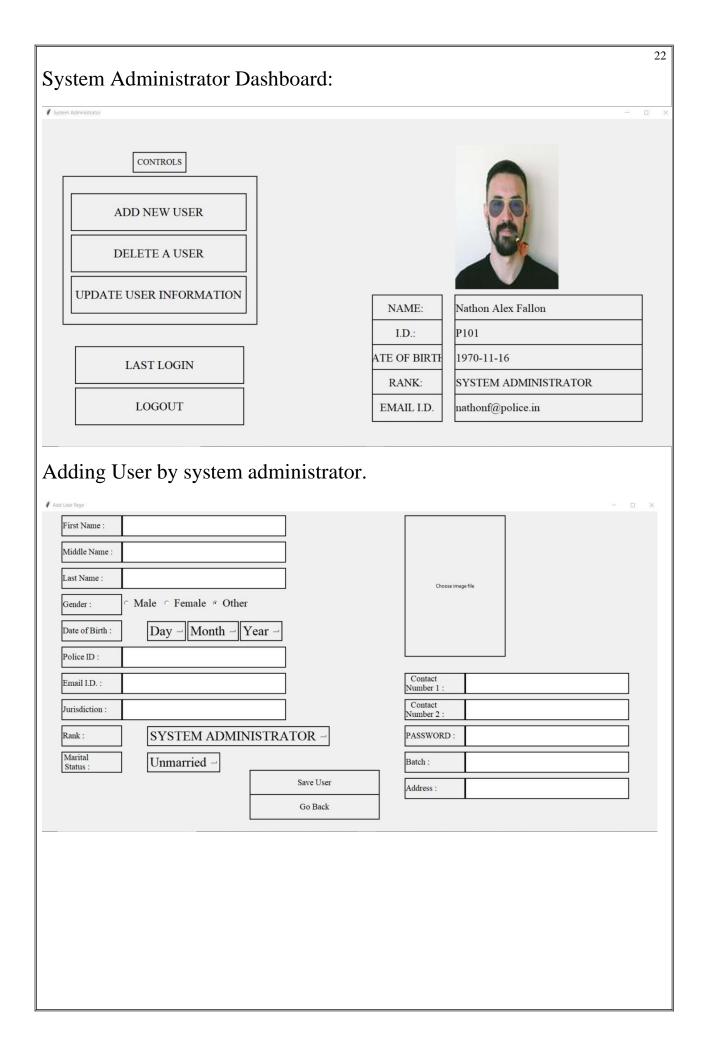
REGISTER

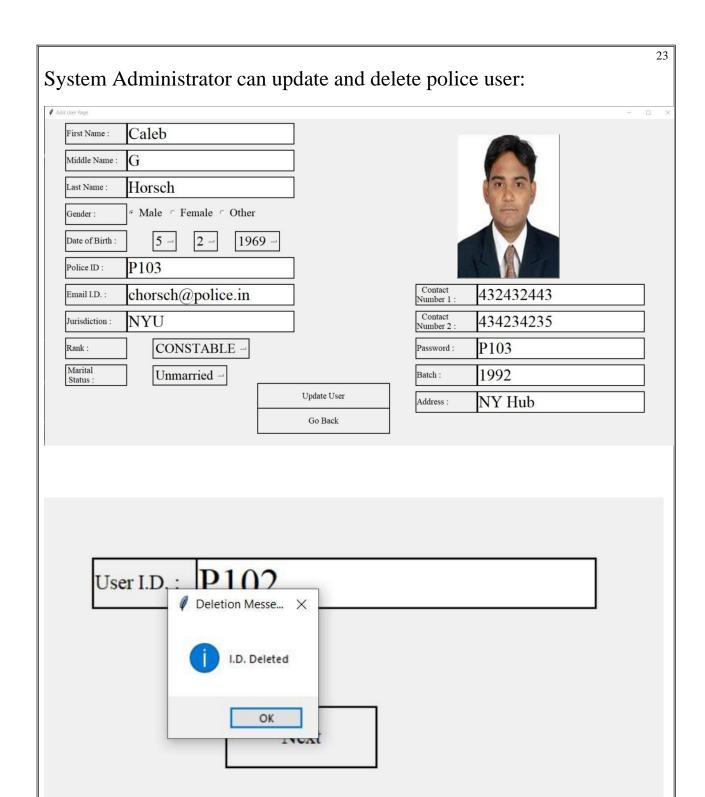
EXIT

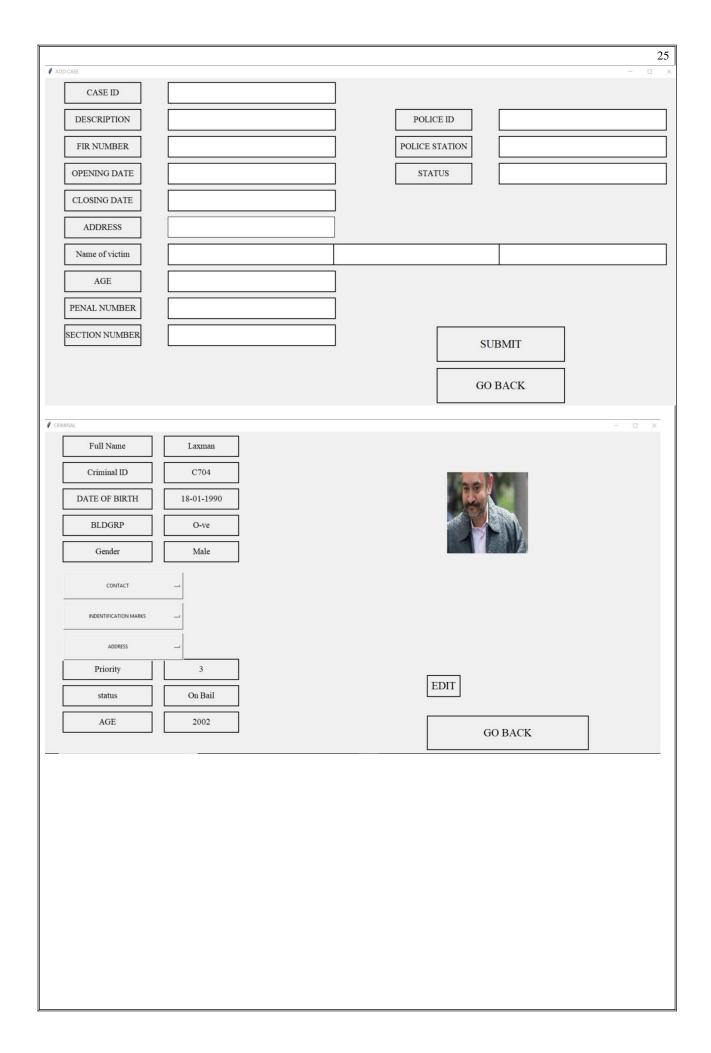
The fields to enter Login ID & password:

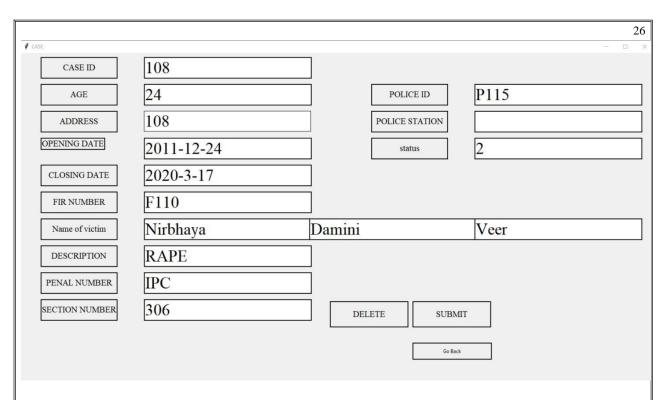




