Step 1:

```
public int find(int p) {
    validate(p);
    int root = p;
    // FIXME
    // END
    if(pathCompression)
         doPathCompression(p);
    while (root != parent[root]) {
         root = parent[root];
    }
    return root;
}
private void mergeComponents(int i, int j) {
   // FIXME make shorter root point to taller one
   if(i != j){
       if(height[i] < height[j]){</pre>
           height[j] += height[i];
          parent[i] = j;
       } else {
           height[i] += height[j];
          parent[j] = i;
```

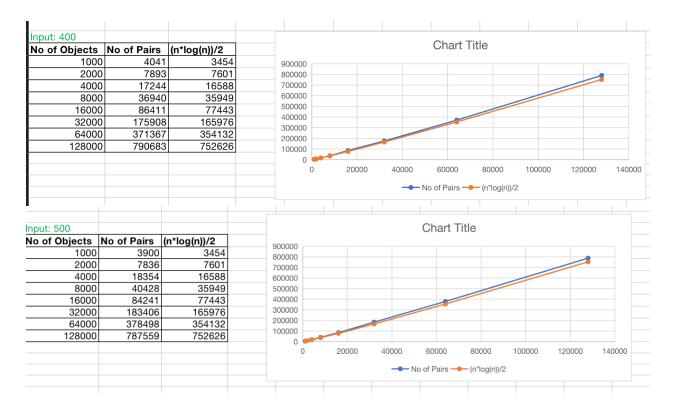
```
private void doPathCompression(int i) {
    // FIXME update parent to value of grandparent
    // END
    while(parent[i] != i){
        parent[i] = parent[parent[i]];
        i = parent[i];
    }
}
```

Step 2:

```
import java.util.Random;
import java.util.Scanner;
public class UFClient {
   public static int cnt(int n) {
       UF_HWQUPC uf = new UF_HWQUPC(n);
       int ct = 0;
            int y = random.nextInt(n);
    public static void main(String[] args) {
        System.out.println("Enter a number");
        Scanner sc = new Scanner(System.in);
        int n = sc.nextInt();
        System.out.println("relationship between m and n");
            int \underline{sum} = 0;
                sum += cnt(i);
            int mean = \underline{sum} / 10;
            System.out.println("total objects " + \underline{i} + ", total pairs " + mean);
```

Run the UFClient.java file. To test the count technique, we may use the terminal to provide a number. To evaluate the link between m and n, we can run more n values and double their values using the doubling method, each 10 times.

Using the approach described above, we mapped the two outputs. First, the command line input was set to 400, and then the command line input was set to 500. Both of their outputs were mapped in a graphical fashion. The following is the output and graphical depiction.



Relationship Conclusion:

We can conclude from above scenario that

Both lines are around the same length. The connection between m and n may be deduced as follows:

$$m = \frac{1}{2} (n \log(n))$$