



Department of Computer Engineering Academic
Term II: 23-24

Class: B.E (Computer), Sem – VI Subject
Name: Artificial Intelligence

Student Name: Siddhesh Pradhan

Roll No: 9632

Practical No:	9
Title:	Eight puzzle game solution by A* algorithm
Date of Performance:	
Date of Submission:	

Rubrics for Evaluation:

Sr. N o	Performance Indicator	Excellent	Good	Below Average	Marks
1	On time Completion & Submission (01)	01 (On Time)	NA	00 (Not on Time)	
2	Logic/Algorithm Complexity analysis (03)	03(Correct)	02(Partial)	01 (Tried)	
3	Coding Standards (03): Comments/indentation/Naming conventions Test Cases /Output	03(All used)	02 (Partial)	01 (rarely followed)	
4	Post Lab Assignment (03)	03(done well)	2 (Partially Correct)	1(submitted)	
Total					

Signature of the Teacher:



Experiment No: 5

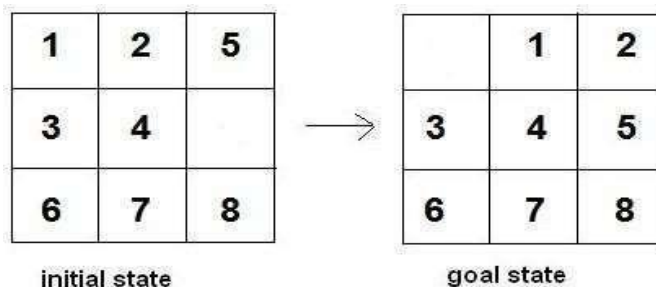
Title: Eight puzzle game solution by A* algorithm

Objective: To study A* algorithm and solutions to 8 puzzle problem using A*

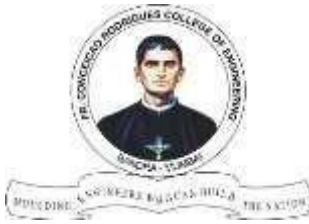
Theory:

The 8-puzzle problem is a puzzle invented and popularized by Noyes Palmer Chapman in the 1870s. It has a set of 3x3 boards having 9 block spaces out of which, 8 blocks are having tiles bearing number from 1 to 8. One space is left blank. The tile adjacent to blank space can move into it. It has to arrange the tiles in a sequence.

The start state is any situation of tiles, and goal state is tiles arranged in a specific sequence. Solution of this problem is reporting of “movement of tiles” in order to reach the goal state. The transition function or legal move is any one tile movement by one space in any direction (i.e., towards left or right or up or down) if that space is blank.



Here the data structure to represent the states can be a 9-element vector indicating the tiles in each board position. Hence, a starting state corresponding to the above configuration will be {1, blank, 4, 6, 5, 8, 2, 3, 7} (there can be various different start positions). The goal state is {1, 2, 3, 4, 5, 6, 7, 8, blank}. Here, the possible movement outcomes after applying a move can be many. They are represented as trees. This tree is called a state space tree. The depth of the tree will depend upon the number of steps in the solution. The part of state space tree of 8-puzzle is shown:

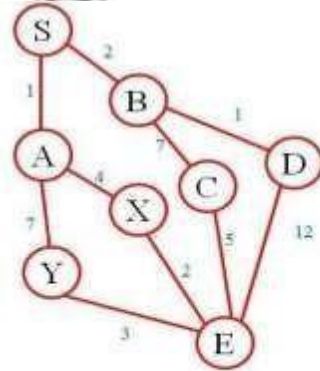


Algorithm:

1. **function** A-STAR-SEARCH (initialState, goalTest)
2. return **SUCCESS** or **FAILURE**: /* Cost $f(n) = g(n) + h(n)$ */
3. frontier = Heap. New(initialState)
4. explored = Set.new()
5. **while not** frontier.isEmpty ();
 - a. state = frontier.deleteMin()
 - b. explored.add(state)
6. **if** goalTest(state):
 - a. return **SUCCESS** (state)
7. **for** neighbor **in** state.neighbours():
 - a. **if** neighbor **not in** frontier U explored:
 - i. frontier.insert(neighbour)
 - b. **else if** neighbor **in** frontier:
 - i. frontier.decreaseKey(neighbour)
8. return **FAILURE**

A* Search algorithm is one of the best and popular technique used in path-finding and graph traversals. It gives the process of plotting an efficiently directed path between multiple points, called nodes. It enjoys widespread use due to its performance and accuracy.