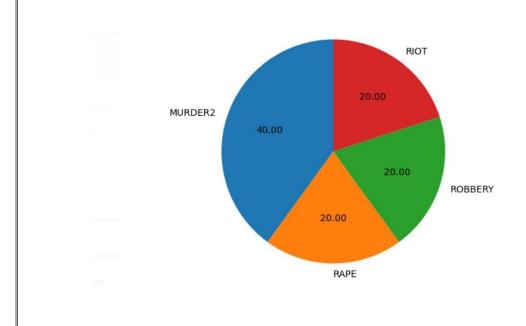




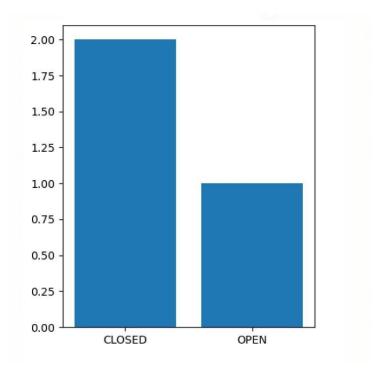


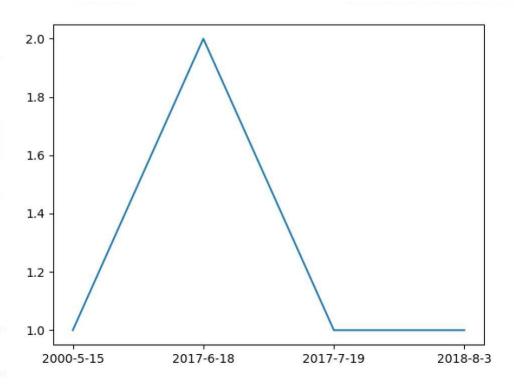


ACP and constable both can visualize data from the database.