Following is an example of A*



• Values for h:
A:5, B:6, C:4, D:15, X:5, Y:8

Expand S

{S,A} $f=1+5=6$

{S,B} $f=2+6=8$

Expand A

{S,B} $f=2+6=8$

{S,A,X} $f=(1+4)+5=10$

{S,A,Y} $f=(1+7)+8=16$

Expand B

{S,A,X} $f=(1+4)+5=10$

{S,B,C} $f=(2+7)+4=13$

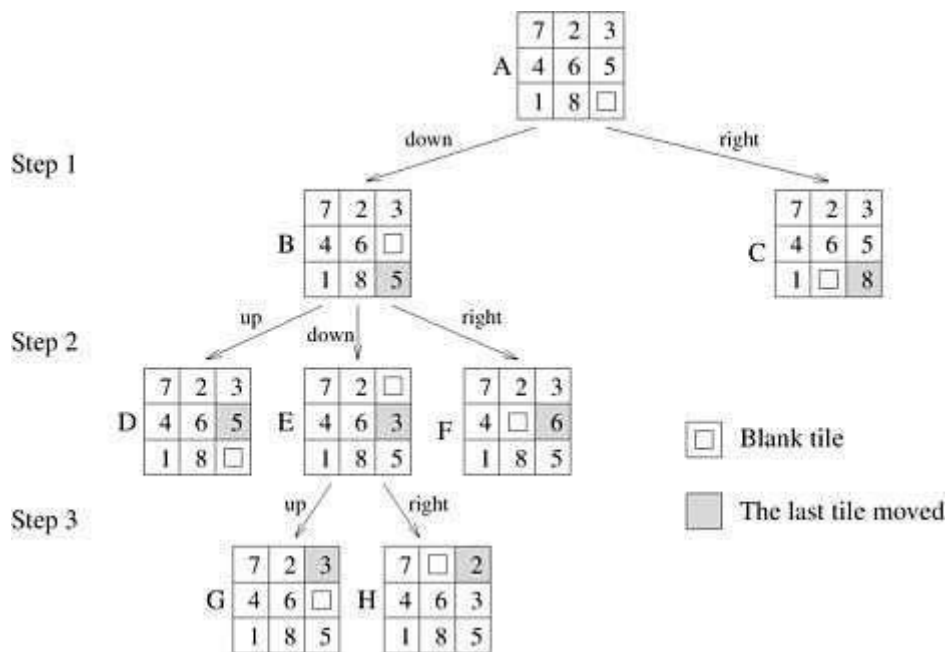
{S,A,Y} $f=(1+7)+8=16$

{S,B,D} $f=(2+1)+15=18$

Expand X

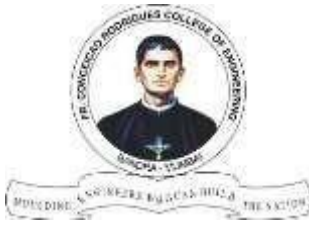
{S,A,X,E} is the best path... (costing 7)

Example 1



Using A* to solve the 8-puzzle problem in a heuristic manner:

- Heuristic 1 (H1): Count the out-of-place tiles, as compared to the goal.



- Heuristic 2 (H2): Sum the distances by which each tile is out of place.
- Heuristic 3 (H3): Multiply the number of required tile reversals by 2.

Analysis of the Evaluation Function:

In developing a good evaluation function for the states in a search space, you are interested in two things:

- $g(n)$: How far is state n from the start state?
- $h(n)$: How far is state n from a goal state?

The first value of, $g(n)$, is important because you often want to find the shortest path. This value can be exactly measured by incorporating a **deep count** into the search algorithm.

The second value, $h(n)$, is important for guiding the search towards the goal. It is an **estimated value** based on your heuristic rules.

Evaluation Function: This gives us the following: $f(n)$
 $= g(n) + h(n)$.

CODE:

```
import numpy as np
import matplotlib.pyplot as plt
import heapq

class State:
    def __init__(self, tiles, blank_position, parent=None, g=0, h=0):
        self.tiles = tiles
        self.blank_position = blank_position
        self.parent = parent
        self.g = g # cost from start to current state
        self.h = h # heuristic cost from current state to goal state

    def __lt__(self, other):
        return (self.g + self.h) < (other.g + other.h)

    def is_goal(self, goal_state):
        return np.array_equal(self.tiles, goal_state.tiles)

    def get_neighbors(self):
```



```
neighbors = []
    i, j = self.blank_position
    m, n = self.tiles.shape
    moves = [(0, 1), (1, 0), (0, -1), (-1, 0)]
    for di, dj in moves:
        ni, nj = i + di, j + dj
        if 0 <= ni < m and 0 <= nj < n:
            neighbor = np.copy(self.tiles)
            neighbor[i, j], neighbor[ni, nj] = neighbor[ni, nj], neighbor[i, j]
            neighbors.append(State(neighbor, (ni, nj), self, self.g + 1, 0)) #
Assuming uniform cost
    return neighbors

def manhattan_distance(state, goal_state):
    return np.sum(np.abs(state.tiles - goal_state.tiles))

def a_star_search(initial_state, goal_test):
    frontier = []
    heapq.heappush(frontier, initial_state)
    explored = set()

    while frontier:
        current_state = heapq.heappop(frontier)
        if goal_test(current_state):
            return current_state
        explored.add(tuple(map(tuple, current_state.tiles)))
        for neighbor in current_state.get_neighbors():
            if tuple(map(tuple, neighbor.tiles)) not in explored:
                heapq.heappush(frontier, neighbor)
                explored.add(tuple(map(tuple, neighbor.tiles)))
    return None

def plot_state(state):
    plt.imshow(state.tiles, cmap='viridis', origin='upper')
    plt.colorbar()
    plt.title('8 Puzzle State')
    plt.show()

def print_solution(solution):
```



```
path = []
while solution:

path.append(solution)
    solution = solution.parent
path.reverse()
for state in path:
    plot_state(state)

def main():
    initial_state = np.array([[1, 2, 3], [0, 4, 6], [7, 5, 8]])
    goal_state = np.array([[1, 2, 3], [4, 5, 6], [7, 8, 0]])

    initial_blank_position = np.where(initial_state == 0)
    initial_blank_position = (initial_blank_position[0][0],
initial_blank_position[1][0])
    initial_state = State(initial_state, initial_blank_position)

    goal_blank_position = np.where(goal_state == 0)
    goal_blank_position = (goal_blank_position[0][0], goal_blank_position[1][0])
    goal_state = State(goal_state, goal_blank_position)

    solution = a_star_search(initial_state, lambda state: state.is_goal(goal_state))

    if solution:
        print("Solution found:")
        print_solution(solution)
    else:
        print("No solution found.")

if __name__ == '__main__':
    main()
```

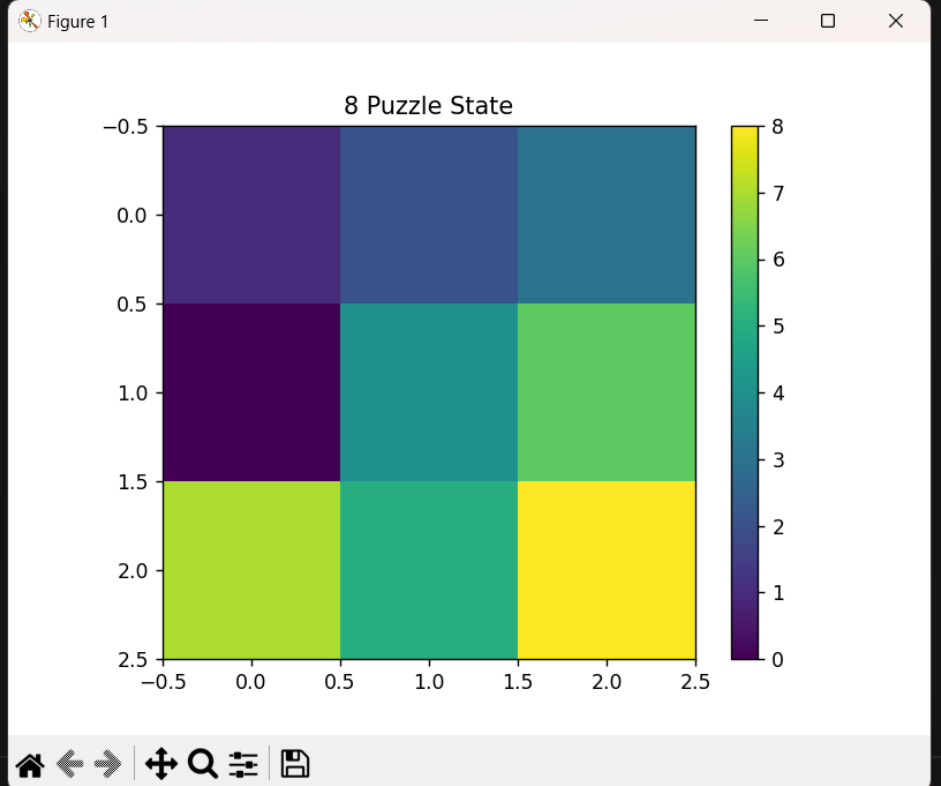


OUTPUT:

```
(base) PS C:\Users\Siddhesh\Desktop> python expt9.py
```

Solution found:

█



Post Lab Assignment:

1. Explain the Time Complexity of the A* Algorithm.
2. What are the limitations of A* Algorithm?
3. Discuss A*, BFS, DFS and Dijkstra's algorithm in detail with examples.



9632

AI Expt 9 Postlab

1. Explain the Time Complexity of the A^* algorithm.

Ans:- Time Complexity of A^* Algorithm

- A^* algorithm time complexity depends on the heuristics quality problem space size.
- Generally expressed as $O(b^d)$, where b is the branching factor and d is the depth of the solution.
- Efficiency improves with a good heuristic function.

2. What are the limitations of A^* algorithm?

Ans:- Limitations of A^* Algorithm

- Can be inefficient or incorrect with a poor heuristic
- Faces challenges with large or dynamic search spaces
- Memory-intensive for large space due to storing explored states.

3. Discuss A^* , BFS, DFS and Dijkstra's algorithm in detail with examples

Ans: A^* : Combines best features of greedy search and Dijkstra's algorithm. Depends on heuristic quality.

BFS: Guarantees shortest path in unweighted graphs. Memory intensive but suitable for small spaces.

DFS: Does not guarantee shortest path. Memory-efficient but may get stuck in infinite loops.

Dijkstra's: Finds shortest path in weighted graphs suitable for non-negative edge weights.