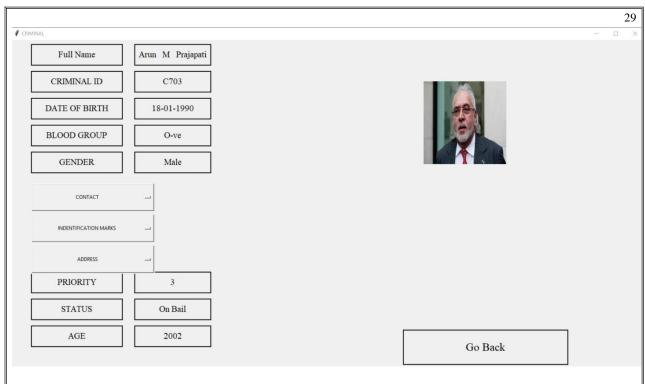




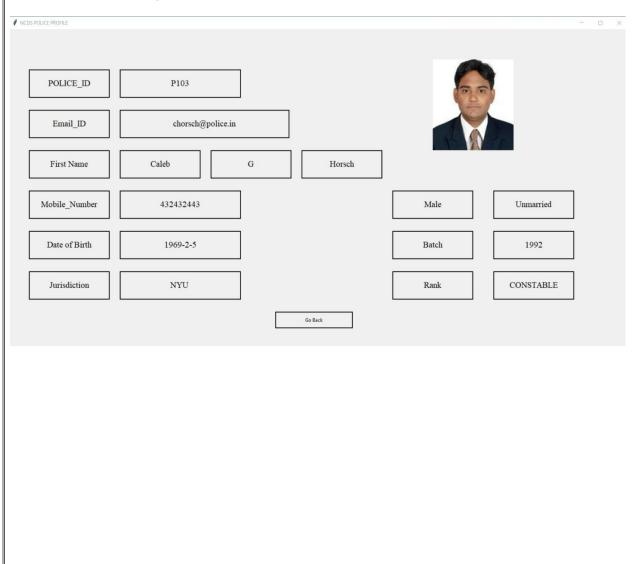


Constable end is similar to ACP, but with restricted access to database.

Meaning they can only view and not alter records.				
	- o x			
CA	ALEB HORSCH			
ACCESS RECORDS DASHBOARD YOUR PROFILE BACK	LAST LOGIN DETAILS LOGOUT SEARCH BY CRAMMAL D CAST SEARCH BY SUBMIT			



Police Profile Page:



DATABASE STRUCTURE

Tables (16)

Name	Туре	Schema	
CASE1		CREATE TABLE CASE1 (CASENO number PRIMARY KEY, PENALCODETYPE text, SECTIONNUMBER number, POLICESTATION text, DESCRIPTION text NOT NULL, OPENDATE text NOT NULL, CLOSEDATE text, COMPLAINT_NO TEXT, FOREIGN KEY (COMPLAINT_NO) REFERENCES COMPLAINT(COMPLAINT_NO))	
CASENO	number	"CASENO" number	
PENALCODETYPE	text	"PENALCODETYPE" text	
SECTIONNUMBER	number	"SECTIONNUMBER" number	
POLICESTATION	text	"POLICESTATION" text	
DESCRIPTION	text	"DESCRIPTION" text NOT NULL	
OPENDATE	text	"OPENDATE" text NOT NULL	
CLOSEDATE	text	"CLOSEDATE" text	
COMPLAINT_NO	TEXT	"COMPLAINT_NO" TEXT	
CASE2		CREATE TABLE CASE2 (CASENO number, POLICEID text, FOREIGN KEY (POLICEID) REFERENCES POLICE(POLICEID), FOREIGN KEY(CASENO) REFERENCES CASE1(CASENO))	
CASENO	number	"CASENO" number	
POLICEID	text	"POLICEID" text	
CASE3		CREATE TABLE CASE3(CASENO number , VFNAME text, VMNAME text, VLNAME text, VAGE number, VADDRESS text, FOREIGN KEY (CASENO) REFERENCES CASE1(CASENO))	
CASENO	number	"CASENO" number	
VFNAME	text	VFNAME" text	
VMNAME	text	'VMNAME" text	
VLNAME	text	VLNAME" text	
VAGE	number	'VAGE" number	
VADDRESS	text	"VADDRESS" text	
CASE4		CREATE TABLE CASE4(CASENO number, FIRNO number, FOREIGN KEY(CASENO) REFERENCES CASE1(CASENO), FOREIGN KEY(FIRNO) REFERENCES CRIME(FIRNO))	
CASENO	number	"CASENO" number	
FIRNO	number	"FIRNO" number	
CIVILIAN1		CREATE TABLE CIVILIAN1 (USERID text PRIMARY KEY CHECK(USERID <> ''), PASSWORD text NOT NULL CHECK(PASSWORD <> ''), FNAME text, MNAME text, LNAME text, DOB text, GENDER text, MARITALSTATUS text, EMAILID text NOT NULL, OCCUPATION text, ADDRESS text, LASTLOGIN text, PHOTO blob)	
USERID	text	"USERID" text CHECK(USERID<>")	
PASSWORD	text	"PASSWORD" text NOT NULL CHECK(PASSWORD<>")	
FNAME	text	"FNAME" text	
MNAME	text	"MNAME" text	
LNAME	text	"LNAME" text	
DOB	text	"DOB" text	
GENDER	text	"GENDER" text	
MARITALSTATUS	text	"MARITALSTATUS" text	
EMAILID	text	"EMAILID" text NOT NULL	
OCCUPATION	text	"OCCUPATION" text	
ADDRESS	text	"ADDRESS" text	
LASTLOGIN	text	"LASTLOGIN" text	
PHOTO	blob	"PHOTO" blob	

text	CREATE TABLE CIVILIAN2 (USERID text , CONTACT number, FOREIGN KEY (USERID) REFERENCES CIVILIAN1(USERID))	
text		
	"USERID" text	
number	"CONTACT" number	
	CREATE TABLE COMPLAINT (COMPLAINT_NO text PRIMARY KEY, PLACEOFCRIME text NOT NULL CHECK(PLACEOFCRIME <> ''), TIMEOFCRIME text, CRIMEDESCRIPTION text, CITY text, POLICESTATION text, STATUS text DEFAULT Registered, VFNAME text, VMNAME text, VLNAME text, AFNAME text, AMNAME text, ALNAME text, USERID text, FOREIGN KEY(USERID) REFERENCES CIVILIAN1(USERID))	
Туре	Schema	
text	"COMPLAINT_NO" text	
text	"PLACEOFCRIME" text NOT NULL CHECK(PLACEOFCRIME<>>")	
text	"TIMEOFCRIME" text	
text	"CRIMEDESCRIPTION" text	
text	"CITY" text	
text	"POLICESTATION" text	
text	"STATUS" text DEFAULT Registered	
text	"VFNAME" text	
text	"VMNAME" text	
text	"VLNAME" text	
	"AFNAME" text	
	"AMNAME" text	
	"ALNAME" text	
	"USERID" text	
LEXI	CREATE TABLE CRIME (FIRNO number PRIMARY KEY, DAMAGEAMOUNT number, INJURED	
	number, DEATHS number, DATEOFCRIME text NOT NULL, PLACEOFCRIME text)	
number	"FIRNO" number	
number	"DAMAGEAMOUNT" number	
number	"INJURED" number	
number	"DEATHS" number	
text	"DATEOFCRIME" text NOT NULL	
text	'PLACEOFCRIME" text	
	CREATE TABLE CRIME2 (FIRNO number, CRIMINALID number, FOREIGN KEY(FIRNO) REFERENCES CRIME(FIRNO), FOREIGN KEY(CRIMINALID) REFERENCES CRIMINAL (CRIMINALID))	
number	"FIRNO" number	
number	"CRIMINALID" number	
	CREATE TABLE CRIME3 (FIRNO number, PENALCODETYPE text, SECTIONNUMBER number, FOREIGN KEY (FIRNO) REFERENCES CRIME(FIRNO))	
number	"FIRNO" number	
text	"PENALCODETYPE" text	
number	"SECTIONNUMBER" number	
	CREATE TABLE CRIMINAL(CRIMINALID number PRIMARY KEY, FNAME text, MNAME text, LNAME text, DOB text, BLOODGROUP text, STATUS text, PRIORITY number, GENDER text, PHOTO BLOB NOT NULL)	
number	"CRIMINALID" number	
text	"FNAME" text	
text	"MNAME" text	
text	"LNAME" text	
text	"DOB" text	
text	"BLOODGROUP" text	
text	"STATUS" text	
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CONCLUSION & FUTURE SCOPE

The above project aimed to remove discrepancies in the record management system of the police department via a centralized database system. Through this project, we successfully demonstrated the fulfillment of our aim and its positive impact it will have on the performance of the Police. Since, police department is spread from villages to metros, a large scale project will be suitable for actual working where multiple users could access the system. Currently, a similar project is implemented in India on pilot basis.

Future Scope of the project involves:

- # Adding AI based Facial Recognition tool to get instant results regarding identity of a criminal and his/her history.
- # More Transparency can be provided to the civilians regarding the working of the police department.
- # Audio samples can also be stored in the Database to match voice samples and same goes to the digital storing of DNA sequence.
- # More Data Analytic features can be added so as to help, reorganize data for direct use by Police and in case to civilians too.

These are just few of the numerous ways in which this project can be further built into something much bigger, magnificent and highly productive.